# Efficient Black-box Checking of Snapshot Isolation in Databases

(Conference VLDB'2024)

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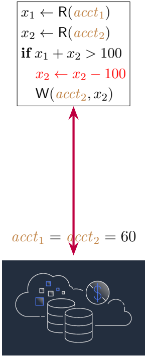
# Transaction and Isolation Level

A transaction is a *group* of operations that is executed atomically.



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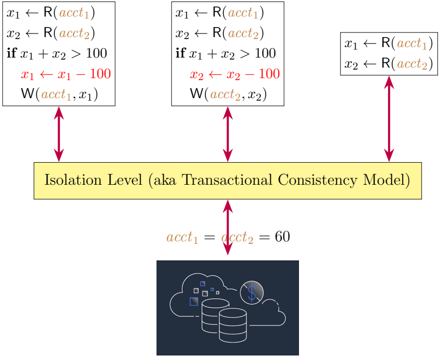

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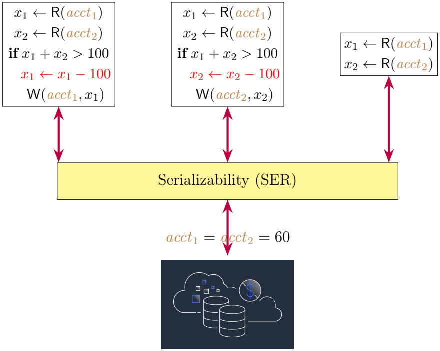
Isolation Level (aka Transactional Consistency Model)

$acct_1 = acct_2 = 60$

The isolation levels specify how they are isolated from each other.

# Serializability (SER)

All transactions appear to execute in some total order.



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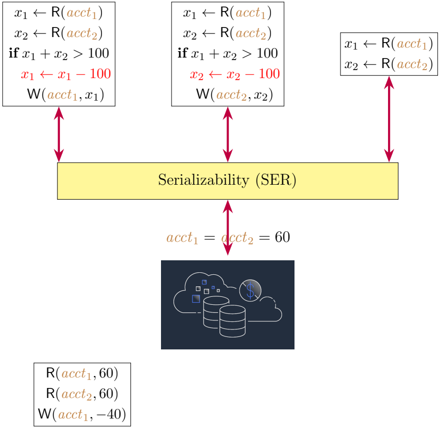

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Serializability (SER)

$acct_1 = acct_2 = 60$

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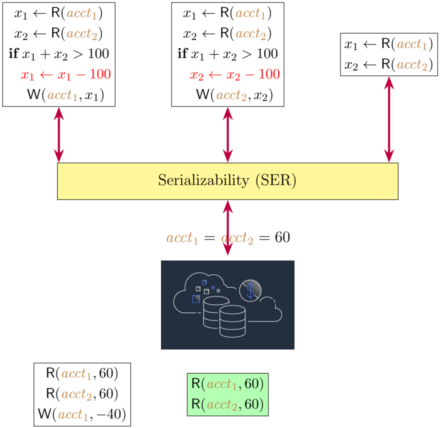

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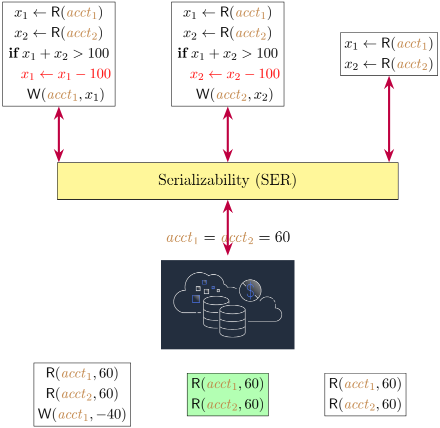

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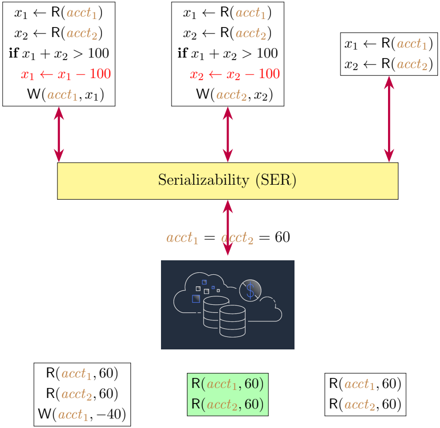

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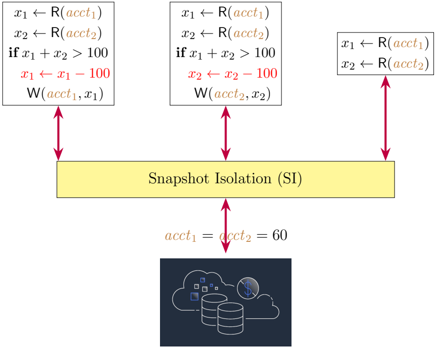
too expensive, especially for distributed transactions

# Snapshot Isolation (SI)



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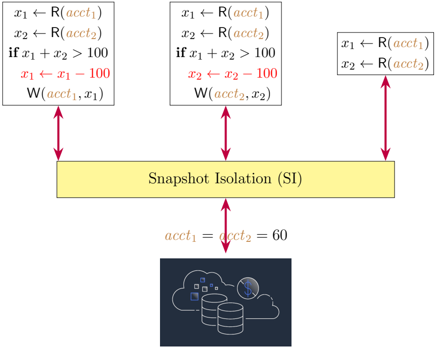

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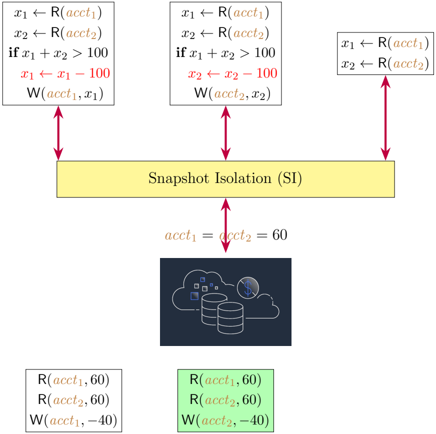

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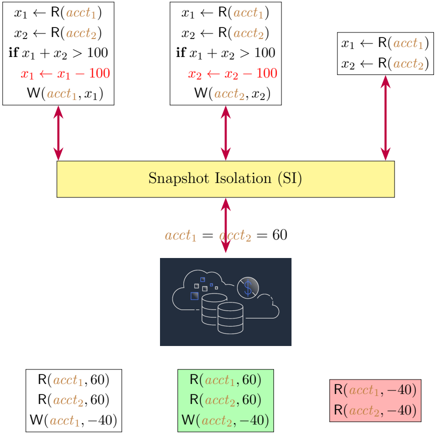

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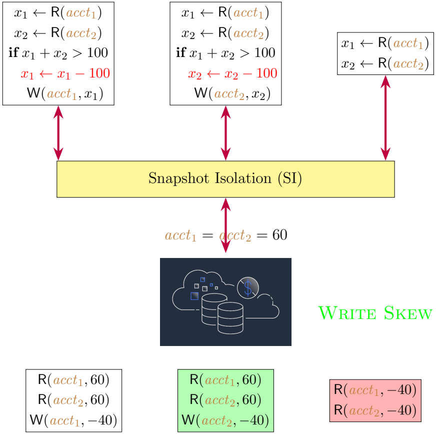

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Write Skew

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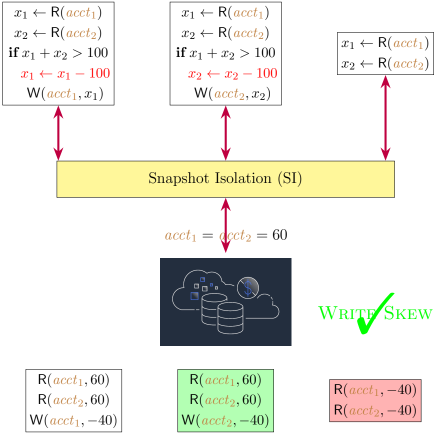

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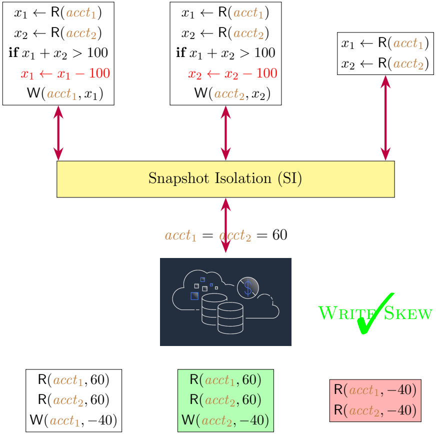

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Snapshot Isolation (SI)

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WRITE SKEW ✓

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$\mathsf{R}(acct_1, -40)$
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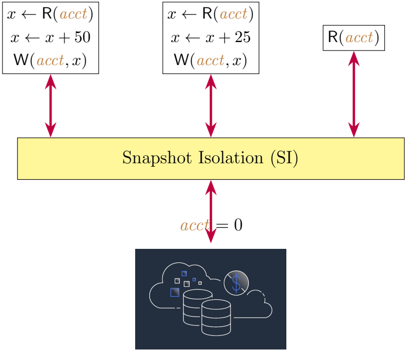
Snapshot Read: Each transaction reads data from a *snapshot* of committed data valid as of the (logical) time the transaction started.

# Snapshot Isolation (SI)



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$\mathsf{W}(acct, x)$

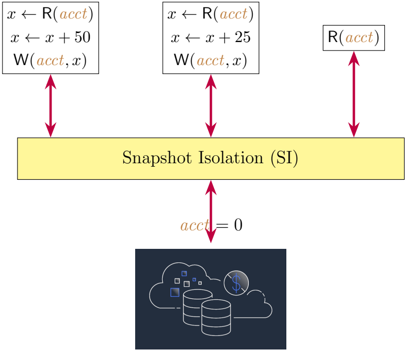
$x \leftarrow \mathsf{R}(acct)$
$x \leftarrow x + 25$
$\mathsf{W}(acct, x)$

$\mathsf{R}(acct)$

Snapshot Isolation (SI)

$acct = 0$

# Snapshot Isolation (SI)



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$\mathsf{W}(acct, x)$

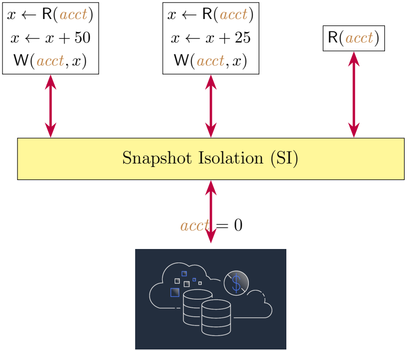
$\mathsf{R}(acct)$

Snapshot Isolation (SI)

$acct = 0$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 50)$

# Snapshot Isolation (SI)



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$\mathsf{R}(acct)$

Snapshot Isolation (SI)

$acct = 0$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 50)$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, ?)$

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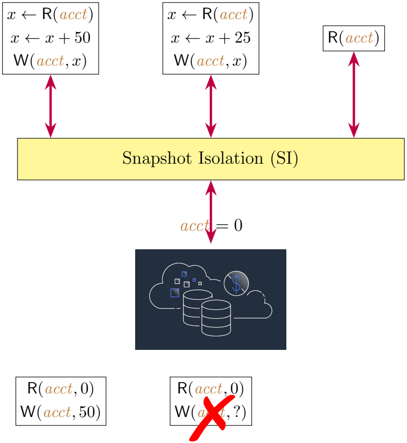

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$x \leftarrow x + 25$
$\mathsf{W}(acct, x)$

$\mathsf{R}(acct)$

Snapshot Isolation (SI)

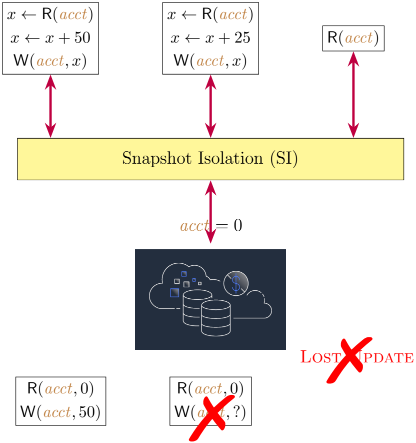
$acct = 0$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 50)$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(\cancel{acct}, ?)$

Snapshot Write: Concurrent transactions cannot write to the same key. One of them must be aborted.

# Snapshot Isolation (SI)



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$\mathsf{R}(acct)$

Snapshot Isolation (SI)

$acct = 0$

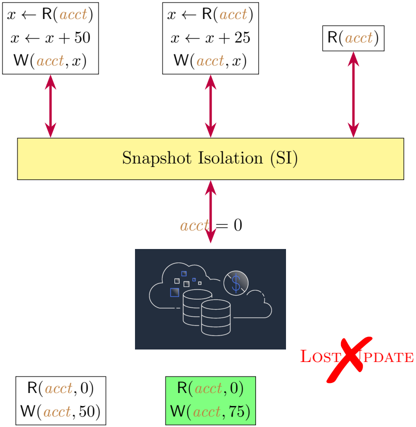
Lost Update

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 50)$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, ?)$

Snapshot Write: Concurrent transactions cannot write to the same key. One of them must be aborted.

# Snapshot Isolation (SI)



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$\mathsf{R}(acct)$

Snapshot Isolation (SI)

$acct = 0$

Lost Update ✗

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 50)$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 75)$

Snapshot Write: Concurrent transactions cannot write to the same key. One of them must be aborted.

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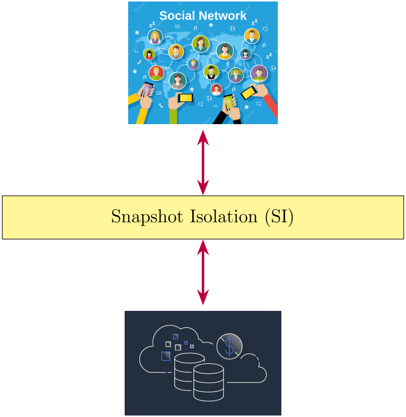

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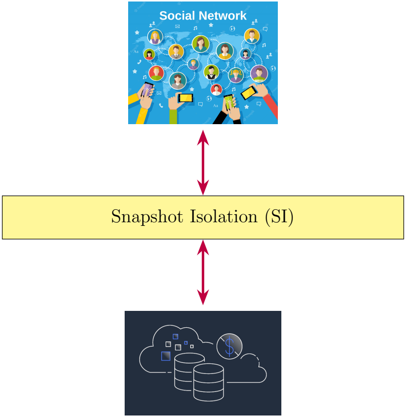

Snapshot Isolation (SI)

$\boxed{\mathsf{W}(x, post)}$

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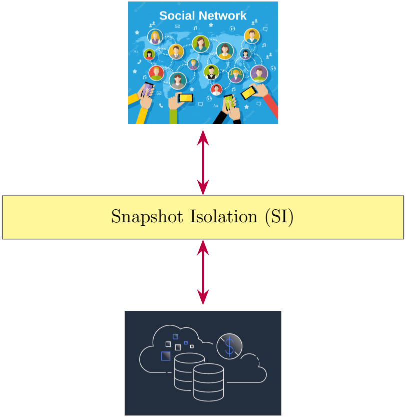
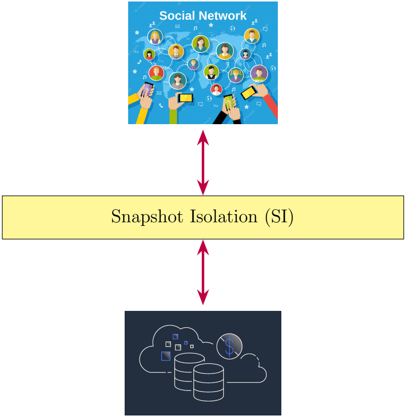


Snapshot Isolation (SI)

$\mathsf{W}(x, post)$

$\mathsf{R}(x, post)$
$\mathsf{W}(y, comment)$

# Snapshot Isolation (SI)



Snapshot Isolation (SI)

| | | |
|---|---|---|
| W($x$, *post*) | R($x$, *post*)<br>W($y$, *comment*) | R($x$, *empty*)<br>R($y$, *comment*) |

# Snapshot Isolation (SI)

Snapshot Isolation (SI)

CAUSALITY VIOLATION

$\mathsf{W}(x, post)$

$\mathsf{R}(x, post)$
$\mathsf{W}(y, comment)$

$\mathsf{R}(x, empty)$
$\mathsf{R}(y, comment)$

# Snapshot Isolation (SI)

Snapshot Isolation (SI)

CAUSALITY VIOLATION ✗

$\mathsf{W}(x, \mathit{post})$

$\mathsf{R}(x, \mathit{post})$
$\mathsf{W}(y, \mathit{comment})$

$\mathsf{R}(x, \mathit{empty})$
$\mathsf{R}(y, \mathit{comment})$ ✗

# Databases and Snapshot Isolation

database logos

Many databases claim to support SI.

# Databases and Snapshot Isolation

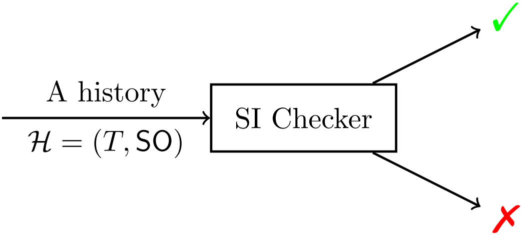
+papers

Databases may fail to provide SI as they claim.

# The SI Checking Problem

Definition (The SI Checking Problem)

The SI checking problem is the decision problem of determing whether a given history $\mathcal{H} = (T, \mathsf{SO})$ satisfies SI?



# The SI Checking Problem

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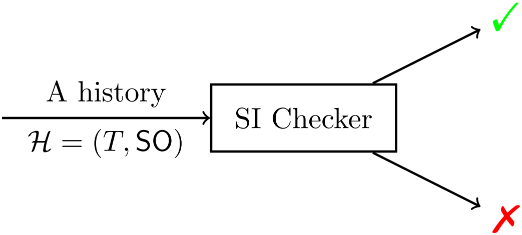
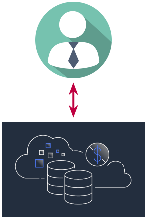


$\mathsf{SO}$ : *session order* among the set $T$ of transactions

# The SI Checking Problem

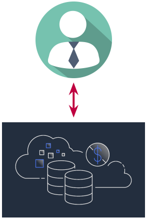
*Black-box checking:* do not rely on database internals



The histories are collected from database logs.

# The SI Checking Problem

*Black-box checking:* do not rely on database internals



W($x$, $post$)

R($x$, $post$)
W($y$, $comment$)

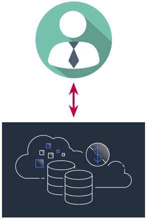
R($x$, $empty$)
R($y$, $comment$)

Lost Update ✗

The histories are collected from database logs.

# The SI Checking Problem

*Black-box checking:* do not rely on database internals



$\mathsf{W}(x, post)$

$\mathsf{R}(x, post)$
$\mathsf{W}(y, comment)$

$\mathsf{R}(x, empty)$
$\mathsf{R}(y, comment)$

LOST UPDATE ✗

$\mathsf{R}(acct_1, 60)$
$\mathsf{R}(acct_2, 60)$
$\mathsf{W}(acct_1, -40)$

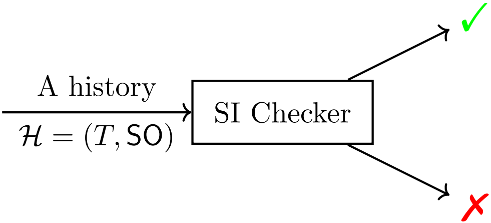
$\mathsf{R}(acct_1, 60)$
$\mathsf{R}(acct_2, 60)$
$\mathsf{W}(acct_2, -40)$

$\mathsf{R}(acct_1, -40)$
$\mathsf{R}(acct_2, -40)$

WRITE SKEW ✓

The histories are collected from database logs.

# The SI Checking Problem



A history

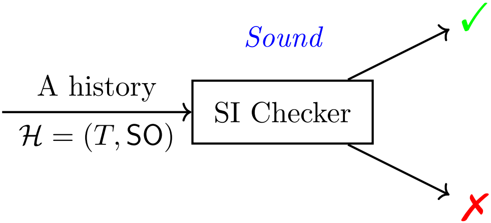
$\mathcal{H} = (T, \mathsf{SO})$

SI Checker

✓

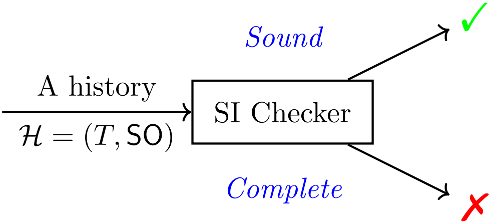
✗

# The SI Checking Problem



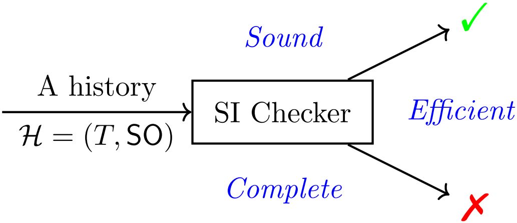
*Sound:* If the checker says ✗, then the history does *not* satisfy SI.

# The SI Checking Problem



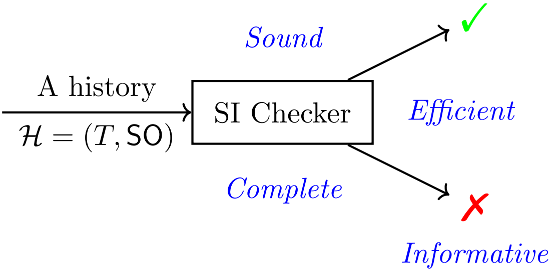
*Complete:* If the checker says ✓, then the history *satisfies* SI.

# The SI Checking Problem



*Efficient:* The checker should *scale* up to large workloads.

# The SI Checking Problem

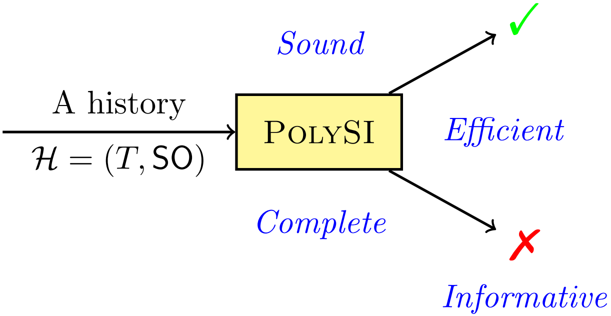


*Informative:* The checker should provide understandable *counterexamples* if it says ✗.

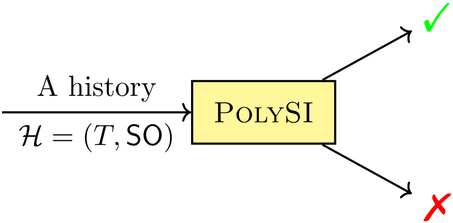
# The SI Checking Problem

related-work

# Contribution: the PolySI Checker



# Contribution: the PolySI Checker



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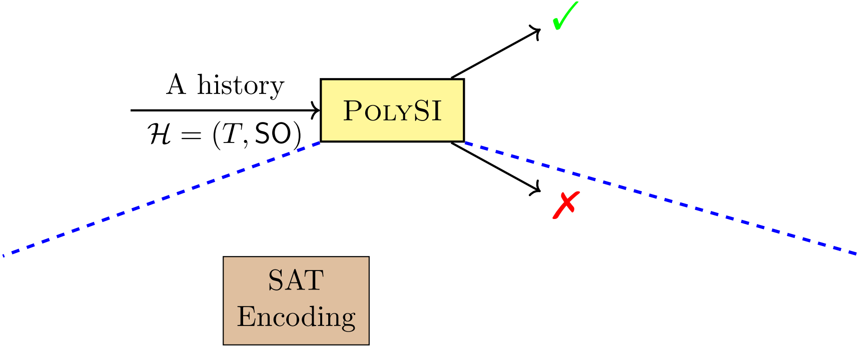
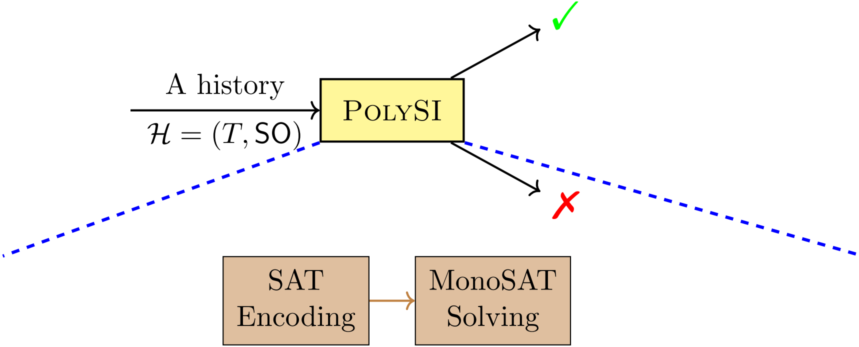


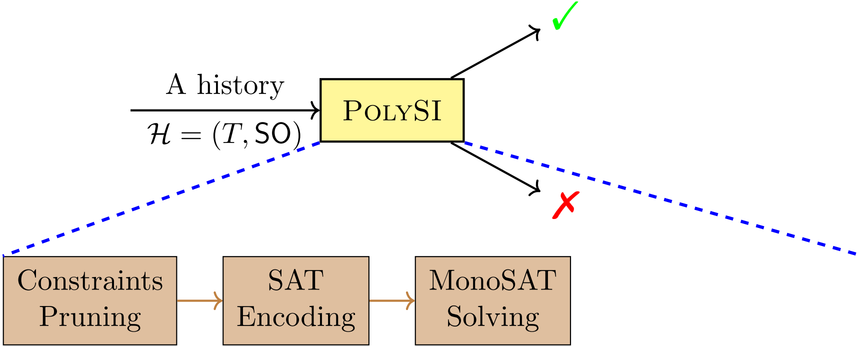
*Sound & Complete:* polygraph-based characterization of SI

# Contribution: the PolySI Checker



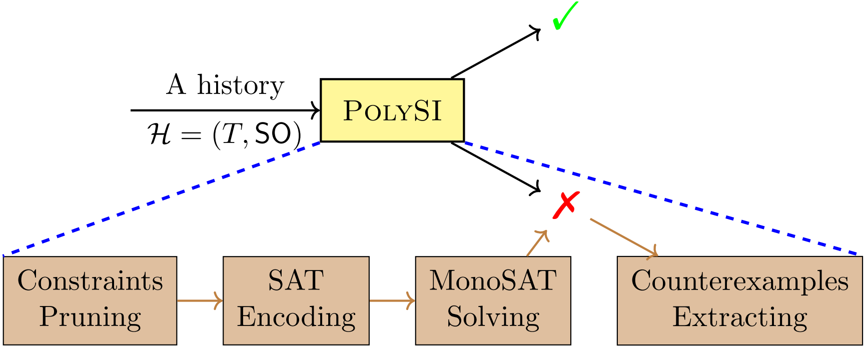
*Efficient:* utilizing MonoSAT solver optimized for graph problems

# Contribution: the PolySI Checker



*Efficient:* domain-specific pruning before encoding

# Contribution: the PolySI Checker



A history

$\mathcal{H} = (T, \mathsf{SO})$

PolySI

✓

✗

Constraints Pruning

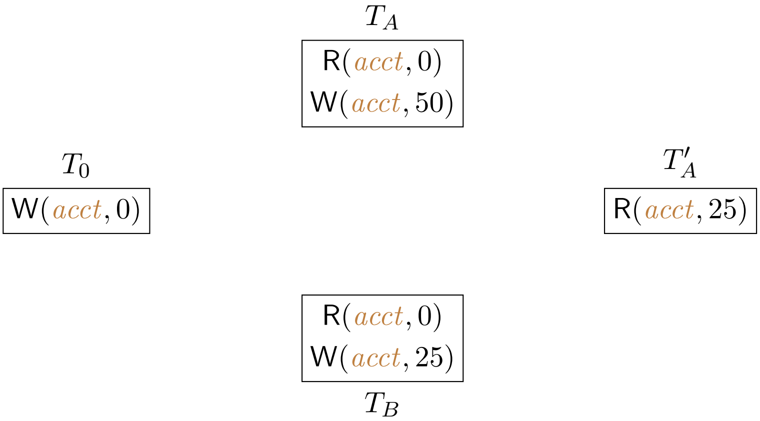
SAT Encoding

MonoSAT Solving

Counterexamples Extracting

*Informative:* extract counterexamples from the unsatisifiable core

# Dependency Graph-based Characterization of SI



$T_A$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 50)$

$T_0$

$\mathsf{W}(acct, 0)$

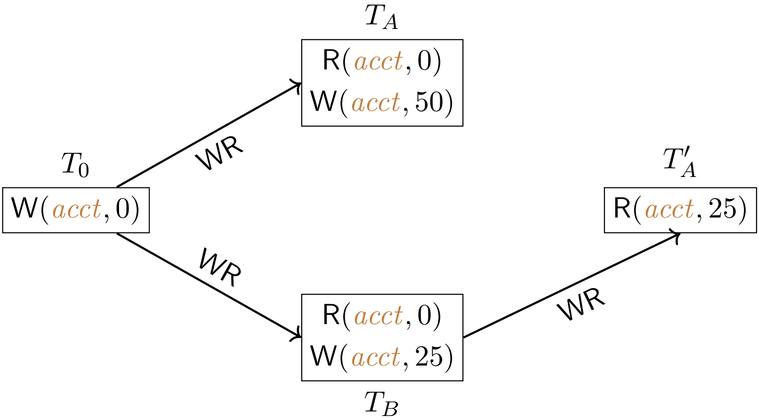
$T'_A$

$\mathsf{R}(acct, 25)$

$\mathsf{R}(acct, 0)$
$\mathsf{W}(acct, 25)$

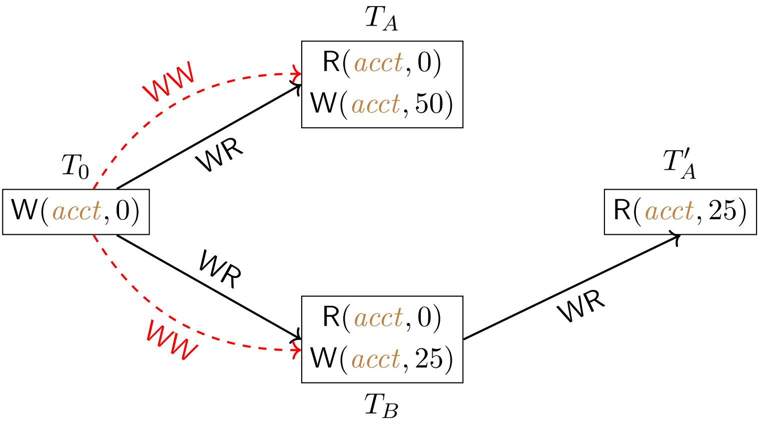
$T_B$

# Dependency Graph-based Characterization of SI



WR: "write-read" dependency capturing the "read-from" relation

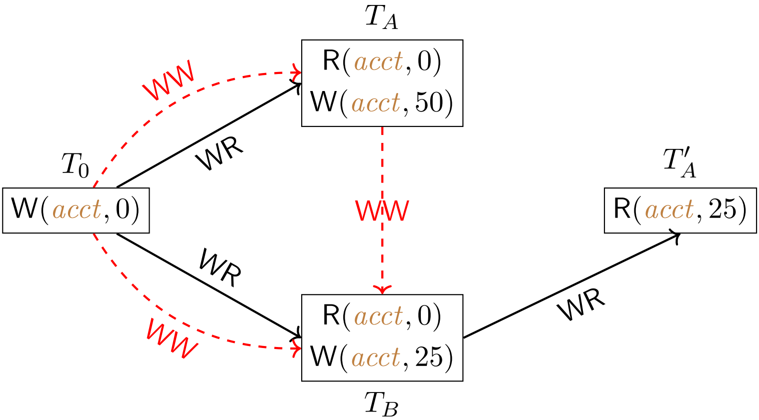
# Dependency Graph-based Characterization of SI



WW: “write-write” dependency capturing the version order

# Dependency Graph-based Characterization of SI

*Suppose that* $T_A \xrightarrow{\textsf{WW}} T_B$



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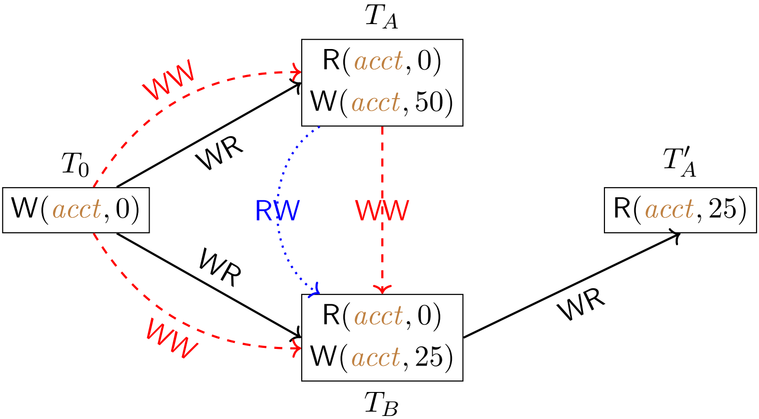
# Dependency Graph-based Characterization of SI

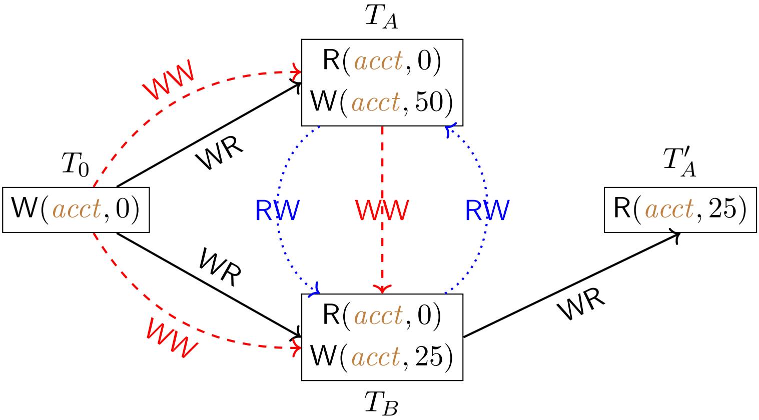
$$T_0 \xrightarrow{\mathsf{WR}} T_A \wedge T_0 \xrightarrow{\mathsf{WW}} T_B \implies T_A \xrightarrow{\mathsf{RW}} T_B$$



RW: “read-write” dependency capturing the overwritten relation

# Dependency Graph-based Characterization of SI

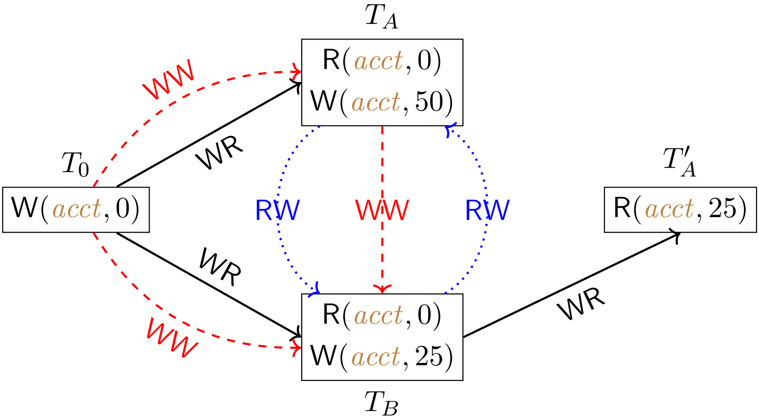
$$T_0 \xrightarrow{\mathsf{WR}} T_B \wedge T_0 \xrightarrow{\mathsf{WW}} T_A \implies T_A \xrightarrow{\mathsf{RW}} T_A$$



RW: “read-write” dependency capturing the overwritten relation

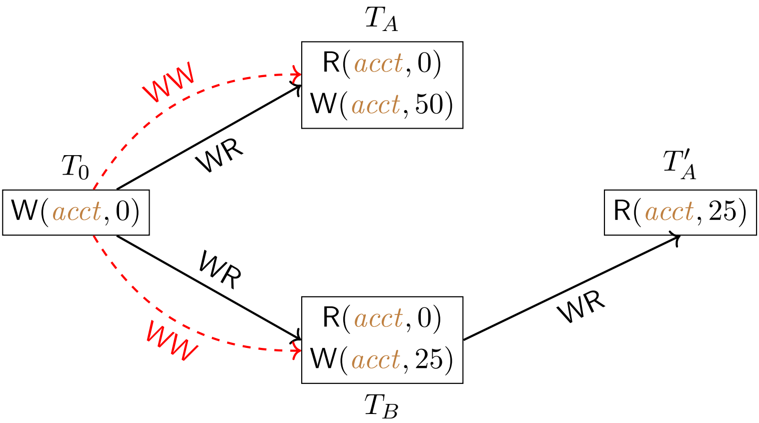
# Dependency Graph-based Characterization of SI

$\boxed{\textit{Suppose that } T_A \xrightarrow{\mathsf{WW}} T_B}$



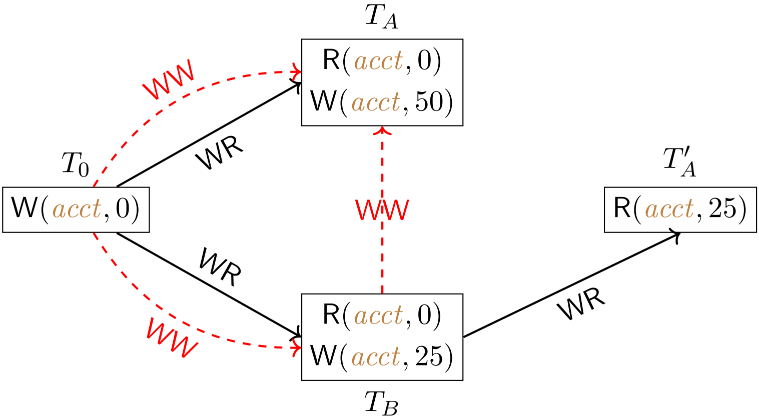
undesired cycle: $T_A \xrightarrow{\mathsf{WW}} T_B \xrightarrow{\mathsf{RW}} T_A$

# Dependency Graph-based Characterization of SI



# Dependency Graph-based Characterization of SI

$\boxed{\textit{Suppose that } T_B \xrightarrow{\mathsf{WW}} T_A}$



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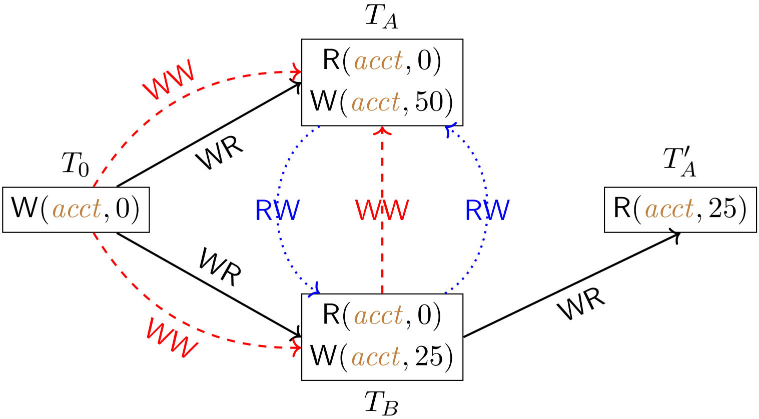
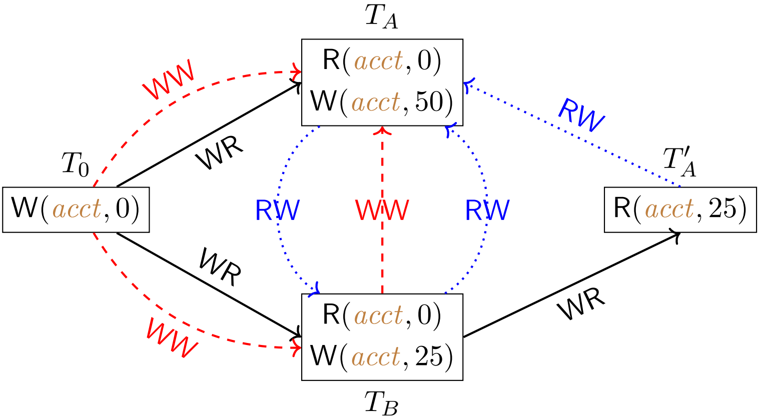
$\boxed{\textit{Suppose that } T_B \xrightarrow{\mathsf{WW}} T_A}$



# Dependency Graph-based Characterization of SI

*Suppose that* $T_B \xrightarrow{\mathsf{WW}} T_A$



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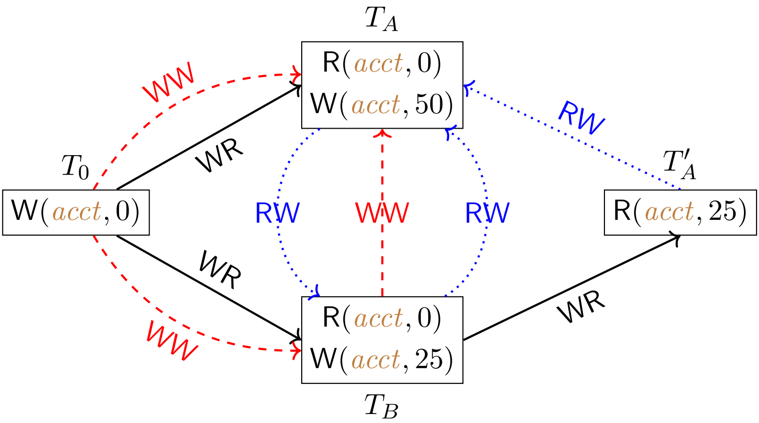
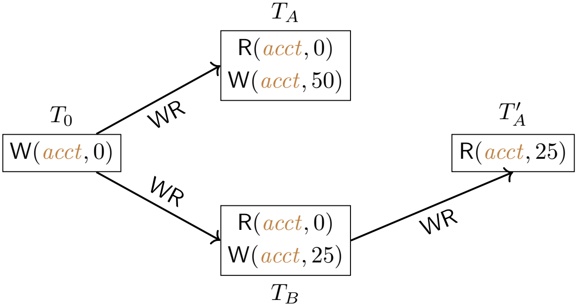
$\boxed{\textit{Suppose that } T_B \xrightarrow{\text{WW}} T_A}$



undesired cycle: $T_B \xrightarrow{\text{WW}} T_A \xrightarrow{\text{RW}} T_B$

# Dependency Graph-based Characterization of SI

We have considered both bases $T_A \xrightarrow{\mathsf{WW}} T_B$ and $T_B \xrightarrow{\mathsf{WW}} T_A$.



Either case leads to an undesired cycle.

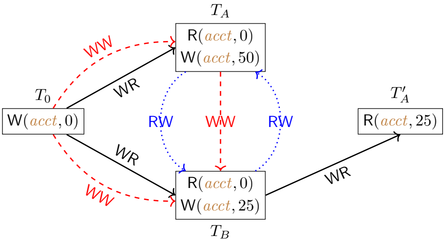
Therefore, it does not satisfy SI.

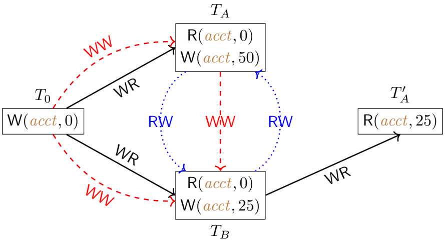
# Dependency Graph-based Characterization of SI

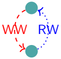
Theorem (Theorem 4.1 of [**AnalysingSI:JACM2018**])

*Informally, a history satisfies SI if only if*
*there exists a dependency graph for it that contains*
*only cycles (if any) with at least two adjacent* RW *edges.*

# Dependency Graph-based Characterization of SI





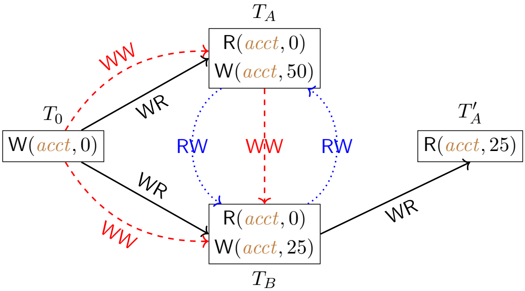
Every possible dependency graph contains an undesired  cycle.

# Dependency Graph-based Characterization of SI

**Theorem (Theorem 4.1 of [AnalysingSI:JACM2018])**

*For a history* $\mathcal{H} = (T, \mathsf{SO})$,

$$\begin{aligned} \mathcal{H} \models \mathrm{SI} \iff \mathcal{H} \models \textsc{Int} \land \\ \exists\, \mathsf{WR}, \mathsf{WW}, \mathsf{RW}.\ \mathcal{G} = (\mathcal{H}, \mathsf{WR}, \mathsf{WW}, \mathsf{RW}) \land \\ (((\mathsf{SO}_{\mathcal{G}} \cup \mathsf{WR}_{\mathcal{G}} \cup \mathsf{WW}_{\mathcal{G}}) \,;\, \mathsf{RW}_{\mathcal{G}}?) \ \textit{is acyclic}). \end{aligned}$$



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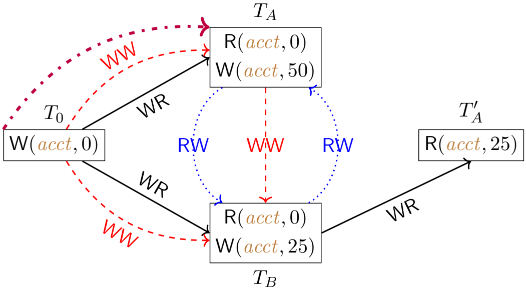

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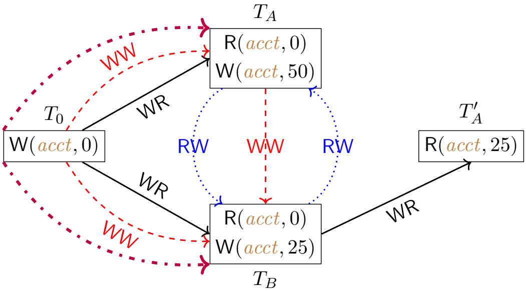

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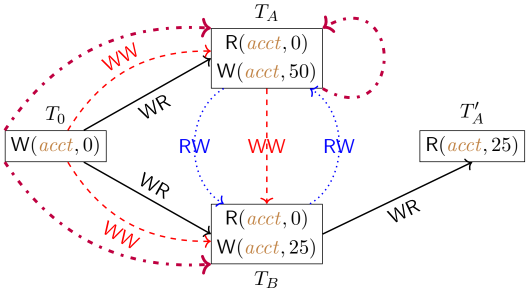

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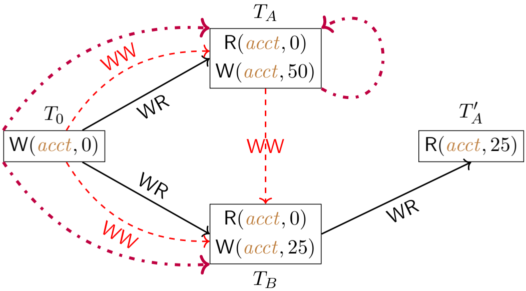
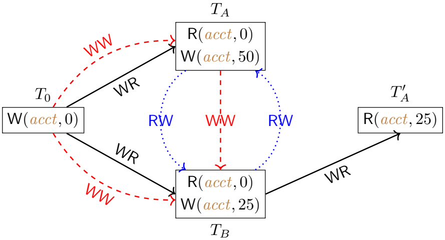
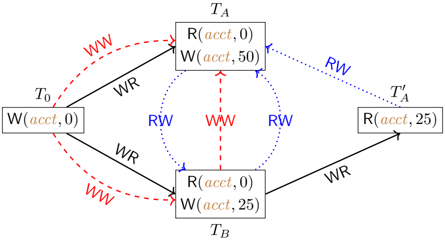


# Dependency Graph-based Characterization of SI

$\mathcal{Q}$ : How to capture and resolve all possible WW dependencies?





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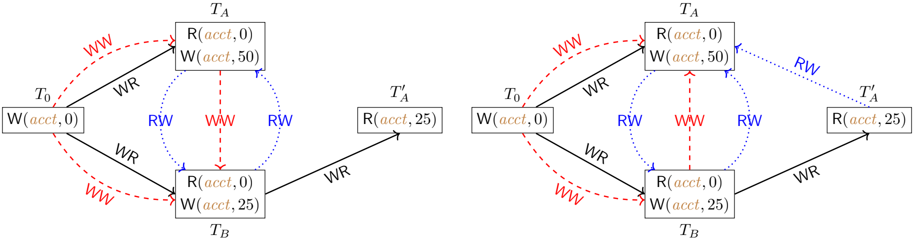
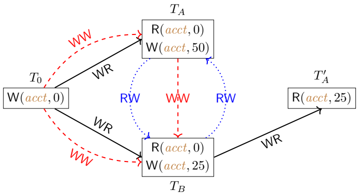
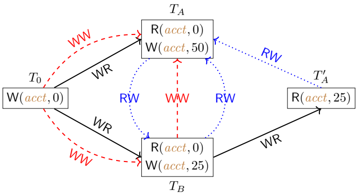


$\mathcal{A}$ : encode them into SAT formulas based on (generalized) polygraphs and solve them using SAT solvers.

# Polygraphs: A Family of Dependency Graphs

Consider the two cases of WW dependencies between $T_A$ and $T_B$.





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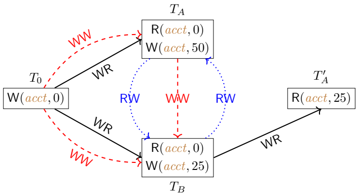

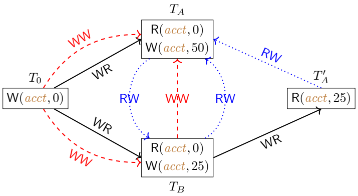

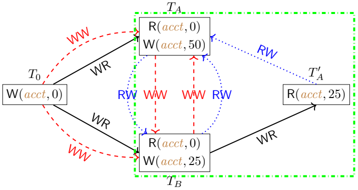

generalized polygraph:

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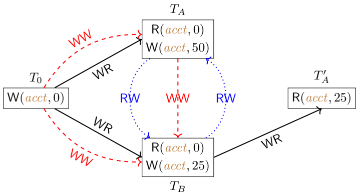
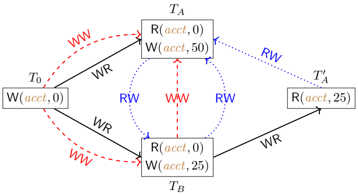
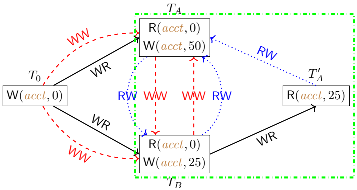


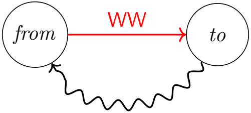


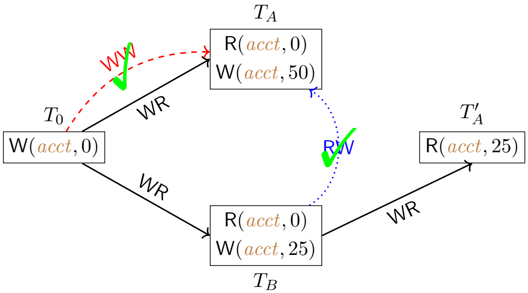


generalized polygraph:

$$\langle \textit{either} \triangleq \{T_A \xrightarrow{\mathsf{WW}} T_B\}, \textit{or} \triangleq \{T_B \xrightarrow{\mathsf{WW}} T_A, T'_A \xrightarrow{\mathsf{RW}} T_A\} \rangle$$

# PolySI: Pruning before Encoding (the WW case)





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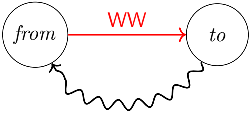
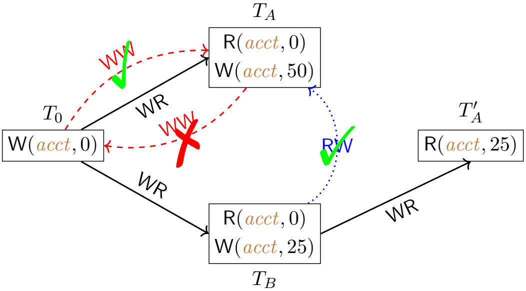
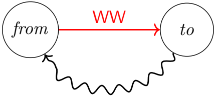
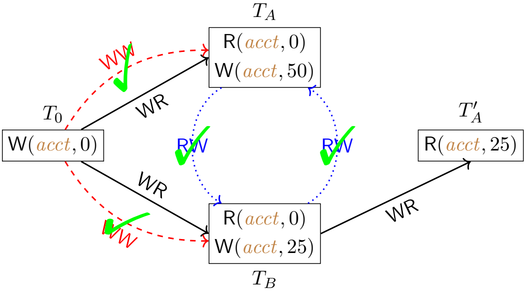




$T_A \xrightarrow{\text{WW}} T_0$ can be pruned due to the $T_A \xrightarrow{\text{WW}} T_0 \xrightarrow{\text{WR}} T_A$ cycle.

# PolySI: Pruning before Encoding (the WW case)





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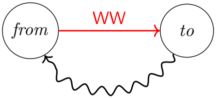
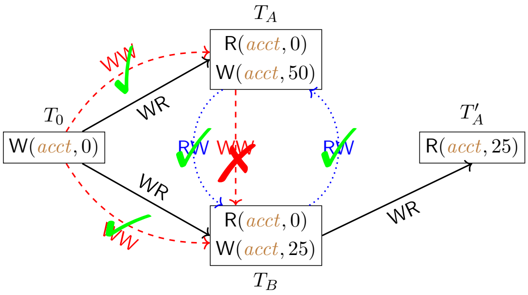
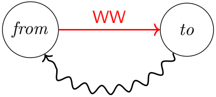
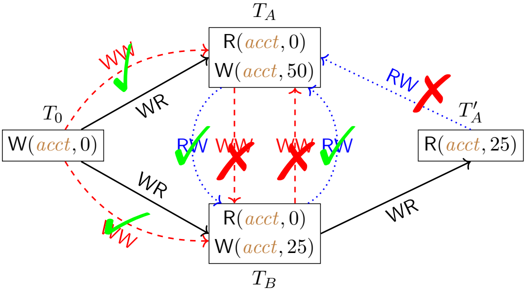




$T_A \xrightarrow{\text{WW}} T_B$ is pruned due to the $T_A \xrightarrow{\text{WW}} T_B \xrightarrow{\text{RW}} T_A$ cycle.

# PolySI: Pruning before Encoding (the WW case)





$T_A \xrightarrow{\text{WW}} T_B$ is pruned due to the $T_A \xrightarrow{\text{WW}} T_B \xrightarrow{\text{RW}} T_A$ cycle.

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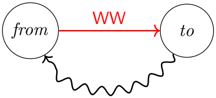
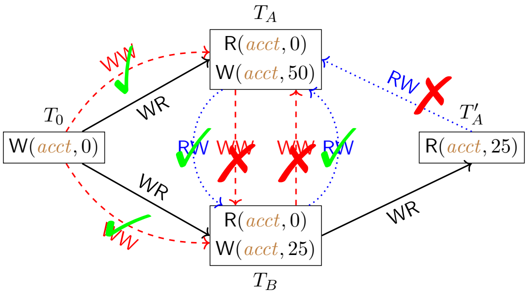


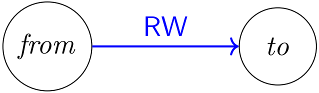


$T_A \xrightarrow{\text{WW}} T_B$ is pruned due to the $T_A \xrightarrow{\text{WW}} T_B \xrightarrow{\text{RW}} T_A$ cycle.

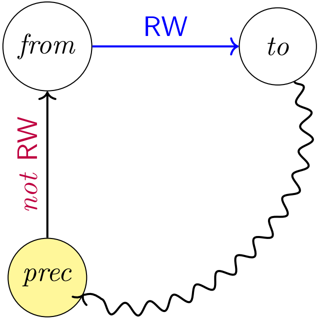
$T_B \xrightarrow{\text{WW}} T_A$ is pruned due to the $T_B \xrightarrow{\text{WW}} T_A \xrightarrow{\text{RW}} T_B$ cycle.

Therefore, we are sure that the history does *not* satisfy SI.

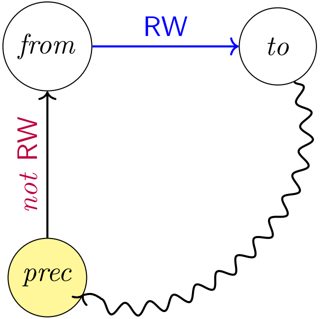
# PolySI: Pruning before Encoding (the RW case)



# PolySI: Pruning before Encoding (the RW case)



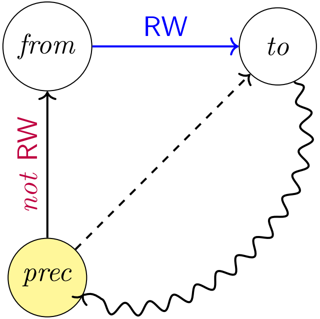
# PolySI: Pruning before Encoding (the RW case)



Theorem (Theorem 4.1 of [**AnalysingSI:JACM2018**])

*Informally, a history satisfies SI if only if there exists a dependency graph for it that contains only cycles (if any) with at least two adjacent* RW *edges.*

# PolySI: Pruning before Encoding (the RW case)



Theorem (Theorem 4.1 of [**AnalysingSI:JACM2018**])

*Informally, a history satisfies SI if only if there exists a dependency graph for it that contains only cycles (if any) with at least two adjacent* RW *edges.*

# PolySI: An Illustrating Example of “Long Fork”

$$T_0 \; \boxed{\mathsf{W}(x, 0)\ \mathsf{W}(y, 0)}$$

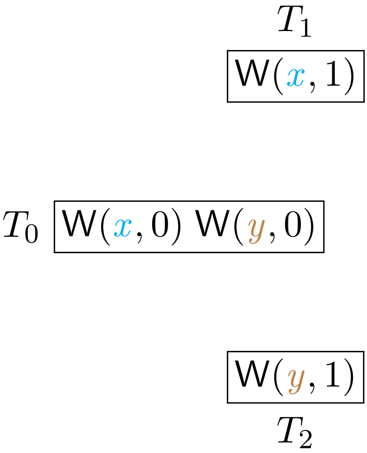
# PolySI: An Illustrating Example of “Long Fork”

$$T_1$$

$$\boxed{\mathsf{W}(x, 1)}$$

$$T_0 \ \boxed{\mathsf{W}(x, 0)\ \mathsf{W}(y, 0)}$$

# PolySI: An Illustrating Example of "Long Fork"



$T_1$

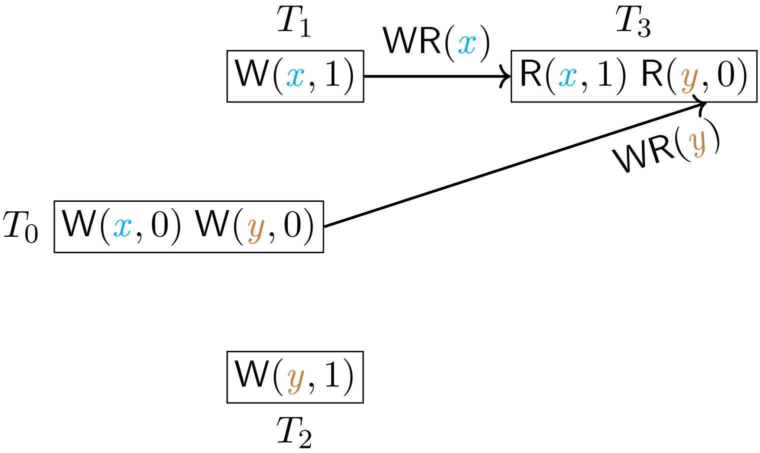
$\boxed{\mathsf{W}(x, 1)}$

$T_0$ $\boxed{\mathsf{W}(x, 0)\ \mathsf{W}(y, 0)}$

$\boxed{\mathsf{W}(y, 1)}$

$T_2$

# PolySI: An Illustrating Example of "Long Fork"



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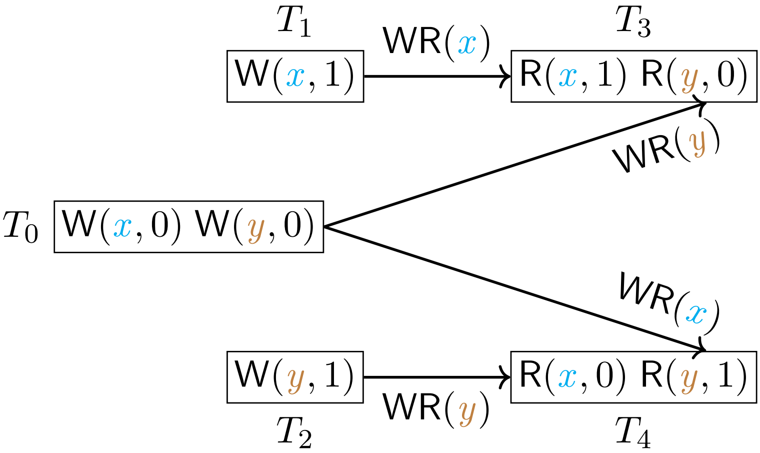

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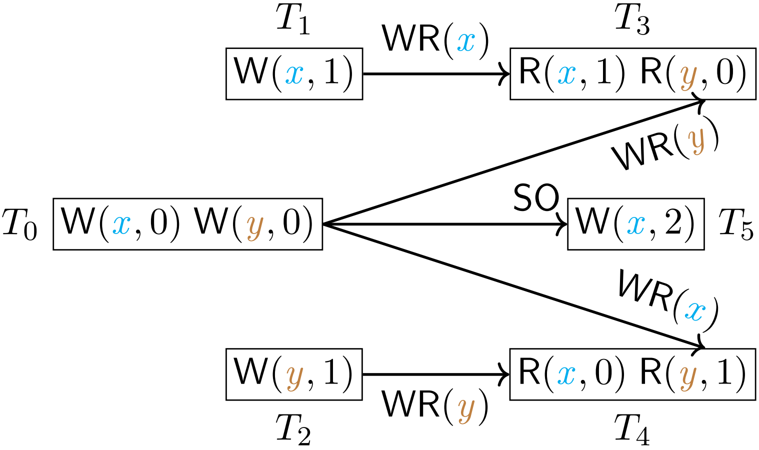
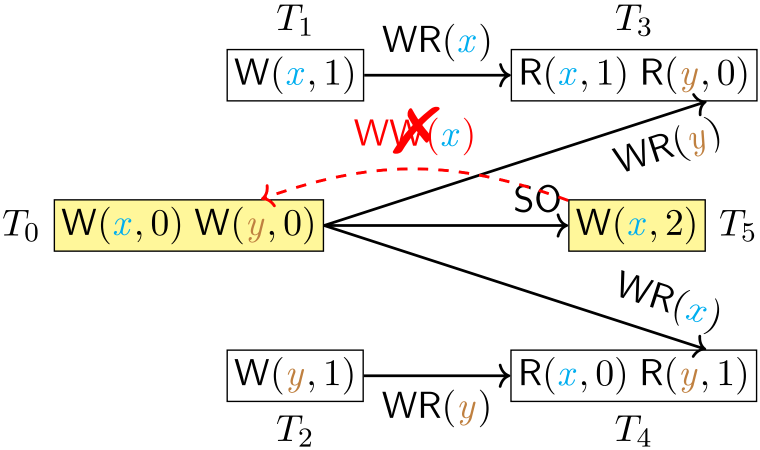


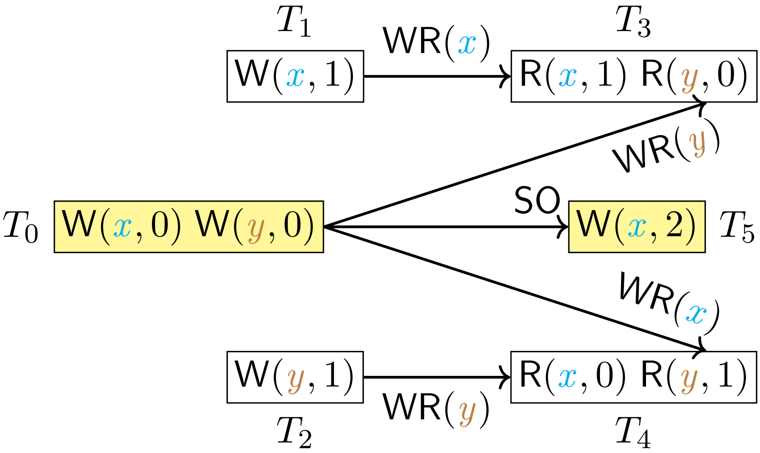
order between $T_0$, $T_1$, and $T_5$ (on $x$) and between $T_0$ and $T_2$ (on $y$)

# PolySI: An Illustrating Example of “Long Fork”



The $T_5 \xrightarrow{\mathsf{WW}(x)} T_0$ case is pruned due to $T_0 \xrightarrow{\mathsf{SO}} T_5 \xrightarrow{\mathsf{WW}(x)} T_0$.

# PolySI: An Illustrating Example of "Long Fork"



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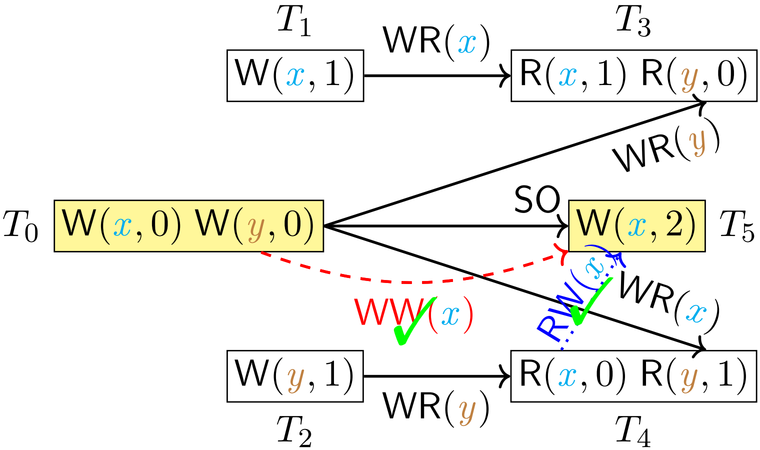
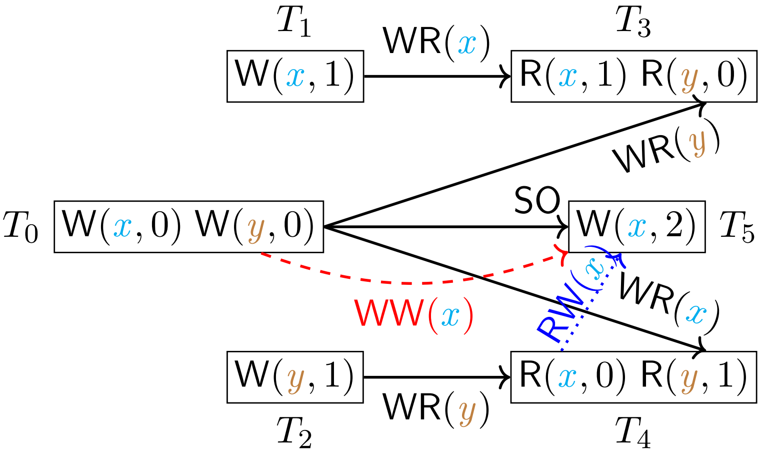


The $T_0 \xrightarrow{\mathsf{WW}(x)} T_5$ case becomes known.

# PolySI: An Illustrating Example of "Long Fork"



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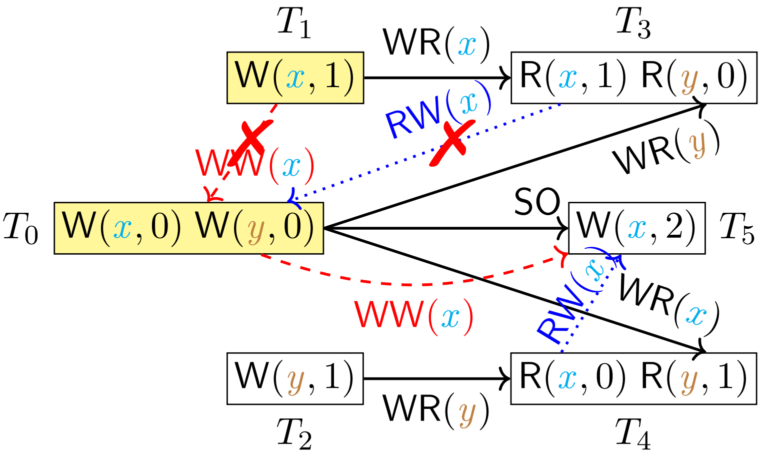
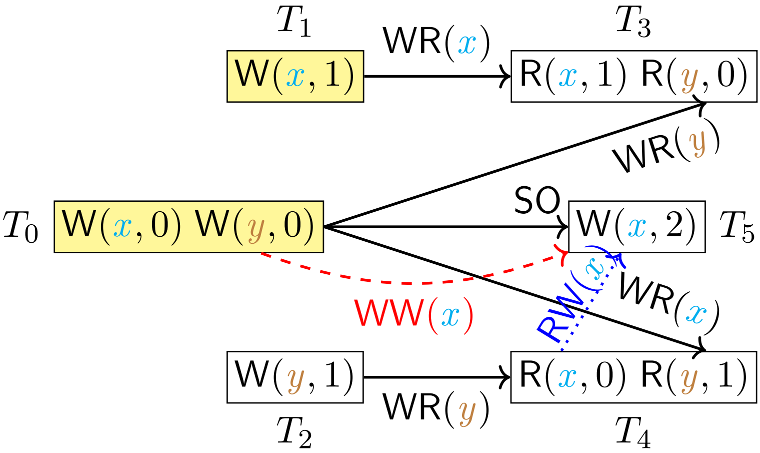


The $T_1 \xrightarrow{\mathsf{WW}(x)} T_0$ case is pruned due to $T_3 \xrightarrow{\mathsf{RW}(x)} T_0 \xrightarrow{\mathsf{WR}(y)} T_3$.

# PolySI: An Illustrating Example of “Long Fork”



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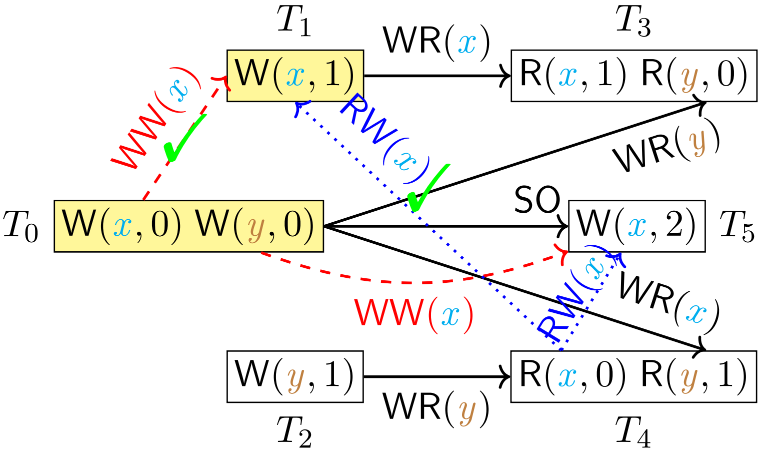
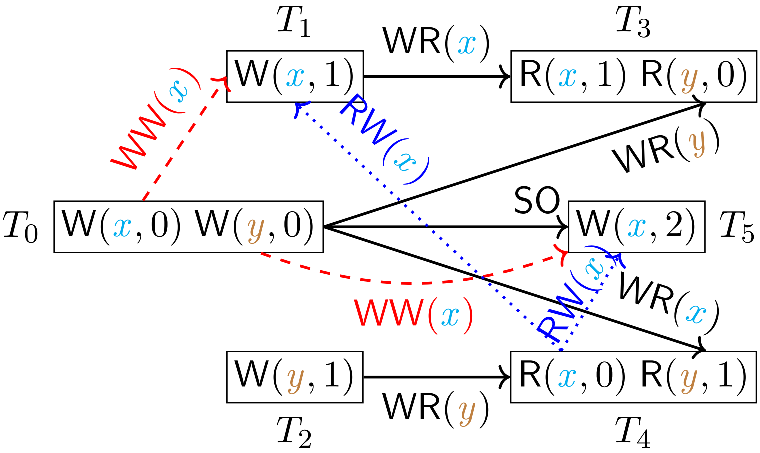
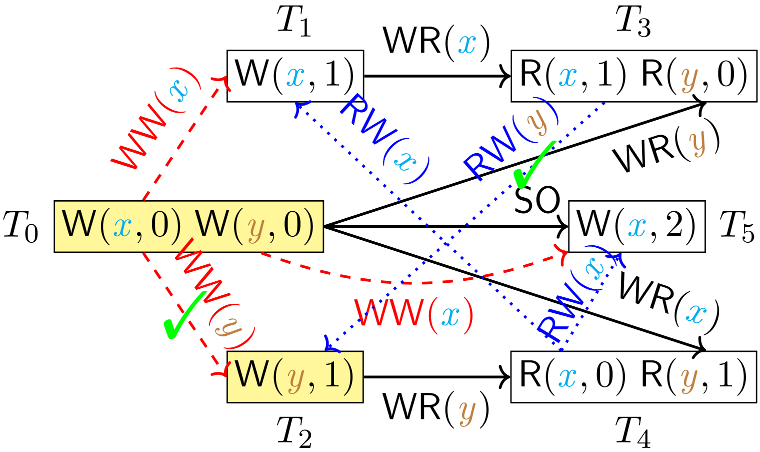


The $T_0 \xrightarrow{\mathsf{WW}(x)} T_1$ case becomes known.

# PolySI: An Illustrating Example of “Long Fork”



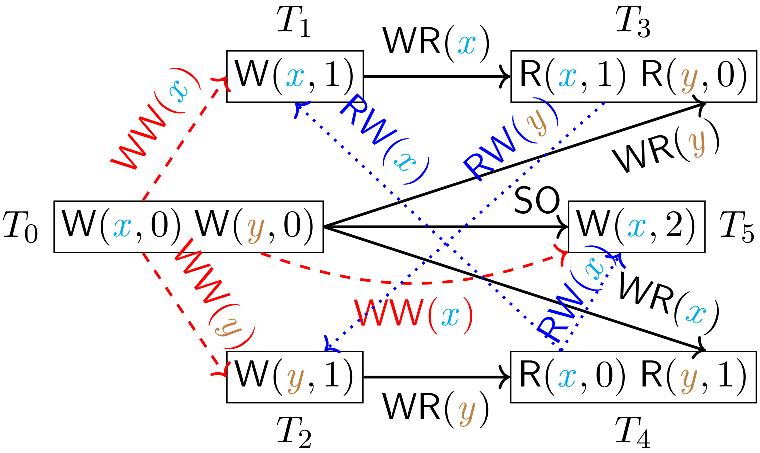
# PolySI: An Illustrating Example of “Long Fork”



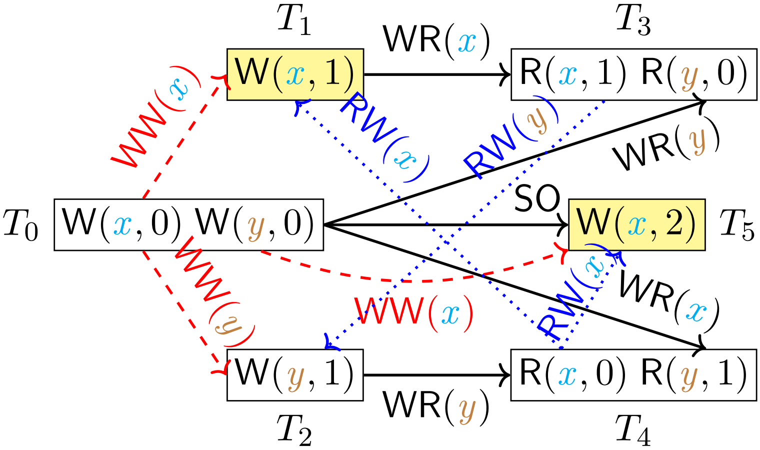
The $T_2 \xrightarrow{\mathsf{WW}(y)} T_0$ case is pruned,

while the $T_0 \xrightarrow{\mathsf{WW}(y)} T_2$ case becomes known.

# PolySI: An Illustrating Example of "Long Fork"



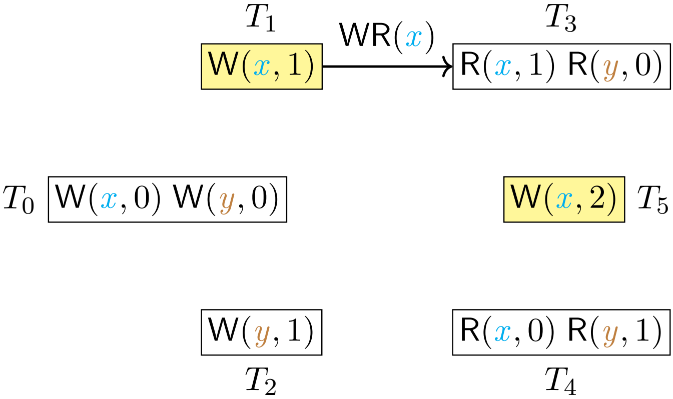
# PolySI: An Illustrating Example of "Long Fork"



The order between $T_1$ and $T_5$ is still uncertain after pruning.

# PolySI: An Illustrating Example of “Long Fork”

$\langle \qquad , \qquad \rangle$



$T_1$ $\mathsf{W}(x, 1)$ $\xrightarrow{\mathsf{WR}(x)}$ $T_3$ $\mathsf{R}(x, 1)$ $\mathsf{R}(y, 0)$

$T_0$ $\mathsf{W}(x, 0)$ $\mathsf{W}(y, 0)$

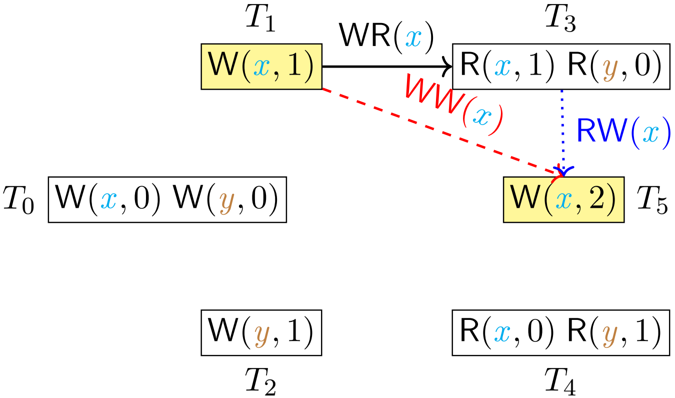
$\mathsf{W}(x, 2)$ $T_5$

$\mathsf{W}(y, 1)$ $T_2$

$\mathsf{R}(x, 0)$ $\mathsf{R}(y, 1)$ $T_4$

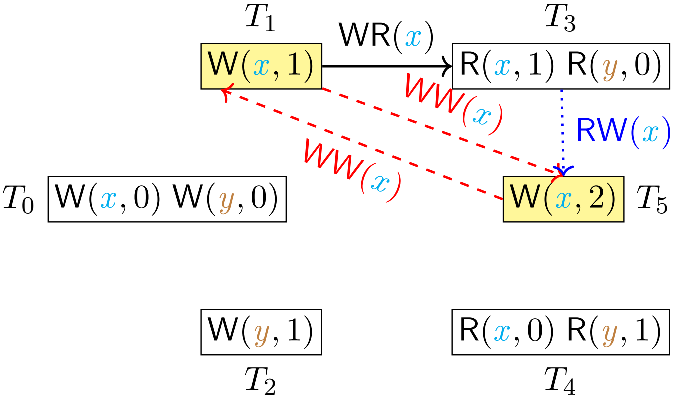
# PolySI: An Illustrating Example of “Long Fork”

$\langle \textit{either} = \{T_1 \xrightarrow{\mathsf{WW}(x)} T_5, T_3 \xrightarrow{\mathsf{RW}(x)} T_5\}, \qquad \rangle$



# PolySI: An Illustrating Example of "Long Fork"

$$\langle either = \{T_1 \xrightarrow{\mathsf{WW}(x)} T_5, T_3 \xrightarrow{\mathsf{RW}(x)} T_5\}, or = \{T_5 \xrightarrow{\mathsf{WW}(x)} T_1\}\rangle$$



# PolySI: An Illustrating Example of "Long Fork"

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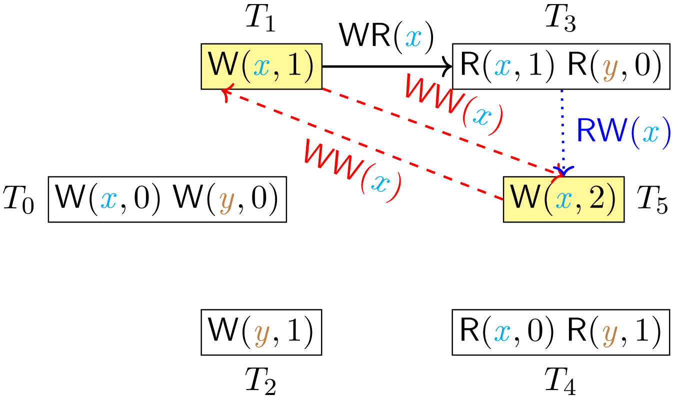
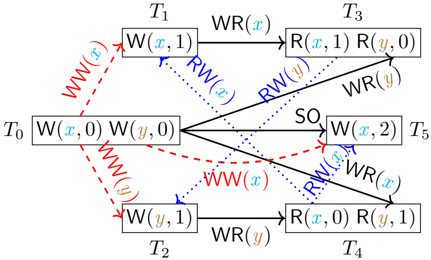
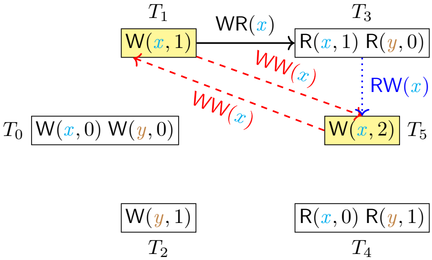


$$(\mathsf{BV}_{1,5} \land \mathsf{BV}_{3,5} \land \neg\mathsf{BV}_{5,1}) \lor (\mathsf{BV}_{5,1} \land \neg\mathsf{BV}_{1,5} \land \neg\mathsf{BV}_{3,5})$$

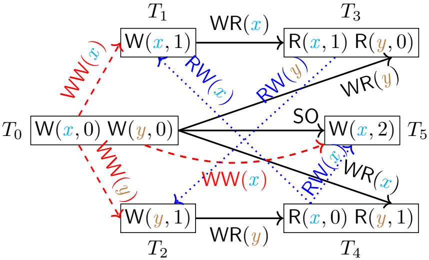
# PolySI: An Illustrating Example of "Long Fork"

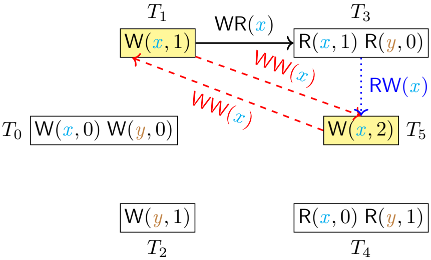




# PolySI: An Illustrating Example of "Long Fork"

$\boxed{((\mathsf{SO}_{\mathcal{G}} \cup \mathsf{WR}_{\mathcal{G}} \cup \mathsf{WW}_{\mathcal{G}}) \; ; \; \mathsf{RW}_{\mathcal{G}}?)}$ *is acyclic.*





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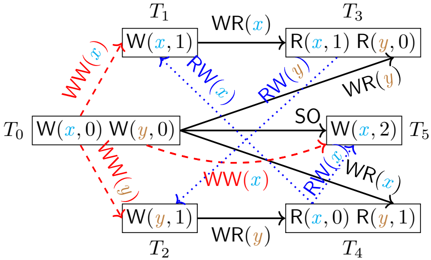

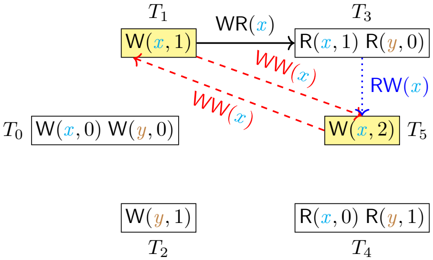

We need to encode the "composition ( ; )" of dependency edges.

# PolySI: An Illustrating Example of “Long Fork”

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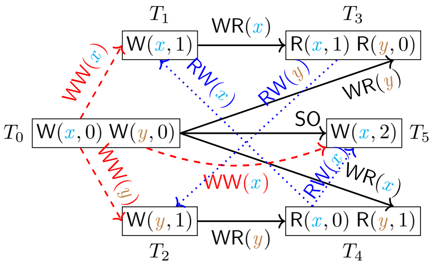

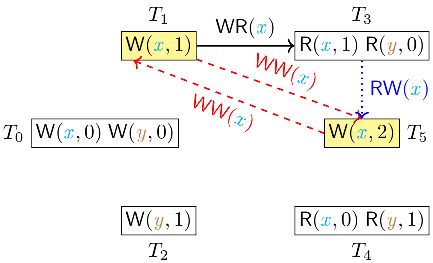

We need to encode the “composition ( ; )” of dependency edges.

$$T_1 \xrightarrow{\mathsf{WR}} T_3 \xrightarrow{\mathsf{RW}} T_2 : \ \mathsf{BV}^{I}_{1,2} = \mathsf{BV}_{1,3} \wedge \mathsf{BV}_{3,2} \quad (I \text{ for the induced graph})$$

# PolySI: An Illustrating Example of "Long Fork"

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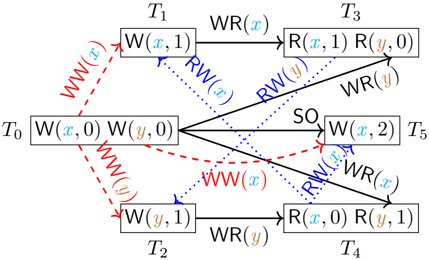
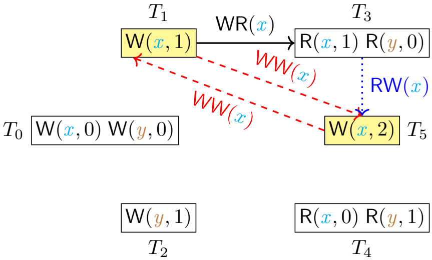




We need to encode the "composition ( ; )" of dependency edges.

$$T_1 \xrightarrow{\mathsf{WR}} T_3 \xrightarrow{\mathsf{RW}} T_2 : \; \mathsf{BV}^{I}_{1,2} = \mathsf{BV}_{1,3} \wedge \mathsf{BV}_{3,2} \quad (I \text{ for the induced graph})$$

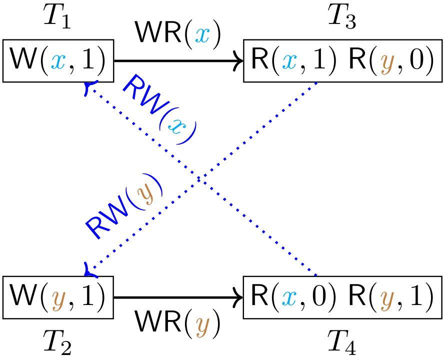
$$T_1 \xrightarrow{\mathsf{WR}} T_3 \xrightarrow{\mathsf{RW}} T_5 : \; \mathsf{BV}^{I}_{1,5} = \mathsf{BV}_{1,3} \wedge \mathsf{BV}_{3,5} \quad (I \text{ for the induced graph})$$

# PolySI: An Illustrating Example of "Long Fork"

Feed the SAT formula into the MonoSAT solver [**MonoSAT:AAAI2015**] optimized for *cycle detection*

Assert that the induced graph $I$ is acyclic.

# PolySI: An Illustrating Example of "Long Fork"



The undesired cycle for "long fork" found by MonoSAT.

# Experimental Evaluation

(1) *Effective:* Can PolySI find SI violations in production databases?

(2) *Informative:* Can PolySI provide understandable counterexamples for SI violations?

(3) *Efficient:* How efficient is PolySI? Is it scalable?

`https://github.com/hengxin/PolySI-PVLDB2023-Artifacts`

# Workloads

Table: Workload parameters and their default values.

| Parameter | Default Value |
|---|---|
| #sess | 20 |
| #txns/sess | 100 |
| #ops/txn | 15 |
| #keys | 10, 000 |
| %reads | 50% |
| distribution | zipfian |

# Benchmarks

RuBis: an eBay-like bidding system

TPC-C: an open standard for OLTP benchmarking

C-Twitter: a Twitter clone

GeneralRH: read-heavy workloads with 95% reads

GeneralRW: medium workloads with 50% reads

GeneralWH: write-heavy workloads with 30% reads

Use a simple database schema of a *two-column table* storing keys and values.

# Finding SI Violations

Table: Reproducing known SI violations.

| Database | GitHub Stars | Kind | Release |
|---|---|---|---|
| CockroachDB | 25.1k | Relational | v2.1.0, v2.1.6 |
| MySQL-Galera | 381 | Relational | v25.3.26 |
| YugabyteDB | 6.7k | Multi-model | v1.1.10.0 |

An extensive collection of 2477 anomalous histories

[**Complexity:OOPSLA2019**; **CockroachDB-bug**; **YugabyteDB-bug**]

# Finding SI Violations

Dgraph: helped the Dgraph team confirm some of their suspicions about their latest release

Table: Detecting new violations.

| **Database** | **GitHub Stars** | **Kind** | **Release** |
|---|---|---|---|
| Dgraph | 18.2k | Graph | v21.12.0 |
| MariaDB-Galera | 4.4k | Relational | v10.7.3 |
| YugabyteDB | 6.7k | Multi-model | v2.11.1.0 |

Galera: confirmed the incorrect claim on preventing “lost updates” for transactions issued on different cluster nodes

# Understanding Violations

# Performance Evaluation

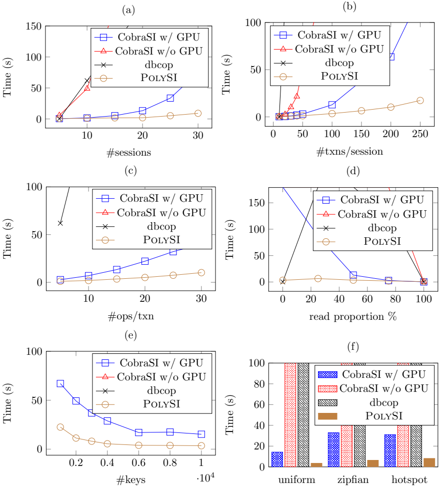
dbcop [**Complexity:OOPSLA2019**]: the state-of-the-art SI checker without using SAT solvers

Cobra [**Cobra:OSDI2020**]: the state-of-the-art SER checker using both MonoSAT and GPU; as a baseline

CobraSI: reducing SI checking to SER checking [**Complexity:OOPSLA2019**] to leverage Cobra with/without GPU

# Performance Evaluation: Runtime

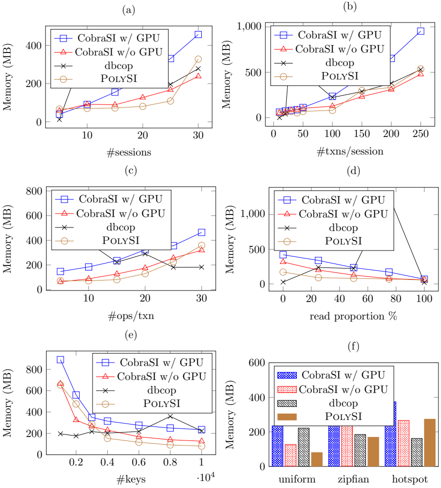
PolySI significantly outperforms the competitors.



All the input histories extracted from PostgreSQL satisfy SI.

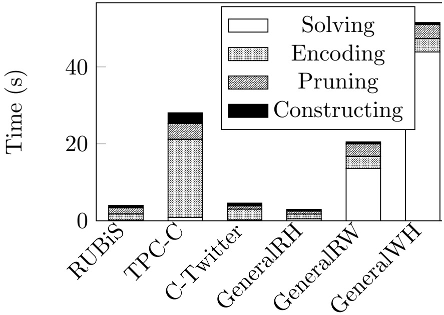
# Performance Evaluation: Memory

PolySI consumes less memory.



# Performance Evaluation: Decomposition

TPC-C incurs more overhead in *encoding* as the number of operations in total is 5x more than the others.



The solving time depends on the remaining constraints and unknown dependencies *after pruning*.

# Performance Evaluation: Pruning

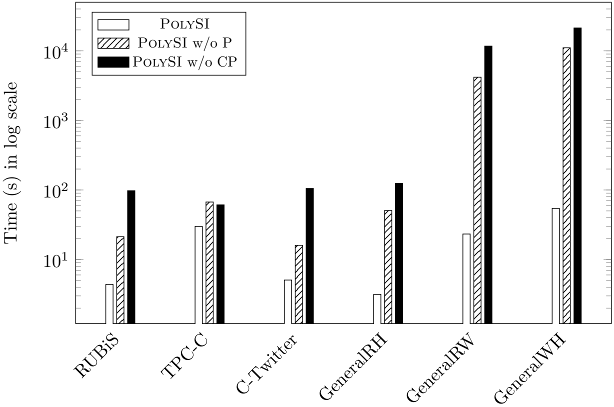
PolySI can effectively prune a huge number of constraints.

| Benchmark | #cons.<br>before P | #cons.<br>after P | #unk. dep.<br>before P | #unk. dep.<br>after P |
|---|---|---|---|---|
| TPC-C | 386k | 0 | 3628k | 0 |
| GeneralRH | 4k | 29 | 39k | 77 |
| RUBiS | 14k | 149 | 171k | 839 |
| C-Twitter | 59k | 277 | 307k | 776 |
| GeneralRW | 90k | 2565 | 401k | 5435 |
| GeneralWH | 167k | 6962 | 468k | 14376 |

TPC-C: read-only transactions + RMW transactions

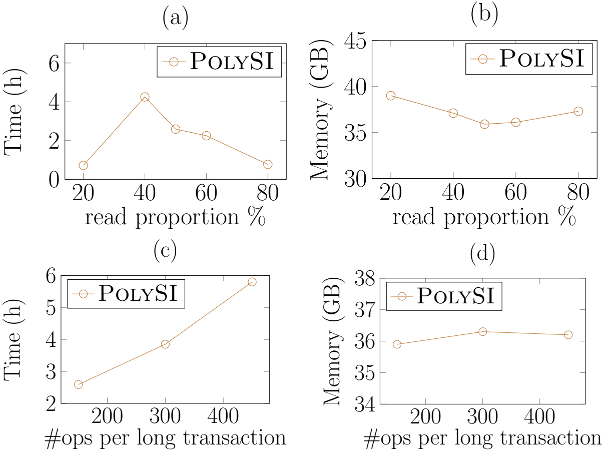
# Performance Evaluation: Differential Analysis

Pruning is crucial to the efficiency of PolySI.

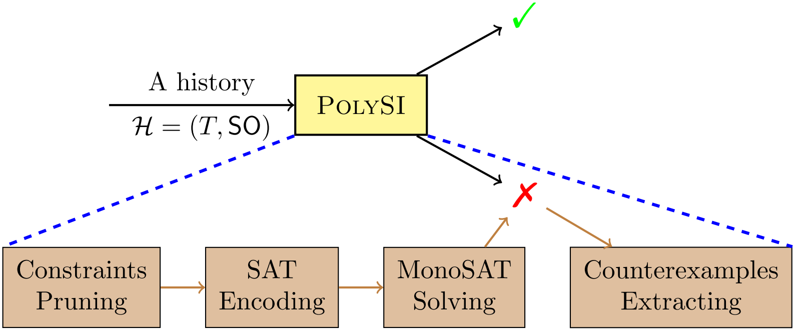


# Performance Evaluation: Scalability

several hours and $35 \sim 40$GB memory for checking 1M transactions



# Conclusion


A history
$\mathcal{H} = (T, \mathsf{SO})$
PolySI
✓
✗
Constraints Pruning
SAT Encoding
MonoSAT Solving
Counterexamples Extracting

# Future Work

PolySI uses MonoSAT as a black-box.

Working on a **theory solver** dedicated to isolation level checking, which is deeply integrated with SAT solvers [**Zord:PLDI2021**].

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