Saikrishna Arcot (edits by M. Hudachek-Buswell)

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Problems with Binary Search Trees

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Problems with Binary Search Trees

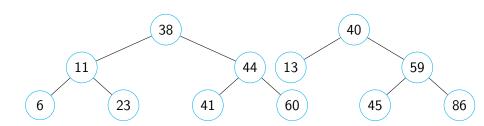
- A binary search tree has, on average, a big-O of $O(\log n)$ for adding, searching, and removing.
- However, in the worst case, all three operations can become O(n) when the tree is unbalanced. This can easily be done by adding the items in sorted order.
- There needs to be a way to keep the tree balanced while adding or removing, and make sure that all three operations have a worst case of O(n).

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- The idea is that whenever you're adding or removing a data item, you check to see if the other nodes are balanced; if they're not balanced, then you make rotations to balance the tree.
- AVL trees are just balanced BSTs; in other words, all AVL trees are also BSTs, but not all BSTs are AVL trees.

Example AVL Trees



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- Searching in an AVL tree is the same as searching in a BST.
- Adding to and removing from an AVL tree has the same basic steps as adding to and removing from a BST, but there is more work you need to do for an AVL tree.
- In addition to each node having a certain height, each node in an AVL tree also has a balance factor.

Balance Factor

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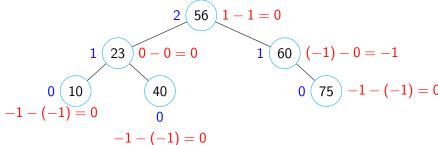
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- The balace factor of a node is calculated by subtracting the height of the left subtree by the height of the right subtree.
- A balance factor of 0 means that the subtree is perfectly balanced and that no rotations should be performed.

 A balance factor between -1 and 1 (inclusive) means that the subtree is considered balanced, but is slighly heavier on the left side (if the balance factor is positive) or the right side (if the balance factor is negative), and that no rotations needs to be done.

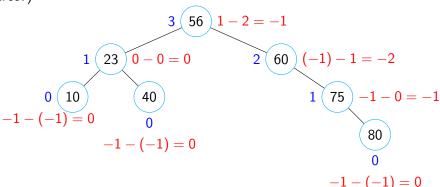
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 - This is because doing any combination of rotations will not make the subtree perfectly balanced.
- Any other balance factor means that the subtree is not balanced and that one or more rotations need to be done.

In the following AVL tree, the height of each node is in blue, and the balance factor of each node is in red.



In the following BST, the height of each node is in blue, and the balance factor of each node is in red. (Why is this not an AVL tree?)



 In AVL trees, to balance a tree after adding or removing something, one or more rotations are done.

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- There are two single rotations and two double rotations to choose from. (Double rotations are just two single rotations put together.)

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- There are two single rotations and two double rotations to choose from. (Double rotations are just two single rotations put together.)
- Which rotations are done (and which rotations can be done) depend on the balance factor of each node.

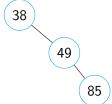
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- In the examples that follow, A, B, C, and D represent subtrees that may (or may not) exist. This is so that you can see where the subtrees move to after the rotation.

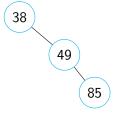
 A left rotation happens when you have a node whose balance factor is -2, and its right child has a balance factor of 0 or -1.

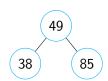
- A left rotation happens when you have a node whose balance factor is -2, and its right child has a balance factor of 0 or -1.
- When a left rotation happens, the top node becomes the left child of the middle node.

Ignoring any child subtrees:

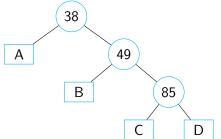


Ignoring any child subtrees:

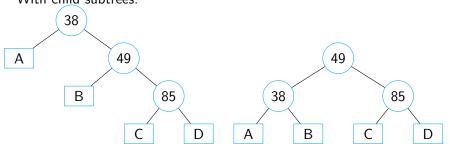




With child subtrees:



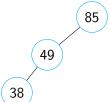
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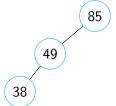
• A right rotation happens when you have a node whose balance factor is 2, and its left child has a balance factor of 0 or 1.

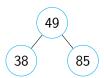
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- When a right rotation happens, the top node becomes the right child of the middle node.

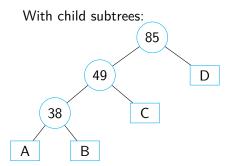
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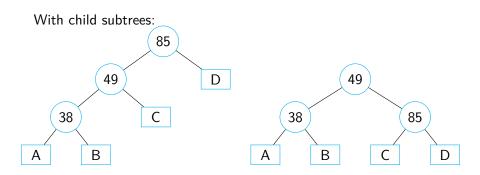
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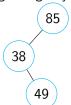
Right Rotation

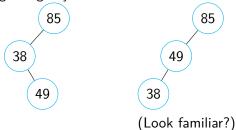


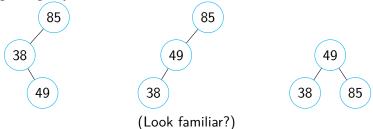
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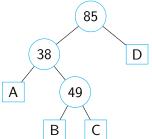
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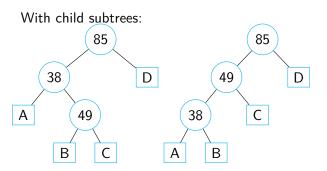


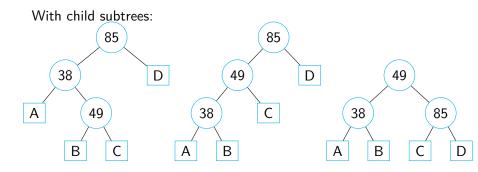




With child subtrees:



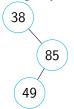




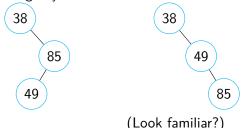
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Ignoring any child subtrees:



(Look familiar?)

Ignoring any child subtrees:

38

85

49

49

49

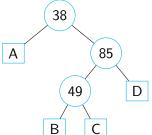
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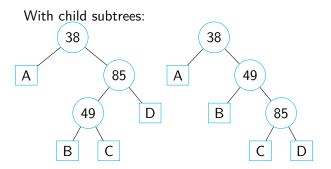
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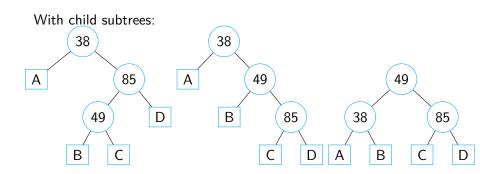
(Look familiar?)

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With child subtrees:







Adding and Removing

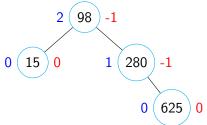
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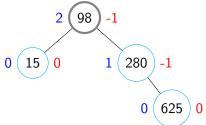
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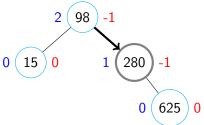
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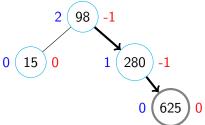
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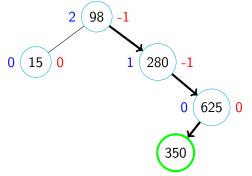
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- If a rotation is done, the heights and balance factors of the nodes involved are updated again.

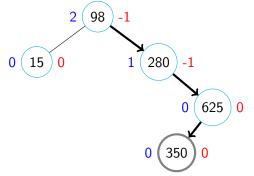


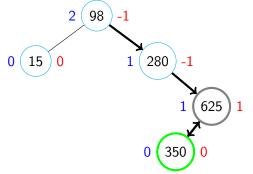


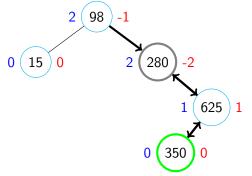


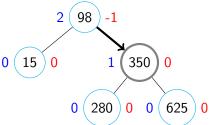


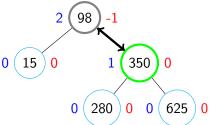


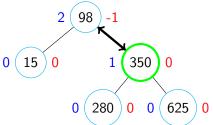




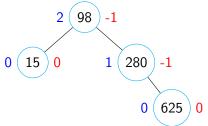




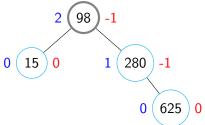




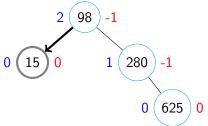
Removing

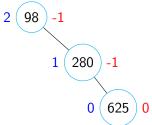


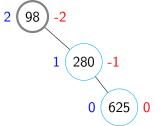
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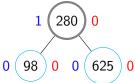


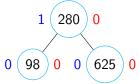
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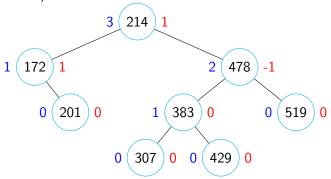


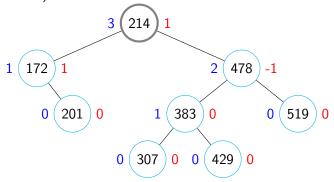


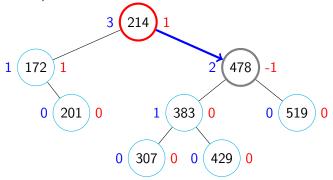


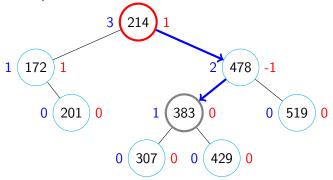


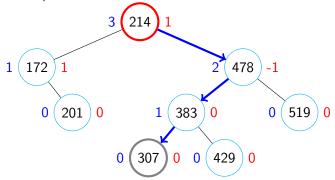


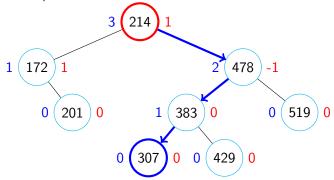


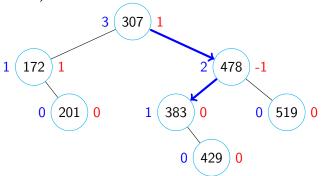


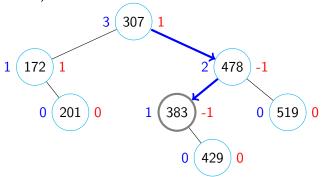


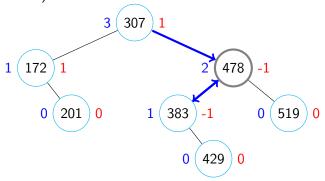


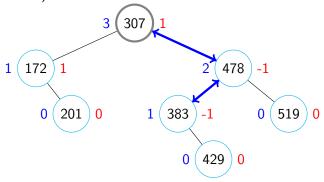


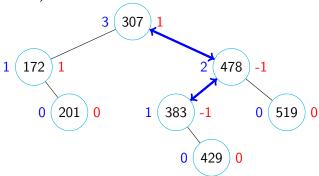












Adding

```
procedure ADD(data, node)
   if node is not valid then
      return new node containing data
   else
      if data = node.data then
          Implementation-defined behavior
      else if data < node.data then
          node.left = Add (data, node.left)
      else
          node.right = Add(data, node.right)
      end if
```

Adding

```
Update heights and balance factors
         if subtree is out of balance then
             Balance subtree
             node \leftarrow new root of subtree
             Update heights and balance factors of node
         end if
         return node
      end if
  end procedure
(Removing is done in a similar way.)
```



Subtree Update Method

Remember this from BSTs?

Subtree Update Method

- Remember this from BSTs?
- This method will make coding an AVL tree easier because you can easily update the left/right child of the parent node after doing a rotation.

Performance

• The average case of insertion, search, and deletion for an AVL tree are all $O(\log n)$.

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- The average case of insertion, search, and deletion for an AVL tree are all $O(\log n)$.
- If the data items are added in sequential order, then the rotations that would be done would prevent the tree from turning into a linked list. Therefore, the worst case is also O(log n).