

Technical Summary For: “A Formalization of Membrane Systems with Dynamically Evolving Structures”

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1 Formal Framework

1.1 Configurations

1. $C = \{(id_1, w_1), \dots, (id_j, w_j), \dots, (id_m, w_m)\} \subseteq (\mathbb{N} \times O^\circ)^*$ - *basic configuration*

$$C = \begin{array}{c} \bigcirc \quad \dots \quad \bigcirc \quad \dots \quad \bigcirc \\ w_1 \quad \quad \quad w_j \quad \quad \quad w_m \\ id_1 \quad \quad \quad id_j \quad \quad \quad id_m \end{array}$$

- $id_j \in \mathbb{N}$ - *cell id*
- $w_j \in O^\circ$ - *cell content*
- (id_j, w_j) - *cell*

2. $\mathbb{C} = \{C \mid size(C) > 0\}$

- $size(C)$ - size of basic configuration C
- \mathbb{C} - set of all basic configurations with size greater than 0.

3. $\mathcal{C} = (L, \rho) = (L = \{(id_1, l_1, w_1), \dots, (id_j, l_j, w_j), \dots, (id_m, l_m, w_m)\}, \rho)$

$$L = \left\{ \begin{array}{c} \bigcirc \quad \dots \quad \bigcirc \quad \dots \quad \bigcirc \\ w_1 \quad \quad \quad w_j \quad \quad \quad w_m \\ id_1 \quad \quad \quad id_j \quad \quad \quad id_m \end{array} \right\}$$

$$\rho \ni \begin{array}{c} \bigcirc \quad \longrightarrow \quad \bigcirc \\ id_k \quad \quad \quad id_{k'} \end{array}$$

- $id_j \in \mathbb{N}$ - *id*
- $w_j \in O^\circ$ - *multiset over O*
- $l_j \in Lab$ - *label*
- (id_j, l_j, w_j) - *labelled cell*
- $L \in (\mathbb{N} \times Lab \times O^\circ)^*$ - *list of labelled cells*
- $\rho \subseteq \mathbb{N} \times \mathbb{N}$ - *relations between cells (ids)*
- $\mathcal{C} = (L, \rho)$ - *configuration*
- $\mathcal{C}_L = L$ and $\mathcal{C}_\rho = \rho$.

- $\bar{\mathcal{C}}_L \in \mathbb{C}$ - *projection* of \mathcal{C}_L as basic configuration.

$$\bar{\mathcal{C}}_L = \left\{ \begin{array}{c} \textcircled{w_1} \cdots \textcircled{w_j} \cdots \textcircled{w_m} \\ id_1 \quad \quad id_j \quad \quad id_m \end{array} \right\}$$

$$4. \mathfrak{C} = \{\mathcal{C} \mid \mathcal{C} = (L, \rho)\}$$

- \mathfrak{C} is the set of all possible configurations.

1.2 Components of a Rule

$$0. r = (\text{Labels}, \rho, \text{Perm}, \text{For}, \text{Rewrite}, \text{Label-Rename}, \text{Delete}, \text{Delete-and-Move}, \text{Generate}, \text{Generate-and-Copy}, \text{Change-Relation})$$

- r is a rule.

$$1. \text{Labels}(r) = (l_1, \dots, l_j, \dots, l_k) \in \text{Lab}^k$$

$$\text{Label}(r) = \begin{array}{c} \textcircled{\quad} \cdots \textcircled{\quad} \cdots \textcircled{\quad} \\ 1 \quad \quad j \quad \quad k \end{array}$$

$$2. \rho(r) \subseteq \mathbb{N}_k \times \mathbb{N}_k$$

- $\mathbb{N}_k = \{1, \dots, j, \dots, k\}$

$$3. \text{Perm}(r) = \{P_1, \dots, P_{j'}, \dots, P_{\bar{p}}\} \subseteq \mathbb{C}_k$$

$$\begin{aligned} \text{Perm}(r) = & \left\{ \begin{array}{l} P_1 = \left\{ \begin{array}{c} \textcircled{p_{(1,1)}} \cdots \textcircled{p_{(1,j)}} \cdots \textcircled{p_{(1,k)}} \\ 1 \quad \quad j \quad \quad k \end{array} \right\} , \dots \\ P_{j'} = \left\{ \begin{array}{c} \textcircled{p_{(j',1)}} \cdots \textcircled{p_{(j',j)}} \cdots \textcircled{p_{(j',k)}} \\ 1 \quad \quad j \quad \quad k \end{array} \right\} , \dots \\ P_{\bar{p}} = \left\{ \begin{array}{c} \textcircled{p_{(\bar{p},1)}} \cdots \textcircled{p_{(\bar{p},j)}} \cdots \textcircled{p_{(\bar{p},k)}} \\ 1 \quad \quad j \quad \quad k \end{array} \right\} \end{array} \right\} \end{aligned}$$

- $\mathbb{C}_k \subseteq \mathbb{C}$ - *basic configurations* with cell ids in \mathbb{N}_k .

- $p_{(j',j)} \in O^\circ$

$$4. \text{For}(r) = \{F_1, \dots, F_{j'}, \dots, F_{\bar{f}}\} \subseteq \mathbb{C}_k$$

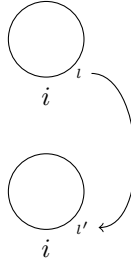
$$\begin{aligned}
For(r) = & \left\{ F_1 = \left\{ \underset{1}{\bigcirc f_{(1,1)}} \cdots \underset{j}{\bigcirc f_{(1,j)}} \cdots \underset{k}{\bigcirc f_{(1,k)}} \right\} , \dots \right. \\
& F_{j'} = \left\{ \underset{1}{\bigcirc f_{(j',1)}} \cdots \underset{j}{\bigcirc f_{(j',j)}} \cdots \underset{k}{\bigcirc f_{(j',k)}} \right\} , \dots \\
& \left. F_{\bar{j}} = \left\{ \underset{1}{\bigcirc f_{(\bar{j},1)}} \cdots \underset{j}{\bigcirc f_{(\bar{j},j)}} \cdots \underset{k}{\bigcirc f_{(\bar{j},k)}} \right\} \right\}
\end{aligned}$$

- $f_{(j',j)} \in O^\circ$

5. $Rewrite(r) = U \rightarrow V$

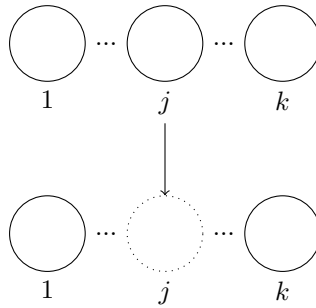
- $U, V \in \mathbb{C}_k$
- $Rewrite(r)$ is a general rewriting rule, rewriting a finite basic configuration U to another finite basic configuration V .

6. $Label-Rename(r) = \{..., (i, l'), ...\} \in (\mathbb{N}_k \times Lab)^*$



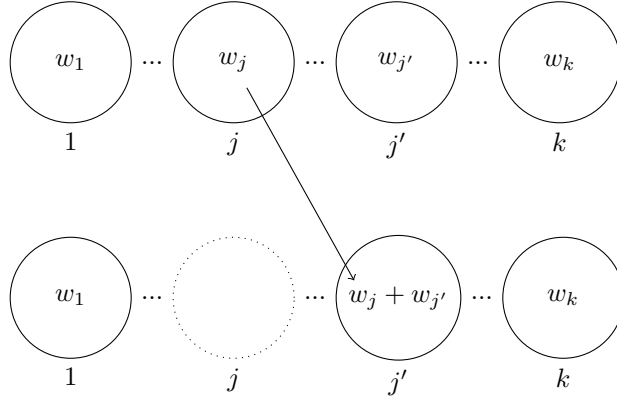
- $i \in \mathbb{N}_k$ - id
- $l' \in Lab$ - new label

7. $Delete(r) = \{..., j, ...\} \in \mathbb{N}_k^*$



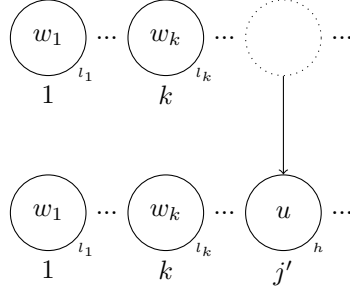
- j - id of cell to be deleted

8. $Delete-Move(r) \in (\mathbb{N}_k \times \mathbb{N}_k)^*$



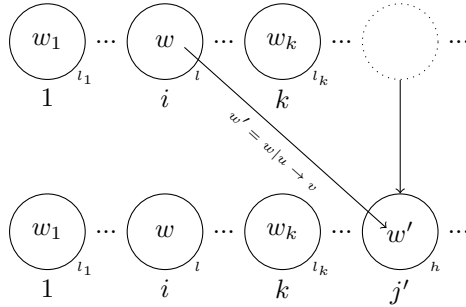
- (j, j') - pair of ids
- j - id of cell to be deleted
- j' - id of cell to receive the multiset
- *Delete-Move*(r)- list of pairs of ids

9. $Generate(r) = \{..., (j', h, u), ...\} \in (\mathbb{N}' \times Lab \times O^\circ)^*$



- j' - primed id - new id
- h - label
- u - multiset

10. $Generate-Copy(r) = \{..., (j', h, i, (u, v)), ...\} \in (\mathbb{N}' \times Lab \times \mathbb{N} \times (O^\circ \times O^\circ))^*$



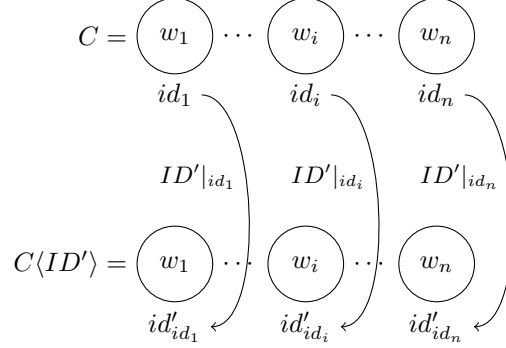
- $j' \in \mathbb{N}'$ - primed id
- $h \in Lab$ - label
- $i \in \mathbb{N}'_k$ - cell id
- $w_1, ..., w, ..., w_k \in O^\circ$ - multisets
- $u, v \in O^\circ$ - multisets
- (u, v) - written as $u \rightarrow v$

11. *Change-Relation*

- *Change-Relation* is a *graph transducer* that updates the relation ρ .
- This transducer should be recursive and it can only add and remove edges.

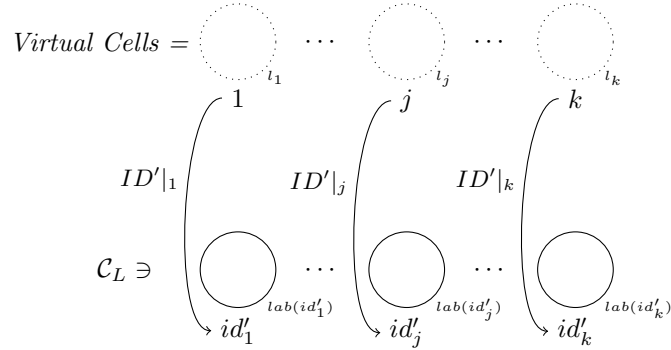
1.3 Eligibility of an Id Vector

1. $ID' = (id'_1, \dots, id'_i, \dots, id'_{n'}) \in \mathbb{N}^{n'}$



- $ID'|_x = id'_x$
- $ID'|_{id_i} = id'_{id_i}$
- $C = \{(id_1, w_n), \dots, (id_i, w_i), \dots, (id_n, w_n)\}$
- $C\langle ID' \rangle = \{(ID'|_{id_1}, w_n), \dots, (ID'|_{id_i}, w_i), \dots, (ID'|_{id_n}, w_n)\}$
- $C\langle ID' \rangle = \{(id'_{id_1}, w_n), \dots, (id'_{id_i}, w_i), \dots, (id'_{id_n}, w_n)\}$
- $|ID'| = n' = \max\{id_i\}$
- $|C| = n$
- $|C| = n \leq n' = |ID'|$

2. $Eligible(ID', r, \mathcal{C})$



- $ID' = (id'_1, \dots, id'_j, \dots, id'_k) \in \mathbb{N}^k$
- r is a k -degree interaction rule
- $r\langle ID' \rangle$ - *instantiated rule* - replace virtual ids with ids in ID'
- $Labels(r) = (l_1, \dots, l_j, \dots, l_k)$
- $\mathcal{C} = (L, \rho)$
- $lab : \mathbb{N} \rightarrow Lab$
- $Eligible(ID', r, \mathcal{C}) \Leftrightarrow :$
 - $\forall x, y \in \{1, \dots, k\} : (x \neq y) \rightarrow (id'_x \neq id'_y)$ [id uniqueness]
 - $\forall x \in \{1, \dots, k\} : l_x = lab(id'_x)$ [label consistency]
 - $\forall x, y \in \{1, \dots, k\} : (x, y) \in \rho(r) \rightarrow (id'_x, id'_y) \in \mathcal{C}_\rho$ [relations consistency]
- $\mathcal{I}_\mathcal{C}(r) = \{ID' \mid Eligible(ID', r, \mathcal{C})\}$

1.4 Applicability of a Multiset of Rules

- $R = \{r_1, \dots, r_i, \dots, r_n\}$ - *multiset* of rules
- $\mathcal{I}_C(r_i) = \{ID'_{i,1}, \dots, ID'_{i,j_i}, \dots, ID'_{i,k_i}\}$
- $RI = \{r_1\langle ID'_{1,j_1} \rangle, \dots, r_i\langle ID'_{i,j_i} \rangle, \dots, r_n\langle ID'_{n,j_n} \rangle\}$ - *multiset of instantiated rules*

$applicable(RI, \mathcal{C}) \Leftrightarrow \forall r_i \in R : \mathcal{I}_C(r_i) \neq \emptyset : \forall r_i\langle ID'_{i,j_i} \rangle \in RI :$

1. $\forall P \in Perm(r_i\langle ID'_{i,j_i} \rangle) \cup \{DP\} : P \subseteq \bar{\mathcal{C}}_L$

- $DP = \{(id_i, u) \mid (j', h, id_i, u \rightarrow v) \in Generate\text{-}and\text{-}Copy(r_i\langle ID'_{i,j_i} \rangle)\}$

$$\begin{aligned}
 Perm(r_i\langle ID'_{i,j_i} \rangle) = & \left\{ \begin{aligned} P_1 = & \left\{ \begin{array}{ccc} \textcircled{p_{(1,1)}} & \cdots & \textcircled{p_{(1,j)}} & \cdots & \textcircled{p_{(1,k)}} \\ ID'_{i,j_i}|_1 & & ID'_{i,j_i}|_j & & ID'_{i,j_i}|_k \end{array} \right\} & , \quad \dots \\ P_{j'} = & \left\{ \begin{array}{ccc} \textcircled{p_{(j',1)}} & \cdots & \textcircled{p_{(j',j)}} & \cdots & \textcircled{p_{(j',k)}} \\ ID'_{i,j_i}|_1 & & ID'_{i,j_i}|_j & & ID'_{i,j_i}|_k \end{array} \right\} & , \quad \dots \\ P_{\bar{p}} = & \left\{ \begin{array}{ccc} \textcircled{p_{(\bar{p},1)}} & \cdots & \textcircled{p_{(\bar{p},j)}} & \cdots & \textcircled{p_{(\bar{p},k)}} \\ ID'_{i,j_i}|_1 & & ID'_{i,j_i}|_j & & ID'_{i,j_i}|_k \end{array} \right\} & \left. \vphantom{\begin{array}{c} P_1 \\ P_{j'} \\ P_{\bar{p}} \end{array}} \right\} \\ DP = & \left\{ \begin{array}{c} \cdots \textcircled{u} \cdots \\ id_i \end{array} \right\} \end{aligned} \right.
 \end{aligned}$$

$$\bar{\mathcal{C}}_L = \left\{ \begin{array}{ccc} \textcircled{w_1} & \cdots & \textcircled{w_j} & \cdots & \textcircled{w_n} \\ id_1 & & id_j & & id_n \end{array} \right\}$$

2. $\forall F \in For(r_i\langle ID'_{i,j_i} \rangle) : F \subseteq \bar{\mathcal{C}}_L$

$$\begin{aligned}
 For(r_i\langle ID'_{i,j_i} \rangle) = & \left\{ \begin{aligned} F_1 = & \left\{ \begin{array}{ccc} \textcircled{f_{(1,1)}} & \cdots & \textcircled{f_{(1,j)}} & \cdots & \textcircled{f_{(1,k)}} \\ ID'_{i,j_i}|_1 & & ID'_{i,j_i}|_j & & ID'_{i,j_i}|_k \end{array} \right\} & , \quad \dots \\ F_{j'} = & \left\{ \begin{array}{ccc} \textcircled{f_{(j',1)}} & \cdots & \textcircled{f_{(j',j)}} & \cdots & \textcircled{f_{(j',k)}} \\ ID'_{i,j_i}|_1 & & ID'_{i,j_i}|_j & & ID'_{i,j_i}|_k \end{array} \right\} & , \quad \dots \\ F_{\bar{f}} = & \left\{ \begin{array}{ccc} \textcircled{f_{(\bar{f},1)}} & \cdots & \textcircled{f_{(\bar{f},j)}} & \cdots & \textcircled{f_{(\bar{f},k)}} \\ ID'_{i,j_i}|_1 & & ID'_{i,j_i}|_j & & ID'_{i,j_i}|_k \end{array} \right\} & \left. \vphantom{\begin{array}{c} F_1 \\ F_{j'} \\ F_{\bar{f}} \end{array}} \right\} \end{aligned} \right.
 \end{aligned}$$

$$3. (\bigcup_{i=1}^n U_i) \subseteq \overline{\mathcal{C}}_L$$

- $Rewrite(r\langle ID'_{i,j_i} \rangle) = U_i \rightarrow V_i$
- $U_i, V_i \in \mathbb{C}$

$$U_i = \left\{ \dots, \underset{id_j}{\bigcirc u}, \dots \right\}$$

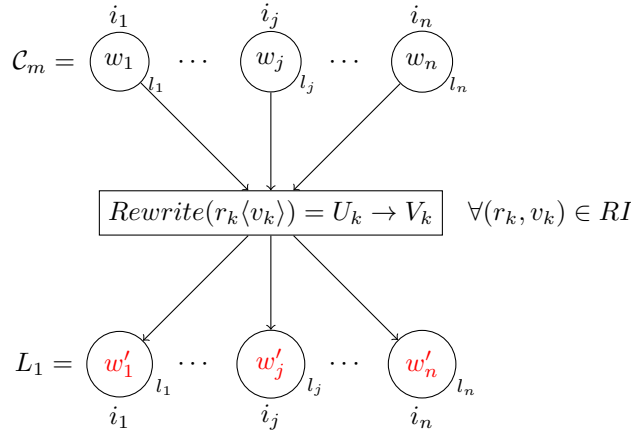
$$4. \forall id_s, i, k : [(id_s, l) \in Label-Rename(r_i\langle ID'_{i,j_i} \rangle) \wedge (id_s, l') \in Label-Rename(r_k\langle ID'_{k,j_k} \rangle)] \rightarrow l = l'$$

$$5. \forall i, k : Change-Relation(r_k)(Change-Relation(r_i)(\mathcal{C}_\rho)) = Change-Relation(r_i)(Change-Relation(r_k)(\mathcal{C}_\rho))$$

- $Applicable(R, \mathcal{C}) = \{RI \mid applicable(RI, \mathcal{C})\}$
- $Applicable(\Pi, \mathcal{C}) = \bigcup_{Applicable(R, \mathcal{C}) \neq \emptyset} Applicable(R, \mathcal{C})$

1.5 Applying a Multiset of Rules

$$1. L_1 = \{(i_1, l_1, w'_1) \dots (i_n, l_n, w'_n)\}$$

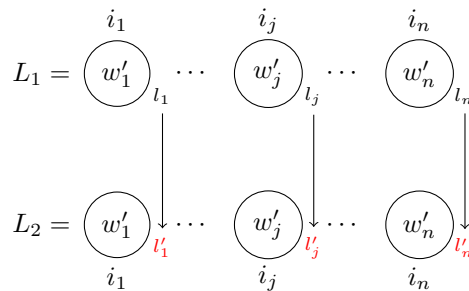


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$$w'_j = w_j - \left(\bigcup_{(r_k, v_k) \in RI} U_{k|j} \right) + \left(\bigcup_{(r_k, v_k) \in RI} V_{k|j} \right)$$

- $RI = \{(r_1, v_1) \dots (r_n, v_n)\}$
- $\mathcal{C} = (\{(i_1, l_1, w_1) \dots (i_n, l_n, w_n)\}, \rho)$
- $Rewrite(r_k\langle v_k \rangle) = U_k \rightarrow V_k$
- $U_k, V_k \in \mathbb{C}_k$

$$2. L_2 = \{(i_1, l'_1, w'_1) \dots (i_n, l'_n, w'_n)\}$$



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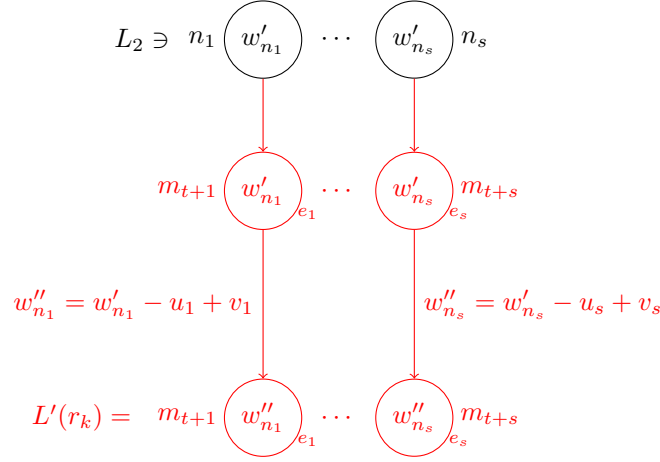
$$l'_j = \begin{cases} e_s, & \text{if } \exists(r_k, v_k) \in RI \text{ such that } (j, e_s) \in \text{Label-Rename}(r_k \langle v_k \rangle) \\ l_j, & \text{otherwise} \end{cases}$$

3. $L_c = L_c(r_1) \cdot \dots \cdot L_c(r_n)$

$$L_c(r_k) = \underbrace{(u_1)_{h_1}}_{m_1} \cdots \underbrace{(u_t)_{h_t}}_{m_t}$$

- $L_c(r_k) = \{(m_1, h_1, u_1) \cdots (m_t, h_t, u_t)\}$
- $\text{Generate}(r_k) = \{(1', h_1, u_1) \cdots (t', h_t, u_t)\}$

4. $L'_c = L'_c(r_1) \cdot \dots \cdot L'_c(r_n)$



- $L'_c(r_k) = \{(m_{t+1}, e_1, w'_{n_1} - u_1 + v_1) \dots (m_{t+s}, e_s, w'_{n_s} - u_s + v_s)\}$
- $\text{Generate-and-Copy}(r_k) = \{(1', e_1, n_1, u_1 \rightarrow v_1) \dots (s', e_s, n_s, u_s \rightarrow v_s)\}$
- $(i_j, l'_j, w'_j) \in L_2$

5. $L_3 = L_2 \cdot L_c \cdot L'_c$

6. $L_4 = \{(i_1, l'_1, w'_1) \cdots (i_n, l'_n, w'_n)\}$

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$$w''_j = w'_j + \left(\bigcup_{\text{last}(i_k)=i_j} w'_k \right)$$

- $(i_j, l'_j, w'_j) \in L_3$
- $\mathbf{p} = (p_1, \dots, p_j, \dots, p_n)$ - vector of “destination” cell ids for multisets from cells that may be deleted.
- $p_j = p(i_j)$ - destination cell for multiset from cell i_j .

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$$p(i_j) = \begin{cases} *, & \text{if } \exists(r_k, v_k) \in RI \text{ such that } i_j \in \text{Delete}(r_k \langle v_k \rangle) \\ e, & \text{if } \exists(r_k, v_k) \in RI \text{ such that } (i_j, e) \in \text{Delete-and-Move}(r_k \langle v_k \rangle) \\ i_j, & \text{otherwise} \end{cases}$$

- $\text{Delete}(r) \in \mathbb{N}_k^*$ is the list of cell ids to be deleted.
- $\text{Delete-and-Move}(r) \in (\mathbb{N}_k \times \mathbb{N}_k)^*$. $(i, j) \in \text{Delete-and-Move}(r)$ means to delete cell i and move its multiset to cell j .
- If $p_k = p(i_k) = i_k$, then cell i_k will not be deleted.
- If $p_k = p(i_k) \neq i_k$ (either $p(i_k) = *$ or $p(i_k) = e$), then cell i_k will be deleted.
- If $p_k = i_k$, then there is a sequence $x_1, \dots, x_{j-1}, \dots, x_m$ of cell ids where:

- $x_1 = p_k$
- $x_j = p(x_{j-1})$ for $2 \leq j \leq m$
- $x_m = z$

There is a chain of “delete-and-move” and z is the last cell id to not be deleted or to be deleted and not moved.

- $last(i_j) = z$
7. $L_5 = \{(i_1, l'_1, w''_1) \dots (i_{n_1}, l'_{n_1}, w''_{n_1})\}$ where $(i_j, l'_j, w''_j) \in L_4$ and $p_j = i_j$.
- L_5 contains L_4 cells with the deleted cells removed.
8. $Apply(RI, \mathcal{C}) = (L_2, \mathcal{C}'_\rho)$
- \mathcal{C}'_ρ is the updated ‘parent’ relation after *CREATE-NODES*, *DELETE-NODES*, and *Change-Relation*($r_k \langle v_k \rangle$) have been computed for all $(r_k, v_k) \in RI$ on \mathcal{C}_ρ .