

Recursive Definitions & Structural Induction

In Week 40, we saw how to define sequences, like the Fibonacci numbers, recursively:

$$f_0 = 0, f_1 = 1$$

$$f_n = f_{n-1} + f_{n-2}, \quad n \geq 2$$

Basis step

Recursive step

We used strong induction to prove that $f_n > \phi^{n-2}$, for $n \geq 3$.

Now, we shall see how structures and sets can be defined recursively.

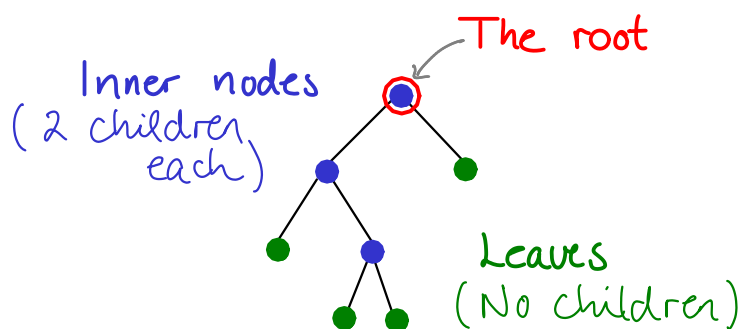
Then, we will see how structural induction can be useful for proving properties of recursively defined sets.

Full binary trees

Trees are a special kind of graphs, i.e., they consist of vertices and edges.

Full binary trees are a special kind of rooted trees in which each vertex has 0 or 2 children.

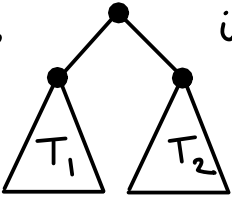
Ex:



Recursive definition of full binary trees (FBTs)

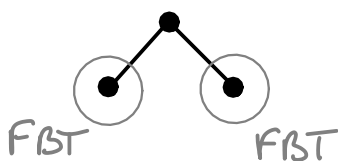
- is a FBT

(Basis step)

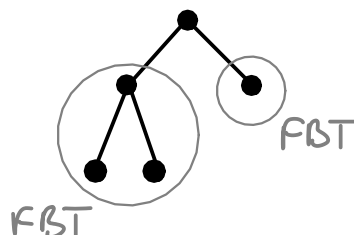
If T_1 and T_2 are FBTs,
then  is a FBT

(Recursive step)

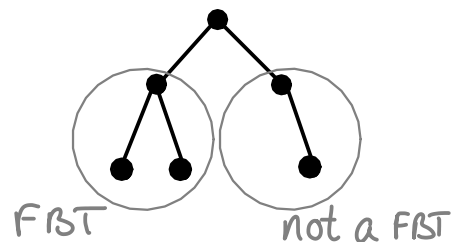
Ex:



✓



✓



✗

An alternative way of phrasing the recursive definition:

Def. 5.3.5: Full Binary Trees

$$S_1 = \{ \bullet \}$$

root

$$S_i = S_{i-1} \cup \{ T_1, T_2 \mid T_1, T_2 \in S_{i-1} \}, i \geq 2$$

left
subtree

right
subtree

Basis step

Recursive step

Thus,

$$S_2 = \{ \bullet, \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array} \}$$

$$S_3 = \{ \bullet, \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array}, \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array}, \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array}, \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array} \}$$

$$S_4 = \{ \text{---} \parallel \text{---}, \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array}, \dots \}$$

$$(|S_4| = 26)$$

Structural Induction

Proving $P(S_i)$, for all $i \geq 1$

Prove $P(S_1)$

Prove that $P(S_i) \Rightarrow P(S_{i+1})$, for all $i \geq 1$

We will use structural induction to prove a relation between the number of vertices and the height (defined on the next page) of a full binary tree.

The height, h , of a rooted tree is defined as:

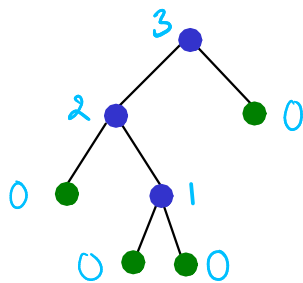
$$h(\bullet) = 0$$

$$h\left(\begin{array}{c} \bullet \\ / \quad \backslash \\ \triangle_{T_1} \quad \triangle_{T_2} \end{array}\right) = 1 + \max\{h(T_1), h(T_2)\}$$

Basis step

Recursive step

Ex:



height ($h=3$):

#edges in longest path from the root to a leaf

How many vertices can we fit in a full binary tree of a given height?

$$h=0: \quad \bullet \quad 1 = 2^0 = 2^1 - 1$$

$$h=1: \quad \begin{array}{c} \bullet \\ / \quad \backslash \\ \bullet \quad \bullet \end{array} \quad 3 = 2^0 + 2^1 = 2^2 - 1$$

$$h=2: \quad \begin{array}{c} \bullet \\ / \quad \backslash \\ \bullet \quad \bullet \\ / \quad \backslash \quad / \quad \backslash \\ \bullet \quad \bullet \quad \bullet \quad \bullet \end{array} \quad 7 = 2^0 + 2^1 + 2^2 = 2^3 - 1$$

$n(T)$: # vertices in T

Theorem 5.3.2

For any full binary tree, T , $n(T) \leq 2^{h(T)+1} - 1$

Ex:

$$\left. \begin{array}{c} \text{Diagram of a full binary tree with 3 nodes (root and two children)} \\ n=3 \\ h=1 \end{array} \right\} \Rightarrow 2^{h+1} - 1 = 3 = n$$

$$\left. \begin{array}{c} \text{Diagram of a full binary tree with 5 nodes (root, two children, and two grandchildren)} \\ n=5 \\ h=2 \end{array} \right\} \Rightarrow 2^{h+1} - 1 = 7 > n$$

Proof by structural induction:

Basis: $S_1 = \{\bullet\}$

$$\left. \begin{array}{c} n=1 \\ h=0 \end{array} \right\} \Rightarrow 2^{h+1} - 1 = 2 - 1 = 1 = n$$

Induction hypothesis ($i \geq 1$):

$$n(T) \leq 2^{h(T)+1} - 1, \text{ for all } T \in \mathcal{S}_i$$

Inductive step ($i \geq 1$):

Let $T \in \mathcal{S}_{i+1} - \mathcal{S}_i$:

$$\left. \begin{array}{c} h(T_1) \left\{ \begin{array}{c} \text{Diagram of } T \\ \text{with root and children } T_1, T_2 \end{array} \right\} h(T_2) \end{array} \right\} h(T) = 1 + \max\{h(T_1), h(T_2)\},$$

by Def. 5.3.6

$$T_1, T_2 \in \mathcal{S}_i, \text{ by Def. 5.3.5}$$

$$\begin{aligned} n(T) &= 1 + n(T_1) + n(T_2) \\ &\leq 1 + (2^{h(T_1)+1} - 1) + (2^{h(T_2)+1} - 1), \text{ by the ind. hyp.} \\ &= 2^{h(T_1)+1} + 2^{h(T_2)+1} - 1 \\ &\leq 2^{h(T)} + 2^{h(T)} - 1, \text{ by Def. 5.3.6} \\ &= 2 \cdot 2^{h(T)} - 1 \\ &= 2^{h(T)+1} - 1 \end{aligned}$$

□

Sequences (Section 2.4)

Ex: (Fibonacci) $0, 1, 1, 2, 3, 5, 8, \dots$

$$f_0 = 0, f_1 = 1$$

$$f_{n+1} = f_n + f_{n-1}, \quad n \geq 2$$

Sequence (Følge)

Ordered set or

Function from (a subset of) \mathbb{N}

Ex: (Harmonic) $1, \frac{1}{2}, \frac{1}{3}, \frac{1}{4}, \dots$

$$f: \mathbb{N}^+ \rightarrow \mathbb{Q}^+, \quad f(n) = \frac{1}{n}$$

or

$$a_n = \frac{1}{n}, \quad n \geq 1$$

$$\{a_n\} = 1, \frac{1}{2}, \frac{1}{3}, \frac{1}{4}, \dots$$

↖ Conflict with set notation

Arithmetic Progressions

(Equidistant terms)

Ex: 1, 3, 5, 7, ...

$$a_n = 1 + 2n, \quad n \in \mathbb{N}$$

or

$$a_0 = 1$$

$$a_n = a_{n-1} + 2, \quad n \geq 1$$

Arithmetic progression

$a, a+d, a+2d, \dots$

or

$$a_n = a + nd, \quad n \in \mathbb{N}$$


Initial term Common difference

Ex. above: $a=1, d=2$

Geometric Progression

(Exponentially increasing or decreasing terms)

Ex: $1, \frac{1}{2}, \frac{1}{4}, \frac{1}{8}, \dots$

$$a_n = \frac{1}{2^n}, \quad n \in \mathbb{N}$$

or

$$a_0 = 1$$

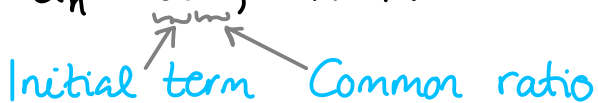
$$a_n = \frac{1}{2} \cdot a_{n-1}, \quad n \geq 1$$

Geometric Progression

a, ar, ar^2, \dots

or

$$a_n = ar^n, \quad n \in \mathbb{N}$$


Initial term Common ratio

Ex above: $a=1, r=\frac{1}{2}$