VE475 Homework 2

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Ex.1 - Simple Questions

1.

$$101 = 5 \times 17 + 16$$

$$17 = 1 \times 16 + 1$$

$$16 = 16 \times 1$$

$$1 = 17 - 1 \times 16$$

$$= 6 \times 17 - 101$$

So the inverse of 17 modulo 101 is 6.

2. To solve $12x \equiv 28 \mod 236$, is equivalent to solve $12x + 236y = 28 \Leftrightarrow 3x + 59y = 7$.

$$59 = 19 \times 3 + 2$$

$$3 = 1 \times 2 + 1$$

$$2 = 2 \times 1$$

$$1 = 3 - 1 \times 2$$

$$= 20 \times 3 - 59$$

$$7 = 140 \times 3 + (-7) \times 59$$

So $x = 140 + 59t, t \in \mathbb{Z}$.

3. From Euclidean algorithm, one can infer that

$$\left\{ \begin{array}{l} \gcd(c,31)=1\\ \gcd(7,\phi(31))=\gcd(7,30)=1 \end{array} \right.$$

Using extended Euclidean algorithm to find p such that $7p \equiv 1 \mod 30$, one can get p = 13 + 30t, $t \in \mathbb{Z}$. Assume t = 0, then the plain text $m \equiv c^p \equiv c^{13} \mod 31$.

4.

$$4883 = 19 \times 257$$

 $4369 = 17 \times 257$

5. Let
$$M = \begin{pmatrix} 3 & 5 \\ 7 & 3 \end{pmatrix}$$
.

M is not invertible modulo p implies that $det(M \mod p) = 0$.

When p > 7, $det(M \mod p) = det(M) = -26$, M is invertible.

Try all $p \le 7$: p = 2, 3, 5, 7, one can get that $det(M \mod 2) = 0$. So when p = 2, M is not invertible.

6. Suppose a and p are coprime, i.e., gcd(a, p) = 1. Then according to Bézout's identity, $\exists r, s \in \mathbb{Z}$ such that ra + sp = 1. Multiply both sides by b, one can get rab + spb = b. Since $p \mid rab, p \mid spb$, then $p \mid b$, i.e., $b \equiv 0 \mod p$.

Similarly, if b and p are coprime, then $a \equiv 0 \mod p$. So at least one of a, b should be divisible by p.

7.

$$2^{2017} \equiv (2^2)^{1008} \times 2$$

 $\equiv (-1)^{1008} \times 2$
 $\equiv 2 \mod 5$

$$2^{2017} \equiv (2^6)^{336} \times 2$$
$$\equiv (-1)^{336} \times 2$$
$$\equiv 2 \mod 13$$

$$2^{2017} \equiv (2^5)^{403} \times 4$$
$$\equiv 1^{403} \times 4$$
$$\equiv 4 \mod 31$$

Since $2015 = 5 \times 13 \times 31$, Chinese Remainder Theorem may be used to calculated $2^{2017} \mod 2015$.

$$(-161) \times 5 + 2 \times 403 = 1$$

 $12 \times 13 + (-1) \times 155 = 1$
 $21 \times 31 + (-10) \times 65 = 1$

$$2^{2017} \equiv 2 \times 2 \times 403 + 2 \times (-1) \times 155 + 4 \times (-10) \times 65$$

 $\equiv -1298$
 $\equiv 717 \mod 2015$

Ex.2 - Rabin Cryptosystem

1. Find two big prime number p and q, and $p, q \equiv 3 \mod 4$. Rabin Cryptosystem takes n = pq as public key, (p, q) as private key.

To encrypt plain text m, cipher text c is calculated as

$$c = m^2 \bmod n$$

To decrypt cipher text, solve the simultaneous congruence equations

$$\begin{cases} m_p \equiv c^{\frac{p+1}{4}} \bmod p \\ m_q \equiv c^{\frac{q+1}{4}} \bmod q \end{cases}$$

Where m_p and m_q are defined as

$$\begin{cases} m \equiv \pm m_p \bmod p \\ m \equiv \pm m_q \bmod q \end{cases}$$

Apply the Chinese Remainder Theorem, one can get four different solutions for the plain text.

- 2. (a) Since at most four different solutions can be found, there is 25% of chance to observe a meaningful message when feeding random numbers.
 - (b) No, since if Eve only has x and public key n, he cannot solve like above. Whether to factorize n, or to solve $m \equiv \sqrt{x} \mod n$ are as the same difficulty.
 - (c) Eve should run a chosen cipher attack. Eve can then use the device to get the four different solutions as $\pm a, \pm b$. Since gcd(|a-b|, n) is a non-trivial factor of n, then Eve can factorize n by calculating |a-b|, which is either p or q.

Ex.3 - CRT

Assume there are x people in the group, then x satisfies that

$$x \equiv 1 \mod 3$$
$$x \equiv 2 \mod 4$$
$$x \equiv 3 \mod 5$$

Since 3, 4, 5 are coprime to each other, one can use Chinese Remainder Theorem to solve x.

$$7 \times 3 + (-1) \times 20 = 1$$

$$4 \times 4 + (-1) \times 15 = 1$$

$$5 \times 5 + (-2) \times 12 = 1$$

$$x \equiv 1 \times (-1) \times 20 + 2 \times (-1) \times 15 + 3 \times (-2) \times 12$$

$$\equiv -122$$

$$\equiv 58 \mod 60$$

Then the two smallest possible x are 58 and 118.