

# CME338 Final Project

## LSLQ Solver

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## 1 Abstract

## 2 Introduction

### 2.1 Notation

$\mathcal{K}_k = \mathcal{K}_k(A^T A, b) = \text{span}(b, A^T A b, \dots, (A^T A)^{k-1} b)$ .

## 3 Derivation of LSLQ

In this section, we derive the short recurrence formulas for LSLQ beginning from the Golub-Kahan process.

### 3.1 Golub-Kahan Process

The Golub-Kahan process is defined by the recurrence defined in Algorithm 1.

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**Algorithm 1** Golub-Kahan Process

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Set  $\beta_1 u_1 = b, \alpha_1 v_1 = A^T u_1$   
for  $k = 1, 2, \dots$  do  
     $\beta_{k+1} u_{k+1} = A v_k - \alpha_k u_k$   
     $\alpha_{k+1} v_{k+1} = A^T u_{k+1} - \beta_{k+1} v_k$   
end for
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After performing  $k$  iterations of this process, we obtain the decompositions

$$A V_k = U_{k+1} B_k \tag{1}$$

$$A^T U_{k+1} = V_{k+1} L_{k+1}^T. \tag{2}$$

where  $U_k = (u_1 | \dots | u_k)$ ,  $V_k = (v_1 | \dots | v_k)$ , and

$$B_k = \begin{pmatrix} \alpha_1 & & & \\ \beta_2 & \alpha_2 & & \\ & \ddots & \ddots & \\ & & \beta_k & \alpha_k \\ & & & \beta_{k+1} \end{pmatrix} \quad L_{k+1} = (B_{k+1} | \alpha_{k+1} e_{k+1}).$$

We can then observe that  $V_k$  is a basis for  $\mathcal{K}_k(A^T A, b)$ , since

$$A^T AV_k = A^T U_{k+1} B_k = V_{k+1} L_{k+1}^T B_k = V_{k+1} \begin{pmatrix} B_k^T B_k \\ \alpha_{k+1} \beta_{k+1} e_k^T \end{pmatrix}.$$

### 3.2 The LSLQ subproblem

In order to solve the linear system or least squares problem  $Ax = b$ , we instead solve the normal equations  $A^T A = A^T b$ . Since  $A^T A$  is symmetric positive semidefinite and we assume that the normal equations are consistent, we may consider applying SYMMLQ to this system. We solve the normal equations iteratively, where at iteration  $k$  we solve the problem

$$\begin{aligned} x_k = \arg \min_{x \in \mathcal{K}_k} \quad & \|x\|_2 \\ \text{s.t.} \quad & A^T r \perp \mathcal{K}_{k-1}. \end{aligned} \tag{3}$$

where  $r = b - Ax$ .

In order to solve this problem, we first note that we may formulate this as an unconstrained problem in a smaller space if we minimize in  $y_k$  and set  $x_k = V_k y_k$ . Then

$$\begin{aligned} 0 = V_{k-1}^T A^T r_k &= V_{k-1}^T A^T (b - Ax_k) \\ &= V_{k-1}^T A^T b - V_{k-1}^T A^T A V_k y_k \\ &= \alpha_1 \beta_1 e_1 - B_{k-1}^T U_k^T A V_k y_k \\ &= \alpha_1 \beta_1 e_1 - B_{k-1}^T L_k y_k. \end{aligned}$$

Thus in order to solve 3, we can solve

$$\begin{aligned} y_k = \arg \min_{y \in \mathbb{R}^k} \quad & \|y\|_2 \\ \text{s.t.} \quad & B_{k-1}^T L_k y = \alpha_1 \beta_1 e_1. \end{aligned} \tag{4}$$

### 3.3 First QR decomposition

We first take the QR factorization of  $Q_k R_k = B_{k-1}$ . Suppose we have the QR factorization of  $Q_{k-1} R_{k-1} = B_{k-2}$ , with

$$R_{k-1} = \begin{pmatrix} \rho_1 & \theta_2 & & \\ & \rho_2 & \ddots & \\ & & \ddots & \theta_{k-1} \\ & & & \rho_{k-2} \end{pmatrix}.$$

Then we may recurse to obtain the factorization of  $B_{k-1}$

$$\begin{aligned}
B_{k-1} &= \left( \begin{array}{cccc|c} \alpha_1 & & & & 0 \\ \beta_2 & \alpha_2 & & & \vdots \\ & \ddots & \ddots & & \vdots \\ & & \beta_{k-2} & \alpha_{k-2} & 0 \\ & & & \beta_{k-1} & \alpha_{k-1} \\ \hline 0 & \dots & \dots & 0 & \beta_k \end{array} \right) = Q_{k-1} \left( \begin{array}{cccc|c} \rho_1 & \theta_2 & & & 0 \\ & \rho_2 & \ddots & & \vdots \\ & & \ddots & \theta_{k-2} & 0 \\ & & & \rho_{k-2} & \theta_{k-1} \\ & & & 0 & \hat{\rho}_{k-1} \\ \hline 0 & \dots & \dots & 0 & \beta_k \end{array} \right) \\
&= Q_{k-1} G_k^{(1)} \left( \begin{array}{cccc|c} \rho_1 & \theta_2 & & & 0 \\ & \rho_2 & \ddots & & \vdots \\ & & \ddots & \theta_{k-2} & \vdots \\ & & & \rho_{k-2} & \theta_{k-1} \\ & & & 0 & \rho_{k-1} \\ \hline 0 & \dots & \dots & 0 & 0 \end{array} \right) = Q_{k-1} G_k^{(1)} \begin{pmatrix} R_{k-1} \\ 0 \end{pmatrix},
\end{aligned}$$

by defining  $Q_k = Q_{k-1} G_k^{(1)}$ , where  $G_k^{(1)}$  is the Givens rotation

$$G_k^{(1)} = \begin{pmatrix} c_1 & s_1 \\ -s_1 & c_1 \end{pmatrix}.$$

Using  $G_{k-1}^{(2)}$  defined in the previous iteration, we have

$$\begin{aligned}
\begin{pmatrix} \theta_{k-1} \\ \hat{\rho}_{k-1} \end{pmatrix} &= \begin{pmatrix} c_1^{(k-1)} & -s_1^{(k-1)} \\ s_1^{(k-1)} & c_1^{(k-1)} \end{pmatrix} \begin{pmatrix} 0 \\ \alpha_{k-1} \end{pmatrix}, \\
\begin{pmatrix} \theta_k \\ \hat{\rho}_k \end{pmatrix} &= \begin{pmatrix} c_1 & -s_1 \\ s_1 & c_1 \end{pmatrix} \begin{pmatrix} 0 \\ \alpha_k \end{pmatrix}, \\
\begin{pmatrix} \rho_{k-1} \\ 0 \end{pmatrix} &= \begin{pmatrix} c_1 & -s_1 \\ s_1 & c_1 \end{pmatrix} \begin{pmatrix} \hat{\rho}_{k-1} \\ \beta_k \end{pmatrix},
\end{aligned}$$

and we therefore obtain the recurrences

$$\hat{\rho}_{k-1} = \frac{\alpha_{k-1} \hat{\rho}_{k-2}}{\rho_{k-2}} = c_1^{(k-1)} \alpha_{k-1}, \quad (5)$$

$$\rho_{k-1} = \sqrt{\hat{\rho}_{k-1}^2 + \beta_k^2}, \quad (6)$$

$$c_1 = \hat{\rho}_{k-1} / \rho_{k-1}, \quad (7)$$

$$s_1 = -\beta_k / \rho_{k-1}, \quad (8)$$

$$\theta_k = \frac{\alpha_k \beta_k}{\rho_{k-1}} = -s_1 \alpha_k. \quad (9)$$

### 3.4 Forward Substitution

With the previous QR decomposition, the system we intend to solve becomes

$$\begin{aligned} \begin{pmatrix} \alpha_1\beta_1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} &= \begin{pmatrix} & 0 \\ B_{k-1}^T B_{k-1} & \vdots \\ & 0 \\ & \alpha_k\beta_k \end{pmatrix} y_k = \begin{pmatrix} & 0 \\ R_k^T R_k & \vdots \\ & 0 \\ & \alpha_k\beta_k \end{pmatrix} y_k \\ &= R_k^T \begin{pmatrix} & 0 \\ R_k & \vdots \\ & 0 \\ & \frac{\alpha_k\beta_k}{\rho_{k-1}} \end{pmatrix} y_k = R_k^T \begin{pmatrix} & 0 \\ R_k & \vdots \\ & 0 \\ & \theta_k \end{pmatrix} y_k. \end{aligned}$$

Define

$$z_k = \begin{pmatrix} \zeta_2 \\ \vdots \\ \zeta_k \end{pmatrix} = \begin{pmatrix} & 0 \\ R_k & \vdots \\ & 0 \\ & \theta_k \end{pmatrix} y_k = \tilde{R}_k y_k \quad (10)$$

so that we have  $R_k^T z_k = \alpha_1\beta_1 e_1$ . As in the Conjugate Gradient method, we obtain a short recurrence for  $\zeta_k$ ,

$$\zeta_k = -\frac{\theta_{k-1}}{\rho_{k-1}} \zeta_{k-1}. \quad (11)$$

### 3.5 Second QR decomposition

Using the recurrence of the previous section, we now need to solve the minimum norm problem

$$\tilde{R}_k y_k = z_k.$$

We accomplish this by taking the QR decomposition of  $\hat{Q}_k \hat{R}_k = \tilde{R}_k^T$ . Suppose we have the QR decomposition from the previous iteration,  $\hat{Q}_{k-1} \hat{R}_{k-1} = \tilde{R}_{k-1}^T$ , with

$$\hat{R}_{k-1} = \begin{pmatrix} \sigma_1 & \eta_2 & & \\ & \sigma_2 & \ddots & \\ & & \ddots & \eta_{k-1} \\ & & & \sigma_{k-2} \end{pmatrix}.$$

Then as was done in the first QR decomposition, we can recurse to obtain a fast update for the second QR decomposition.

$$\begin{aligned}
\tilde{R}_k^T &= \left( \begin{array}{cccc|c} \rho_1 & & & & 0 \\ \theta_2 & \rho_2 & & & \vdots \\ & \ddots & \ddots & & \vdots \\ & & \theta_{k-2} & \rho_{k-2} & 0 \\ & & & \theta_{k-1} & \rho_{k-1} \\ \hline 0 & \dots & \dots & 0 & \theta_k \end{array} \right) = \hat{Q}_k \left( \begin{array}{cccc|c} \sigma_1 & \eta_2 & & & 0 \\ & \sigma_2 & \ddots & & \vdots \\ & & \ddots & \eta_{k-2} & \vdots \\ & & & \sigma_{k-2} & \eta_{k-1} \\ & & & 0 & \hat{\sigma}_{k-1} \\ \hline 0 & \dots & \dots & 0 & \theta_k \end{array} \right) \\
&= \hat{Q}_k G_k^{(2)} \left( \begin{array}{cccc|c} \sigma_1 & \eta_2 & & & 0 \\ & \sigma_2 & \ddots & & \vdots \\ & & \ddots & \eta_{k-2} & \vdots \\ & & & \sigma_{k-2} & \eta_{k-1} \\ & & & 0 & \sigma_{k-1} \\ \hline 0 & \dots & \dots & 0 & 0 \end{array} \right)
\end{aligned}$$

We define  $\hat{Q}_k = \hat{Q}_{k-1} G_k^{(2)}$ , where  $G_k^{(2)}$  is the Givens rotation

$$G_k^{(2)} = \begin{pmatrix} c_2 & s_2 \\ -s_2 & c_2 \end{pmatrix}.$$

Using  $G_{k-1}^{(2)}$  defined in the previous iteration, we have

$$\begin{aligned}
\begin{pmatrix} \eta_{k-1} \\ \hat{\sigma}_{k-1} \end{pmatrix} &= \begin{pmatrix} c_2^{(k-1)} & -s_2^{(k-1)} \\ s_2^{(k-1)} & c_2^{(k-1)} \end{pmatrix} \begin{pmatrix} 0 \\ \rho_{k-1} \end{pmatrix}, \\
\begin{pmatrix} \eta_k \\ \hat{\sigma}_k \end{pmatrix} &= \begin{pmatrix} c_2 & -s_2 \\ s_1 & c_2 \end{pmatrix} \begin{pmatrix} 0 \\ \rho_k \end{pmatrix}, \\
\begin{pmatrix} \sigma_{k-1} \\ 0 \end{pmatrix} &= \begin{pmatrix} c_2 & -s_2 \\ s_2 & c_2 \end{pmatrix} \begin{pmatrix} \hat{\sigma}_{k-1} \\ \theta_k \end{pmatrix}.
\end{aligned}$$

We then obtain the following recurrences,

$$\eta_{k-1} = -s_2^{(k-1)} \rho_{k-1} \quad (12)$$

$$\hat{\sigma}_{k-1} = \frac{\rho_{k-1} \hat{\sigma}_{k-2}}{\sigma_{k-2}} = c_2^{(k-1)} \rho_{k-1}, \quad (13)$$

$$\sigma_{k-1} = \sqrt{\hat{\sigma}_{k-1}^2 + \theta_k^2}, \quad (14)$$

$$c_2 = \hat{\sigma}_{k-1} / \sigma_{k-1}, \quad (15)$$

$$s_2 = -\theta_k / \sigma_{k-1}, \quad (16)$$

$$\eta_k = \frac{\rho_k \theta_k}{\sigma_{k-1}} = -s_2 \rho_k. \quad (17)$$

$$(18)$$

### 3.6 Recurrence for $x_k$

We now derive a fast recurrence for  $x_k$  using the second QR decomposition. From the second QR decomposition, we have

$$\hat{R}_k^T \hat{Q}_k^T y_k = z_k.$$

Define

$$\hat{R}_k^T \hat{z}_k = z_k \quad (19)$$

$$\hat{Q}_k \begin{pmatrix} I_{k-1} \\ 0 \end{pmatrix} \hat{z}_k = y_k, \quad \hat{z}_k = \begin{pmatrix} \hat{\zeta}_2 \\ \vdots \\ \hat{\zeta}_k \end{pmatrix} \quad (20)$$

$$W_k = V_k \hat{Q}_k = (w_2^{(k)} | \dots | w_k^{(k)}). \quad (21)$$

With these definitions, we have,

$$W_k \begin{pmatrix} I_{k-1} \\ 0 \end{pmatrix} \hat{z}_k = V_k \hat{Q}_k \begin{pmatrix} I_{k-1} \\ 0 \end{pmatrix} \hat{z}_k = V_k y_k = x_k,$$

and so

$$x_k = W_{k-1} \begin{pmatrix} I_{k-2} \\ 0 \end{pmatrix} \hat{z}_{k-1} + w_{k-1}^{(k)} \hat{\zeta}_k.$$

The recursion for  $\hat{z}_k$  is similar to that of  $z_k$ , since it is a similar triangular solve via forward substitution, where we obtain

$$\hat{\zeta}_k = \frac{1}{\sigma_{k-1}} (\zeta_k - \eta_{k-1} \hat{\zeta}_{k-1}). \quad (22)$$

To get the recursion for  $W_k$ , we observe that

$$W_k = V_k \hat{Q}_k \quad (23)$$

$$= (V_{k-1} | v_k) \begin{pmatrix} \hat{Q}_{k-1} & \\ & 1 \end{pmatrix} \begin{pmatrix} I_{k-2} & \\ & G_k^{(2)} \end{pmatrix} \quad (24)$$

$$= (W_{k-1} | v_k) \begin{pmatrix} I_{k-2} & \\ & G_k^{(2)} \end{pmatrix} \quad (25)$$

$$= (w_1^{(k-1)} | \dots | w_{k-2}^{(k-1)} | w_{k-1}^{(k-1)} | v_k) \begin{pmatrix} I_{k-2} & \\ & G_k^{(2)} \end{pmatrix}. \quad (26)$$

Then we see that the first  $k-2$  columns of  $W_{k-1}$  and  $W_k$  are equal to each other, and so the only update that is required is

$$\begin{pmatrix} w_{k-1}^{(k)} & w_k^{(k)} \end{pmatrix} = \begin{pmatrix} w_{k-1}^{(k-1)} & v_k \end{pmatrix} \begin{pmatrix} c_2 & s_2 \\ -s_2 & c_2 \end{pmatrix}. \quad (27)$$

Although we compute both  $w_{k-1}^{(k)}$  and  $w_k^{(k)}$ , we need only  $w_{k-1}^{(k)}$  in order to compute  $x_k$ , while  $w_k^{(k)}$  is necessary for the computation of  $W_{k+1}$ .

We summarize this procedure in Algorithm 2.

## 4 Norms and Stopping Criteria

Here we will derive recurrences for computing estimates of  $\|r_k\|$ ,  $\|A^T r_k\|$ ,  $\|A\|$  and  $\text{cond}(A)$ .

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**Algorithm 2** LSLQ

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$$\beta_1 u_1 = b, \alpha_1 v_1 = A^T u_1$$

$$\beta_2 u_2 = A v_1 - \alpha_1 u_1$$

$$\alpha_2 v_2 = A^T u_2 - \beta_2 v_1$$

$$\rho_2 = \sqrt{\alpha_1^2 + \beta_2^2}$$

$$c_1^{(2)} = \alpha_1 / \rho_2, \quad s_1^{(2)} = \beta_2 / \rho_2$$

$$\theta_2 = \alpha_2 \beta_2 / \rho_2$$

$$\zeta_2 = \alpha_1 \beta_1 / \rho_2$$

$$\hat{\sigma}_2 = \rho_2$$

$$\sigma_2 = \sqrt{\hat{\sigma}_2^2 + \theta_2^2}$$

$$c_2^{(2)} = \hat{\sigma}_2 / \sigma_2, \quad s_2^{(2)} = -\theta_2 / \sigma_2$$

$$\hat{\zeta}_2 = \zeta_2 / \sigma_2$$

$$\begin{pmatrix} w_1^{(2)} & w_2^{(2)} \end{pmatrix} = \begin{pmatrix} v_1 & v_2 \end{pmatrix} \begin{pmatrix} c_2^{(2)} & -s_2^{(2)} \\ s_2^{(2)} & c_2^{(2)} \end{pmatrix}$$

$$x_k = \hat{\zeta}_2 w_1^{(2)}$$

**for**  $k = 3, \dots$  **do**

$$\beta_k u_k = A v_{k-1} - \alpha_{k-1} u_{k-1}$$

$$\alpha_k v_k = A^T u_k - \beta_k v_{k-1}$$

$$\hat{\rho}_{k-1} = c_1^{(k-1)} \alpha_{k-1}$$

$$\rho_{k-1} = \sqrt{\hat{\rho}_{k-1}^2 + \beta_k^2}$$

$$c_1^{(k)} = \hat{\rho}_{k-1} / \rho_{k-1}, \quad s_1^{(k)} = \beta_k / \rho_{k-1}$$

$$\theta_k = -s_k^{(k)} \alpha_k$$

$$\zeta_k = -\zeta_{k-1} \theta_{k-1} / \rho_{k-1}$$

$$\eta_{k-1} = -s_2^{(k-1)} \rho_{k-1}$$

$$\hat{\sigma}_{k-1} = c_2^{(k-1)} \rho_{k-1}$$

$$\sigma_{k-1} = \sqrt{\hat{\sigma}_{k-1}^2 + \theta_k^2}$$

$$c_2^{(k)} = \hat{\sigma}_{k-1} / \sigma_{k-1}, \quad s_2^{(k)} = -\theta_k / \sigma_{k-1}$$

$$\hat{\zeta}_k = (\zeta_k - \eta_{k-1} \hat{\zeta}_{k-1}) / \sigma_{k-1}$$

$$\begin{pmatrix} w_{k-1}^{(k)} & w_k^{(k)} \end{pmatrix} = \begin{pmatrix} w_{k-1}^{(k-1)} & v_k \end{pmatrix} \begin{pmatrix} c_2^{(k)} & -s_2^{(k)} \\ s_2^{(k)} & c_2^{(k)} \end{pmatrix}$$

$$x_k = x_{k-1} + \hat{\zeta}_k w_{k-1}^{(k)}$$

**end for**

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#### 4.1 Recurrence for $y_k$

Here we derive a recurrence for the last two entries of  $y_k$ , which will be used in the estimates of the norm quantities in which we are interested. From equation 20, we have that

$$\begin{aligned}
y_k &= \hat{Q}_k \begin{pmatrix} I_{k-1} \\ 0 \end{pmatrix} \begin{pmatrix} \hat{\zeta}_2 \\ \vdots \\ \hat{\zeta}_k \end{pmatrix} \\
&= \begin{pmatrix} G_{k-1}^{(2)} & \\ & 1 \end{pmatrix} \begin{pmatrix} I_{k-2} & \\ & G_k^{(2)} \end{pmatrix} \begin{pmatrix} \hat{\zeta}_2 \\ \vdots \\ \hat{\zeta}_k \\ 0 \end{pmatrix} \\
&= \begin{pmatrix} \times & \cdots & \times \\ \times & \cdots & \times \\ & \ddots & \vdots \\ & s_{k-1}^{(2)} & -c_{k-1}^{(2)} c_k^{(2)} \\ & 0 & s_k^{(2)} \end{pmatrix} \begin{pmatrix} \hat{\zeta}_2 \\ \vdots \\ \hat{\zeta}_k \end{pmatrix}.
\end{aligned}$$

Thus we can obtain the last 2 entries of  $y_k$  from

$$\begin{pmatrix} \psi_1^{(k)} \\ \psi_2^{(k)} \end{pmatrix} = \begin{pmatrix} s_{k-1}^{(2)} & -c_{k-1}^{(2)} c_k^{(2)} \\ 0 & s_k^{(2)} \end{pmatrix} \begin{pmatrix} \hat{\zeta}_{k-1} \\ \hat{\zeta}_k \end{pmatrix}. \quad (28)$$

#### 4.2 Recurrence for $\|r_k\|$

We observe the following relationships based on the equations

$$\tilde{R}_k y_k = z_k \quad (29)$$

$$R_{k+1}^T z_{k+1} = \alpha_1 \beta_1 e_1 \quad (30)$$

$$Q_{k+1} R_{k+1} = B_k. \quad (31)$$

Taking the transpose of the second equation and defining  $q^{(k+1)} = \beta_1 Q_{k+1}^T e_1$ , we see that

$$R_{k+1}^T Q_{k+1}^T = B_k^T \quad (32)$$

$$R_{k+1}^T q^{(k+1)} = \alpha_1 \beta_1 e_1. \quad (33)$$

We can obtain a recurrence for  $q^{(k+1)}$  by observing that

$$\begin{aligned}
q^{(k+1)} &= \beta_1 Q_{k+1}^T e_1 \\
&= \beta_1 \begin{pmatrix} I_{k-2} & \\ & (G_{k+1}^{(1)})^T \end{pmatrix} \begin{pmatrix} Q_k^T & \\ & 1 \end{pmatrix} e_1 \\
&= \begin{pmatrix} I_{k-2} & \\ & (G_{k+1}^{(1)})^T \end{pmatrix} \begin{pmatrix} q^{(k)} \\ 0 \end{pmatrix} = \begin{pmatrix} q_1^{(k)} \\ \vdots \\ q_{k-1}^{(k)} \\ q_k^{(k)} c_1^{(k+1)} \\ -q_k^{(k)} s_1^{(k+1)} \end{pmatrix}.
\end{aligned}$$



Now, since  $R_{k+1}^T$  is nonsingular, we have that

$$q^{(k+1)} = \begin{pmatrix} z_{k+1}^{(k)} \\ -q_k^{(k)} s_1^{(k+1)} \end{pmatrix}. \quad (34)$$

We now use this relationship to obtain a short recurrence for  $\|r_k\|$ . Thus

$$\begin{aligned} \|r_k\| &= \|b - Ax_k\| \\ &= \|U_{k+1}(\beta_1 e_1 - B_k y_k)\| \\ &= \|Q_{k+1}^T(\beta_1 e_1 - B_k y_k)\| \\ &= \left\| \beta_1 Q_{k+1}^T e_1 - \begin{pmatrix} R_{k+1} \\ 0 \end{pmatrix} y_k \right\| \\ &= \left\| \beta_1 Q_{k+1}^T e_1 - \begin{pmatrix} R_{k+1} \\ 0 \end{pmatrix} y_k \right\| \\ &= \left\| \beta_1 Q_{k+1}^T e_1 - \begin{pmatrix} z_k \\ \rho_k \psi_2^{(k)} \\ 0 \end{pmatrix} \right\| \\ &= \left\| \begin{pmatrix} 0 \\ -q_k^{(k)} c_1^{(k+1)} - \rho_k \psi_2^{(k)} \\ q_k^{(k)} s_1^{(k+1)} \end{pmatrix} \right\|. \end{aligned}$$

Since we require only the last 2 entries of  $q^{(k+1)}$  for which we have a fast recurrence, the computation of  $\|r_k\|$  can be achieved in  $O(1)$  flops. Note that in order to estimate the residual at iteration  $k$ , we need values computed at iteration  $k+1$ .

### 4.3 Recurrence for $\|A^T r_k\|$

We can obtain an estimate of  $\|A^T r_k\|_2$  with  $O(1)$  flops. We have

$$\begin{aligned} \|A^T r_k\| &= \|A^T(b - Ax_k)\| \\ &= \|V_{k+1}(\alpha_1 \beta_1 e_1 - L_{k+1}^T B_{k+1} y_k)\| \\ &= \|\alpha_1 \beta_1 e_1 - L_{k+1}^T B_{k+1} y_k\|. \end{aligned}$$

We note that

$$\begin{aligned} L_{k+1}^T B_{k+1} y_k &= \left( \begin{array}{ccc|c} & & & 0 \\ & B_{k-1}^T & & \vdots \\ & & & 0 \\ \hline 0 & \cdots & \alpha_k & \beta_{k+1} \\ 0 & \cdots & 0 & \alpha_{k+1} \end{array} \right) \begin{pmatrix} L_k^T \\ \beta_{k+1} e_k^T \end{pmatrix} y_k \\ &= \begin{pmatrix} B_{k-1}^T L_k \\ \alpha_k \beta_k e_{k-1}^T + (\alpha_k^2 + \beta_{k+1}^2) e_k^T \\ \alpha_{k+1} \beta_{k+1} e_k^T \end{pmatrix} y_k \\ &= \begin{pmatrix} \alpha_1 \beta_1 e_1 \\ (\alpha_k \beta_k e_{k-1} + (\alpha_k^2 + \beta_{k+1}^2) e_k)^T y_k \\ \alpha_{k+1} \beta_{k+1} e_k^T y_k \end{pmatrix}. \end{aligned}$$

Thus we have

$$\|\alpha_1\beta_1e_1 - L_{k+1}^TB_{k+1}y_k\| = \left\| \begin{pmatrix} 0 \\ (\alpha_k\beta_k e_{k-1} + (\alpha_k^2 + \beta_{k+1}^2)e_k)^T y_k \\ \alpha_{k+1}\beta_{k+1}e_k^T y_k \end{pmatrix} \right\|,$$

and so we need only the last 2 entries of  $y_k$  which we can obtain in  $O(1)$  flops as described in Section 4.1.

#### 4.4 Estimate of $\|A\|$ and $\text{cond}(A)$

As in LSMR, we may estimate  $\|A\|$  by using  $\|B_k\|_F$ . In order to estimate  $\text{cond}(A)$ , we note that  $\text{cond}(A)^2 = \text{cond}(A^T A)$ , and that we may estimate  $\text{cond}(A^T A)$  by  $\text{cond}(B_k^T B_k)$ . Since  $B_k^T B_k = R_{k+1}^T R_{k+1}$ , we can estimate  $\text{cond}(A)$  by  $\text{cond}(R_k)$  at each iteration.  $\text{cond}(R_k)$  may be estimated by the largest and smallest entries on its diagonal, that is,  $\max \sigma_i / \min \sigma_i$ .