

# Mathematical Logic\_Assignment\_2

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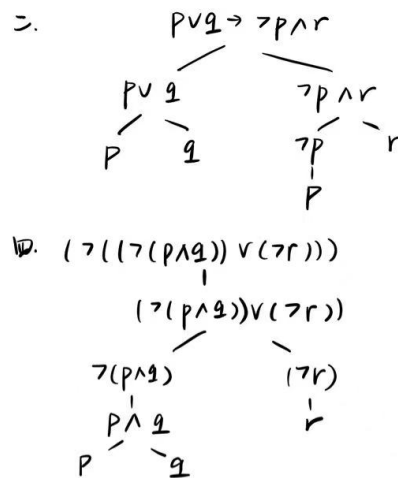
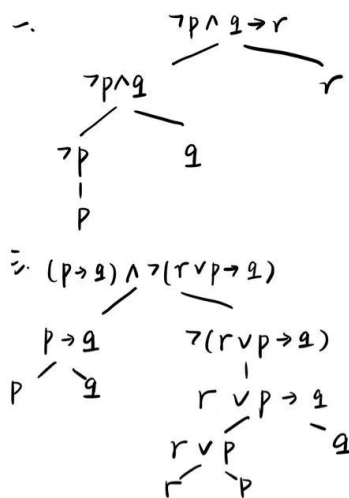
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# 1 Which of the following are well formed propositional formulas?

- It is not a well formed propositional formula.
- It is a well formed propositional formula.
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## 2 Parse Tree



## 3 Propositional Formalization

- $(D \rightarrow (B \wedge C))$
- $((A \wedge B) \rightarrow (\neg D))$
- $(D \leftrightarrow (\neg(C \wedge A)))$
- $((D \wedge (\neg C)) \rightarrow A)$
- $((\neg D) \rightarrow C) \vee (D \rightarrow (\neg B))$

- $((\neg(B \wedge C)) \wedge D) \rightarrow A$
- $((((A \wedge B) \wedge C) \leftrightarrow (\neg D)) \vee (((\neg A) \wedge (\neg B)) \rightarrow ((\neg C) \rightarrow (\neg D))))$

## 4 Properties of formulas I

① **Base Case:** For the simplest propositional language, which contains only one item.  $m = 1$  and  $n = 0$ . Thus,  $m = n + 1$ . Satisfied.

② **Inductive Step:** Suppose that a propositional language A contains k atoms and satisfies  $m = n + 1$ .

Let us consider a propositional language that contains  $(k + 1)$  atoms. we can write it as  $B = A \star P$  ( $\star$  includes  $\wedge \vee \rightarrow \leftrightarrow$ ) and P stands for a new atom. The total number of logical connectives in B is obviously  $(k - 1 + 1)$ .

Therefore, for B,  $(m_B = k + 1)$  and  $(n_B = n_A + 1 = k - 1 + 1 = k)$ .

Therefore,  $(m_B = n_B + 1)$ .

$(m = n + 1)$  stays true for all propositional language.

## 5 Properties of formulas II

① If  $\mathbf{UV}$  is a formula in propositional language. Then  $\mathbf{V}$  is the suffix of  $\mathbf{UV}$ , which contains more closing brackets than opening brackets.

② If  $\mathbf{VW}$  is a formula in propositional language. Then  $\mathbf{V}$  is the prefix of  $\mathbf{VW}$ , which contains more opening brackets than closing brackets.

③  $\mathbf{v}$  cannot satisfy the statements above at the same time.

In conclusion,  $\mathbf{UV}$  and  $\mathbf{VW}$  cannot both be formulas in propositional language.