

Dynamic Connectivity Problem

The **dynamic connectivity** problem involves determining if there is a path between two objects in a set of objects that can be connected or disconnected along time. This problem is frequently addressed by using algorithms that supports the following operations:

- **Union:** connects two objects;
- **Find:** verify if two objects are connected;

REAL LIFE APPLICATIONS: social networks, image pixels, electric circuits,

The **connections** between objects must attend to natural and intuitive properties named as **equivalence relations**:

- **Reflexive:** p is connected to p .
- **Symmetric:** if p is connected to q , then q is connected to p .
- **Transitive:** if p is connected to q and q is connected to r , then p is connected to r .

Connected components: maximal set of objects that are mutually connected.

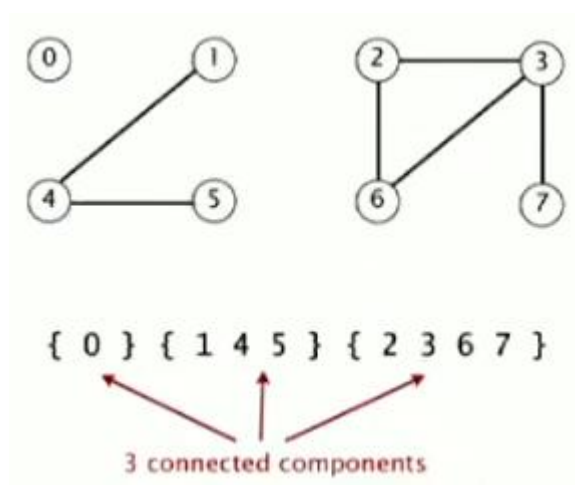


Image 1 – Connected components.

Quick Find Algorithm

The **quick find algorithm** is one way to solve the **dynamic connectivity** problem with fast **find** operations but slow **union** operations.



Data Structure: uses an *array* $id[]$, where $id[i]$ stores the identifier of the group which element i belongs to.

Two elements are connected if they have the same id .

Operations:

- **Find/Connected:** checks if $id[p] == id[q]$. (constant time-complexity: $O(1)$).
- **Union:** to connect p and q , changes all elements with $id[p]$ to $id[q]$. (linear time-complexity: $O(N)$).

Quick Union Algorithm

The **quick union algorithm** is another way of solving the **dynamic connectivity** problem with fast **union** operations but slow **find** operations.

Data Structure: uses an *array* $id[]$, but now $id[i]$ stores the parent of i in a tree structure. The **root** of an element is found by scaling up to parents until it finds an element that is its own parent ($id[i] == i$).

*Two elements are connected if they have the same **root**.*

Operations:

- **Find/Connected:** compares the **roots** of p and q . (tree height proportional time-complexity: $O(N)$ worst-case scenario.)
- **Union:** connects the **root** of p to the root of q . (tree height proportional time-complexity: $O(N)$ worst-case scenario.)

Improving Quick-Union

There are faster ways to implement the **quick-union** algorithms:

Weighted Quick-Union:

- Avoid tall trees.
 - Keep track of number of objects in each tree.
 - Balance by linking root of smaller tree to root of larger tree.
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Module 2 Assignment

Percolation Problem:

Given a $N \times N$ matrix, check if there exists a path of full-open nodes between any open node in first row and any open node in last row. **Model a Percolation system.**

```
public class Percolation {  
  
    private static final int TOP = 0;  
  
    private final boolean[][] opened;  
  
    private final int size;  
  
    private final int bottom;  
  
    private int openSites;  
  
    private final WeightedQuickUnionUF qf;  
  
    public Percolation(int n) {  
  
        if (n <= 0) {  
  
            throw new IllegalArgumentException();  
  
        }  
  
        size = n;  
  
        bottom = size * size + 1;  
  
        qf = new WeightedQuickUnionUF(size * size + 2);  
  
        opened = new boolean[size][size];  
  
        openSites = 0;  
  
    }  
  
    public void open(int row, int col) {  
  
        checkException(row, col);  
  
        opened[row - 1][col - 1] = true;  
  
        ++openSites;  
  
        if (row == 1) {  
  
            qf.union(getQuickFindIndex(row, col), TOP);  
  
        }  
  
        if (row == size) {  
  
            qf.union(getQuickFindIndex(row, col), bottom);  
  
        }  
  
    }  
  
}
```



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```
}

if (row > 1 && isOpen(row - 1, col)) {

    qf.union(getQuickFindIndex(row, col), getQuickFindIndex(row - 1, col));

}

if (row < size && isOpen(row + 1, col)) {

    qf.union(getQuickFindIndex(row, col), getQuickFindIndex(row + 1, col));

}

if (col > 1 && isOpen(row, col - 1)) {

    qf.union(getQuickFindIndex(row, col), getQuickFindIndex(row, col - 1));

}

if (col < size && isOpen(row, col + 1)) {

    qf.union(getQuickFindIndex(row, col), getQuickFindIndex(row, col + 1));

}

}

private void checkException(int row, int col) {

    if (row <= 0 || row > size || col <= 0 || col > size) {

        throw new IllegalArgumentException();

    }

}

public boolean isOpen(int row, int col) {

    checkException(row, col);

    return opened[row - 1][col - 1];

}

public int numberOfOpenSites() {

    return openSites;

}

public boolean isFull(int row, int col) {

    if ((row > 0 && row <= size) && (col > 0 && col <= size)) {

        return qf.find(TOP) == qf.find(getQuickFindIndex(row, col));

    }

}
```



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```
        else throw new IllegalArgumentException();  
    }  
  
    private int getQuickFindIndex(int row, int col) {  
        return size * (row - 1) + col;  
    }  
  
    public boolean percolates() {  
        return qf.find(TOP) == qf.find(bottom);  
    }  
}
```

Notes:

- **Public:** public access by other classes; **Private:** restrict access to its own class;
- **Static:** does not need an object's instance to be accessed;
- **Final:** constant value;
- **Void:** non-returning method;
- **int/String/boolean/char/double:** specifies the data type of variable or returning method.