Xevolver



Vectorization-aware Loop Optimization with User-defined Code Transformations

Re-Emergence of Vector Architectures (REV-A Workshop) Sept 5, 2017@Honoluli, Hawaii

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Background

Computer Architecture Design

HPC Application Development

Make Common Case Fast

- What is the fast case of vector architecture?
 - → Vectorized long loops!
 - Another system/compiler/programmer may require/prefer a different loop structure.

Separate vectorization-awareness from application codes

How Is Code Modified?



- Bad News -- Messy
 - Vectorization-aware code modifications are scattered over a code
- Good News -- Repetitive
 - Same (or similar) code modifications are required many times

Manual code modifications can be replaced with a smaller number of mechanical code transformations.

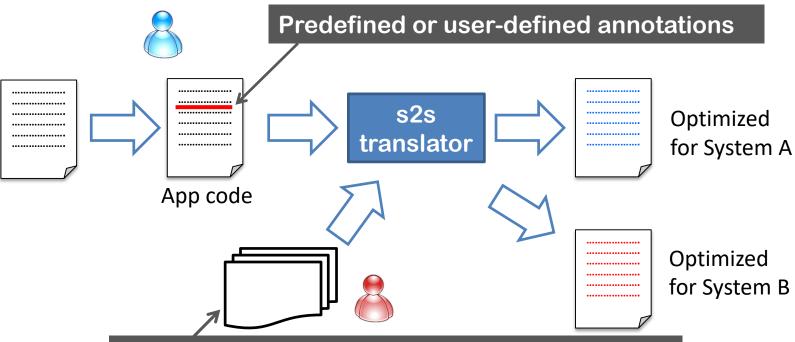
→ Express application-specific and/or system-specific code modifications as mechanical code transformations

Xevolver Framework



Various transformations are required for replacing <u>arbitrary code modifications</u>.

- = cannot be expressed by combining predefined transformations.
- → Xevolver : a framework for custom code transformations

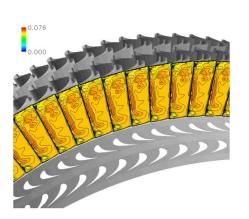


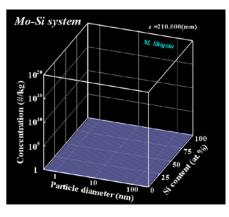
Translation rules

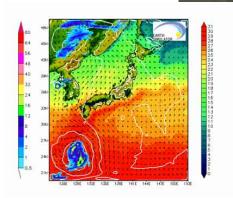
- Define the code transformation of each annotation
- Different systems can use different rules
- Users can define their own code transformations

Case Studies with Real Applications

- Real-world applications originally developed for NEC SX-9 have been ported to OpenACC.
 - Numerical Turbine (Yamamoto et al@Tohoku-U)
 - Nano-Powder Growth Simulation (Shigeta@Osaka-U)
 - MSSG-A (Takahashi et al@JAMSTEC)







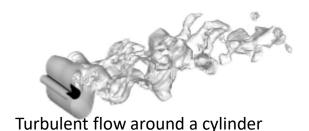
Is our approach effective for vectorization-aware loop optimizations?

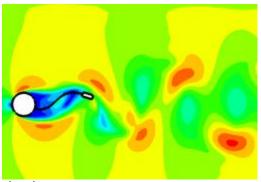
- There are clear patterns in such code modifications!
 - = Xevolver should be helpful for expressing vectorization-awareness.

This Work

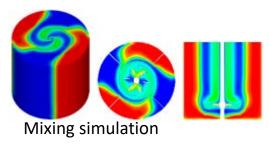


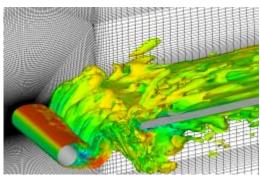
- Case Study: The FASTEST incompressible flow solver
 - Originally developed in the 1990's at FAU, Germany
 - SPPEXA ExaFSA version developed at Technical University Darmstadt
 - http://www.fnb.tu-darmstadt.de/forschung_fnb/software_fnb/software_fnb.en.jsp
 - Once written for old vector machines
 - → Some kernels already have their default and vector versions





Fluid-structure interactions





3D flow around cylinder-plate configuration



What is the difference btwn default and vector versions?



- Loop collapse and its variants
 - Vector versions have different loop structures to enable vectorization and increase their loop lengths.

```
real a(n,n,n),b(n,n,n)

do k=2,nkm

do i=2,nim

do j=2,njm

a(j,i,k) = a(j,i,k)+b(j,i,k)

end do

end do

end do

end do
```

```
real a(n,n,n),b(n,n,n)

do kij=0,(nke-1)*(nie-1)*(nje-1)-1

k=kij/((nie-1)*(nje-1))+2

i=mod(kij,(nie-1)*(nje-1))/(nje-1)+2

j=mod(kij,(nje-1))+2

a(j,i,k) = a(j,i,k)+b(j,i,k)

end do
```

What is the difference btwn default and vector versions?



Loop collapse and its variants

eal a(n,n,n),b(n,n,n)

 Vector versions have different loop structures to enable vectorization and increase their loop lengths.

Standard version

```
Transform
               o k=2, nkm
                do i=2, nim
                  do j=2, njm
                    a(j,i,k) = a(j,i,k) + b(j,i,k)
                  end do
                end do
                                    The compiler used in this case study
              end do
                                    does not collapse this loop structure.
                                                     Vector version
              real a(n,n,n),b(n,n,n)
              do kij=0, (nke-1) * (nie-1) * (nje-1) -1
                k=kij/((nie-1)*(nje-1))+2
                 i=mod(kij, (nie-1)*(nje-1))/(nje-1)+2
                 j=mod(kij,(nje-1))+2
                 a(j,i,k) = a(j,i,k) + b(j,i,k)
              end do
```

How to Describe Code Transformation

- One common way to explain a code transformation is to show before-and-after versions
 - e.g. Loop unrolling (screen capture of en.wikipedia.org)

A procedure in a computer program is to delete 100 items from a collection. This is normally accomplished by mea overhead of the loop requires significant resources compared to those for the *delete(x)* loop, unwinding can be use

Normal loop	After loop unrolling			
	int x; for (x = 0; x < 100; x += 5)			
int x;	1			
for (x = 0; x < 100; x++)	delete(x);			
{	delete(x + 1);			
delete(x);	delete(x + 2);			
}	delete(x + 3);			
	delete(x + 4);			
	}			

As a result of this modification, the new program has to make only 20 iterations, instead of 100. Afterwards, only 20 decrease in the loop administration overhead. To produce the optimal benefit, no variables should be specified in the referencing.

Users do not need any special knowledge (e.g. XML and AST) to describe a code transformation.



Simple Loop Collapse Rule



```
!$xev tgen var(nke, nje, nie) exp
   !$xev tgen list(flxc_stl1) stmt
   !$xev tgen trans stmt src begin
                                            Original Loop Structure
   !$xev vector flxc begin
       do k=2, nke
5
         do i=2, nie
            do j=2, nje
    !$xev tgen stmt(flxc stl1)
            end do
          end do
10
        end do
   e^{n\alpha} ao !$xev vector All statements in the loop body are copied.
11
12
   !$xev tgen trans stmt src end
13
                                                   Transformed Loop Structure
14
   !$xev tgen trans stmt dst begin
15
   !$xev vector flxc begin
16
   #include "vdi nodep.h"
17
   do kij=0, (nke-1) * (nie-1) * (nje-1)-1
18
     k=kij/((nie-1)*(nje-1))+2
19
     i=mod(kij,(nie-1)*(nje-1))/(nje-1)+2
20
     j=mod(kij/(nje-1))+2
21
    !$xev tgen stmt(flxc_stl1)
22
       end do
23
   !$xev vector flxc end
   !$xev tgen trans stmt dst end
```

Code patterns are written in the rule → easily customizable for individual cases





Collapsing imperfectly-nested loops

```
!$xev tgen list(vint_stl1) stmt
   !$xev tgen trans src begin
   !$xev vector vint begin
                                 Original Loop Structure
     do kcq=1, nkmq
       kfa=2*kca-1
                                  There is a statement between loops.
       do icg=1, nimg
         ifg=2*icg-1
         do jcg=1, njmg
           jfg=2*jcg-1
   !$xev tgen stmt(vint_stl1)
         end do
12
       end do
     end do
13
   !$xev vector vint end
   !$xev tgen trans src end
   !$xev tgen trans dst begin
                                                    Transformed Loop Structure
   !$xev vector vint begin
   #include "vdirnodep.h"
                                                    Any statements can be inserted to avoid
     do kijcg=0,nkmg*nimg*njmg-1
20
       kcq=kijcq/(nimq*njmq)+1
21
                                                    changing the code behaviors.
       icq=mod(kijcq, (nimq*njmq))/njmq+1
       jcg=mod(kijcg,njmg)+1
24
       kfq=2*kcq-1
25
       ifg=2*icg-1
26
       jfg=2*jcg-1
   !$xev tgen stmt(vint_stl1)
     end do
   !$xev vector vint end
   !$xev tgen trans dst end
```

Experimental Setup







	NEC SX-ACE	Intel Xeon E5-2695v2
Peak Performance [Gflop/s]	256/socket, 64/core	230.4/socket, 19.2/core
Number of cores	4	12
Cache size	1MB/core	L2:256KB/core, L3:30MB/socket
Memory bandwidth [GB/s]	256	59.7
Vector length/ SIMD width (double)	256	4
Compiler	NEC SX/Fortran 111	Intel Compiler 16.0.3

User-defined Code Transformations



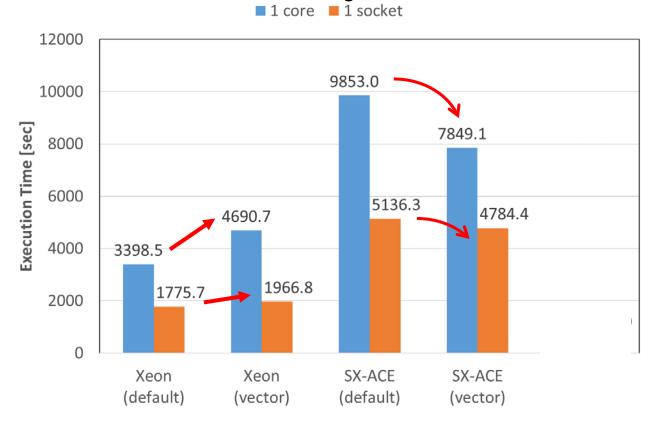
- 19 kernels of the original FASTEST code have their default and vector versions.
 - The difference between the two versions is expressed as a code transformation.
- 10 rules are defined for converting the 19 kernels
 - Most of them are for collapsing loop nests.
 - A simple loop collapse rule is customized for defining other special rules.
 - Some rules can be reused for converting multiple loops.

name	calcdp	calcp	celp2	flxc	vint	flxcoa	sipsol1	sipsol2	sipsol4	celuvw1
# loops	1	5	2	3	1	3	1	1	1	1

- Improving code maintainability
 - Only the default version is kept, and the vector version is removed from the code.
 - The default version of each kernel is converted to its vector version using Xevolver when it is compiled for vector systems.

Performance Evaluation Results

- 19 kernels already have their vector versions
 - →The SX performance increases by using vector versions.
 - = The other kernels may not be vectorized.



Obstacles for Vectorization

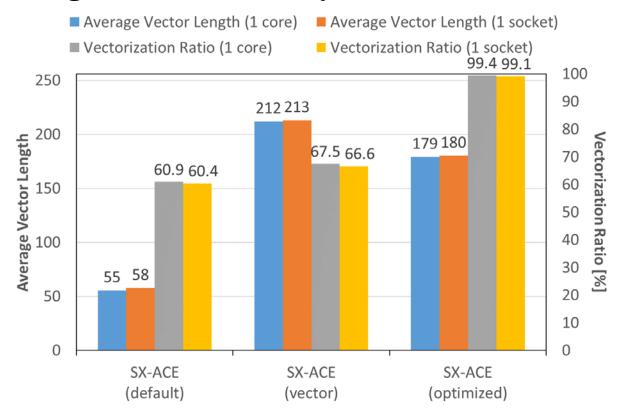


```
program main
2
   do i=1, ni
                          This loop nest is not vectorized.
     do j=1, n j
4
     call foo (x, j, mode)
6
     enddo j
   enddo i
9
   end program main
10
11
   subroutine foo(x,itr,iflag)
                                    The value of x is undefined
12
     integer x
13
                                    if iflag is neither 1 nor 2.
     integer itr
14
     integer iflag
15
16
                                     This part is removed to
     if (iflag == 1) then
17
      x = itr
18
                                     prevent undefined variables.
     else if (iflag
19
      x = it.r*2
20
     end if
21
   end subroutine foo
```

Vectorization Ratio

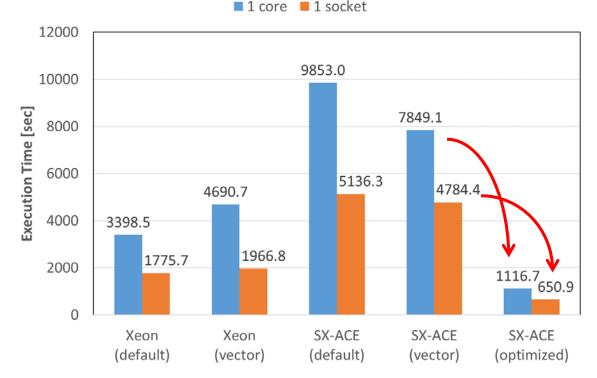


- Most loops are vectorized by minor modifications.
 - Obstacles for vectorization are simply removed.
 - System-specific optimizations are not applied to avoid degrading the maintainability.



Finally Achieved Performance

- Most kernels are now vectorized
 - = Significant SX performance improvement.
 - → Vectorization ratio is very important for SX-ACE.
 - → Average vector length is less important.



Conclusions



- Xevolver framework
 - Vectorization-aware loop optimizations are separated from application codes.
 - Application developers can maintain the original code
 - System-specific optimizations are defined in an external file
- The right one of

code refactoring or code transformation

should be used for the right purpose.

- Xevolver provides the latter option.
- An application code should be refactored if the refactoring could improve the maintainability.
- → We need to further explore the best practices of code optimizations with user-defined code transformations.

Acknowledgements



- We would like to thank <u>Kenta Yamaguchi</u> and <u>Yoichi</u>
 <u>Shimomura</u> for their contributions to this work <u>.</u>
- This work was partially supported by JST Post-Peta CREST, DFG SPPEXA ExaFSA project, and Grant-in-Aid for Scientific Research(B) 16H02822.

Xevolver	SEARCH
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Xevolver with some sample translation rules is online available at http://xev.sc.cc.tohoku.ac.jp.

Your feedbacks (and bug reports) are welcome!