

## Walkthrough

### !!! Spoilers Ahead !!!

This walkthrough contains **spoilers and detailed instructions**.

If you want detailed instructions, please proceed to the next page.

## Analysis of the Dropper: procmon64 Step 1

1. **Right-click** the dropper and examine its **properties/description**.
  2. Rename the dropper as **w.exe**
  3. Start **procmon64** and add a filter where the **Process Name equals w.exe**.
  4. **Execute** the dropper
  5. Filter the **procmon64** log to identify **where the files are dropped and how they are executed**.
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## Analysis of the Dropped Files

### Analyze the msedgewebview4.exe file

1. **Open in pe-bear:**
  - Examine the **imports**, **strings**, and **sections**. Does anything look suspicious?
2. **Open in IDAFree:**
  - **Understand the general layout:**
    - Find the function that **reads and loads the file**.
    - Find the function that **loops over the encrypted file content** and determine **where the encryption key is stored**.
  - **Understand how the encryption is implemented:**
    - Analyze the **XOR loop** (which suggests a **Pseudo-Random Number Generator (PRNG)**).
    - Find the **encryption key** (which serves as the **seed for the PRNG**).

### Analyze the SearchHost.bin file

1. **Decrypt the bytecode** using `decrypt_bytecode.py` with the **correct key (seed)**.
  2. **Check if it matches a well-known architecture** using `find_shellcode_arch.py`:
    - **What heuristic** does `find_shellcode_arch.py` implement?
  3. **Open the shellcode in ghidra** (specifying the **correct architecture**):
    - **Analyze the API resolution function:**
      - Use `hash_x65599_exports.py` to **identify which APIs are imported** (e.g., `kernelbase!VirtualProtect`, `kernel32!CreateThread`).
      - Determine **how the API is used**.
    - **Analyze the main loop function**.
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## Toward the Second Stage: procmon64 Step 2

### 1. Continue monitoring with procmon64:

- Identify the **APIs that are invoked** by analyzing the **stack trace**.
- Look for `GetComputerName` and `GetUserNameW` in the stack trace.

### 2. Create a user:

```
net user <username> <pass> /add
```

### 3. Add the user to the local administrators group (optional):

```
net localgroup administrators <username> /add
```

### 4. copy the dropped files into `c:\windows\temp` and execute the dropper under the new user's credentials:

```
cd c:\windows\temp
Start-Process powershell.exe -ArgumentList "-Command & { Start-Process msedgewebview4.exe
```

### 5. Attach to the process with windbg (activate all user processes view in windbg - requires elevation):

- Place a **breakpoint** on `kernelbase!VirtualProtect`.
- Wait for the breakpoint to **hit**. Look at the stack trace with `kb`. **Identify the parameters**; we expect `r8` to equal `0x40` (**Read-Write-Execute - RWX**).
- **Execute `VirtualProtect`** and open **SystemInformer**. Find the process and look for the **RWX allocated memory regions**. **Dump the content** and **trim the dump** properly using a hex editor.
- Greetings! You've found **Stage 2**!

## Analysis of the stage-2

### Stage-2 Malware Analysis

This section details the analysis of `stage-2.dll`, focusing on static and initial dynamic examination to understand its loading mechanism and primary function.

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### Static Analysis and Initial Triage

#### 1. File Triage with PE-bear:

- Load `stage-2.dll` into a PE viewer like **PE-bear**.
- Examine the **PE Headers** and sections.
- Analyze **Import Address Table (IAT)** and **Export Address Table (EAT)** to identify its dependencies and entry points.
- Review embedded **Strings** for immediate indicators (e.g., file paths, domain names, API calls, custom error messages).

## 2. Disassembly and Initial Code Review (IDA-Free):

- Open `stage-2.dll` in a disassembler (e.g., **IDA-Free**). Initial inspection suggests it is a **COM DLL** based on its exported functions.
  - **Review Standard COM Exports:**
    - Examine `DllRegisterServer` and `DllUnregisterServer` to confirm their conventional structure or look for malicious deviations.
  - **Identify Loader Invocation:**
    - Note that the **stage-1 loader** specifically invokes **Ordinal 1**, which corresponds to the export `DllGetClassObject`.
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## Reflective Loading Mechanism Discovery

### 3. Analysis of `DllGetClassObject` (Ordinal 1)

- **Unconventional implementation:** Inspect the control flow of `DllGetClassObject`. It does **not** follow the standard COM pattern.
- **Key indicators of reflective loading:** Note the presence of magic constants such as `0x6A4ABC5B`, `0x3CFA685D`, and `0x534C0AB8`.
- **Verification:** A quick search (Google or an LLM) shows these constants are commonly associated with reflective loader implementations (e.g., Metasploit-style reflective DLL injection). This strongly suggests a reflective loader is present.

### 4. Reflective loader execution flow

- The reflective loader maps the DLL into memory and transfers execution to its internal entry point.
  - The reflective loader eventually calls `DllEntryPoint`, which is the **true main** for the stage-2 payload.
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## Payload Execution Flow

### 5. `DllEntryPoint` analysis

- **Thread creation:** `DllEntryPoint` creates a new thread to run a function we label `start_function`. The loader waits on that thread. This detaches payload execution from the loader and helps evade analysis.

### 6. `start_function` dissection

- **API resolution:** One of the first tasks of `start_function` is a PEB walk to dynamically resolve required APIs. This avoids static imports and complicates static analysis.
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## Key derivation and 3DES decryption chain

7. Important function: `sub_7FFC86652680`

- `sub_7FFC86652680` takes as input a `fLink` (a DLL pointer) and a pointer to `&unk_7FFC86655370` (a constant).
- Nested calls reveal this repeated pattern:

```
LOBYTE(n8_2) = 101;
v24 = sub_7FFC86651500(v16, *a2, n8_2);
LOBYTE(n8_3) = 100;
v26 = sub_7FFC86651500(v24, a2[1], n8_3);
LOBYTE(n8_4) = 101;
v28 = sub_7FFC86651500(v26, a2[2], n8_4);
n8_10 += 8;
```

This structure is distinctive and suggests a multi-pass block cipher application.

8. Byte table used by `sub_7FFC86651500` performs mathematical transformations using the following byte array:

```
.byte_7FFC86655000: db 3Ah, 32h, 2Ah, 22h, 1Ah, 12h, 0Ah, 2, 3Ch, 34h, 2Ch
.byte_7FFC8665500B: db 24h, 1Ch, 14h, 0Ch, 4, 3Eh, 36h, 2Eh, 26h, 1Eh, 16h
.byte_7FFC86655016: db 0Eh, 6, 40h, 38h, 30h, 28h, 20h, 18h, 10h, 8, 39h, 31h
.byte_7FFC86655022: db 29h, 21h, 19h, 11h, 9, 1, 3Bh, 33h, 2Bh, 23h, 1Bh, 13h
```

Correlating the constants and call structure reveals DES-like S-box / permutation tables. The calling sequence implements 3DES (Triple DES) in ECB mode (no IV). Block size is 8 bytes (64 bits).

9. Key size and origin The key length is 24 bytes, matching the size of a global variable passed to:

```
sub_7FFC86652D10(byte_7FFC86655A80);
```

`sub_7FFC86652D10` generates a random value and performs operations on a specific DLL; therefore the key can be recovered:

- *statically* (by reversing the generation routine)
- *dynamically* (by using the debugger to read the generated bytes).

To speed up the process, this walkthrough uses the dynamic approach (debugger).

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## Debugger setup & dynamic key recovery (IDA)

10. **Debugger setup (IDA):** In IDA's Debugger options, set the default local debugger. In Process options, set the target application to:

```
C:\Windows\System32\rundll32.exe
```

As parameters, pass the renamed sample DLL and a dummy exported ordinal, e.g.:

`C:\Users\user\2stage\2stage.dll, #2`

Place a breakpoint on the first instruction of the DLL main (press **F2** on that instruction).

11. **Run and advance to start\_function:** Launch the debugger (green arrow) and wait the execution breaks at DLL main.
12. **Key discovery (dynamic approach):** Open `sub_7FFC86652D10` and set the IP to its start (use **Ctrl+N** or “Set IP”) Follow the execution till an instruction immediately after it writes the global key (e.g., `byte_7FFC86655A80`).

The key is derived from opcodes of a randomly chosen export (e.g., from `ntdll`). In our sample the extracted 24-byte key was:

`0000DEADB8D18B4C017FDEAD082504F6C32EDEAD050F0375` (hex representation)

Read memory at the global key address and dump 24 bytes. Format them as hex: that is your 3DES key (`K1||K2||K3`).

13. **Confirming the decryption chain:** Cross-reference the decryption routine (`sub_7FFC86652680` and callers). All `unk_*` values referenced by callers correspond to encrypted blobs in the stage-2 payload.

Use CyberChef or any 3DES tool:

Algorithm: Triple DES (3DES)

Mode: ECB (no IV)

Key: 24-byte key recovered (hex)

Input format: Hex (raw ciphertext), e.g. `281a5261479d3d42087515567ec409c3`

Decrypt all encrypted blobs. The decrypted outputs reveal stage-2 strings and behavior. One buffer at `0x00007FFC86655450` should be converted to code in IDA (press **C**).