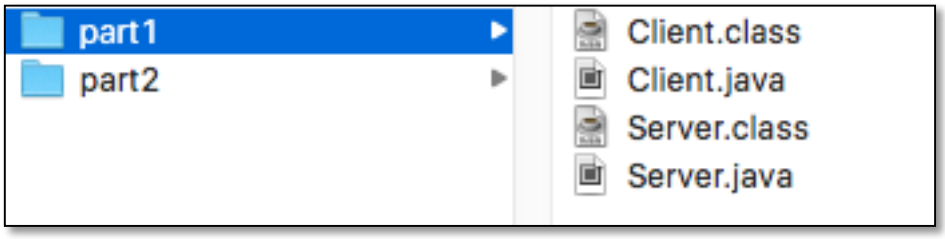
Report Document of CS655 PA1

Part1

Directory:



Compilation instruction:

Client.java and Server.java are included in a package called "part1", they can be compiled with command:

java Client.java
java Server.java

When we want to run the program, cd to the upper directory, and use:

java part1.Client
java part1.Server

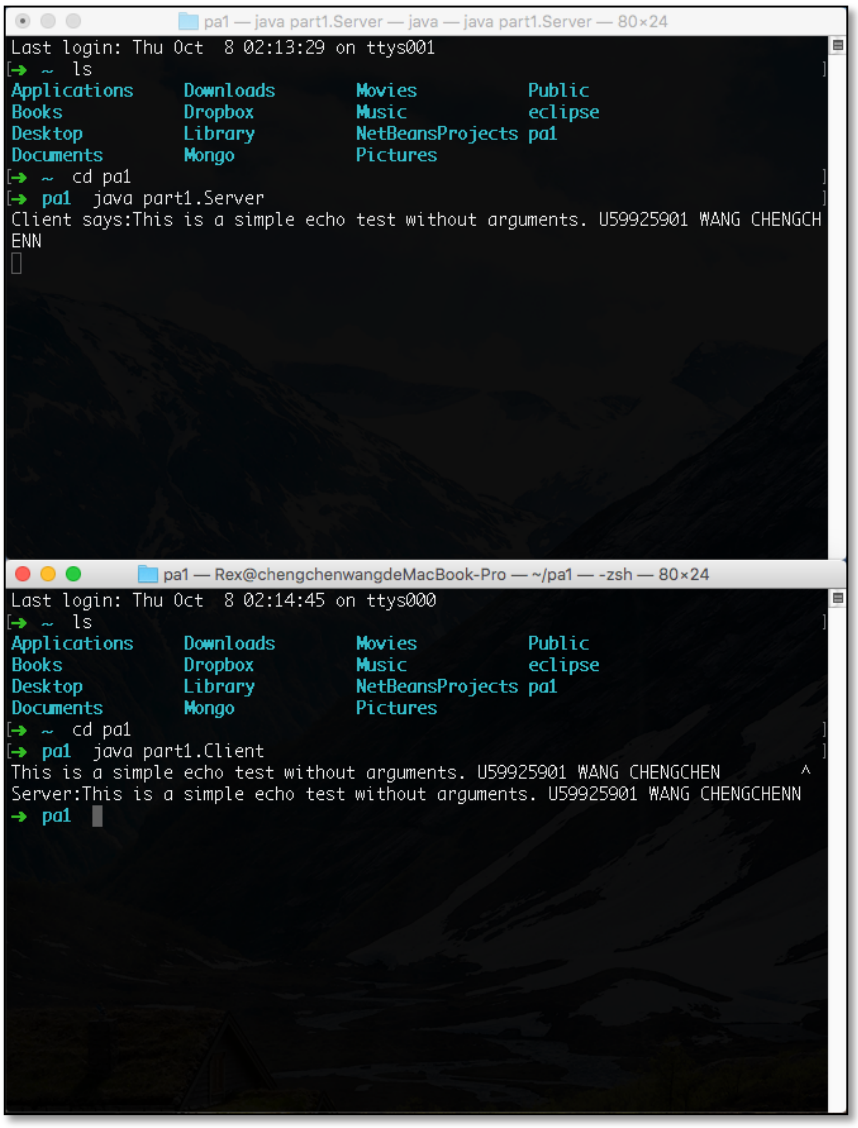
Argument Specification:

When using client, and user wants to specify Server IP address and port number, just run the client with command arguments like (use whitespace between IP and port):

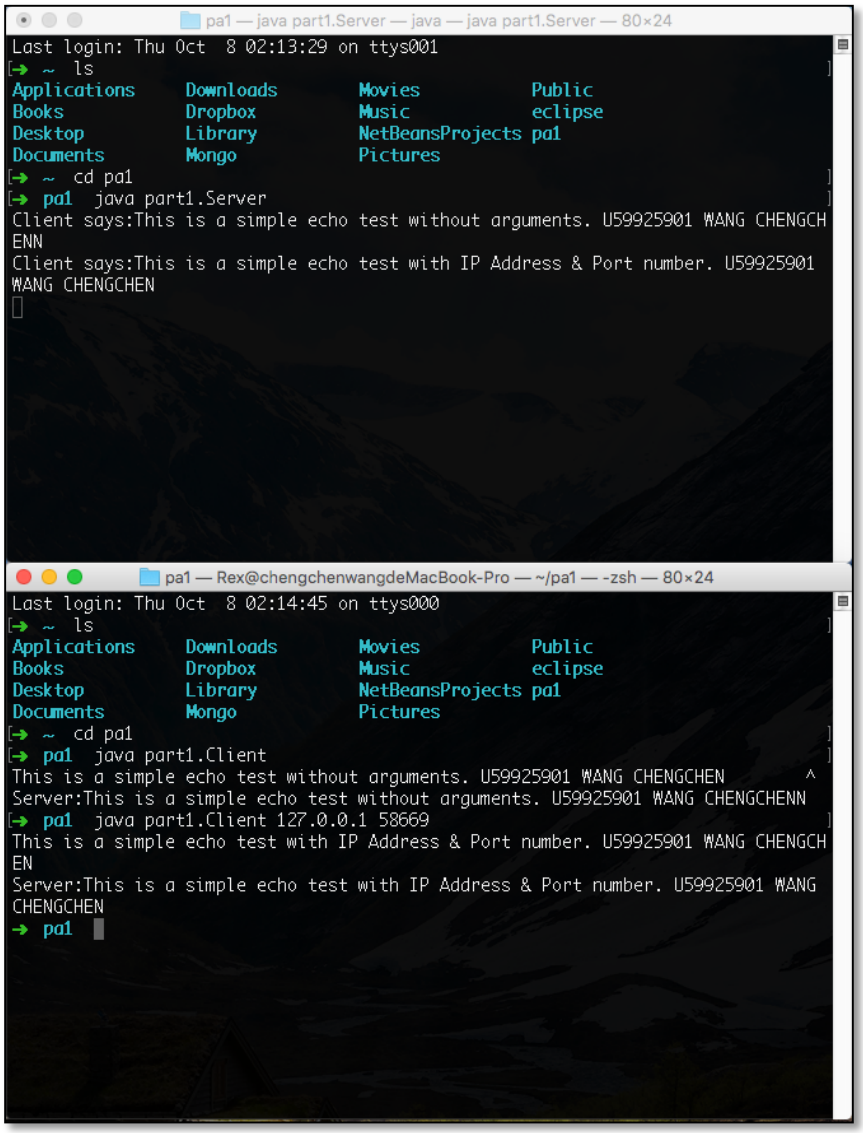
java part1.Client 192.168.1.1 58669

Testing Results:

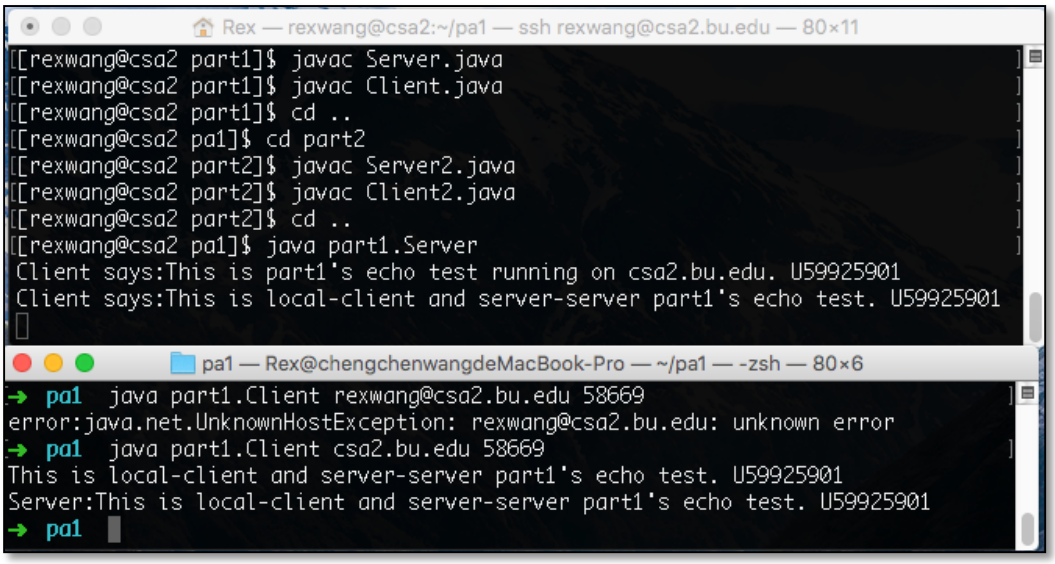
(1) Run client and server both locally without specifying IP and Port:



(2) Run client and server both locally with my specifying IP and Port (127.0.0.1 58669):



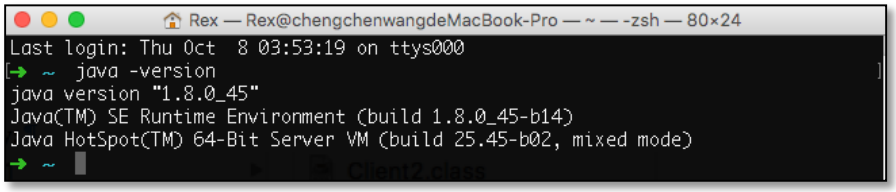
(3) Run client locally and server on csa2@bu.edu with my specifying Port number 58669 (client runs on a different machine than server):



Part2

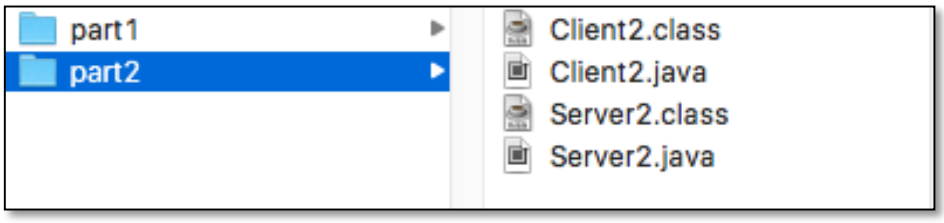
Experiment Machine:
Type: MacBook Pro 13.3 retina with Mac OS 10.11

Java: 1.8



Network: Comcast Service

Directory:



Compilation instruction:

Client2.java and Server2.java are included in a package called "part2", they can be compiled with command:

javac Client2.java

javac Server2.java

When we want to run the program, cd to the upper directory, and use:

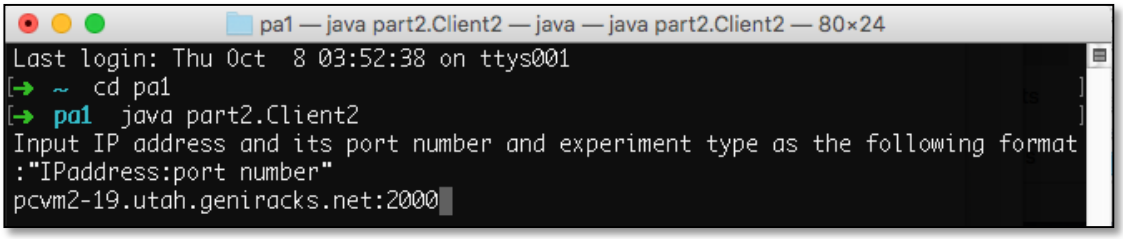
java part2.Client2

java part2.Server2

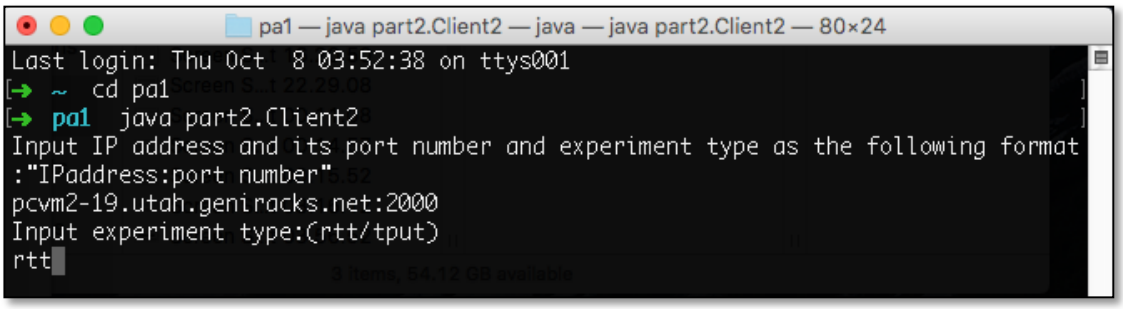
Argument Specification:

When using Client2, it required you to input several arguments to finish the experiment:

Input Server IP address: port number like: pcvm2-19.utah.geniracks.net:2000

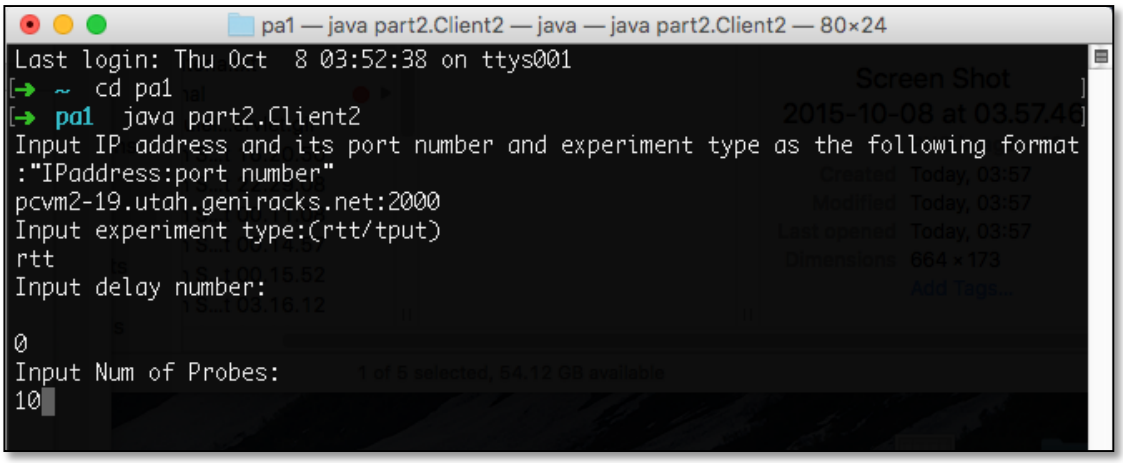


Input experiment type like: "rtt" or "tput" (without quotation mark and must be either rtt or tput)

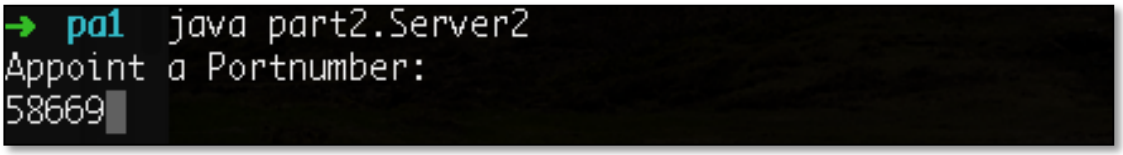


Input delay number like: 0

Input number of probes (at least 10 times) like: 10



In another hand, when using Server2, it required you to appoint a port number to run on, like:



Expected results of Client2:

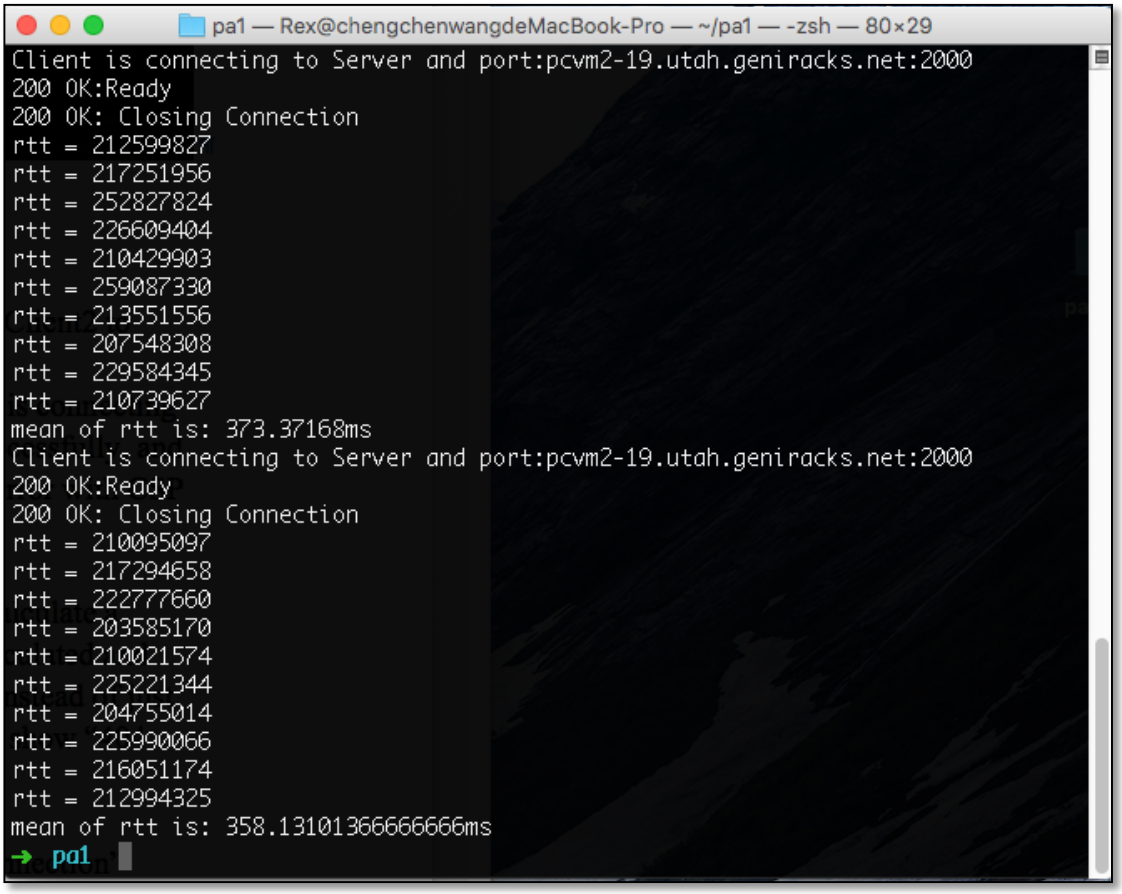
For each size of packet:

During the CSP-MP-CTP method experiment, I designed and expected Client2 to show information to tell the status of the connection:

In CSP, the Client2 is supposed to show: "200 OK: Ready" and “Client is connecting to Server and port:pcvm2-19.utah.geniracks.net:2000” if it connects successfully, and show “404 ERROR: Invalid Connection Setup Message” if there is an error with CSP message.

During the MP, the Client2 will show rtt values for every probes, and calculate a average of rtt/throughput value. (Note that this information will be calculated and showed after ctp, because we must calculate the whole round trip time instead of just mp's trip time) If there is something wrong with the mp message, it will show "404 ERROR: Invalid Measurement Message".

During the CTP, the Client2 is supposed to show "200 OK: Closing Connection" normally, and show "404 ERROR: Invalid Connection Termination Message" from server response, if there is something wrong with ctp message.



Design of program-Client2:

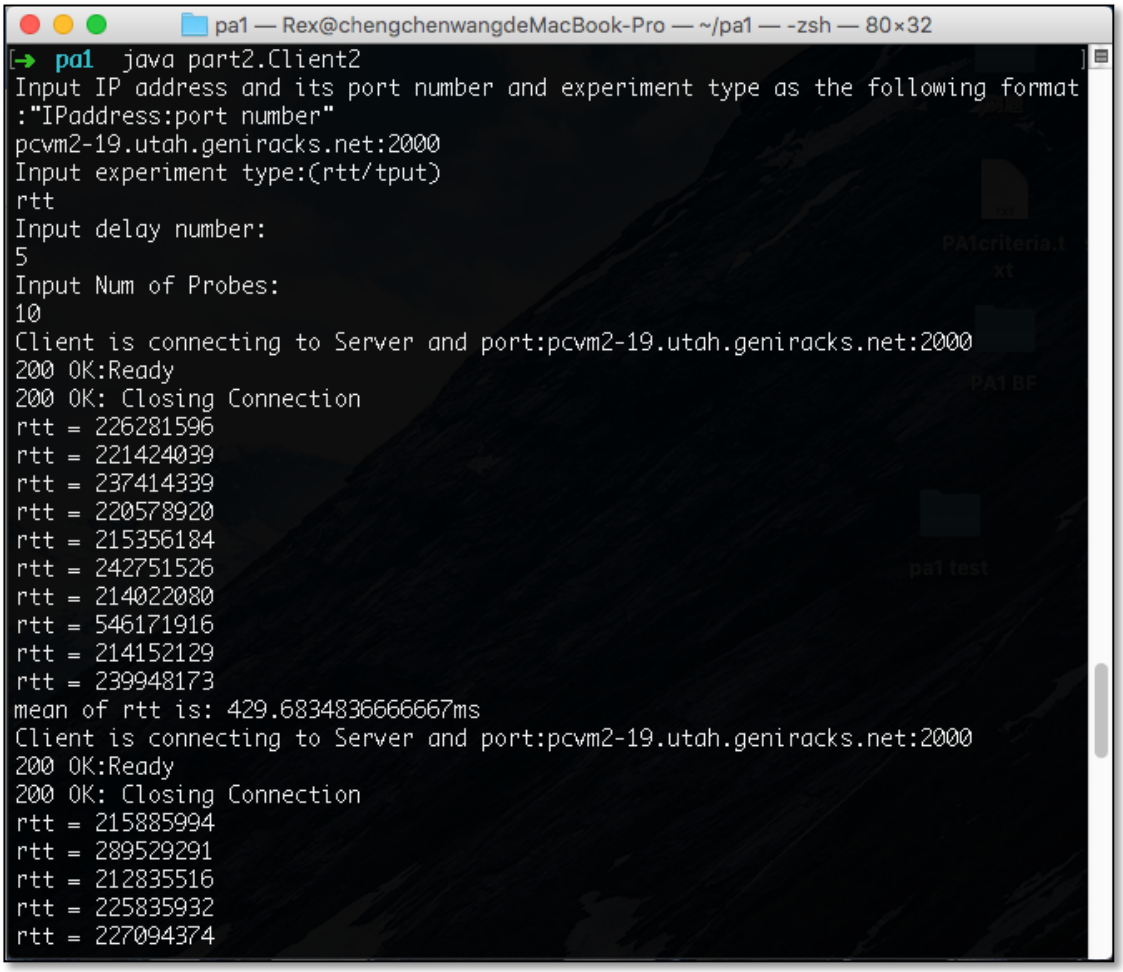
Client2 first define initializes some basic arguments, and then use method Scanner and Bufferedreader to scan the users' inputs, which provides an interactive interface to specify Sever IP address and port number, this experiment's type (rtt/ tput), server delay, number of probes (at least 10). Then Client2 uses these specified arguments to set up a socket connection with the server and start CSP-MP-CTP method process, using bufferedReader to read messages form server, and using printWriter with flush() to send messages to the server. During this process, whenever Client2 get a "404 ERROR" from server, it prints the message out and close the socket right away. During MP phase, it notes down the rtt values for each probe, after the CTP phase, Client2 calculate the mean rtt value or the mean throughput value, determined by what type of this experiment is.

Design of program-Server2:

After users entering port number to run on, the Server2 establish a server socket by using method ServerSocket. And in a "while(true)" loop, which means Server2 continuously accepting connections from clients, Server2 examines CSP messages/ MP messages /CTP messages by splitting the messages by "whitespace". It examines protocol type, number of probes and so on. When it comes to sequence number, Server2 uses a loop to self add the sequence number and examines it with the sequence number accepting from clients. During the process, each error type of message cause a socket close and Server2 prints the type of error out.

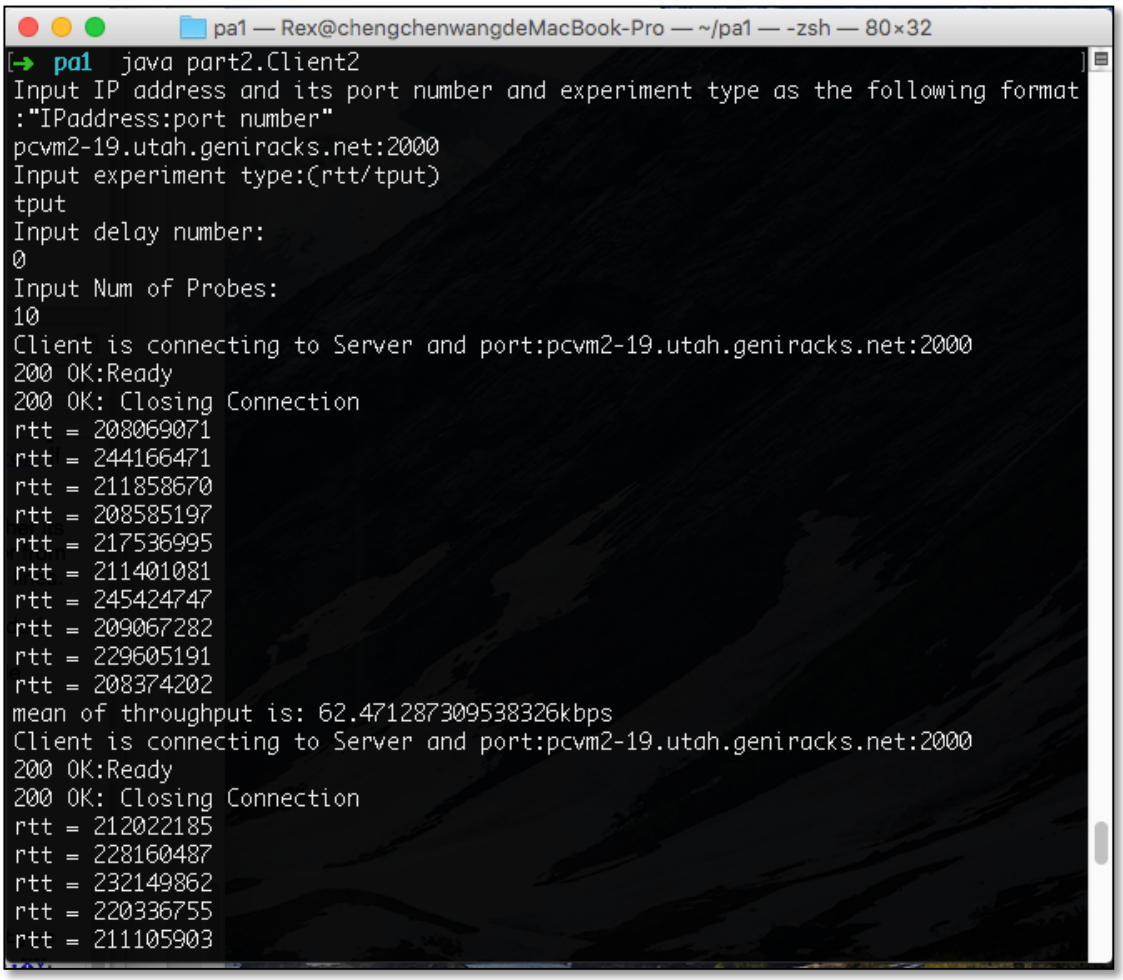
Description and results of Experiment:

For RTT: using **pcvm2-19.utah.geniracks.net:2000** as the Server, and use experiment type **rtt**, from **1byte to 100, 200, 400, 800, 1000 bytes**, for each size, the number of probes is **10**, so there will be 10 rtt values for each size. Calculate each size's **average rtt value**, so we get 60 results and 6 average values. Varying delay number from **0 to 10 to 100**, and we can get different groups of average rtt values:



RTT Average(ms)				
size(bytes)\delay	0	1	10	100
1	351.327	363.14403	381.28389	538.41334
100	366.3712	380.34703	386.00675	531.83799
200	352.5901	372.09824	374.64523	534.81884
400	361.9832	363.19053	395.191	514.05933
800	373.3717	379.38718	379.11662	523.21387
1000	358.131	407.81461	376.39269	529.79634

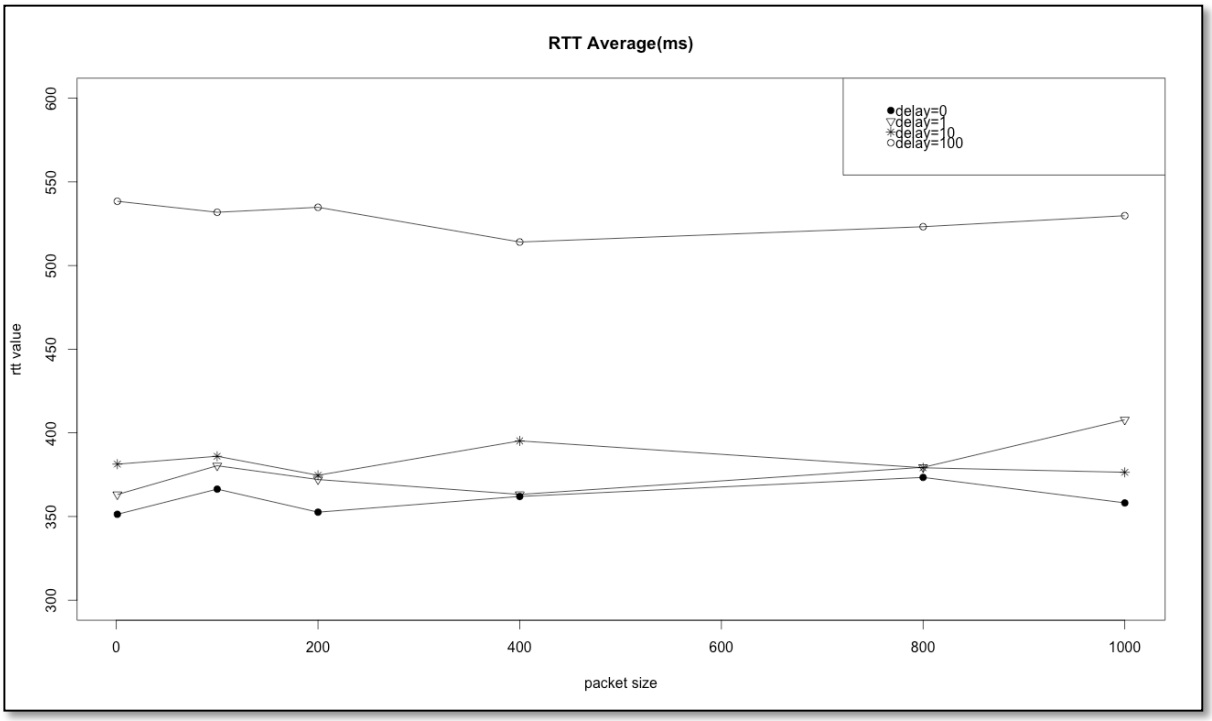
For Tput: using **pcvm2-19.utah.geniracks.net:2000** as the Server, and use experiment type **tput**, from **1k byte to 2k, 4k, 8K, 16K 32K bytes**, for each size, the number of probes is **10**, so there will be 10 tput values for each size. Calculate each size's average tput value (using size/rtt, and then convert the unit of value to kbps), so we get 60 results and 6 average values. Varying delay number from 0 to 10 to 100, and we can get different groups of average tput values:



Tput(Throughput) Average(kbps)				
size(K bytes)\delay	0	1	10	100
1	62.47129	62.736546	52.046581	43.546032
2	125.3883	121.09124	119.84405	85.884168
4	259.1951	250.36827	246.67293	156.97092
8	515.9116	493.81537	480.78444	320.07234
16	1039.965	1006.3825	912.62213	698.79173
32	1963.322	1991.5436	1879.2058	1379.2273

Graphs and analysis:

RTT:



Plotted by R language:

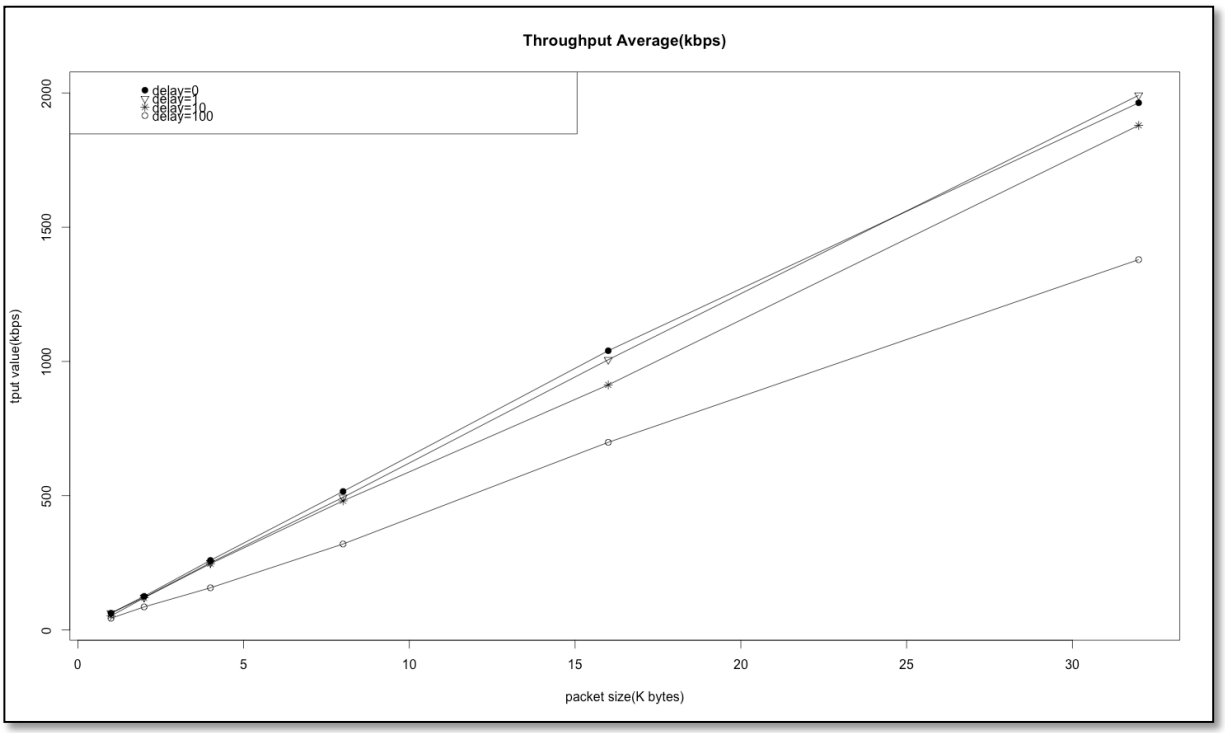
```
plot(x,y,main="RTT Average(ms)",xlab="packet size",ylab="rtt value",ylim=c(300,600),pch=19)
legend("topright",pch=c(19,25,8,21),legend=c("delay=0","delay=1","delay=10","delay=100"),x.intersp=0.1,y.intersp = 0.2)
points(x,y2,main="RTT Average(ms)",ylim=c(300,600),pch=25)
points(x,y3,main="RTT Average(ms)",ylim=c(300,600),pch=8)
points(x,y4,main="RTT Average(ms)",ylim=c(300,600),pch=21)

lines(x,y,main="RTT Average(ms)",ylim=c(300,600))
lines(x,y2,main="RTT Average(ms)",ylim=c(300,600))
lines(x,y3,main="RTT Average(ms)",ylim=c(300,600))
lines(x,y4,main="RTT Average(ms)",ylim=c(300,600))
```

According to the image above, when delay=0, we find that different packet sizes have nearly no effects on mean rtt values, when delay=1/10/100, we also find that different packet sizes have nearly no effects on mean rtt values. But as the delay increasing from 1 to 100, the mean rtt values are increasingly as well. So we know that even though increasing the server delay merely increases the processing time at the server,

it causes the feedback delay, which observed by us, has an effect similar to increasing the path's propagation.

Tput:



Plotted by R language:

```
plot(x,y1,main="Throughput Average(kbps)",xlab="packet size(K bytes)",ylab="tput value(kbps)",ylim=c(40,2000),pch=19)
legend("topleft",pch=c(19,25,8,21),legend=c("delay=0","delay=1","delay=10","delay=100"),x.intersp=0.1,y.intersp = 0.3)
points(x,y2,main="Throughput Average(kbps)",ylim=c(40,2000),pch=25)
points(x,y3,main="Throughput Average(kbps)",ylim=c(40,2000),pch=8)
points(x,y4,main="Throughput Average(kbps)",ylim=c(40,2000),pch=21)

lines(x,y1,main="Throughput Average(kbps)",ylim=c(40,2000))
lines(x,y2,main="Throughput Average(kbps)",ylim=c(40,2000))
lines(x,y3,main="Throughput Average(kbps)",ylim=c(40,2000))
lines(x,y4,main="Throughput Average(kbps)",ylim=c(40,2000))
```

According to the image above, when delay-0, as the packet size is increasing, we find that the mean throughput values are also increasing. It's not a liner function, but the increasing speed is really high and nearly equals to liner. So we can know that when we increase a packet size, the path's throughput will be correspondingly increase. Besides, when delay increases from 0 to 100, the increase speed of throughput somewhat decreases, as the image shows, but the overall trend does not change.