Optimal Rates of (Locally) Differentially Private Heavy-tailed Multi-Armed Bandits

Youming Tao *1 Yulian Wu *2 Peng Zhao 3 Di Wang 2

Abstract

In this paper we investigate the problem of stochastic multi-armed bandits (MAB) in the (local) differential privacy (DP/LDP) model. Unlike previous results that assume bounded/sub-Gaussian reward distributions, we focus on the setting where each arm's reward distribution only has (1+v)-th moment with some $v \in (0,1]$. In the first part, we study the problem in the central ϵ -DP model. We first provide a near-optimal result by developing a private and robust Upper Confidence Bound (UCB) algorithm. Then, we improve the result via a private and robust version of the Successive Elimination (SE) algorithm. Finally, we establish the lower bound to show that the instance-dependent regret of our improved algorithm is optimal. In the second part, we study the problem in the ϵ -LDP model. We propose an algorithm that can be seen as locally private and robust version of SE algorithm, which provably achieves (near) optimal rates for both instancedependent and instance-independent regret. Our results reveal differences between the problem of private MAB with bounded/sub-Gaussian rewards and heavy-tailed rewards. To achieve these (near) optimal rates, we develop several new hard instances and private robust estimators as byproducts, which might be used to other related problems.

1. INTRODUCTION

Multi-Armed Bandits (MAB), and its general form, bandit learning, have already been studied for more than half

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a century, starting from (Thompson, 1933) and (Robbins, 1952). They find numerous applications in many areas such as medicine (Gutiérrez et al., 2017), finance (Shen et al., 2015), social science (Nakayama et al., 2017), and clinical research (Press, 2009). However, due to the existence of sensitive data and their distributed nature in many applications, it is often challenging to preserve the privacy of such data.

To preserve the privacy of these sensitive data, Differential Privacy (DP) (Dwork et al., 2006) has received a great deal of attention and now has established itself as a de facto notion of privacy for data analysis. Over the past decade, differentially private bandit learning has been extensively studied from various setups including classical stochastic MAB (Mishra & Thakurta, 2015; Tossou & Dimitrakakis, 2016; Sajed & Sheffet, 2019; Ren et al., 2020; Kalogerias et al., 2020), combinatorial semi-bandits (Chen et al., 2020), and contextual bandits (Shariff & Sheffet, 2018; Hannun et al., 2019; Malekzadeh et al., 2020; Zheng et al., 2020).

However, these problems are still not well-understood. For example, all of the previous results and methods need to assume that the rewards are sampled from some bounded (or sub-Gaussian) distributions to guarantee the DP property. However, such assumptions may not hold when designing decision-making algorithms for complicated realworld systems. In particular, previous papers have shown that the rewards or the interactions in such systems often lead to heavy-tailed and power law distributions (Dubey & Pentland, 2019), such as modeling stock prices (Bradley & Taqqu, 2003), preferential attachment in social networks (Mahanti et al., 2013), and online behavior on websites (Kumar & Tomkins, 2010). Thus, it is necessary to develop new methods to deal with these heavy-tailed rewards in the private bandit learning.

To address the above issue, in this paper, we focus on the most fundamental bandit model, *i.e.*, multi-armed bandits, with heavy-tailed rewards. We conduct a comprehensive and the first study on MAB with heavy-tailed rewards in both central and local DP models. Our contributions are summarized as follows.

• In the first part (Section 2), we consider the problem in the central ϵ -DP model. Specifically, we first pro-

^{*}Equal contribution ¹School of Computer Science, Shandong University, Qingdao, China ²Division of Computer, Electrical and Mathematical Sciences and Engineering, King Abdullah University of Science and Technology, Jeddah, Saudi Arabia ³School of Artificial Intelligence, Nanjing University, Nanjing, China. Correspondence to: Di Wang <di.wang@kaust.edu.sa>.

pose a method based on a robust version of the Upper Confidence Bound (UCB) algorithm, and also design a new mechanism that could be seen as an adaptive version of the Tree-based mechanism (Dwork et al., 2010). To further improve the result, we then develop a private and robust version of the Successive Elimination (SE) algorithm and show that the (expected) regret bound is improved by a factor of $\log^{1.5+\frac{1}{v}} T$, where T is the number of rounds. Moreover, we establish the lower bound and show that the instance-dependent regret bound of $O\left(\frac{\log T}{\epsilon}\sum_{\Delta_a>0}\left(\frac{1}{\Delta_a}\right)^{\frac{1}{v}}+\max_a\Delta_a\right)$ achieved by our second algorithm is optimal (up to poly($\log \log \frac{1}{\Delta_a}$) factors), where Δ_a is the mean gap.

- In the second part (Section 3), we study the problem in the ϵ -LDP model. We first develop a LDP version of the SE algorithm which achieves an instance-dependent regret bound $\begin{array}{ll} \text{of} & O\left(\frac{\log T}{\epsilon^2}\sum_{\Delta_a>0}\left(\frac{1}{\Delta_a}\right)^{\frac{1}{v}}+\max_a\Delta_a\right) \quad \text{and} \quad \text{an} \\ \widetilde{O}\left(\left(\frac{K}{\epsilon^2}\right)^{\frac{v}{1+v}}T^{\frac{1}{1+v}}\right) \quad \text{instance-independent} \quad \text{bound.} \end{array}$ Then, we show that the above instance-dependent regret bound is optimal and the instance-independent regret bound is near-optimal (up to poly(log T)factors).
- · All of our results also reveal the differences between the problem of private MAB with bounded/sub-Guassian rewards and that with heavy-tailed rewards. To achieve these (near) optimal results, we develop several new hard instances, mechanisms and private robust estimators as byproducts, which could be used to other related problems, such as private contextual bandits (Shariff & Sheffet, 2018) or private reinforcement learning (Vietri et al., 2020).

2. DP HEAVY-TAILED MAB

In the classical setting where the rewards follow some bounded distributions, the most commonly used approach is using the Tree-based mechanism to privately calculate the sum of rewards and then modify the Upper Confidence Bound (UCB) algorithm (Auer et al., 2002), such as (Mishra & Thakurta, 2015; Tossou & Dimitrakakis, 2016). However, their methods cannot be directly generalized to the heavytailed setting, since now the reward is unbounded. Thus, the most natural idea is to first preprocess the rewards to make them bounded and then use the Tree-based mechanism and UCB algorithm, see Algorithm 1 for details.

Lemma 1 ((Adaptive) Tree-based Mechanism). Given a stream σ such that $\sigma(t) \in [-B_t, B_t]$ for $\forall t \in [T]$, where B_t is non-decreasing with t, we want to privately and continually release the sum of the stream $S(t) \triangleq \sum_{i=1}^{t} \sigma(i)$ for each $t \in [T]$. Tree-based Mechanism (Algorithm 2) outputs

Algorithm 1 DP Robust Upper Confidence Bound

Input: time horizon T, parameters ϵ, v, u .

- 1: Create an empty tree $tree_a$ for each arm $a \in [K]$.
- 2: Initialize pull number $n_a \leftarrow 0$ for each arm $a \in [K]$. 3: Denote B_n as $(\frac{\epsilon un}{\log^{1.5}T})^{1/(1+v)}$ for any $n \in \mathbb{N}^+$. 4: **for** $t = 1, \dots, K$ **do**
- Pull arm t and observe a reward x_t .
- 6: Update the pull number $n_t \leftarrow n_t + 1$.
- Truncate the reward by $\widetilde{x}_t \leftarrow x_t \cdot \mathbb{I}_{|x_t| < B_n}$. 7:
- 8: Insert \tilde{x}_t into $tree_t$.
- 9: end for
- 10: **for** t = K + 1, ..., T **do**
- Obtain $\widehat{S}_a(t)$ for each $a \in [K]$ via Algorithm 2. 11:

$$a_t = \argmax_a \frac{\widehat{S}_a(t)}{n_a} + 18u^{\frac{1}{1+v}} (\frac{\log(2t^4)\log^{1.5 + \frac{1}{v}}T}{n_a \epsilon})^{\frac{v}{1+v}}$$

and observe the reward x_t .

- 13: Update the pull number $n_{a_t} \leftarrow n_{a_t} + 1$.
- Truncate the reward by $\widetilde{x}_t \leftarrow x_t \cdot \mathbb{I}_{|x_t| \leq B_{n_0}}$. 14:
- 15: Insert \widetilde{x}_t into $tree_{a_t}$.
- 16: **end for**

Algorithm 2 (Adaptive) Tree-based Mechanism

Input: time horizon T, privacy budget ϵ , a stream σ . **Output:** A private version $\hat{S}(t)$ for $S(t) = \sum_{i=1}^{t} \sigma(i)$ at

each $t \in [T]$ 1: Initialize each α_i and noisy $\widehat{\alpha}_i$ to 0.

- 2: $\epsilon' \leftarrow \epsilon / \log T$.
- 3: **for** t = 1, ..., T **do**
- Express t in binary form: $t = \sum_{j} \text{Bin}_{j}(t) \cdot 2^{j}$.
- $i \leftarrow \min\{j : \operatorname{Bin}_j(t) \neq 0\}.$
- 6:
- $\alpha_i \leftarrow \sum_{j < i} \alpha_j + \sigma(t).$ for $j = 0, \dots, i-1$ do 7:
- $\alpha_j \leftarrow 0, \, \widehat{\alpha}_j \leftarrow 0.$ 8:
- 9:
- 10:
- $\begin{array}{l} \widehat{\alpha}_i \leftarrow \alpha_i + \operatorname{Lap}(2B_t/\epsilon'). \\ \widehat{S}(t) \leftarrow \sum_{j: \operatorname{Bin}_j(t) = 1} \widehat{\alpha}_j. \end{array}$ 11:
- 12: **end for**

an estimation $\widehat{S}(t)$ for S(t) at each $t \in [T]$ such that $\widehat{S}(t)$ preserves ϵ -differential privacy and guarantees the following noise bound with probability at least $1 - \delta$ for any $\delta > 0$,

$$\left| \widehat{S}(t) - S(t) \right| \le \frac{2B_t}{\epsilon} \cdot \log^{1.5} T \cdot \log \frac{1}{\delta}. \tag{1}$$

When $B_t = B$, Algorithm 2 will be the same as the original one. Theorem 1 presents the privacy guarantee of overall algorithm (Algorithm 1 and Algorithm 2).

Theorem 1. For any $\epsilon > 0$, the overall algorithm (Algorithm 1 and Algorithm 2) is ϵ -differentially private.

In fact, the $\widehat{S}_a(t)/n_a$ term in step 12, which is denoted by $\widehat{\mu}_a(n_a,t)$, could be seen as a robust and private estimator of the mean μ_a after total n_a pulls of arm a till time t. Our selection strategy in step 12 is based on the following estimation error between $\widehat{\mu}_a(n_a,t)$ and μ_a , which is also a key lemma that will be used to bound the regret of Algorithm 1.

Lemma 2. In Algorithm 1, for a fixed arm a and t, we have the following estimation error with probability at least $1-t^{-4}$,

$$\widehat{\mu}_{a}(n_{a}, t) \leq \mu_{a} + 18u^{\frac{1}{1+v}} \left(\frac{\log(2t^{4}) \log^{1.5 + \frac{1}{v}} T}{n_{a} \epsilon} \right)^{\frac{v}{1+v}}.$$
(2)

We have the following instance-dependent regret bound by the proof of Theorem 1 in (Auer et al., 2002).

Theorem 2. Under our assumptions, for any $0 < \epsilon \le 1$ the instance-dependent expected regret of Algorithm 1 satisfies

$$\mathcal{R}_T \le O\left(\sum_{a:\Delta_a>0} \left(\frac{\log^{2.5 + \frac{1}{v}} T}{\epsilon} \left(\frac{u}{\Delta_a}\right)^{\frac{1}{v}} + \Delta_a\right)\right). \quad (3)$$

Compared with the $O(\sum_{a:\Delta_a>0} [\log T(\frac{u}{\Delta_a})^{\frac{1}{v}} + \Delta_a])$ optimal rate of the regret in the non-private version (Bubeck et al., 2013), we can see that there is an additional factor of $\frac{\log^{1.5+\frac{1}{v}}T}{\epsilon}$ in the private case. Thus, a natural question arises here is whether it is possible to further improve the regret. We answer this question affirmatively by designing an optimal algorithm, see Algorithm 3 for details.

Theorem 3. For any $\epsilon > 0$, Algorithm 3 is ϵ -differentially private.

Theorem 4 (DP Upper Bound). If we set $\beta = \frac{1}{T}$ in Algorithm 3, then for sufficiently large T and any $\epsilon \in (0, 1]$, the instance-dependent expected regret of Algorithm 3 satisfies

$$\mathcal{R}_T \le O\left(\frac{u^{\frac{1}{1+v}}\log T}{\epsilon} \sum_{\Delta_a > 0} \left(\frac{1}{\Delta_a}\right)^{\frac{1}{v}} + \max_a \Delta_a\right). \tag{4}$$

Moreover, the instance-independent expected regret of Algorithm 3 satisfies

$$\mathcal{R}_T \le O\left(u^{\frac{v}{(1+v)^2}} \left(\frac{K \log T}{\epsilon}\right)^{\frac{v}{1+v}} T^{\frac{1}{1+v}}\right), \quad (5)$$

where the $O(\cdot)$ -notation omits $\log \log \frac{1}{\Delta_a}$ terms.

From Theorem 4 we can see that compared with the regret bound $O(\frac{\log^{2.5}T}{\epsilon}\sum_{\Delta_a>0}(\frac{1}{\Delta_a})^{\frac{1}{v}})$ in Theorem 2, we achieve an improved bound of $O(\frac{\log T}{\epsilon}\sum_{\Delta_a>0}(\frac{1}{\Delta_a})^{\frac{1}{v}})$. We show the instance-dependent regret bound presented in Theorem 4 is **optimal**. The lower bound of the instance-independent regret is still unclear, and we leave it as an open problem.

Algorithm 3 DP Robust Successive Elimination

```
Input: confidence \beta, parameters \epsilon, v, u.
  1: S \leftarrow \{1, \cdots, K\}
  2: Initialize: t \leftarrow 0, \tau \leftarrow 0.
  3: repeat
  4:
               \tau \leftarrow \tau + 1.
              Set \bar{\mu}_a = 0 for all a \in \mathcal{S}.
              \begin{split} & r \leftarrow 0, \, D_{\tau} \leftarrow 2^{-\tau}, \\ & R_{\tau} \leftarrow \left[ u^{\frac{1}{v}} \left( \frac{24^{(1+v)/v} \log(4|\mathcal{S}|\tau^{2}/\beta)}{\epsilon D_{\tau}^{(1+v)/v}} \right) + 1 \right], \\ & B_{\tau} \leftarrow \left( \frac{uR_{\tau}\epsilon}{\log(4|\mathcal{S}|\tau^{2}/\beta)} \right)^{1/(1+v)}. \end{split}
  8:
               while r < R_{\tau} do
  9:
10:
                    r \leftarrow r + 1.
                    for a \in \mathcal{S} do
11:
12:
                          t \leftarrow t + 1.
                          Sample a reward x_{a,r}.
13:
14:
                          \widetilde{x}_{a,r} \leftarrow x_{a,r} \cdot \mathbb{I}_{\{|x_{a,r}| \leq B_{\tau}\}}
15:
16:
               For each a \in \mathcal{S}, compute \bar{\mu}_a \leftarrow (\sum_{l=1}^{R_{\tau}} \tilde{x}_{a,l})/R_{\tau}.
17:
               Set \widetilde{\mu}_a \leftarrow \overline{\mu}_a + \operatorname{Lap}(\frac{2B_{\tau}}{R_{-\epsilon}}) for all a \in \mathcal{S}.
18:
              \widetilde{\mu}_{\max} \leftarrow \max_{a \in \mathcal{S}} \widetilde{\mu}_{a}.
err_{\tau} \leftarrow u^{1/(1+v)} \left(\frac{\log(4|\mathcal{S}|\tau^{2}/\beta)}{R_{\tau}\epsilon}\right)^{v/(1+v)}.
19:
20:
21:
               for all viable arm a do
22:
                     if \widetilde{\mu}_{\max} - \widetilde{\mu}_a > 12err_{\tau} then
23:
                           Remove arm a from S.
24:
                     end if
               end for
25:
26: until |S| = 1
27: Pull the arm in S in all remaining T - t rounds.
```

Theorem 5 (DP Instance-dependent Lower Bound). There exists a heavy-tailed K-armed bandit instance with $u \leq 1$, $\mu_a \leq \frac{1}{6}$ and $\Delta_a \in (0,\frac{1}{12})$, such that for any ϵ -DP ($0 < \epsilon \leq 1$) algorithm $\mathcal A$ whose expected regret is at most $T^{\frac{3}{4}}$, we have

$$\mathcal{R}_T \ge \Omega \left(\frac{\log T}{\epsilon} \sum_{\Delta_a > 0} \left(\frac{1}{\Delta_a} \right)^{\frac{1}{v}} \right). \tag{6}$$

Remark 1. In the MAB with bounded rewards case, it has been shown that the optimal rate of the expected rate is $O(\frac{K\log T}{\epsilon} + \sum_{\Delta_a>0} \frac{\log T}{\Delta_a})$ (Sajed & Sheffet, 2019). Compared with the optimal rate $O(\frac{\log T}{\epsilon} \sum_{\Delta_a>0} (\frac{1}{\Delta_a})^{\frac{1}{v}})$ in the heavy-tailed case, we can see there is a huge difference. First, the dependency on $\frac{1}{\Delta_a}$ now becomes to $(\frac{1}{\Delta_a})^{\frac{1}{v}}$. Secondly, the price of privacy in the bounded rewards case is an additional term of $O(\frac{K\log T}{\epsilon})$ compared with the non-private rate, while in the heavy-tailed case, there is an additional factor of $\frac{1}{\epsilon}$ compared with the non-private one.

3. LDP HEAVY-TAILED MAB

Unlike the Tree-based mechanism in the central model, the LDP version of the UCB algorithm will introduce a huge amount of error to estimate the mean. To achieve a better utility, we propose an ϵ -LDP version of the SE algorithm, see Algorithm 4 for details. The following theorems provide the privacy and utility guarantees for Algorithm 4.

Theorem 6. For any $\epsilon > 0$, Algorithm 4 is ϵ -local differentially private.

Theorem 7 (LDP Upper Bound). Set $\beta=\frac{1}{T}$ in Algorithm 4. For any $\epsilon\in(0,1]$ and sufficiently large T, the instance-dependent expected regret of Algorithm 4 satisfies

$$\mathcal{R}_T \le O\left(\frac{u^{\frac{2}{v}}\log T}{\epsilon^2} \sum_{\Delta_a > 0} \left(\frac{1}{\Delta_a}\right)^{\frac{1}{v}} + \max_a \Delta_a\right). \quad (7)$$

Moreover, the instance-independent expected regret of Algorithm 4 satisfies

$$\mathcal{R}_T \le O\left(u^{\frac{2}{1+v}} \left(\frac{K \log T}{\epsilon^2}\right)^{\frac{v}{1+v}} T^{\frac{1}{1+v}}\right), \tag{8}$$

where the $O(\cdot)$ -notations omit $\log \log \frac{1}{\Delta_a}$ terms.

In the following, we derive both instance-dependent and instance-independent lower bounds for heavy-tailed MAB in the ϵ -LDP model.

Theorem 8 (LDP Instance-dependent Lower Bound). There exists a heavy-tailed K-armed bandit instance with $u \le 1$ and $\Delta_a \triangleq \mu_1 - \mu_a \in (0, \frac{1}{5})$, such that for any ϵ -LDP $(0 < \epsilon \le 1)$ algorithm whose regret $\le o(T^\alpha)$ for any $\alpha > 0$, the regret satisfies

$$\liminf_{T \to \infty} \frac{\mathcal{R}_T}{\log T} \ge \Omega\left(\frac{1}{\epsilon^2} \sum_{\Delta_a > 0} (\frac{1}{\Delta_a})^{\frac{1}{v}}\right).$$

Remark 2. The attained bound in Theorem 7 is optimal. Compared with the optimal rate $O(\frac{1}{\Delta^{\frac{1}{v}}})$ in the non-private case, we can see the price of privacy is an additional factor of $\frac{1}{\epsilon^2}$, which is similar to other MAB with bounded/sub-Gaussian rewards problems in the LDP model (Zhou & Tan, 2021; Ren et al., 2020).

Theorem 9 (LDP Instance-independent Lower Bound). There exists a heavy-tailed K-armed bandit instance with the (1+v)-th bounded moment of each reward distribution is bounded by 1. Moreover, if T is large enough, for any the ϵ -LDP algorithm $\mathcal A$ with $\epsilon \in (0,1]$, the expected regret must satisfy

$$\mathcal{R}_T \ge \Omega\left(\left(\frac{K}{\epsilon^2}\right)^{\frac{v}{1+v}} T^{\frac{1}{1+v}}\right).$$

Algorithm 4 LDP Robust Successive Elimination

```
Input: Confidence \beta, parameters \epsilon, v, u.
  1: \mathcal{S} \leftarrow \{1, \cdots, K\}
  2: Initialize: t \leftarrow 0, \tau \leftarrow 0.
  3: repeat
  4:
                \tau \leftarrow \tau + 1.
                Set \bar{\mu}_a = 0 for all a \in \mathcal{S}.
  5:
              r \leftarrow 0, D_{\tau} \leftarrow 4^{-\tau}.
R_{\tau} \leftarrow \left[ u^{\frac{2}{v}} \left( \frac{28^{2(1+v)/v} \log(8|\mathcal{S}|\tau^{2}/\beta)}{\epsilon^{2} D_{\tau}^{2(1+v)/v}} \right) + \log\left(\frac{8|\mathcal{S}|\tau^{2}}{\beta}\right) \right].
B_{\tau} \leftarrow \left( \frac{u\sqrt{R_{\tau}\epsilon}}{\sqrt{\log(8|\mathcal{S}|\tau^{2}/\beta)}} \right)^{1/(1+v)}.
  8:
  9:
                while r < R_{\tau} do
10:
                      r \leftarrow r + 1.
                      for a \in \mathcal{S} do
11:
                            t \leftarrow t + 1.
12:
                            Sample a reward x_{a,r} for each arm a \in \mathcal{S}.
13:
                            \widetilde{x}_{a,r} \leftarrow x_{a,r} \cdot \mathbb{I}_{\{|x_{a,r}| \leq B_{\tau}\}}.
\widehat{x}_{a,r} \leftarrow \widetilde{x}_{a,r} + \operatorname{Lap}(\frac{2B_{\tau}}{\epsilon})
14:
15:
                      end for
16:
17:
                end while
                For each a \in \mathcal{S}, compute \bar{\mu}_a \leftarrow (\sum_{l=1}^{R_{\tau}} \widehat{x}_{a,l})/R_{\tau}.
18:
               \begin{split} \bar{\mu}_{\max} &\leftarrow \max_{a \in \mathcal{S}} \bar{\mu}_i. \\ err_{\tau} &\leftarrow u^{1/(1+v)} \big(\frac{\sqrt{\log(8|\mathcal{S}|\tau^2/\beta)}}{R_{\tau\epsilon}}\big)^{v/(1+v)}. \\ \text{for all viable arm } a \text{ do} \end{split}
19:
20:
21:
22:
                      if \widetilde{\mu}_{\max} - \widetilde{\mu}_a > 14err_{\tau} then
                            Remove arm a from S.
23:
24:
                      end if
                end for
25:
26: until |S| = 1
27: Pull the arm in S in all remaining T - t rounds.
```

Remark 3. From Theorem 9, we can see the upper bound (8) of Algorithm 4 is nearly optimal. However, compared with instance-independent lower bound, there is still a poly $(\log T)$ factor gap. We conjecture this factor could be removed by using some more advanced robust estimator, such as the estimator in (Lee et al., 2020) and we will leave it as an open problem. For MAB with bounded rewards in the LDP model, (Basu et al., 2019) shows that its instance-dependent regret bound is always at least $\Omega(\frac{\sqrt{KT}}{\epsilon})$, *i.e.*, there is an additional factor of $\frac{1}{\epsilon}$ compared with the nonprivate case. However, for heavy-tailed MAB, compared with the lower bound of $\Omega(K^{\frac{v}{1+v}}T^{\frac{1}{1+v}})$ in the non-private case, from Theorem 9 we can observe that the difference is a factor of $(\frac{1}{\epsilon^2})^{\frac{v}{1+v}}$. Thus, combining with Remark 1, we can conclude that heavy-tailed MAB and bounded MAB are quite different in both central and local differential privacy models.

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