

Exercise 3

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Problem 1

$$\min_{\mathbf{x}} \mathbf{c}^T \mathbf{x} \quad \text{st. } \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq 0$$

KKT conditions:

$$\mathbf{A}^T \boldsymbol{\lambda} + \mathbf{s} = \mathbf{c} \quad (1)$$

$$\mathbf{A}\mathbf{x} = \mathbf{b} \quad (2)$$

$$\mathbf{x} \geq 0 \quad (3)$$

$$\mathbf{s} \geq 0 \quad (4)$$

$$x_i s_i = 0 \quad (5)$$

- a) Calculation of Newton direction assumes that the Hessian of $\mathbf{c}^T \mathbf{x}$ is invertible, but

$$\nabla^2 \mathbf{c}^T \mathbf{x} = 0 \quad (\text{zero-matrix})$$

Since it is linear, and thus $(\nabla^2 \mathbf{c}^T \mathbf{x})^{-1}$ doesn't exist.

b) The objective function is linear so it is clearly convex.

$$\text{Proof: } f(\alpha x + (1-\alpha)y)$$

$$= C^T (\alpha x + (1-\alpha)y)$$

$$= \alpha C^T x + (1-\alpha) C^T y$$

$$= \alpha f(x) + (1-\alpha) f(y)$$

$$\leq \alpha f(x) + (1-\alpha) f(y)$$

□

To show that the feasible set Ω is convex we'll consider two points $x, y \in \Omega$ and the line connecting them $\alpha x + (1-\alpha)y$, $\alpha \in [0, 1]$. If this line is fully contained in Ω , then Ω is convex.

$$\begin{aligned} A(\alpha x + (1-\alpha)y) &= \alpha Ax + (1-\alpha)Ay \\ &= \alpha b + (1-\alpha)b \\ &= b \quad \checkmark \end{aligned}$$

Since $x \geq 0$, $y \geq 0$, $\alpha \geq 0$ and $(1-\alpha) \geq 0$ we can say

$$\alpha x + (1-\alpha)y \geq 0$$

✓

This shows that the line is feasible and thus that the set is convex.

$$c) \max_{\lambda} b^T \lambda \text{ s.t. } A^T \lambda \leq c$$

We first rephrase it as a minimization problem with equality constraints

$$\min_{\lambda} -b^T \lambda \text{ s.t. } c - A^T \lambda - s = 0, s \geq 0$$

The Lagrangian is given by

$$\mathcal{L}(\lambda, s, x) = -b^T \lambda + x^T (c - A^T \lambda - s) + x^T s$$

Optimality requires $\nabla_{\lambda} \mathcal{L} = 0, x \geq 0, x^T s = 0$

$$\nabla_{\lambda} \mathcal{L} = 0 \Rightarrow 0 = -b + Ax$$

$$\Leftrightarrow Ax = b$$

$$x^T s = 0 \Leftrightarrow x_i s_i = 0 \text{ for } i=1, \dots, n$$

Summarizing these requirements:

$$A^T \lambda + s = c$$

$$Ax = b$$

$$x \geq 0$$

$$s \geq 0$$

$$x_i s_i = 0$$

(KKT-1)

d) $C^T x^* = (A^T \lambda^* + S^*)^T x^*$

$$= (\lambda^* A + S^{*T}) x^*$$
$$= \underbrace{\lambda^* A x^*}_{\lambda^* b} + S^{*T} x^*$$
$$= \lambda^* b + 0$$
$$= (b^T \lambda^*)^T$$
$$= b^T \lambda^* \quad (\text{scalar})$$

They are equal

e) A point x is feasible if there exists a subset B of the indices, such that

- $|B| = m$
- For $i \notin B$ we have $x_i = 0$
- The matrix B defined by the i 'th columns of A , where $i \notin B$ is non-singular.

f) If a_i is the i 'th row of A , then we can write the equality constraints as

$$C_i(x) = a_i^T x - b_i = 0$$

$$\Rightarrow \nabla C_i = a_i^T$$

As long as $x_i \neq 0$ for all i , the active set is given by the equality constraints. Then the linear independence of the rows of A guarantees that the LICQ holds.

Problem 2

a) $x_1 = x_A = \underbrace{\text{amount}}$ of A ^{in 100kg}

$x_2 = x_B = \underline{\quad}$ 11 - B

$f(x) = \text{"profit"}$

$$= \frac{3}{2}x_1 + x_2 = -C^T x, \quad C = \begin{pmatrix} -3/2 \\ -1 \end{pmatrix}$$

Constraints:

$$0 \leq R_I \leq 8$$

$$0 \leq R_{II} \leq 15$$

where

$$R_I = 2 \cdot x_1 + 1 \cdot x_2$$

$$R_{II} = 1 \cdot x_1 + 3 \cdot x_2$$

By adding slack variables $x_3, x_4 \geq 0$ we can enforce the equality constraint.

This constraint can be stated as

$$Ax = b, \quad A = \begin{pmatrix} 2 & 1 & 1 & 0 \\ 1 & 3 & 0 & 1 \end{pmatrix}, \quad b = \begin{pmatrix} 8 \\ 15 \end{pmatrix}$$

Also have $x \geq 0$ since we can't have a net loss of product.

We can now state it as a maximization LP:

$$\max_{\mathbf{x}} -\mathbf{c}^T \mathbf{x}, \quad \mathbf{A}\mathbf{x} = \mathbf{b}, \quad \mathbf{x} \geq \mathbf{0}$$

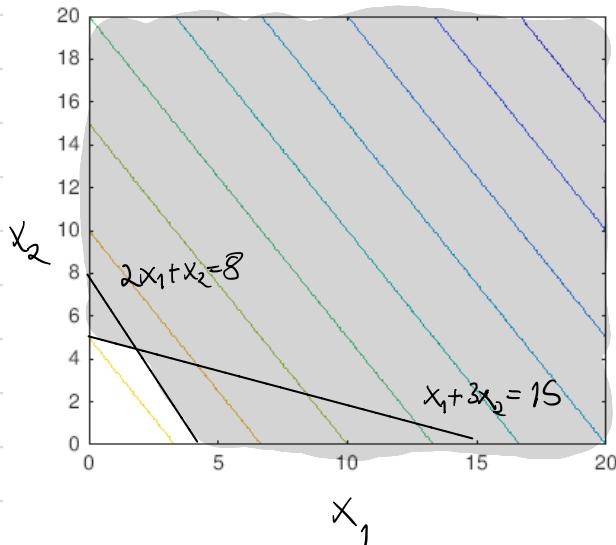
Flipping it a standard form LP we get

$$\min_{\mathbf{x}} \mathbf{c}^T \mathbf{x} \text{ s.t. } \mathbf{A}\mathbf{x} = \mathbf{b}, \quad \mathbf{x} \geq \mathbf{0}$$

where

$$\mathbf{c} = \begin{pmatrix} -3 \\ 2 \\ -1 \end{pmatrix}, \quad \mathbf{A} = \begin{pmatrix} 2 & 1 \\ 1 & 3 \end{pmatrix}, \quad \mathbf{b} = \begin{pmatrix} 8 \\ 15 \end{pmatrix}$$

b)



c) Running the simplex algorithm gives solution

$$A: x_1 = 1.8 \quad (\cdot 1000 \text{ kg})$$

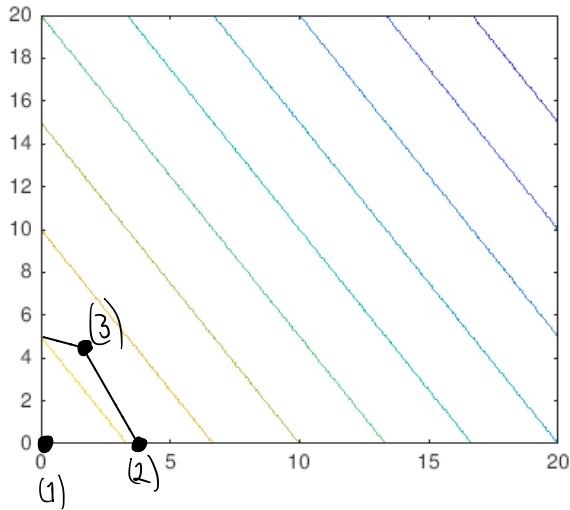
$$B: x_2 = 4.4 \quad (\cdot 1000 \text{ kg})$$

The constraint $x \geq 0$ is not active.

The solution is at a point where both the slack variables are zero, so it is at the intersection of the constraints

$$R_I = 8, R_{II} = 15$$

d)



e)

Problem 3

$$\min_x \quad q(x) = \frac{1}{2} x^T G x + x^T c$$

$$\text{s.t.} \quad a_i^T x = b_i \quad i \in \mathcal{E}$$

$$a_i^T x \geq b_i \quad i \in \mathcal{I}$$

a)

$$A(x^*) = \mathcal{E} \cup \{i \mid i \in \mathcal{I}, a_i^T x^* = b_i\}$$

$$b) \quad \mathcal{L}(x, \lambda) = q(x) - \sum \lambda_i (a_i^T x - b_i)$$

$$\Rightarrow \nabla_x \mathcal{L}(x, \lambda) = Gx + c - \sum \lambda_i a_i$$

For $i \notin A(x^*)$ we know that $\lambda_i = 0$, thus

$$\lambda_i c_i(x^*) = \lambda_i (a_i^T x^* - b_i) = 0 \quad \text{for all } i \in \mathcal{E} \cup \mathcal{I}$$

This allows us to simplify the $\nabla_x \mathcal{L}$ at x^*

$$\nabla_x \mathcal{L}(x^*, \lambda^*) = Gx^* + c - \sum_{i \in A(x^*)} \lambda_i^* a_i$$

The inactive inequality constraints must satisfy

$$a_i^T x^* > b_i$$

The active constraints must satisfy

$$a_i^T x^* = b_i$$

If the solution is optimal the Lagrangian multipliers must be non-negative for inequality constraints.

$$\lambda_i^* \geq 0, i \in \mathbb{Z}$$

If the solution is optimal, then $\nabla_x \mathcal{L}(x^*, \lambda^*) = 0$, so

$$G x^* + c = \sum_{i \in A(x^*)} \lambda_i^* a_i$$