CME338 Final Project LSLQ Solver

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Abstract 1

2 Introduction

Notation

$$\mathcal{K}_k = \mathcal{K}_k(A^T A, b) = \operatorname{span}(b, A^T A b, \dots, (A^T A)^{k-1} b).$$

Derivation of LSLQ 3

In this section, we derive the short recurrence formulas for LSLQ beginning from the Golub-Kahan process.

3.1Golub-Kahan Process

The Golub-Kahan process is defined by the recurrence defined in Algorithm 1.

Algorithm 1 Golub-Kahan Process

Set
$$\beta_1 u_1 = b$$
, $\alpha_1 v_1 = A^T u_1$
for $k = 1, 2, ...$ do
 $\beta_{k+1} u_{k+1} = A v_k - \alpha_k u_k$
 $\alpha_{k+1} v_{k+1} = A^T u_{k+1} - \beta_{k+1} v_k$
end for

After performing k iterations of this process, we obtain the decompositions

$$AV_k = U_{k+1}B_k \tag{1}$$

$$AV_k = U_{k+1}B_k
 A^T U_{k+1} = V_{k+1}L_{k+1}^T.
 (2)$$

where $U_k = (u_1 | \dots | u_k), V_k = (v_1 | \dots | v_k),$ and

$$B_k = \begin{pmatrix} \alpha_1 \\ \beta_2 & \alpha_2 \\ & \ddots & \ddots \\ & & \beta_k & \alpha_k \\ & & & \beta_{k+1} \end{pmatrix} \qquad L_{k+1} = (B_{k+1} | \alpha_{k+1} e_{k+1}).$$

We can then observe that V_k is a basis for $\mathcal{K}_k(A^TA, b)$, since

$$A^{T}AV_{k} = A^{T}U_{k+1}B_{k} = V_{k+1}L_{k+1}^{T}B_{k} = V_{k+1}\begin{pmatrix} B_{k}^{T}B_{k} \\ \alpha_{k+1}\beta_{k+1}e_{k}^{T} \end{pmatrix}.$$

3.2 The LSLQ subproblem

In order to solve the linear system or least squares problem Ax = b, we instead solve the normal equations $A^TA = A^Tb$. Since A^TA is symmetric positive semidefinite and we assume that the normal equations are consistent, we may consider applying SYMMLQ to this system. We solve the normal equations iteratively, where at iteration k we solve the problem

$$x_k = \underset{x \in \mathcal{K}_k}{\operatorname{arg \, min}} \quad \|x\|_2$$

s.t. $A^T r \perp \mathcal{K}_{k-1}$. (3)

where r = b - Ax.

In order to solve this problem, we first note that we may formulate this as an unconstrained problem in a smaller space if we minimize in y_k and set $x_k = V_k y_k$. Then

$$0 = V_{k-1}^{T} A^{T} r_{k} = V_{k-1}^{T} A^{T} (b - Ax_{k})$$

$$= V_{k-1}^{T} A^{T} b - V_{k-1}^{T} A^{T} A V_{k} y_{k}$$

$$= \alpha_{1} \beta_{1} e_{1} - B_{k-1}^{T} U_{k}^{T} A V_{k} y_{k}$$

$$= \alpha_{1} \beta_{1} e_{1} - B_{k-1}^{T} L_{k} y_{k}.$$

Thus in order to solve 3, we can solve

$$y_k = \underset{y \in \mathbb{R}^k}{\operatorname{arg \, min}} \quad ||y||_2$$

s.t.
$$B_{k-1}^T L_k y = \alpha_1 \beta_1 e_1.$$
 (4)

3.3 First QR decomposition

We first take the QR factorization of $Q_k R_k = B_{k-1}$. Suppose we have the QR factorization of $Q_{k-1}R_{k-1} = B_{k-2}$, with

$$R_{k-1} = \begin{pmatrix} \rho_1 & \theta_2 & & & \\ & \rho_2 & \ddots & & \\ & & \ddots & \theta_{k-1} \\ & & & \rho_{k-2} \end{pmatrix}.$$

Then we may recurse to obtain the factorization of B_{k-1}

by defining $Q_k = Q_{k-1}G_k^{(1)}$, where $G_k^{(1)}$ is the Givens rotation

$$G_k^{(1)} = \begin{pmatrix} c_1 & s_1 \\ -s_1 & c_1 \end{pmatrix}.$$

Using $G_{k-1}^{(2)}$ defined in the previous iteration, we have

$$\begin{pmatrix} \theta_{k-1} \\ \hat{\rho}_{k-1} \end{pmatrix} = \begin{pmatrix} c_1^{(k-1)} & -s_1^{(k-1)} \\ s_1^{(k-1)} & c_1^{(k-1)} \end{pmatrix} \begin{pmatrix} 0 \\ \alpha_{k-1} \end{pmatrix},$$

$$\begin{pmatrix} \theta_k \\ \hat{\rho}_k \end{pmatrix} = \begin{pmatrix} c_1 & -s_1 \\ s_1 & c_1 \end{pmatrix} \begin{pmatrix} 0 \\ \alpha_k \end{pmatrix},$$

$$\begin{pmatrix} \rho_{k-1} \\ 0 \end{pmatrix} = \begin{pmatrix} c_1 & -s_1 \\ s_1 & c_1 \end{pmatrix} \begin{pmatrix} \hat{\rho}_{k-1} \\ \beta_k \end{pmatrix},$$

and we therefore obtain the recurrences

$$\hat{\rho}_{k-1} = \frac{\alpha_{k-1}\rho_{k-2}}{\rho_{k-2}} = c_1^{(k-1)}\alpha_{k-1}, \tag{5}$$

$$\rho_{k-1} = \sqrt{\hat{\rho}_{k-1}^2 + \beta_k^2},$$

$$c_1 = \hat{\rho}_{k-1}/\rho_{k-1},$$

$$s_1 = -\beta_k/\rho_{k-1},$$
(8)

$$c_1 = \hat{\rho}_{k-1}/\rho_{k-1},\tag{7}$$

$$s_1 = -\beta_k/\rho_{k-1}, \tag{8}$$

$$\theta_k = \frac{\alpha_k \beta_k}{\rho_{k-1}} = -s_1 \alpha_k. \tag{9}$$

3.4 Forward Substitution

With the previous QR decomposition, the system we intend to solve becomes

$$\begin{pmatrix} \alpha_1 \beta_1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \begin{pmatrix} B_{k-1}^T B_{k-1} & 0 \\ \vdots \\ 0 \\ \alpha_k \beta_k \end{pmatrix} y_k = \begin{pmatrix} R_k^T R_k & \vdots \\ 0 \\ \alpha_k \beta_k \end{pmatrix} y_k$$
$$= \begin{pmatrix} R_k^T R_k & \vdots \\ 0 \\ \alpha_k \beta_k \end{pmatrix} y_k$$
$$= \begin{pmatrix} R_k^T R_k & 0 \\ \vdots \\ 0 \\ \alpha_k \beta_k \end{pmatrix} y_k$$
$$= \begin{pmatrix} R_k^T R_k & 0 \\ \vdots \\ 0 \\ \beta_{k-1} \end{pmatrix} y_k$$

Define

$$z_{k} = \begin{pmatrix} \zeta_{2} \\ \vdots \\ \zeta_{k} \end{pmatrix} = \begin{pmatrix} R_{k} & 0 \\ \vdots \\ 0 \\ \theta_{k} \end{pmatrix} y_{k} = \tilde{R}_{k} y_{k}$$
 (10)

so that we have $R_k^T z_k = \alpha_1 \beta_1 e_1$. As in the Conjugate Gradient method, we obtain a short recurrence for ζ_k ,

$$\zeta_k = -\frac{\theta_{k-1}}{\rho_{k-1}} \zeta_{k-1}. \tag{11}$$

3.5 Second QR decomposition

Using the recurrence of the previous section, we now need to solve the minimum norm problem

$$\tilde{R}_k y_k = z_k.$$

We accomplish this by taking the QR decomposition of $\hat{Q}_k\hat{R}_k = \tilde{R}_k^T$. Suppose we have the QR decomposition from the previous iteration, $\hat{Q}_{k-1}\hat{R}_{k-1} = \tilde{R}_{k-1}^T$, with

$$\hat{R}_{k-1} = \begin{pmatrix} \sigma_1 & \eta_2 & & \\ & \sigma_2 & \ddots & \\ & & \ddots & \eta_{k-1} \\ & & & \sigma_{k-2} \end{pmatrix}.$$

Then as was done in the first QR decomposition, we can recurse to obtain a fast update for the second QR decomposition.

$$\tilde{R}_{k}^{T} = \begin{pmatrix} \rho_{1} & & & & 0 \\ \theta_{2} & \rho_{2} & & & \vdots \\ & \ddots & \ddots & & \vdots \\ & \theta_{k-2} & \rho_{k-2} & 0 \\ & & \theta_{k-1} & \rho_{k-1} \\ \hline 0 & \cdots & \cdots & 0 & \theta_{k} \end{pmatrix} = \hat{Q}_{k} \begin{pmatrix} \sigma_{1} & \eta_{2} & & 0 \\ & \sigma_{2} & \ddots & & \vdots \\ & & \ddots & \eta_{k-2} & \vdots \\ & & \sigma_{k-2} & \eta_{k-1} \\ \hline 0 & \cdots & \cdots & 0 & \theta_{k} \end{pmatrix}$$

$$= \hat{Q}_{k}G_{k}^{(2)} \begin{pmatrix} \sigma_{1} & \eta_{2} & & 0 \\ & \sigma_{2} & \ddots & & \vdots \\ & & \ddots & \eta_{k-2} & \vdots \\ & & & \sigma_{k-2} & \eta_{k-1} \\ \hline & & & & \sigma_{k-2} & \eta_{k-1} \\ \hline & & & & & 0 & \sigma_{k-1} \\ \hline & & & & & & 0 \end{pmatrix}$$

We define $\hat{Q}_k = \hat{Q}_{k-1}G_k^{(2)}$, where $G_k^{(2)}$ is the Givens rotation

$$G_k^{(2)} = \begin{pmatrix} c_2 & s_2 \\ -s_2 & c_2 \end{pmatrix}.$$

Using $G_{k-1}^{(2)}$ defined in the previous iteration, we have

$$\begin{pmatrix} \eta_{k-1} \\ \hat{\sigma}_{k-1} \end{pmatrix} = \begin{pmatrix} c_2^{(k-1)} & -s_2^{(k-1)} \\ s_2^{(k-1)} & c_2^{(k-1)} \end{pmatrix} \begin{pmatrix} 0 \\ \rho_{k-1} \end{pmatrix},$$

$$\begin{pmatrix} \eta_k \\ \hat{\sigma}_k \end{pmatrix} = \begin{pmatrix} c_2 & -s_2 \\ s_1 & c_2 \end{pmatrix} \begin{pmatrix} 0 \\ \rho_k \end{pmatrix},$$

$$\begin{pmatrix} \sigma_{k-1} \\ 0 \end{pmatrix} = \begin{pmatrix} c_2 & -s_2 \\ s_2 & c_2 \end{pmatrix} \begin{pmatrix} \hat{\sigma}_{k-1} \\ \theta_k \end{pmatrix}.$$

We then obtain the following recurrences,

$$\eta_{k-1} = -s_2^{(k-1)} \rho_{k-1} \tag{12}$$

$$\hat{\sigma}_{k-1} = \frac{\rho_{k-1}\hat{\sigma}_{k-2}}{\sigma_{k-2}} = c_2^{(k-1)}\rho_{k-1}, \tag{13}$$

$$\sigma_{k-1} = \sqrt{\hat{\sigma}_{k-1}^2 + \theta_k^2}, \tag{14}$$

$$c_2 = \hat{\sigma}_{k-1}/\sigma_{k-1}, \tag{15}$$

$$s_2 = -\theta_k/\sigma_{k-1}, \tag{16}$$

$$\eta_k = \frac{\rho_k \theta_k}{\sigma_{k-1}} = -s_2 \rho_k. \tag{17}$$

(18)

3.6 Recurrence for x_k

We now derive a fast recurrence for x_k using the second QR decomposition. From the second QR decomposition, we have

$$\hat{R}_k^T \hat{Q}_k^T y_k = z_k.$$

Define

$$\hat{R}_k^T \hat{z}_k = z_k \tag{19}$$

$$\hat{Q}_k \begin{pmatrix} I_{k-1} \\ 0 \end{pmatrix} \hat{z}_k = y_k, \qquad \hat{z}_k = \begin{pmatrix} \hat{\zeta}_2 \\ \vdots \\ \hat{\zeta}_k \end{pmatrix}$$
 (20)

$$W_k = V_k \hat{Q}_k = (w_2^{(k)} | \dots | w_k^{(k)}). \tag{21}$$

With these definitions, we have,

$$W_k \begin{pmatrix} I_{k-1} \\ 0 \end{pmatrix} \hat{z}_k = V_k \hat{Q}_k \begin{pmatrix} I_{k-1} \\ 0 \end{pmatrix} \hat{z}_k = V_k y_k = x_k,$$

and so

$$x_k = W_{k-1} \begin{pmatrix} I_{k-2} \\ 0 \end{pmatrix} \hat{z}_{k-1} + w_{k-1}^{(k)} \hat{\zeta}_k.$$

The recursion for \hat{z}_k is similar to that of z_k , since it is a similar triangular solve via forward substitution, where we obtain

$$\hat{\zeta}_k = \frac{1}{\sigma_{k-1}} (\zeta_k - \eta_{k-1} \hat{\zeta}_{k-1}). \tag{22}$$

To get the recursion for W_k , we observe that

$$W_k = V_k \hat{Q}_k \tag{23}$$

$$= (V_{k-1}|v_k) \begin{pmatrix} \hat{Q}_{k-1} & \\ & 1 \end{pmatrix} \begin{pmatrix} I_{k-2} & \\ & G_k^{(2)} \end{pmatrix}$$
 (24)

$$= (W_{k-1}|v_k) \begin{pmatrix} I_{k-2} & & \\ & G_k^{(2)} \end{pmatrix}$$
 (25)

$$= (w_1^{(k-1)}|\dots|w_{k-2}^{(k-1)}|w_{k-1}^{(k-1)}|v_k) \begin{pmatrix} I_{k-2} & & \\ & G_k^{(2)} \end{pmatrix}.$$
 (26)

Then we see that the first k-2 columns of W_{k-1} and W_k are equal to each other, and so the only update that is required is

$$\begin{pmatrix} w_{k-1}^{(k)} & w_k^{(k)} \end{pmatrix} = \begin{pmatrix} w_{k-1}^{(k-1)} & v_k \end{pmatrix} \begin{pmatrix} c_2 & s_2 \\ -s_2 & c_2 \end{pmatrix}.$$
(27)

Although we compute both $w_{k-1}^{(k)}$ and $w_k^{(k)}$, we need only $w_{k-1}^{(k)}$ in order to compute x_k , while $w_k^{(k)}$ is necessary for the computation of W_{k+1} .

We summarize this procedure in Algorithm 2.

Algorithm 2 LSLQ

end for

$$\begin{array}{l} \widehat{\beta_1 u_1} = b, \ \alpha_1 v_1 = A^T u_1 \\ \beta_2 u_2 = A v_1 - \alpha_1 u_1 \\ \alpha_2 v_2 = A^T u_2 - \beta_2 v_1 \\ \rho_2 = \sqrt{\alpha_1^2 + \beta_2^2} \\ c_1^{(2)} = \alpha_1/\rho_2, \quad s_1^{(2)} = \beta_2/\rho_2 \\ \theta_2 = \alpha_2 \beta_2/\rho_2 \\ \end{array}$$

$$\begin{array}{l} \widehat{\zeta_2} = \rho_2 \\ \sigma_2 = \sqrt{\widehat{\sigma}_2^2 + \theta_2^2} \\ c_2^{(2)} = \widehat{\sigma}_2/\sigma_2, \quad s_2^{(2)} = -\theta_2/\sigma_2 \\ \end{array}$$

$$\begin{array}{l} \widehat{\zeta_2} = (c_1) \\ \widehat{\zeta_2} = (c_2) \\ \widehat{\zeta_2} =$$