

Monads

Principles of Functional Programming

Martin Odersky

Monads

Data structures with map and flatMap seem to be quite common.

In fact there's a name that describes this class of a data structures together with some algebraic laws that they should have.

They are called *monads*.

What is a Monad?

A monad M is a parametric type M[T] with two operations, flatMap and unit, that have to satisfy some laws.

```
extension [T, U](m: M[T])
  def flatMap(f: T => M[U]): M[U]

def unit[T](x: T): M[T]
```

In the literature, flatMap is also called bind. It can be an extension method, or be defined as a regular method in the monad class M.

Examples of Monads

- ► List is a monad with unit(x) = List(x)
- Set is monad with unit(x) = Set(x)
- Option is a monad with unit(x) = Some(x)
- Generator is a monad with unit(x) = single(x)

Monads and map

map can be defined for every monad as a combination of flatMap and unit:

Note: andThen is defined function composition in the standard library.

```
extension [A, B, C](f: A => B)
  @infix def andThen(g: B => C): A => C =
    x => g(f(x))
```

Monad Laws

```
To qualify as a monad, a type has to satisfy three laws:
Associativity:
    m.flatMap(f).flatMap(g) == m.flatMap(f(_).flatMap(g))
Left unit
    unit(x).flatMap(f) == f(x)
Right unit
    m.flatMap(unit) == m
```

Checking Monad Laws

Let's check the monad laws for Option.

Here's flatMap for Option:

```
extension [T](xo: Option[+T])
  def flatMap[U](f: T => Option[U]): Option[U] = xo match
    case Some(x) => f(x)
    case None => None
```

Checking the Left Unit Law

```
Need to show: Some(x).flatMap(f) == f(x)

Some(x).flatMap(f)
```

Checking the Left Unit Law

```
Need to show: Some(x).flatMap(f) == f(x)
Some(x).flatMap(f)

== Some(x) match
    case Some(x) => f(x)
    case None => None
```

Checking the Left Unit Law

Checking the Right Unit Law

```
Need to show: opt.flatMap(Some) == opt
    opt.flatMap(Some)
```

Checking the Right Unit Law

```
Need to show: opt.flatMap(Some) == opt
    opt.flatMap(Some)

== opt match
    case Some(x) => Some(x)
    case None => None
```

Checking the Right Unit Law

```
Need to show: opt.flatMap(f).flatMap(g) ==
opt.flatMap(f(_).flatMap(g))
       opt.flatMap(f).flatMap(g)
       (opt match { case Some(x) \Rightarrow f(x) case None \Rightarrow None })
             match { case Some(v) => g(v) case None => None }
       opt match
          case Some(x) =>
            f(x) match { case Some(y) => g(y) case None => None }
          case None =>
            None match { case Some(y) \Rightarrow g(y) case None \Rightarrow None }
```

```
== opt match
    case Some(x) =>
    f(x) match { case Some(y) => g(y) case None => None }
    case None => None
```

```
== opt match
    case Some(x) =>
        f(x) match { case Some(y) => g(y) case None => None }
    case None => None

== opt match
    case Some(x) => f(x).flatMap(g)
    case None => None
```

```
opt match
  case Some(x) =>
    f(x) match { case Some(y) => g(y) case None => None }
  case None => None
opt match
  case Some(x) \Rightarrow f(x).flatMap(g)
  case None => None
opt.flatMap(x \Rightarrow f(x).flatMap(g))
```

```
opt match
   case Some(x) =>
     f(x) match { case Some(v) => g(v) case None => None }
   case None => None
opt match
   case Some(x) \Rightarrow f(x).flatMap(g)
   case None => None
 opt.flatMap(x \Rightarrow f(x).flatMap(g))
 opt.flatMap(f(_).flatMap(g))
```

Significance of the Laws for For-Expressions

We have seen that monad-typed expressions are typically written as for expressions.

What is the significance of the laws with respect to this?

1. Associativity says essentially that one can "inline" nested for expressions:

```
for
    y <- for x <- m; y <- f(x) yield y
    z <- g(y)
    yield z

== for x <- m; y <- f(x); z <- g(y)
    yield z</pre>
```

Significance of the Laws for For-Expressions

2. Right unit says:

```
for x <- m yield x
== m</pre>
```

3. Left unit does not have an analogue for for-expressions.