

# Game Theory Supplementary Exercises 1

1a

	L	C	R
T	<u>4</u>	1	2
M	<u>4</u>	0	<u>3</u>
B	1	<u>1</u>	1
	2	-1	2

Best responses underlined

At

Player i's action  $a_i$  is rationalisable, by definition, if it is a best response against some potentially correlated mix of other players' actions. Equivalently, player i's action is rationalisable, by definition, if it survives iterated elimination of ~~strict~~ actions that are not best responses against some potentially correlated mix of other players' actions.

$L \geq R$ , so  $R$  is never a best response

By inspection,  $L \geq R$ , then by definition of strict dominance and BR,  $R$  is never a BR.

Likewise,  $T \geq B$ , so  $B$  is never a BR.

Each  $\neq$  Then, each remaining action is a BR to some action of the other player, so iterated elimination of never-BR terminates here. T, M, L, and C only are rationalisable.

Since CKR implies players play only rationalisable actions, we consider the reduced game involving only rationalisable actions

	L	C
T	<u>4</u>	1
M	<u>4</u>	<u>1</u>

Best responses underlined

By inspection, (T, L) and (M, C) are the only pure NE where players play mutual best responses.

By inspection there are no hybrid NE since no player is ever indifferent between his actions when the other plays some fixed strategy. Consider

at  $s^*$  P1 plays T with probability  $p$  and M with probability  $1-p$  and P2 plays L with probability  $q$  and C with probability  $1-q$  for  $p, q \in (0, 1)$ . If  $s^*$  is a mixed NE, then  $\pi_1(T, s^*) = \pi_1(M, s^*)$ ,

$$4p = q + (1-q), q = \frac{1}{4} \text{ and}$$

$$\pi_2(L, s^*) = \pi_2(C, s^*),$$

$$4p = p + (1-p), p = \frac{1}{4}.$$

So the only mixed NE is  $s^* = (\frac{1}{4}T + \frac{3}{4}M, \frac{1}{4}L + \frac{3}{4}C)$

b

	L	C	R
T	<u>4</u>	1	2
M	<u>4</u>	0	<u>3</u>
B	1	<u>1</u>	1
	2	-1	2

Best responses underlined

By inspection,  $L \geq R$  and  $T \geq B$ .

Since  $R$  is never a BR

we consider the reduced game involving only strategies that survive IED.

	L	C
T	<u>4</u>	1
M	<u>4</u>	<u>0</u>

By inspection, (T, L) and (M, C) are the only pure NE where players play mutual best responses. Consider some strategy profile  $s^*$  such that P1 mixes T and M, i.e.  $s^*_1 = (pT + (1-p)M)$ ,  $p \in (0, 1)$ . If  $s^*$  is an NE, P1 has no profitable deviation, so  $\pi_1(T, s^*_2) = \pi_1(M, s^*_2)$ ,  $4p = 1q$  where

$$s^*_2 = (qL + (1-q)C) \text{ for } q \in [0, 1], \text{ so } q = 0 \text{ and } s^*_2 = C$$

If  $s^*$  is an NE, P2 has no profitable deviation, so  $\pi_2(C, s^*_1) \geq \pi_2(L, s^*_1)$ ,  $p \geq 4p$ ,  $p \leq 0$ .

This contradicts the supposition that  $p \in (0, 1)$ , so by reductio, there is no NE where P1 mixes T and M. By symmetry, there is no NE where P2 mixes L and C.

A mixed NE does not exist because each player has a weakly dominant strategy action, so any mixed strategy fails to the deviation which reallocates probability mass

2a

	L	C	R
T	<u>15</u>	11	12
M	10	<u>16</u>	<u>8</u>
B	2	1	<u>5</u>
	12	20	5
	11	19	4

Best responses underlined

Player i's action  $a_i$  is rationalisable, by definition, if it is a best response against some potentially correlated mix of other players' actions.

Equivalently, player i's action  $a_i$  is rationalisable if it survives iterated elimination of actions that are not best responses to any potentially correlated mix of other players' yet uneliminated actions.

By inspection  $M \geq B$ , so  $B$  is by definition of strict dominance and BR,  $B$  is never a best response against any potentially correlated mix of P2's strategies. Similarly,  $R \geq C$  so  $C$  is never



a best response against any of P1's strategies so eliminate B and C. then the game reduces to

	L	R
T	15	12
M	10	8
	2	5

By inspection, each remaining strategy is a best response to some strategy of the other player, so iterated elimination stops here.

T, M, L, and R only are rationalisable.

Let  $s^*$  be the unique NE. By definition of strict dominance and NE, only strategies that survive ED are played in NE, so B and C are not played in NE. Then, given that  $s^*$  is fully mixed,  $s^* = (pT + (1-p)M, qR + (1-q)L)$  for some  $p, q \in (0, 1)$ .

By definition of NE, at  $s^*$ , there is no profitable deviation, so  $\pi_1(T, s^*) = \pi_1(M, s^*)$  and  $\pi_2(L, s^*) = \pi_2(R, s^*)$  since otherwise some player can profitably deviate by reallocating probability mass to the action with higher payoff and  $\pi_1(T, s^*) = \pi_1(M, s^*) \geq \pi_1(B, s^*)$  and  $\pi_2(L, s^*) = \pi_2(R, s^*) \geq \pi_2(C, s^*)$  since otherwise some player can profitably deviate by reallocating probability mass to the action with the higher payoff.

$$10q + 8(1-q) = 12q + 5(1-q) \Rightarrow 1q = \frac{4}{3}(1-q)$$

$$3 - 3q = 2q, 5q = 3, q = \frac{3}{5}$$

$$15p + 2(1-p) = 12p + 5(1-p) \Rightarrow 11p = 1(1-p)$$

$$3p = 3(1-p), 6p = 3, p = \frac{1}{2}$$

$s^* = (\frac{1}{2}T + \frac{1}{2}M, \frac{3}{5}L + \frac{2}{5}R)$  is the unique NE.

Each correlated equilibrium of this game

$$P = \begin{pmatrix} P_{TL} & P_{TC} & P_{TR} \\ P_{ML} & P_{MC} & P_{MR} \\ P_{BL} & P_{BC} & P_{BR} \end{pmatrix}$$

is a joint distribution that satisfies the incentive compatibility constraints

P1 finds it optimal to play T when so instructed, to M when so instructed, and to play B when so instructed. Similarly for P2.

Since  $M \succ B$ , player P1 never finds it optimal to play B when so instructed, since it is always better to play M. Likewise, P2 never finds it optimal to play C since  $R \succ C$ , so each correlated equilibrium never instructs P1 to play B and never instructs P2 to play C.

$P_{TC} = P_{MC} = P_{BL} = P_{BC} = P_{BR} = 0$ . So each correlated equilibrium reduces to

$$P = \begin{pmatrix} P_{TL} & P_{TR} \\ P_{ML} & P_{MR} \end{pmatrix}$$

The incentive constraints are P1 finds it optimal to play T when so instructed

$$10P_{TL} + 8P_{TR} \geq 12P_{TL} + 5P_{TR}$$

P1 finds it optimal to play M when so instructed

$$12P_{ML} + 5P_{MR} \geq 10P_{ML} + 8P_{MR}$$

For player P2,

$$15P_{TL} + 2P_{ML} \geq 12P_{TL} + 5P_{ML}$$

$$12P_{TR} + 5P_{MR} \geq 15P_{TR} + 2P_{MR}$$

Simplifying,

$$3P_{TR} \geq 2P_{TL}, P_{TR} \geq \frac{2}{3}P_{TL}$$

$$2P_{ML} \geq 3P_{MR}, P_{ML} \geq \frac{3}{2}P_{MR}$$

$$3P_{TL} \geq 3P_{ML}, P_{TL} \geq P_{ML}$$

$$3P_{MR} \geq 3P_{TR}, P_{MR} \geq P_{TR}$$

Let  $P_{TL} = x$ , then one solution is

$$P_{TL} = x, P_{TR} = \frac{2}{3}x, P_{MR} = \frac{2}{3}x, P_{ML} = x$$

$$P_{TL} + P_{TR} + P_{ML} + P_{MR} = 1, \frac{10}{3}x = 1, x = \frac{3}{10}$$

$$P_{TL} = \frac{3}{10}, P_{TR} = \frac{2}{10}, P_{ML} = \frac{3}{10}, P_{MR} = \frac{2}{10}$$

This corresponds to the mixed NE  $(\frac{1}{2}T + \frac{1}{2}M, \frac{3}{5}L + \frac{2}{5}R)$

By inspection of the simplified incentive compatibility constraints,

$$\uparrow P_{TL} \text{ from the given solution} \Rightarrow \uparrow P_{TR} \Rightarrow \uparrow P_{MR} \Rightarrow \uparrow P_{ML}$$

$$\Rightarrow \uparrow (P_{TL} + P_{TR} + P_{ML} + P_{MR}) \Rightarrow 1. \text{ So by reductio there is}$$

no solution where  $P_{TL} >$

$$P_{TL} > \frac{3}{10} \Rightarrow P_{TR} > \frac{2}{10} \Rightarrow P_{MR} > \frac{2}{10} \Rightarrow P_{ML} > \frac{3}{10},$$

$$\text{collectively} \Rightarrow (P_{TL} + P_{TR} + P_{ML} + P_{MR}) > 1$$

$$P_{TL} < \frac{3}{10} \Rightarrow P_{TR} < \frac{2}{10} \Rightarrow P_{MR} < \frac{2}{10} \Rightarrow P_{ML} < \frac{3}{10},$$

$$\text{collectively} \Rightarrow (P_{TL} + P_{TR} + P_{ML} + P_{MR}) < 1$$

By reductio definition of a correlated equilibrium,

$$P_{TL} + P_{TR} + P_{ML} + P_{MR} = 1$$

By reductio, there is no correlated equilibrium where  $P_{TL} \neq \frac{3}{10}$ . Similarly, there is no correlated equilibrium where  $P_{TR} \neq \frac{2}{10}$ ,  $P_{ML} \neq \frac{3}{10}$ , or  $P_{MR} \neq \frac{2}{10}$ .

The unique correlated equilibrium is unique and coincides with the unique Nash equilibrium.

This correlated equilibrium both maximises and minimises each player's the sum of players' payoffs.

$$\text{state space } \Omega = \{w, x, y, z\}$$

Probability distribution over  $\Omega$ ,  $\pi$  such that  $\pi(w) = \pi(y) = \frac{3}{10}$ ,  $\pi(x) = \pi(z) = \frac{2}{10}$

Partition for player P1

$$\mathcal{P}_1 = \{\{w, x\}, \{y, z\}\}$$

Partition for P2

$$\mathcal{P}_2 = \{\{w, y\}, \{x, z\}\}$$

Strategy for P1

$$\sigma_1 \text{ such that } \sigma_1(w) = \sigma_1(x) = T, \sigma_1(y) = \sigma_1(z) = M$$

Strategy for P2

$$\sigma_2 \text{ such that } \sigma_2(w) = \sigma_2(y) = L, \sigma_2(x) = \sigma_2(z) = R$$



3a Set of players  $N = \{1, 2\}$

Set of actions of each player  $i$

$$A_i = \mathbb{R}^+ \quad B_i = \mathbb{R}^+$$

Payoff of each player  $i$

$$u_i(b_i, b_{-i}) = \begin{cases} -b_i & \text{if } b_i < b_{-i} \\ \frac{a}{2} - b_i & \text{if } b_i = b_{-i} \\ a - b_i & \text{if } b_i > b_{-i} \end{cases}$$

Let  $BR_i(b_{-i})$  denote player  $i$ 's best response given player  $-i$  plays  $b_{-i}$

Suppose  $b_{-i} = a$ , then

$$u_i(b_i = 0, b_{-i}) = 0$$

$$u_i(b_i \in (0, b_{-i}), b_{-i}) = -b_i < 0$$

$$u_i(b_i = b_{-i}, b_{-i}) = \frac{a}{2} - b_{-i} = -\frac{a}{2} < 0$$

$$u_i(b_i > b_{-i}, b_{-i}) = a - b_{-i} < 0$$

$$\text{So } BR_i(b_{-i} = a) = 0$$

Suppose  $b_{-i} > a$ , then

$$u_i(b_i = 0, b_{-i}) = 0$$

$$u_i(b_i \in (0, b_{-i}), b_{-i}) = -b_i < 0$$

$$u_i(b_i = b_{-i}, b_{-i}) = \frac{a}{2} - b_{-i} < 0$$

$$u_i(b_i > b_{-i}, b_{-i}) = a - b_{-i} < 0$$

$$\text{So } BR_i(b_{-i} > a) = 0$$

Suppose  $b_{-i} < a$ , then

$b_i < b_{-i}$  fails to deviation  $b'_i = b_{-i} + \epsilon$  for sufficiently small  $\epsilon > 0$  since  $u_i(b_i < b_{-i}, b_{-i}) = -b_i \leq 0$  and  $u_i(b'_i = b_{-i} + \epsilon, b_{-i}) = a - b'_i = (a - b_{-i}) - \epsilon > 0$  for  $\epsilon < (a - b_{-i})$

$b_i = b_{-i}$  fails to a similar deviation since

$$u_i(b_i = b_{-i}, b_{-i}) = \frac{a}{2} - b_{-i} < u_i(b'_i = b_{-i} + \epsilon, b_{-i}) = a - b_{-i} - \epsilon \text{ for } \epsilon < \frac{a}{2}$$

$b_i > b_{-i}$  fails to deviation  $b'_i = b_i + (b_i - b_{-i}/2)$  since  $u_i(b_i > b_{-i}, b_{-i}) = a - b_i < u_i(b'_i, b_{-i}) = a - b'_i$  where  $b'_i > b_i$  and  $b_{-i} < b'_i < b_i$ .

$$\text{So } BR_i(b_{-i} < a) = \emptyset$$

$$BR_i(b_{-i}) = \begin{cases} 0 & \text{if } b_{-i} \geq a \\ \emptyset & \text{if } b_{-i} < a \end{cases}$$

By definition of NE and BR, players play mutual best responses at  $NE^*$ .

$b_1 = BR_1(b_2)$  Suppose that some strategy pure strategy profile  $b^*$  is a NE, then

$$b_1^* = BR_1(b_2^*) = \begin{cases} 0 & \text{if } b_2^* \geq a \\ \emptyset & \text{if } b_2^* < a \end{cases}$$

By definition of NE and strategy profile,  $b_1^*, b_2^* \neq \emptyset$ , so  $b_1^* = b_2^* = 0$ , then  $b_2^* < a$ , then  $b_1^* = \emptyset$ . By reductio, there is no pure NE  $b^*$ .

player

b Suppose there is a mixed NE where each firm  $i$  mixes a finite number of actions. Let  $\bar{b}_i$  be the highest bid on which firm  $i$  places positive probability.

Suppose that  $\bar{b}_1 \neq \bar{b}_2$ ,  $\bar{b}_1 \neq \bar{b}_2$ , then  $\bar{b}_i > \bar{b}_j$  for some  $i, j$ . Then player  $i$  can profitably deviate by reallocating probability mass from  $\bar{b}_i$  to  $\bar{b}'_i \in (\bar{b}_j, \bar{b}_i)$  since  $u_i(\bar{b}_i, b_j) = a - \bar{b}_i < u_i(\bar{b}'_i, b_j) = a - \bar{b}'_i$  for all  $b_j$ , so  $E(u_i(\bar{b}_i, b_j)) < E(u_i(\bar{b}'_i, b_j))$ .

Suppose instead that  $\bar{b}_1 = \bar{b}_2$ . Then either player  $i$  can profitably deviate by reallocating probability mass from  $\bar{b}_i$  to  $\bar{b}'_i = \bar{b}_i + \epsilon$  for sufficiently small  $\epsilon > 0$ . This reallocation yields an increase in expected payoff from a greater ~~then player~~ on the probability  $\bar{p}_j$  that player  $j$  plays  $\bar{b}_j$  under the prior strategy profile ~~then where previously there was a  $\bar{p}_i \bar{p}_j$~~  probability that  $b_i = \bar{b}_i = \bar{b}_j = b_j$  so  $u_i(b_i, b_j) = \frac{a}{2} - \bar{b}_i$ , under the deviated strategy profile, there is a  $\bar{p}_i \bar{p}_j$  probability that  $b_i = \bar{b}'_i > \bar{b}_j = b_j$  so  $u_i(b_i, b_j) = a - \bar{b}'_i = a - \bar{b}_i - \epsilon$ .

Suppose further that under the candidate equilibrium,  $\bar{b}_i$  is played with probability  $\bar{p}_i$  and  $\bar{b}_j$  is played with probability  $\bar{p}_j$ . ~~the change in  $u_i$  when~~ Under the proposed deviation, when  $i$  plays  $\bar{b}'_i$  instead of  $\bar{b}_i$  with probability  $\bar{p}_i$ , ~~is payoff decreases by  $\epsilon$  but~~  $u_i(\bar{b}'_i, b_j) = a - \bar{b}'_i = u_i(\bar{b}_i, b_j) - \epsilon$  for  $b_j \neq \bar{b}_j$  and  $u_i(\bar{b}'_i, \bar{b}_j) = u_i(\bar{b}_i, \bar{b}_j) + \frac{a}{2} - \epsilon$  for  $b_j = \bar{b}_j$ . So ~~if~~ is expected payoff increases if  $\epsilon < \bar{p}_j \frac{a}{2}$ . Informally, this deviation is profitable for  $i$  if the cost of the higher bid is smaller than the ~~value of~~ value of winning the prize with greater probability.

So if each player mixes over a finite number of actions, some player has a profitable deviation, this is not a NE.

< Suppose there is a symmetric mixed NE,  $b^*$ , ~~where then by the result~~ where each player  $i$  plays mixed strategy  $b_i$  which is a probability over distribution over  $B_i$ . Let  $F(b)$  denote the common cdf of  $b_1^*$  and  $b_2^*$ .

There are no gaps in the support of  $F(b)$ . Suppose that there is a gap  $[\bar{b}, \bar{b}]$  in the support of  $F(b)$ , then i.e. ~~no~~ players never play an action  $b_i \in [\bar{b}, \bar{b}]$ . Then, each player can profitably deviate by reallocating probability mass from ~~actions to actions~~ "just above" this interval ~~the~~  $\bar{b}$  to actions "just above"  $\bar{b}$  since this reallocation has no effect on a player's the probability of each player's making the winning bid but reduces the expected cost of bidding.



By definition of NE, at  $b^*$ , there is no profitable deviation, so ~~each the payoff to each~~ for each player, the payoff of each action in the support of  $F(b)$  must be equal, since otherwise ~~some~~ some player can profitably deviate by reallocating probability mass from actions in this support with lower expected payoff to actions with higher payoff, i.e.

$E(u_i(b_i, b_{-i}^*)) = F(b_i)a - b_i$  is constant for all  $b_i$  in the support of  $F(b)$ .

~~0 is in the~~

0 is in the support of  $F(b)$ . Suppose that 0 is not in the support of  $F(b)$ , then given that there are no gaps in the support of  $F(b)$ , the support of  $F(b)$  is ~~some~~  $[b, b]$  for some  $0 < b < \infty \in \mathbb{R}^+$ . Then, each player can profitably deviate by reallocating probability mass from actions ~~at the bottom of this~~  $b_i \in [b, b + \varepsilon]$  to  $b_i = 0$  since the former actions have payoff  $-b < 0$  in the limit as  $\varepsilon$  becomes small while  $b_i = 0$  has payoff 0.

$$0 = F(b_i) F(0)a - 0 = F(b_i)a - b_i \geq 0,$$

$$b_i = F(b_i) F(0)a - 0 = F(b_i)a - b_i \geq 0,$$

$F(b_i)$  corresponds to the cdf of  $u[0, a]$ .

At  $b^*$ , each player plays the mixed strategy  $b_i^*$  which is a uniform probability distribution over  $[0, a]$ . Each player's expected payoff is 0. In first-price all-pay auctions, ~~buyers~~ (at least in the two-player case), all surplus is captured by the seller.