

Game theory problem set 5

$$1a \quad u_A(x_A, m_A) = x_A + m_A$$

$$u_B(x_B, m_B) = 3x_B + m_B$$

$$x_A + x_B = 100 \Rightarrow x_B = 100 - x_A$$

$$m_A + m_B = 200 \Rightarrow m_B = 200 - m_A$$

where u_C , x_C , and m_C denote C's utility, C's share of farmland in acres and C's share of cash respectively for $C \in \{A, B\}$.

$$X = \{(x_A, m_A) : x_A \in [0, 100], m_A \in [0, 200]\},$$

$$D = (x_A = 50, m_A = 100)$$

$$U = \{(u_A(x_A, m_A), u_B(x_B = 100 - x_A, m_B = 200 - m_A)) : (x_A, m_A) \in X\}$$

$$d = (u_A(x_A = 50, m_A = 100) = 150, u_B(x_B = 100 - x_A, m_B = 200 - m_A) = 250)$$

$$= u_B(x_B = 50, m_B = 100) = 250$$

where X is the set of possible agreements, D is the disagreement outcome, U is the set of payoff pairs corresponding to X and d is the payoff pair corresponding to D .

Suppose $x_A = 100, m_A = 200$. Then $u_A =$

Let

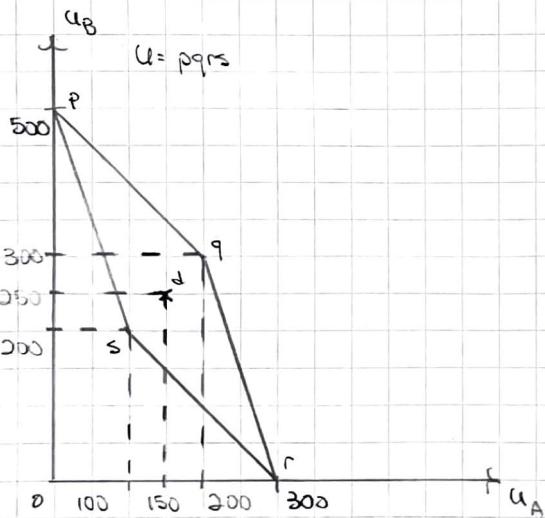
$$u_A(x_A = 100, m_A = 200) = 300, u_B(x_A = 100, m_A = 200) = 0$$

$$u_A(x_A = 100, m_A = 0) = 100, u_B(x_A = 100, m_A = 0) = 200$$

$$u_A(x_A = 0, m_A = 200) = 200, u_B(x_A = 0, m_A = 200) = 300$$

$$u_A(x_A = 0, m_A = 0) = 0, u_B(x_A = 0, m_A = 0) = 500$$

All other $x \in X$ are weighted averages of these possible agreements. Then, because utilities are linear, ~~weighted~~ for all other $x \in X$, the payoff pair corresponding to x is a weighted average of these four possible agreements.



b Let $F(U, d) = u$ be a bargaining solution that satisfies the Nash axioms.

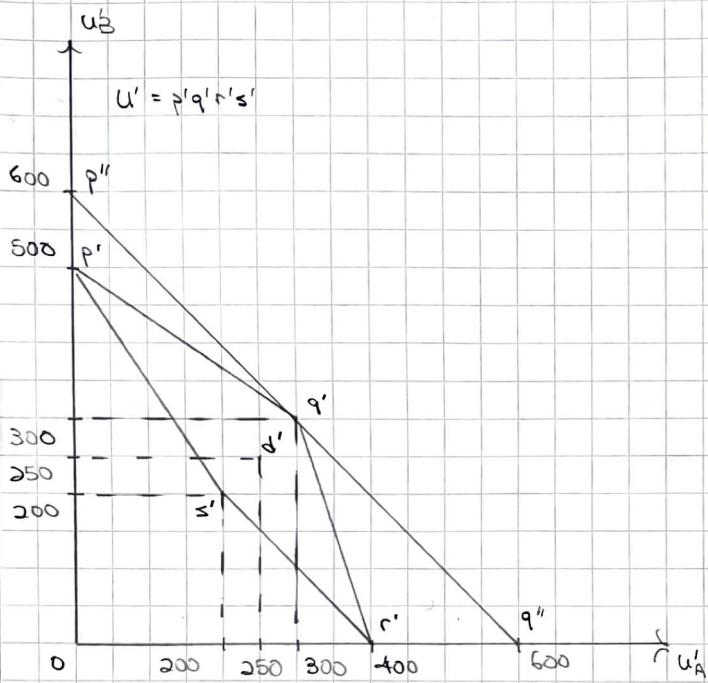
$$(1) \quad u_A'(x_A, m_A) = 2u_A(x_A, m_A) - 100, \text{ and } u_B'(x_B, m_B) = u_B(x_B, m_B),$$

$$(2) \quad \{(u_A'(x_A, m_A), u_B'(x_B = 100 - x_A, m_B = 200 - m_A)) : (x_A, m_A) \in X\},$$

$$\text{and } d = \{(u_A'(x_A = 50, m_A = 100), u_B'(x_B = 100 - 50 = 50, m_B = 200 - 50 = 150))\}$$

= C

Let $U_A(x_A, m_A) = u_A(x_A, m_A) + 100$, $U_B(x_B, m_B) = u_B(x_B, m_B)$,
 $U' = \{(U_A(x_A, m_A), U_B(x_B = 100 - x_A, m_B = 200 - m_A)) : (x_A, m_A) \in X\}$,
 $d' = (U_A, U_B)|_{d} = (250, 250)$



Consider the bargaining problem (U'', d') , where $U'' = \{(U_A, U_B) : U_A + U_B \in [0, 600]\}$. U'' is represented by area $O_p''q''$. By inspection, $U'' \subset U'$.

Suppose that $(U'_1, U'_2) \in U'$, then $U'_1 + U'_2 \in [0, 600]$, then $(U'_1, U'_2) \in U''$. Since $d'_1 = 250 = d'_2$, so the bargaining problem (U'', d') is symmetric.

Let $U'' = F(U', d')$. By symmetry, $U'' = (U_A'', U_B'')$ is such that $U_A'' = U_B''$. Suppose for reductio that $U_A'' + U_B'' > 600$. Let $\epsilon = 600 - U_A'' - U_B''$. Then $U_A'' + \frac{\epsilon}{2} + U_B'' + \frac{\epsilon}{2} = 600$, $\Rightarrow (U_A'' + \frac{\epsilon}{2}, U_B'' + \frac{\epsilon}{2}) \in U''$ and $U_A'' + \frac{\epsilon}{2} > U_A''$ and $U_B'' + \frac{\epsilon}{2} > U_B''$. By weak Pareto efficiency, $U'' \neq F(U', d')$. By reductio $U_A'' + U_B'' = 600$. Solving simultaneously, $U_A'' = U_B'' = 300$. By independence of irrelevant alternatives, $U'' = F(U', d') = U'' = (300, 300)$. By invariance to equivalent payoff representations, $U'' = F(U', d') = (U_A - 100, U_B) = (200, 300)$. This corresponds to the distribution $x_A = 0, m_A = 200, x_B = 100, m_B = 0$.

< Suppose that A is the divider. Denote the portions chosen by A as $f(x_A, m_A)$ and (x_i, m_i) and $(x_2 = 100 - x_1, m_2 = 200 - m_1)$. In the second stage subgame, B's maximisation problem is
 $\max_{i \in \{1, 2\}} U_B(x_i, m_i) = 3x_i + m_i$.

If B chooses (x_1, m_1) $\Rightarrow 3x_1 + m_1 \leq 3x_2 + m_2 = 3(100 - x_1) + 200 - m_1$,
 $3x_1 + m_1 \leq 500 - 3x_1 - m_1$,
 $3x_1 + m_1 \geq 250$

In the first stage, A's payoff reduces to
 $U_A(x_1, m_1) = x_1 + m_1$ iff $3x_1 + m_1 \leq 250$, $(100 - x_1) + (200 - m_1)$
otherwise

Suppose, without loss of generality, that A intends to induce B to choose (x_2, m_2) . Then A's maximisation problem is $\max_{x_1, m_1} x_1 + m_1$ subject to $3x_1 + m_1 \leq 250$, $x_1 \in [0, 100]$, and $m_1 \in [0, 200]$.

By inspection

By inspection, A chooses $(x_1 = 50/3, m_1 = 200)$, then $(x_2 = 250/3, m_2 = 0)$, and B (weakly) prefers (x_2, m_2) to (x_1, m_1) .

The SPE is the strategy profile such that A's strategy is to offer $(x_1 = 50/3, m_1 = 200)$ and $(x_2 = 250/3, m_2 = 0)$ and B's strategy is to choose (x_1, m_1) iff $3x_1 + m_1 > 250$.

Suppose that B is the divider. In the second stage subgame, A's maximisation problem is $\max_{x_1, m_1} \sum_{i \in \{1, 2\}} u_A(x_i, m_i) = x_1 + m_1$, if A chooses (x_1, m_1) & $u_A(x_1, m_1) = x_1 + m_1 > u_A(x_2, m_2) = x_2 + m_2 = (100 - x_1) + (200 - m_1)$ $x_1 + m_1 > 150$. In the first stage, B's payoff function reduces to $u_B(x_1, m_1) = 3x_1 + m_1$, iff $x_1 + m_1 \leq 150$, $3(100 - x_1) + (200 - m_1)$ otherwise. Suppose wlog that B intends to induce A to choose (x_2, m_2) . Then B's maximisation problem is $\max_{x_1, m_1} 3x_1 + m_1$ subject to $x_1 + m_1 \leq 150$, $x_1 \in [0, 100]$, $m_1 \in [0, 200]$. By inspection, $x_1 = 100$, $m_1 = 50$, then $(x_2 = 100 - x_1 = 0, m_2 = 200 - m_1 = 150)$, and A (weakly) prefers (x_2, m_2) to (x_1, m_1) . The SPE is the strategy profile such that B's strategy is to offer $(x_1 = 100, m_1 = 50)$ and $(x_2 = 0, m_2 = 150)$, and A's strategy is to choose (x_1, m_1) iff $x_1 + m_1 > 150$.

d) Suppose that A is the divider, then at SPE, ~~$U_A = 50/3 + 200 = 250$~~
 $U_A = 50/3 + 200 \neq U_B = 250$

Suppose that B is the divider, then at SPE,
 $U_A = 150$, $U_B = 350$

Expected Utilities

$$U_A = \frac{50/3 + 200 + 150}{2} = 1100/6, \quad U_B =$$

$$U_B = (350 + 250 + 350)/2 = 300$$

This solution is not Pareto optimal because it is Pareto-dominated by $(x_1, m_1) = (0, 200)$, where $U_A = 200$ and $U_B = 300$.

Graphically, each of the "pure" divide and choose outcomes is some point on the pqr frontier that does not coincide with q since each first mover divides portions such as to advantage itself. By the convexity of U , any weighted average of the "pure" outcomes lies within the frontier, and is not Pareto optimal.

2a) If $U' \subseteq U$ and $d' = d$ and $F(U, d) \in U'$, then $F(U') = F(U, d)$

Consider $U' = \{(u_1, u_2) : u_1, u_2 \in \mathbb{R}\}$, $U = U' \cup \{(0, 10)\}$,
and $d = d' = (0, 0)$. Then $u_1^+ = u_2^+ = u_1^- = 1$ and $u_2^- = 10$.

$\max_u u_1/u_1^+$ subject to $u_1/u_1^+ = u_2/u_2^+$ and $(u_1, u_2) \in U$

coincides with

$\max_u u_1$ subject to $u_1 = u_2/10$ and $(u_1, u_2) \in U$

~~= max u_1 subject to~~

~~etc~~

= $\max_u u_1$ subject to $(u_1, 10u_1) \in U$

= 1/10

$$(u_1^*, u_2^*) = (1/10, 1)$$

$\max_u u_1/u_1^+$ subject to $u_1/u_1^+ = u_2/u_2^+$ and $(u_1, u_2) \in U'$

coincides with

$\max_u u_1$ subject to $u_1 = u_2$ and $(u_1, u_2) \in U'$

= $\max_u u_1$ subject to $(u_1, u_1) \in U'$

= 1

$$(u_1^*, u_2^*) = (1, 1)$$

Consequently, KS solution does not satisfy IA.

b) $u_1(x_1) = x_1$

$u_2(x_2) = x_2^p$ for $0 < p < 1$

~~etc~~

$$U = \{(u_1(x_1), u_2(x_2)) : x_1 + x_2 = 1\}$$

$$u_1^* = \max_u u_1 \text{ st } (u_1, u_2) \in U = 1$$

$$u_2^* = \max_u u_2 \text{ st } (u_1, u_2) \in U = 1$$

$\max_u u_1/u_1^+$ st $u_1/u_1^+ = u_2/u_2^+, (u_1, u_2) \in U$

~~= max u_1 st u_1 = u_2, (u_1, u_2) \in U~~

~~= max u_1 st (u_1, u_1) \in U~~

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Suppose that $u_1/u_1^+ = u_2/u_2^+$, then $u_1 = u_2, x_1 = x_2^p$,
 $x_1 = (1-x_2)^p, x_1^{1/p} + x_2 - 1 = 0 \Rightarrow x_2 > x_1$ since $0 < p < 1$

and $x_1, x_2 \in [0, 1]$

$\max_u u_1/u_1^+ \text{ st } u_1/u_1^+ = u_2/u_2^+, (u_1, u_2) \in U$

~~= max u_1 st u_1 = u_2, (u_1, u_2) \in U~~

~~= max x_1 st x_1^{1/p} + x_2 - 1 = 0, x_1 \in [0, 1]~~

~~etc~~

Suppose that $u_1/u_1^+ = u_2/u_2^+$ and $(u_1, u_2) \in U$

Suppose for reductio that $x_2 = 1$, then $x_1 = 1$, then

$x_1 = 1$, then $(u_1, u_2) = (1, 1) \notin U$. By reductio, $x_2 \neq 1$

then $x_2 \in [0, 1]$ ~~but so~~ since consequently,

since given $0 < p < 1$, $x_1 > x_2$.

c) Let $U' = \{(u_1(x_1), u_2(x_2)) : x_1 + x_2 = 1\}$, ~~$u_1^{**} = u_2^{**} = 1$~~

$$u_1^{**} = s, u_2^{**} = s^p$$

Suppose that $u_1^{**}/u_1^+ = u_2^{**}/u_2^+$ and $(u_1^{**}, u_2^{**}) \in U'$

$$s^p/x_1^p = s^p/(1-x_1)^p \Rightarrow (x_1/s)^p = (x_2/s)^p$$

$$x_1/s = (x_2/s)^p \Rightarrow 1 - x_2/s = (x_2/s)^{p-1}$$

$$\max u_i/u_1 \text{ subject to } u_i/u_1 = u_2/u_2^*, \quad (u_i, u_2) \in U$$

$$= \max x_1/s \text{ st } x_1/s = x_2/s^{\rho}, \quad (x_1, x_2) \in U$$

$$= \max x_1/s \text{ st } x_2^{\rho} = x_2 s^{-1}, \quad (x_1, x_2) \in U$$

$$= \max x_1/s \text{ st } (x_1, x_2 s^{-1}) \in U$$

coincides with

$$\max x_1 \text{ st } (x_2 s^{1-\rho}, x_2^{\rho}) \in U$$

$$= \max x_1 \text{ st } x_2 s^{1-\rho} + x_2 = s$$

$$\rightarrow \max x_1$$

$$\max x_2 s^{1-\rho} \text{ st } (x_2 s^{1-\rho}, x_2^{\rho}) \in U$$

$$= \max x_2 s^{1-\rho} \text{ st } x_2 s^{1-\rho} + x_2 = s$$

$$= \max x_2 s^{1-\rho} \text{ st } x_2 = \frac{s}{1+s^{1-\rho}}$$

$$u_i^*/u_1^* = u_2^*/u_2^* \text{ and } (u_i^*, u_2^*) \in U$$

$$x_1/s = x_2^{\rho}/s^{\rho}$$

$$x_1 = x_2^{\rho} s^{1-\rho}$$

$$(x_1, x_2^{\rho}) \in U$$

~~$$(x_1, x_2^{\rho} s^{1-\rho}, x_2^{\rho}) \in U$$~~

$$x_2^{\rho} s^{1-\rho} + x_2 = s$$

$$x_2^{\rho} s^{1-\rho} + x_2 s^{-1} = 1$$

~~$$(x_2/s)^{\rho} + (x_2/s)^{-1} = 1$$~~

Differentiating implicitly w.r.t s

$$-\rho x_2^{\rho-1} s^{-\rho-1} + \rho x_2^{\rho-1} \cancel{\frac{dx_2}{ds}} s^{-\rho}$$

$$1 - x_2 s^{-2} + \cancel{\frac{dx_2}{ds}} s^{-1} = 0$$

$$(\rho x_2^{\rho-1} s^{-\rho} + s^{-1}) \cancel{\frac{dx_2}{ds}} = \rho x_2^{\rho-1} s^{-\rho-1} + x_2 s^{-2}$$

$$\text{Given } \cancel{\frac{dx_2}{ds}} > 0, \quad 0 < \rho < 1, \quad x_2 > 0, \quad s > 0,$$

$$\cancel{\frac{dx_2}{ds}} > 0$$

$$\frac{dx_1}{ds} = \rho x_2^{\rho-1} s^{1-\rho} > 0$$

$$\frac{dx_1}{ds} > 0$$

consequently for ~~x_1 and x_2~~ each increase

when ~~s~~ increases from 1 to some $s > 1$

- 3a In the second stage subgame, given that P1 chose x_3 ,
 P2 chooses A (accept) or R (reject) to maximise
 $u_2 = \begin{cases} x_3^p - sx_3^p = x_3^p - s(1-x_3)^p & \text{if } a_2 = A \\ 0 & \text{if } a_2 = R \end{cases}$

where a_2 denotes P2's action. P2 chooses A iff
 $x_3^p - sx_3^p \geq 0, x_3^p \geq sx_3^p, x_3 \geq s^{1/p}x_1, x_3 \geq s^{1/p}(1-x_2)$
 $(1+s^{1/p})x_3 \geq s^{1/p}, x_3 \geq s^{1/p}/(1+s^{1/p})$. Then P1's payoff function in the first stage reduces to
 $u_1(x_3) = \begin{cases} s^{1/p} & \text{if } x_3 \geq s^{1/p}/(1+s^{1/p}) \\ 0 & \text{otherwise} \end{cases}$

then P1's maximisation problem is $\max_{x_3} u_1(x_3)$ which has solution $x_3 = s^{1/p}/(1+s^{1/p})$.

The SPE is the strategy profile such that P1's strategy is to offer $x_3 = s^{1/p}/(1+s^{1/p})$ and P2's strategy is to A iff $x_3 \geq s^{1/p}/(1+s^{1/p})$ and R otherwise.

- b Supposing that s is known to P2, P2's payoff function, maximisation problem, and best response function are unchanged from (a). P2 chooses A iff
 $x_2 \geq s^{1/p}(1-x_3), s^{1/p} \leq x_2/(1-x_3), s \leq (x_2/(1-x_3))^p$.

P1's expected payoff function reduces to

$$\pi_1(x_2) = (1-x_2)^p P(s \leq (x_2/(1-x_3))^p) + \bar{P}(s > (x_2/(1-x_3))^p)$$

$$= (1-x_2)^p F((x_2/(1-x_3))^p)$$

$$= (1-x_2)^p (x_2/(1-x_3))^{p^2}$$

P1's maximisation problem is $\max_{x_2} \pi_1(x_2)$ which coincides with $\max_{x_2} (1-x_2)(x_2/(1-x_3))^{p^2}$ since $0 < p < 1$ so raising to the power p^2 is a strictly increasing monotonic transformation.

$$\begin{aligned} \text{FOC: } & \frac{\partial}{\partial x_2} (1-x_2)(x_2/(1-x_3))^{p^2} \\ &= \frac{\partial}{\partial x_2} x_2^2/(1-x_2) \\ &\leftarrow \cancel{2x_2/(1-x_2)} + \cancel{x_2^2/(1-x_2)^2} = 0 \\ &\cancel{2x_2x_3^2} \cancel{2x_2/(1-x_2)} + x_2^2 = 0 \\ &\equiv 2x_2 - x_2^2 = 0 \end{aligned}$$

$$\begin{aligned} & \cancel{x_2 = 0 \text{ (reject) or } x_2 = 1 \text{ (offer)}} \\ \text{FOC: } & \frac{\partial}{\partial x_2} x_2^2/(1-x_2)^{p^2} \\ &= \cancel{(2p-2)x_2} \cancel{2px_2^{2p-1}}/(1-x_2)^p + x_2^{2p} (-p)(1-x_2)^{p-1}(-1) \\ &= 2px_2^{2p-1}/(1-x_2)^p + p^2x_2^{2p}/(1-x_2)^{p+1} = 0 \\ &\cancel{2x_2^{p-1}} \cancel{2/p} \cancel{x_2} \\ &\cancel{-x_2/(1-x_2)} \\ &\leftarrow 2/(1-x_2)^p + x_2/(1-x_2)^{p+1} = 0 \\ &\leftarrow x_2^2/(1-x_2)^p = 0 \\ &x_2 = 0 \text{ (reject)} \end{aligned}$$

There are no interior solutions. The constraints are $x_2 \in [0, 1]$ and $F(s) \in [0, 1] \Rightarrow (x_2/(1-x_3))^{p^2} \in [0, 1] \Rightarrow x_2 \in [0, 1/2]$. The SPE is such that P1's strategy is to offer $x_2 = 1/2$ and P2's strategy is to choose A iff $s \leq (x_2/(1-x_3))^p$. At SPE, P2 chooses A with probability 1.

c Expected total utility $U = (1-x_2)^p + E(x_2^p - s(1-x_2)^p)$

$$\begin{aligned} &= x_2^p + (1-E(s))(1-x_2)^p = x_2^p + (1-2/3)(1-x_2)^p \\ &= (1/2)^p + 1/3(1/2)^p = 1/3(1/2)^p \end{aligned}$$