

**Objectives: algorithm analysis; traversing mazes****Up next: MP7: due Monday;** all MPs final grading next Wednesday

1. **Work with a neighbor:** How many seconds will each problem take? It takes 5s to ask each person a question and get a response, 10s for two people to swap positions. All participants are sitting in one row.

b. Find the smallest amount of change carried by one person.

Variables used: ?

Dataset size 10 20 40

Best case :

Worst case :

c. Partition (divide) the group of people into two subgroups: "more-obnoxious-than Chapman" group and less-obnoxious-than-Chapman group. Variables used: ?

Dataset size 10 20 40

Best case :

Worst case :

2. Write an expression for the **worst-case** running time of each algorithm.  $t(N) = \dots$ . Define any constants you need.

```
public static boolean fool(int N) {      t(N)= ?
    int i = N * 2;
    return (i*i + N*N )>1000;
}
```

```
public static void foo2(int N) {
    int i= 2000;

    while( i <= N ){

        i += 2;

    }
}
```

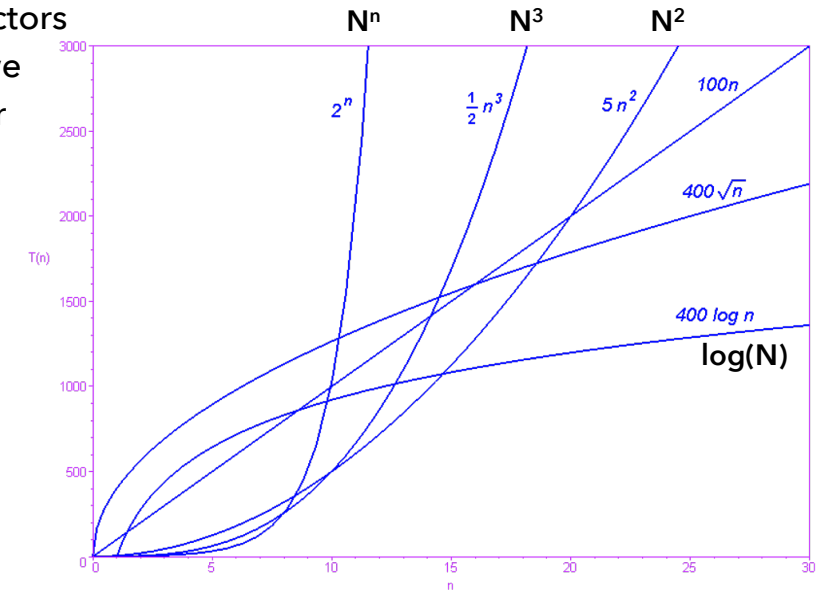
Each take 1 time unit:

arithmetic operations  
assignment ( = )  
boolean comparison  
function activation / return  
array element assignment/  
access  
variable declaration

**3. Algorithm analysis - big idea:****Evaluation and comparison of algorithms.**

**Method:** Characterize the behavior of an algorithm by suppressing constant factors and smaller order terms/constants - we only care when N is large (lower order terms become increasing irrelevant)

Is this all we should ever care about?

**4. What's in a name?**

time complexity / analysis

algorithmic complexity / analysis

asymptotic complexity / analysis

**5. Why is Big-O the way to go?**

- 1) Abstract enough to ignore language and compiler details
- 2) Specific enough to allow differentiation and comparison between different algorithms and implementations, especially with large data inputs

Applications: sorting, searching, multiplying two integers, etc.

**6. What is the algorithmic complexity for the iterative linear search?****7. What is the algorithmic complexity for the recursive linear search?**

8. Write an expression for the **worst-case** running time of this algorithm.  
 $t(N) = \dots$  Define any constants you need.

```
public static int maxFunc(int[][] m) {
// assume a matrix with all dimensions N
int max = 0;
for (int i = 0; i < m.length; i++)
  for (int j = 0; j < m[i].length; j++)
    if (max < m[i][j]) max = m[i][j];
}
return max;
```

What if  $i++$  and  $j++$  above were  $i+=2$  and  $j+=2$ ?

9. Write an expression for the **worst-case** running time of each algorithm.  
 $t(N) = \dots$  Define any constants you need.

```
public static int foo3(int N, int[] data) {
// assume N < data.length-1
return data[N] * 5;
}

public static int foo4(int[] data) {
int best = 0;

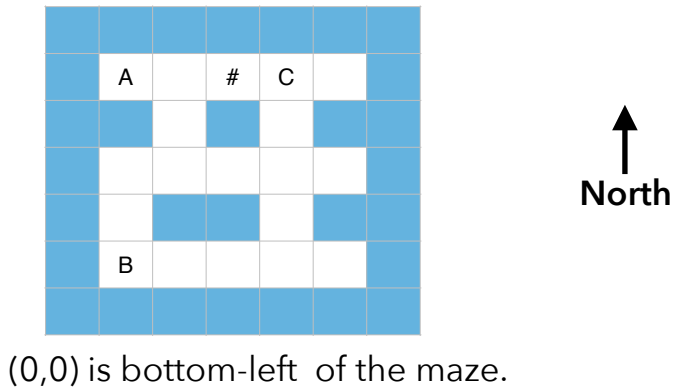
// N is data.length
for(int i=1; i<data.length; i++) {

    if(data[best] > data[i])
        best = i;
}
return best;
}
```

Each take 1 time unit:

arithmetic operations  
assignment ( = )  
boolean comparison  
function activation / return  
array element assignment/  
access  
variable declaration

10. Searching Mazes: recursive algorithm  
From starting points A,B,C, which (x,y) positions will be checked by the code below?



```
Pseudo-code
String explore(x,y,flagX,flagY,wall,...) {
  if(wall[x][y]) return null; // No path here
  if(x == flagX && y == flagY) return ""; // Found

  String goNorth = explore(    ,    , ...) ??
  String goEast  = explore(    ,    , ...) ??
  String goSouth = explore(    ,    , ...) ??
  String goWest  = explore(    ,    , ...) ??
  // There are more elegant implementations

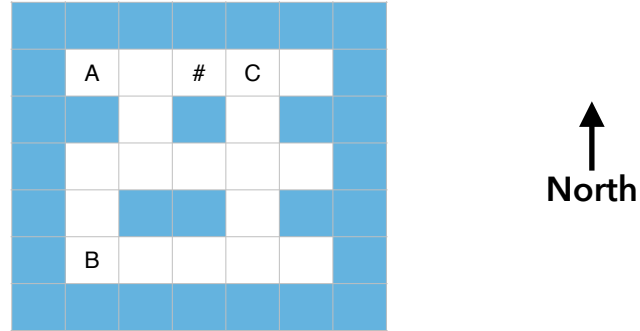
  if(goNorth != null) goNorth = "N" + goNorth;
  if(goEast  != null) goEast  = "E" + goEast ;
  if(goSouth != null) goSouth = "S" + goSouth;
  if(goWest  != null) goWest  = "W" + goWest ;

  if(goNorth == null && goEast  == null
    && goSouth == null && goWest==null ) return null;
  // No path 😞

  String shortestViablePath = ... pick shortest non-null path
  return shortestViablePath;
}
```

## 6. Searching Mazes: recursive algorithm

From starting points A,B,C, which (x,y) positions will be checked by the code below?



(0,0) is bottom-left of the maze.

*Pseudo-code*

```
String explore(x,y,flagX,flagY,wall,...) {
    if(wall[x][y]) return null; // No path here
    if(x == flagX && y == flagY) return ""; // Found
```

```
    String goNorth = explore(    ,    , ...) ??
    String goEast  = explore(    ,    , ...) ??
    String goSouth = explore(    ,    , ...) ??
    String goWest  = explore(    ,    , ...) ??
```

```
// There are more elegant implementations
```

```
    if(goNorth != null) goNorth = "N" + goNorth;
    if(goEast  != null) goEast  = "E" + goEast ;
    if(goSouth != null) goSouth = "S" + goSouth;
    if(goWest  != null) goWest  = "W" + goWest ;
```

```
if(goNorth == null && goEast  == null
   && goSouth == null && goWest==null ) return null;
```

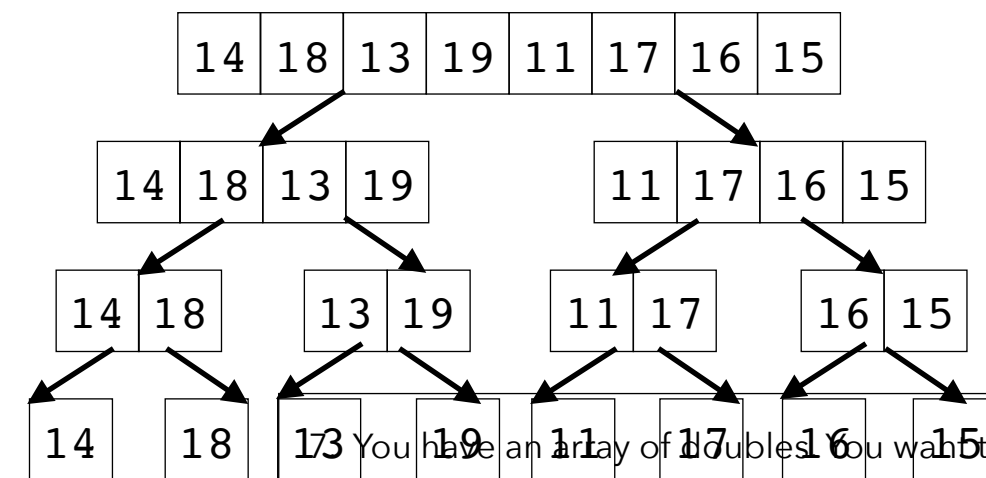
```
// No path 😞
```

```
String shortestViablePath = ... pick shortest non-null path
return shortestViablePath;
```

```
}
```

1. Warm up: Think about the iterative algorithm for linear search.  
Now, write a recursive pseudo-code method for doing linear search:

2. How many levels does the merge sort activation diagram have (roughly)? Why?



You have an array of doubles you want to search between in  
'hi'. Write a recursive method to find the largest product of two  
values. e.g. findPair({ 1.0 , 1.0 , 7.5 , 4.0 , 4.1 ,  
of the array (7.5, 1.420), which has 2 problems

3a. At each level of the tree (j=0, 1, 2, ...), how many subproblems  
are there (as a function of N)?

```
public static double findPair(double[] array, int i, int j) {
```

3b. At each level of the tree (j=0, 1, 2, ...), how many values are in  
the array passed into the recursive activation (as a function of N)?

Write a FORWARD recursive method to find the first index of the  
product of two neighboring values. e.g. findPair({ 1.0 ,  
4.0 , 4.1 , 3.5 }, 0, 5) returns 2 because 7.5x4.0=30.0 is the  
of two neighboring values.

4. Write an expression for the **worst-case** running time of each algorithm.  $t(N) = \dots$  Define any constants you need.

```
public static int foo3(int N, int[] data) {
    // assume N < data.length-1

    return data[N] * 5;
}
```

```
public static int foo4(int[] data) {
    int best = 0;

    // N is data.length
    for(int i=1; i<data.length; i++) {

        if(data[best] > data[i])
            best = i;
    }
    return best;
}
```

#### 5. QuickSort introduction

12	14	11	16	18	17	13	15
----	----	----	----	----	----	----	----

```
static void quickSort(int[] data, int lo, int hi) {
    if (hi > lo) {

        int pivot = ?

        int newPivotIndex = ?

        quickSort(data, lo, newPivotIndex - 1);
        quickSort(data, newPivotIndex + 1, hi);
    }
}
```

#### 6. QuickSort, partitioning...

```
static int partition(int[] data, int lo, int hi, int pivotIndex)
{
    // Move the pivot out of the way; for now we'll put
    // it at the start of the list and ignore it until the end.

    // Start working in, from both L and R ends of the list

    // The pivot will need to go to the left of the final
    // boundary if the last value is larger than the
    // pivot value.

}
```

#### 5. QuickSort summary:

12	14	11	16	18	17	13	15
----	----	----	----	----	----	----	----

How does quick sort differ from merge sort? better? worse?

2. **Merge Sort:**

12	14	11	16	18	15	13	17
----	----	----	----	----	----	----	----

```
static void mergeSort(int[] data, int lo, int hi) {
    if (lo >= hi) return;
    int mid = (lo + hi) / 2;

    mergeSort(data, lo, mid);
    mergeSort(data, mid+1, hi);

    int size = hi - lo + 1;
    int[] temp = new int[size];

    merge(data, temp, lo, mid+1, hi);

    for (int i = 0; i < size; i++) data[i+lo] = temp[i];
}

public static void merge(int[] a, int []tempArray,
                        int lower, int mid, int upper){
    int tempIndex=0;
    int leftLo = lower;
    int leftHi = mid-1;
    int rightLo = mid;
    int rightHi = upper;

    while (leftLo <= leftHi && rightLo <= rightHi) {
        if (a[leftLo] < a[rightLo]) temp[tempIndex++] = a[leftLo++];
        else temp[tempIndex++] = a[rightLo++];
    }

    while (leftLo <= leftHi) temp[tempIndex++] = a[leftLo++];
    while (rightLo <= rightHi) temp[tempIndex++] = a[rightLo++];

    for (int i = 0; i < tempIndex; i++) a[lower+i] = temp[i];
}
```