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CS-210 Functional programming

Date: 08.12.2021

Duration: 25 minutes (dry run).

The real exam will last 90 minutes.

# 1

# Ada Lovelace

SCIPER: 1000001

Wait for the start of the exam before turning to the next page. This document is printed double sided, 8 pages. Do not unstaple.













- This is a closed book exam. No electronic devices allowed.
- Place on your desk: your student ID, writing utensils place all other personal items below your desk or on the side.
- You each have a different exam. For technical reasons, **do use black or blue pens for the MCQ part, no pencils!** Use white corrector if necessary.
- **Your Time:** All points are not equal: we do not think that all exercises have the same difficulty, even if they have the same number of points.

This dry run contains 4 multiple choice questions worth 4 points each, 1 true/false questions worth 2 points and 1 open questions worth 12 points, for a total of **30 points**.

The real exam will last 90 minutes and will have a total of **100 points**:

- 16 multiple choice questions with a single correct answer: +4 for the correct answer, 0 otherwise.
- 6 true/false questions: +2 for the correct answer, -1 for a wrong answer, 0 if left unanswered.
- 2 open questions worth 12 points each.

- **Your Attention:** The exam problems are precisely and carefully formulated, some details can be subtle. Pay attention, because if you do not understand a problem, you cannot obtain full points.
- **Stay Functional:** You are strictly forbidden to use return statements, mutable state (vars) and mutable collections in your solutions.
- The last page of this exam contains an appendix. Do not detach this page.

Respectez les consignes suivantes   Observe this guidelines   Beachten Sie bitte die unten stehenden Richtlinien		
choisir une réponse   select an answer Antwort auswählen	ne PAS choisir une réponse   NOT select an answer NICHT Antwort auswählen	Corriger une réponse   Correct an answer Antwort korrigieren
  		 
ce qu'il ne faut <b>PAS</b> faire   what should <b>NOT</b> be done   was man <b>NICHT</b> tun sollte		
     		



## First part: single choice questions

Each question has **exactly one** correct answer. Marking the box corresponding to the correct answer (and only that one) will get you 4 points. Otherwise, you will get 0 points. There are no negative points.

Given the following lemmas, holding for all types  $A$ ,  $x: A$ ,  $b1: \text{Bool}$ ,  $b2: \text{Bool}$ ,  $p: A \Rightarrow \text{Bool}$ ,  $xs: \text{List}[A]$  and  $ys: \text{List}[A]$ :

(FORALLNIL)  $\text{nil.forall}(p) == \text{True}$

(FORALLCONS)  $(x :: xs).\text{forall}(p) == p(x) \ \&\& \ xs.\text{forall}(p)$

(EXISTSNIL)  $\text{nil.exists}(p) == \text{False}$

(EXISTSCONS)  $(x :: xs).\text{exists}(p) == p(x) \ || \ xs.\text{exists}(p)$

(NEGFALSE)  $!\text{False} == \text{True}$

(NEGOR)  $!(b1 \ || \ b2) == !b1 \ \&\& \ !b2$

(NEGAND)  $!(b1 \ \&\& \ b2) == !b1 \ || \ !b2$

(NEGINVOLUTIVE)  $!!b1 == b1$

Let us prove the following lemma for all  $l: \text{List}[A]$  and all  $p: A \Rightarrow \text{Bool}$ :

(LISTNEGEXISTS)  $!l.\text{exists}(x \Rightarrow !p(x)) == l.\text{forall}(p)$

We prove it by induction on  $l$ .

*Base case:*  $l$  is  $\text{Nil}$ .

Therefore, we need to prove:

$!\text{Nil}.\text{exists}(x \Rightarrow !p(x)) == \text{Nil}.\text{forall}(p)$

**Question 1** Starting from the left hand-side  $(!\text{Nil}.\text{exists}(x \Rightarrow !p(x)))$ , what exact sequence of lemmas should we apply to get the right hand-side  $(\text{Nil}.\text{forall}(p))$ ?

- ☐ NEGINVOLUTIVE, FORALLNIL, EXISTSNIL
- ☐ FORALLNIL, NEGFALSE, EXISTSNIL
- ☐ NEGFALSE, EXISTSNIL, FORALLNIL,
- ☐ NEGFALSE, FORALLNIL, EXISTSNIL
- ☐ EXISTSNIL, NEGINVOLUTIVE, FORALLNIL
- ☐ EXISTSNIL, NEGFALSE, FORALLNIL
- ☐ FORALLNIL, NEGINVOLUTIVE, EXISTSNIL
- ☐ NEGINVOLUTIVE, EXISTSNIL, FORALLNIL



*Induction step:* let  $l = x :: xs$ .

Therefore, we need to prove:

$$\neg (x :: xs).exists(x \Rightarrow \neg p(x)) \iff (x :: xs).forall(p)$$

Our induction hypothesis is that for  $xs$ :

$$(IH) \neg xs.exists(x \Rightarrow \neg p(x)) \iff xs.forall(p)$$

**Question 2** Starting from the left hand-side  $\neg (x :: xs).exists(x \Rightarrow \neg p(x))$ , what exact sequence of lemmas should we apply to get the right hand-side  $(x :: xs).forall(p)$ ?

- ☐ EXISTSCONS, NEGFALSE, IH, NEGAND, FORALLCONS
- ☐ EXISTSCONS, NEGOR, NEGINVOLUTIVE, IH, FORALLCONS
- ☐ EXISTSCONS, IH, NEGINVOLUTIVE, NEGAND, FORALLCONS
- ☐ EXISTSCONS, NEGINVOLUTIVE, IH, NEGOR, FORALLCONS
- ☐ EXISTSCONS, NEGAND, NEGINVOLUTIVE, IH, FORALLCONS
- ☐ EXISTSCONS, NEGFALSE, IH, NEGOR, FORALLCONS
- ☐ EXISTSCONS, IH, NEGAND, NEGFALSE, FORALLCONS
- ☐ EXISTSCONS, NEGINVOLUTIVE, NEGOR, IH, FORALLCONS
- ☐ EXISTSCONS, NEGINVOLUTIVE, IH, NEGAND, FORALLCONS
- ☐ EXISTSCONS, NEGFALSE, NEGOR, IH, FORALLCONS
- ☐ EXISTSCONS, NEGFALSE, NEGAND, IH, FORALLCONS
- ☐ EXISTSCONS, IH, NEGOR, NEGFALSE, FORALLCONS
- ☐ EXISTSCONS, NEGOR, NEGFALSE, IH, FORALLCONS
- ☐ EXISTSCONS, NEGAND, NEGFALSE, IH, FORALLCONS
- ☐ EXISTSCONS, NEGINVOLUTIVE, NEGAND, IH, FORALLCONS
- ☐ EXISTSCONS, IH, NEGFALSE, NEGOR FORALLCONS



Church booleans are a representation of booleans in the lambda calculus. The Church encoding of true and false are functions of two parameters:

Church encoding of `tru`: `t => f => t`

Church encoding of `fls`: `t => f => f`

**Question 3** What does the following function implement?

```
1 b => c => b c fls
```

- ☐ not (b and c)
- ☐ b or c
- ☐ not b
- ☐ not (b xor c)
- ☐ b and c
- ☐ not c

**Question 4** What should replace ??? so that the following function computes not (b and c)?

```
1 b => c => b ??? (not b)
```

- ☐ (not b)
- ☐ (not c)
- ☐ tru
- ☐ fls
- ☐ b
- ☐ c

**Second part: yes/no questions**

The answer of each question is *either* “**Yes**”, *either* “**No**”. Marking the box corresponding to the correct answer (and only that one) will get you 2 points. Not marking anything will get you 0 points. Otherwise, you will get -1 points

**Question 5** Is “type-directed programming” a language mechanism that infers types from values?

☐

Yes

☐

No

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### Third part, open questions

**Question 5:** *This question is worth 12 points.*

☐ 0 ☐ 1 ☐ 2 ☐ 3 ☐ 4 ☐ 5 ☐ 6 ☐ 7 ☐ 8 ☐ 9 ☐ 10 ☐ 11 ☐ 12

Monoids can be represented by the following type class:

```
1 trait Monoid[T]:  
2   extension (x: T) def combine (y: T): T  
3   def unit: T
```

Additionally the three following laws should hold for all `Monoid[M]` and all `m1, m2, m3: M`:

(ASSOCIATIVITY) `a.combine(b).combine(c) == a.combine(b.combine(c))`

(LEFT UNIT) `unit.combine(a) == a`

(RIGHT UNIT) `a.combine(unit) == a`

Write a `Monoid` implementation for pairs of arbitrary types `(A, B)` as a **given**, using `Monoid[A]` and `Monoid[B]`.





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## Appendix: Scala Standard Library Methods

Here are some methods from the Scala standard library that you may find useful. If `xs` is a `List[A]` then:

- `xs.head`: `A` returns the first element of the list. Throws an exception if the list is empty.
- `xs.tail`: `List[A]` returns the list `xs` without its first element. Throws an exception if the list is empty.
- `x :: (xs: List[A]): List[A]` prepends the element `x` to the left of `xs`, returning a `List[A]`.
- `xs ++ (ys: List[A]): List[A]` appends the list `ys` to the right of `xs`, returning a `List[A]`.
- `xs.apply(n: Int): A`, or `xs(n: Int): A` returns the `n`-th element of `xs`. Throws an exception if there is no element at that index.
- `xs.drop(n: Int): List[A]` returns a `List[A]` that contains all elements of `xs` except the first `n` ones. If there are less than `n` elements in `xs`, returns the empty list.
- `xs.filter(p: A => Boolean): List[A]` returns all elements from `xs` that satisfy the predicate `p` as a `List[A]`.
- `xs.flatMap[B](f: A => List[B]): List[B]` applies `f` to every element of the list `xs`, and flattens the result into a `List[B]`.
- `xs.foldLeft[B](z: B)(op: (B, A) => B): B` applies the binary operator `op` to a start value and all elements of the list, going left to right.
- `xs.foldRight[B](z: B)(op: (A, B) => B): B` applies the binary operator `op` to a start value and all elements of the list, going right to left.
- `xs.foreach[U](f: (A) => U): Unit` applies `f` to each element for its side effects.
- `xs.map[B](f: A => B): List[B]` applies `f` to every element of the list `xs` and returns a new list of type `List[B]`.
- `xs.max[A](using ord: Ordering[A]): A` finds the largest element of the list `xs`.
- `xs.min[A](using ord: Ordering[A]): A` finds the smallest element of the list `xs`.
- `xs.isEmpty`: `Boolean` returns **true** if the list has zero element, **false** otherwise.
- `xs.nonEmpty`: `Boolean` returns **true** if the list has at least one element, **false** otherwise.
- `xs.reduce[A](op: (A, A) => A): A` reduces the elements of `xs` using the specified associative binary operator.
- `xs.reduceLeft[A](op: (A, A) => A): A` applies a binary operator to all elements of `xs`, going left to right.
- `xs.reduceRight[A](op: (A, A) => A): A` applies a binary operator to all elements of `xs`, going right to left.
- `xs.reverse`: `List[A]` reverses the elements of `xs`.
- `xs.size`: `Int` returns the number of elements `xs`.
- `xs.sorted[A](using ord: Ordering[A]): List[A]` sorts `xs` according to an `Ordering`.
- `xs.take(n: Int): List[A]` returns a `List[A]` containing the first `n` elements of `xs`. If there are less than `n` elements in `xs`, returns these elements.
- `xs.zip(ys: List[B]): List[(A, B)]` zips elements of `xs` and `ys` in a pairwise fashion. If one list is longer than the other one, remaining elements are discarded. Returns a `List[(A, B)]`.

The trait `Ordering` contains a single abstract method. If `ord` is an `Ordering`, then:

- `ord.compare(x: T, y: T): Int` returns an integer whose sign communicates how `x` compares to `y`.