5. Find the equivalent partition of the following machine.

	Next State,z		
Present State	X = 0	$\overline{X} = 1$	
A	C, 1	D, 0	
В	D, 1	E, 0	
\mathbf{C}	B, 1	E, 1	
D	B, 1	A, 0	
${ m E}$	D, 1	B, 1	

Solution: The partitions are

 $P_0 = (ABCDEF)$

 $P_1 = (AD)(BCE)$ (Depending on o/p for i/p 1)

 $P_2 = (AD)(BE)(C)$ (For i/p 0, the next state of C goes to another set)

 $P_3 = (A)(D)(BE)(C)$ (For i/p 0, the next state of A and D goes to a different set)

 $P_4 = (A)(D)(BE)(C)$

As P_3 and P_4 are the same, P_3 is the equivalent partition.

Minimization: We know that the equivalent partition is unique. So,

 $P_4 = (A)(D)(BE)(C)$ is the unique combination. Here, every single set represents one state of the minimized machine.

Let us rename these partitions for simplification.

Rename (A) as S_1 , (BE) as S_2 , (C) as S_3 , and (D) as S_4 .

The minimized machine becomes

	Next State,z		
Present State	X = 0	$\overline{X} = 1$	
$S_1(A)$	$S_3, 1$	$S_4, 1$	
$S_2(\mathrm{BE})$	$S_4, 1$	$S_2, 0$	
$S_3(\mathbf{C})$	$S_2, 1$	$S_2, 1$	
$S_4(D)$	$S_2, 1$	$S_1, 0$	

6. Simplify the following incompletely specifi ed machine.

	Next State,z			
Present State	$\overline{I_1}$	I_2	I_3	
A	D, 1	E,1	_	
В	B, 0	E,-	C,-	
\mathbf{C}	C, -	C,0	В, –	
D	B, 0	D,-	E,-	
E	_	B,0	A, –	

Solution: Put a temporary state T in the next state place, where the next states are not specifi ed. If the output is not mentioned, there is no need to put any output.

As a temporary state T is considered, T is put in the present state column with the next state T for all inputs with no output. The simplified machine becomes

	Next State,z			
Present State	$\overline{I_1}$	I_2	I_3	
A	D, 1	E,1	T,-	
В	B, 0	E,-	C,-	
\mathbf{C}	C, -	C,0	B, –	
D	B, 0	D,-	E,-	
${ m E}$	T, -	$_{\rm B,0}$	A, –	
${ m T}$	T, 0	T,0	T,0	

7. Minimize the following incompletely specifi ed machine.

	Next State,z			
Present State	$\overline{I_1}$	I_2	$\overline{I_3}$	
A	A, 1	D, -	$\overline{\mathrm{C},-}$	
В	A, –	D, –	E,-	
\mathbf{C}	E, 0	A, 1	-,-	
D	E, –	A, 1	-,-	
E	E, 0	—, —	C,-	

Solution: In an incompletely specifi ed machine, all the next states or all the outputs or both are not mentioned. We can minimize an incompletely specifi ed machine by using the merger graph and compatible graph method. From the merger graph, we have to find the compatible pair and from that compatible pair

and implied pair we have to construct compatible graph. From the compatible graph, we have to find the closed partition. And, the closed partitions are the minimized states of the machine.

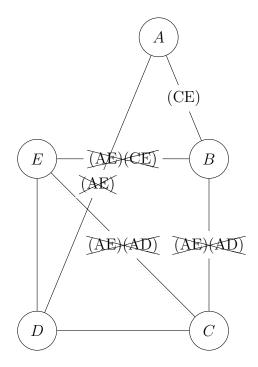
Merger Graph: The machine consists of fi ve states. So, the merger graph consists of fi ve nodes, named A, B, C, D, and E.

The outputs of A and B do not differ, and so there is an arc between A and B. For input I3, the next state conflicts. so the arc is an interrupted arc and in the interrupted portion the conflicting next state pair (CE) is placed.

For states A and C, the output conflicts, and so there is no arc between A and C. For the states C and D, the outputs as well as next states do not conflict. So, an uninterrupted arc is placed between C and D.

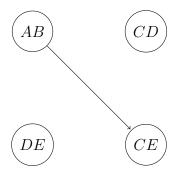
By this process, the merger graph is constructed.

The merger graph is as follows As there is no arc between A and E, the arc between (AD), (BE), (BD), and (BC) are also crossed.



Compatible Graph: (CE) is the implied pair for (AB). A directed arc is drawn from (AB) to (CE). The compatible graph consists of four vertices (AB), (CD), (DE), and (CE).

A subgraph of a compatibility graph is said to be closed cover for the machine, if for every vertex in the subgraph, all outgoing edges and their terminal vertices also belong to the subgraph, and every state of the machine is covered by at least one vertex of the subgraph.



For the constructed compatible graph, (AB), (CD), and (CE) form a closed covering.

The states of the minimized machine are (AB) and (CDE).

Rename (AB) as S_1 and (CDE) as S_2 .

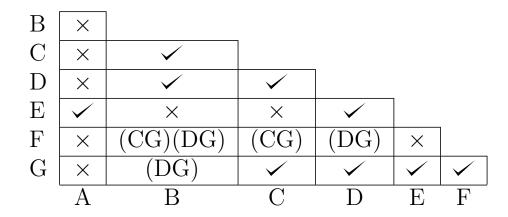
The minimized machine is

	Next State,z		
Present State	$\overline{I_1}$	I_2	$\overline{I_3}$
$\overline{S_1}$	$S_1, 1$	$S_2, -$	$\overline{S_2, -}$
S_2	$S_2, 0$	$S_1, 1$	$S_2, -$

8. Construct a compatible graph for the following incompletely specifi ed machine.

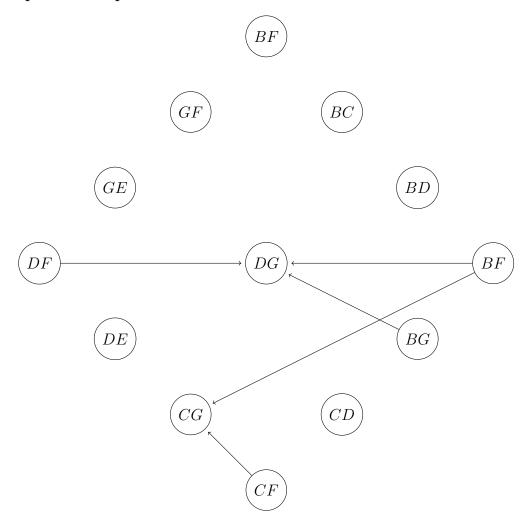
	Next State,z			
Present State	$\overline{I_1}$	I_2	I_3	I_4
A	A, O_1	E, O_2	, _	A, O_2
В	-, -	C, O_3	B, O_1	D, O_4
\mathbf{C}	A, O_1	C, O_3	-, -	-, -
D	A, O_1	-, -	-, -	D, O_4
${ m E}$	-, -	E, O_2	F, O_1	-, -
\mathbf{F}	-, -	G, O_3	F, O_1	G, O_4
G	A, O_1	-, -	-, -	G, O_4

Solution: To construct a compatible graph, we have to fi rst fi nd compatible pairs. To fi nd the compatible pairs, we need to construct a merger graph. But this machine has seven states, and so it is difficult to construct a merger graph for this machine. So, we have to construct a merger table to fi nd compatible pairs.



Compatible pairs are (AE), (BC). (BD), (CD), (CG), (DE), (DG), (EG), (GF), (BF), (BG), (CF) and (DF). If (CG) and (DG) are compatible, then (BF) is compatible. If (DG) is compatible, then (BG) is compatible. If (CG) is compatible, then (CF) is compatible.

Compatible Graph



The closed partitions are (AE), (BF), (BG), (CF), (CG), (DF), and (DG). The states of the minimized machine are (AE), (BCDFG). Rename them as S_1 , S_2 .

The minimized machine is

	Next State,z			
Present State	$\overline{I_1}$	I_2	I_3	$\overline{I_4}$
S_1	S_1, O_1	S_1, O_2	S_2, O_1	S_1, O_2
S_2	S_1, O_1	S_2, O_3	S_2, O_1	S_2, O_4

9. Consider the following machine M_1

	Next State,z			
Present State	$\overline{I_1}$	I_2	I_3	I_4
A		C, 1	E, 1	B, 1
В	E, 0	F, 1	_	
\mathbf{C}	F, 0	F, 1	_	
D	_		B, 1	_
${ m E}$	_	F, 0	A, 0	D, 1
F	C, 0		B, 0	C, 1

- a) Construct a merger table for M_1 .
- b) Find the set of compatibles.
- c) Draw a compatibility graph for M_1 . Describe the procedure used by you.
- d) Obtain a closed covering of M_1 .
- e) Construct a minimized machine M_1 * of M_1 .