Implementing a Simple Programming Language

Functional Programming (CS-210)

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```
Example program:
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    def fact = (n => (if n then (* n (fact (- n 1))) else 1))
        (fact 6)
)
evaluates to:
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  def square = (x \Rightarrow (*x x))
  (twice square 3)
evaluates to:
```

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Example program:
  def fact = (n => (if n then (* n (fact (- n 1))) else 1))
  (fact 6)
evaluates to: 720
  def twice = (f \Rightarrow x \Rightarrow (f (f x)))
  def square = (x \Rightarrow (*x x))
  (twice square 3)
evaluates to: 81
```

Program Representation: Abstract Syntax Trees

```
(def twice = (f \Rightarrow x \Rightarrow (f (f x)))
 def square = (x \Rightarrow (* x x))
 (twice square 3))
\approx
val defs : DefEnv = Map[String, Expr](
  "twice" -> Fun("f", Fun("x",
                             Call(N("f"). Call(N("f"). N("x")))).
  "square" -> Fun("x", BinOp(Times, N("x"), N("x"))))
val expr = Call(Call(N("twice"), N("square")), C(3))
```

- ▶ We represent a program using *expression tree* called Abstract Syntax Tree (AST)
- Our implementation is an interpreter, which traverses AST to produce the result
- We discuss later briefly how to convert an input file into an abstract syntax tree; more on that in the course Computer Language Processing (CS-320) next year

Growing a Language and Its Interpreter

- 101 Language of arithmetic and *if* expressions
- 102 Absolute value and its *desugaring*
- 103 Recursive functions implemented using substitutions
- 104 **Environment** instead of substitutions
- 105 *Higher-order* functions using substitutions
- 106 Higher-order functions using environments
- 107 **Nested recursive** definitions using environments

101. Language of arithmetic and *if* expressions: Trees

Integer constants combined using arithmetic operations and the if conditional

How to describe such trees?

101. Language of arithmetic and *if* expressions: Trees

```
Integer constants combined using arithmetic operations and the if conditional
val expr1 = BinOp(Times, C(6), C(7)) // 6 * 7
val cond1 = BinOp(LessEq, expr1, C(50)) // expr1 <= 50</pre>
val expr2 = IfNonzero(cond1, C(10), C(20)) // if (cond1) 10 else 20
How to describe such trees?
enum Expr
  case C(c: BigInt)
                                               // integer constant
  case BinOp(op: BinOps, e1: Expr, e2: Expr) // binary operation
  case IfNonzero(cond: Expr. trueE: Expr. falseE: Expr)
enum BinOps
  case Plus. Minus. Times. Power. LessEq
```

101. Language of arithmetic and *if* expressions: Printing

```
def str(e: Expr): String = e match
  case C(c) => c.toString
  case BinOp(op, e1, e2) =>
    s"(${strOp(op)} ${str(e1)} ${str(e2)})" // string interpolation
  case IfNonzero(cond, trueE, falseE) =>
    s"(if ${str(cond)} then ${str(trueE)} else ${str(falseE)})"
def strOp(op: BinOps): String = op match
  case Plus => "+"
  case Minus => "-"
  case Times => "*"
 case Power => "^"
 case LessEq => "<="</pre>
> str(IfNonzero(BinOp(LessEq, C(4), C(50)), C(10), C(20)))
(if (<= 4 50) then 10 else 20)
```

101. Language of arithmetic and *if* expressions: Interpreting

```
def eval(e: Expr): BigInt = e match
 case C(c) \Rightarrow c
  case BinOp(op, e1, e2) =>
    evalBinOp(op)(eval(e1), eval(e2))
  case IfNonzero(cond, trueE, falseE) =>
    if eval(cond) != 0 then eval(trueE) else eval(falseE)
def evalBinOp(op: BinOps)(x: BigInt, y: BigInt): BigInt = op match
  case Plus => x + v
  case Minus => x - y
  case Times => x * v
  case Power => x.pow(y.toInt)
  case LessEq => if (x <= y) 1 else 0
> eval(IfNonzero(BinOp(LessEq, C(4), C(50)), C(10), C(20)))
10
```

102. Absolute Value and Its *Desugaring*: Trees

```
enum Expr
  case C(c: BigInt)
  case BinOp(op: BinOps, e1: Expr, e2: Expr)
  case IfNonzero(cond: Expr, trueE: Expr, falseE: Expr)
  case AbsValue(arg: Expr) // new case
```

How to extend evaluator to work with absolute value as well? Two approaches:

- add a case to the interpreter (exercise)
- transform (desugar) trees to reduce them to previous cases

 $Syntactic\ sugar = extra language\ constructs\ that\ are\ not\ strictly\ necessary\ because\ they\ can be\ expressed\ in\ terms\ of\ others\ (they\ make\ the\ language\ sweeter\ to\ use)$

Desugaring = automatically eliminating syntactic sugar by expanding constructs

102. Desugaring Absolute Value: Idea

By definition of absolute value, we would like this equality to hold:

How to write desugar function that eliminates all occurrences of AbsValue?

102. Desugaring Absolute Value: Idea

By definition of absolute value, we would like this equality to hold:

How to write desugar function that eliminates all occurrences of AbsValue?

Replace (recursively) each subtree AbsValue(x) with its definition.

102. Desugaring Absolute Value: Code

```
def desugar(e: Expr): Expr = e match
  case C(c) => e
  case BinOp(op, e1, e2) =>
       BinOp(op. desugar(e1), desugar(e2))
  case IfNonzero(cond, trueE, falseE) =>
       IfNonzero(desugar(cond). desugar(trueE). desugar(falseE))
  case AbsValue(arg) =>
    val x = desugar(arg)
    IfNonzero(BinOp(LessEq. x. C(0)).
              BinOp(Minus, C(0), x).
              x )
```

102. Desugaring Absolute Value: Example Run

```
def show(e: Expr): Unit =
  println("original:")
  println(str(e))
  val de = desugar(e)
  println("desugared:")
  println(str(de))
  println(" ~~> " + eval(de) + "\n")
show(AbsValue(BinOp(Plus.C(10).C(-50))))
original:
(abs (+ 10 -50))
desugared:
(if (<= (+ 10 -50) 0) then (- 0 (+ 10 -50)) else (+ 10 -50))
  ~~> 40
```