Combinatorial Multi-armed Bandits &

 α, β - Approximation

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Presentation Overview

- 1. Multi armed bandit (MAB)
- 2. Combinatorial multi armed bandit (CMAB)
- 3. General CMAB Framework
- 4. Main theorem
- 5. Comparing to classical MAB

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Multi armed bandit

- m arms
- reward
- optimal arm
- regret
- tradeoff between exploration and exploitation

real world application

- online advertising scenario
- website contains a set of web pages
- advertiser can select at most k web pages

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- super arms: any set of arms
- offline computation oracle computes the optimal super arm
- (α, β) -approximation oracle
- $\alpha, \beta < 1$
- α : oracle outputs a super arm whose expected reward is at least α fraction of the optimal expected reward
- β : success probability
- (α, β) -approximation regret

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definition of CMAB framework

- lack m m arms associated with a set of random variables $X_{i,t}$ for $1 \leq i \leq m$ and $t \geq 1$
- super arms : every set of arms
- $igcup X_{i,t}$: a random variable with bounded support on [0,1]
- $igcup X_{i,t}$: random outcome of the i-th arm in its t-th trial
- $igoplus \mu = (\mu_1, \mu_2, ..., \mu_m)$: vector of expectations of all arms
- T_{i,t}: number of times the outcome of arm i is revealed after the first t rounds
- $igcup R_t(S)$: non-negative random variable denoting the reward of super arm s

$\mathbb{E}[R_t(S)]$: a function of only the set of arms S and the expectation vector

General CMAB Framework

- $r_{\mu} = \mathbb{E}\left[R_t(S)\right]$
- Monotonicity : if for all $i \in [m]$, $\mu_i \leq \mu'_i$, we have $r_{\mu}(S) \leq r_{\mu'}(S)$ for all S
- Bounded smoothness: there exists a strictly increasing function f(.) such that, we have $|r_{\mu}(S) - r_{\mu'}(S)| \leq f(\Lambda)$ if $\max_{i \in S} |\mu_i - \mu_i'| \leq \Delta$

CMAB algorithm

- Algorithm A selects S_t^A
- maximize $\mathbb{E}_{S,R}\left[\sum_{t=1}^{n}R_{t}(S_{t}^{A})\right]=\mathbb{E}_{S}\left[\sum_{t=1}^{n}r_{\mu}(S_{t}^{A})\right]$

- (α, β) -approximation oracle : $\mathbb{P}[r_{\mu}(S) \geq \alpha.opt_{\mu}] \geq \beta$
- $\mathsf{Reg}_{\mu,\alpha,\beta}^A(n) = n.\alpha.\beta.\mathsf{opt}_{\mu} \mathbb{E}_{\mathcal{S}}\left[\sum_{t=1}^n r_{\mu}(\mathcal{S}_t^A)\right]$

CMAB algorithm

empirical mean :

$$\hat{\mu}_i = (\sum_{i=1}^s X_{i,j})/s$$

adjusted mean :

$$\bar{\mu}_i = \hat{\mu}_i + \sqrt{\frac{3\ln(t)}{2T_i}}$$

- 1: For each arm i, maintain: (1) variable T_i as the total number of times arm i is played so far; (2) variable µ̂_i as the mean of all outcomes X_{i,*}'s of arm i observed so far.
- 2: For each arm i, play an arbitrary super arm $S \in \mathcal{S}$ such that $i \in S$ and update variables T_i and $\hat{\mu}_i$.
- 3: $t \leftarrow m$.
- 4: while true do
- 5: $t \leftarrow t + 1$.
- 6: For each arm i, set $\bar{\mu}_i = \hat{\mu}_i + \sqrt{\frac{3 \ln t}{2T_i}}$.
- 7: $S = \operatorname{Oracle}(\bar{\mu}_1, \bar{\mu}_2, \dots, \bar{\mu}_m).$
- 8: Play S and update all T_i 's and $\hat{\mu}_i$'s.
- 9: end while

 ${\bf Algorithm~1:~CUCB~with~computation~oracle}$

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Preliminaries

- Super arm S is bad if $r_u(S) < \alpha.opt_u$
- $S_B = \{S | r_\mu(S) < \alpha.opt_\mu\}$
- $\Delta_{max}^{i} = \alpha . opt_{u} min\{r_{u}(S)|S \in S_{B}, i \in S\}$
- $\Delta_{max} = \mathsf{max}_{i \in [m]} \, \Delta_{max}^i$
- $\Delta_{min} = \min_{i \in [m]} \Delta_{min}^i$

Theorem

Theorem

The (α, β) -approximation regret of the CUCB algorithm in n rounds using an (α, β) -approximation oracle is at most

$$\sum_{i \in [m]} \int_{\Delta^i} \left(\frac{6 \ln(n) \Delta^i_{min}}{\left(f^{-1} \left(\Delta^i_{min}\right)\right)^2} + \int_{\Delta^i_{min}}^{\Delta^i_{max}} \frac{6 \ln(n)}{\left(f^{-1} (x)\right)^2} dx \right) + \left(\frac{\pi^2}{3} + 1\right) .m. \Delta_{max}$$

simplified form :

$$\left(\frac{6\ln(n)}{\left(f^{-1}(\Delta_{min})\right)^2} + \frac{\pi^2}{3} + 1\right).m.\Delta_{max}$$

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Comparing to classical MAB

- every super arm is a simple arm
- f(x) = x
- $\Delta_{max}^{i} = \Delta_{min}^{i}$ $\Delta^{i} = \max_{i \in [m]} \mu_{i} \mu_{i}$
- the regret bound of the classical MAB

$$\sum_{i \in [m], \Delta^i > 0} \frac{6 \mathit{ln}(n)}{\Delta^i} + \left(\frac{\pi^2}{3} + 1\right).m.\Delta_{\mathit{max}}$$

lacksquare the coefficient of the regret bound in (Auer et al., 2002) : $\sum_{i\in [m],\Delta^i>0}rac{8}{\Delta^i}$