

# Tutorial 6

## COMP 355: Introduction to Theoretical Computer Science

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1 Theorem

2 Application

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1 Theorem

2 Application

# Pumping Lemma

## Definition

If  $L$  is a regular language, then there is a number  $p$  (the pumping length) where, if  $s$  is any string in  $L$  of length at least  $p$ , then  $s$  may be divided into 3 pieces,  $s = xyz$ , satisfying the following conditions:

- 1 for each  $i \geq 0$ ,  $xy^iz \in L$
- 2  $|y| > 0$
- 3  $|xy| \leq p$

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# Pumping Lemma

## Question 1

Prove  $L = \{ww \mid w \in \Sigma^*\}$  is not regular.

## Question 2

Is the language  $L = \{w_1 w_2 \mid w_1, w_2 \in \{a, b\}^*, |w_1| = |w_2|\}$  regular?

## Question 3

Prove  $L = \{0^i1^j \mid i > j \geq 0\}$  is not regular.



# Pumping Lemma

## Question 4

Consider the language  $L = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and if } i = 1 \text{ then } j = k\}$

- 1 Show that  $L$  is not regular.
- 2 Show that  $L$  acts like a regular language in the pumping lemma. In other words, give a pumping length  $p$ , and demonstrate that  $L$  satisfies the 3 conditions of the pumping lemma for this value of  $p$ .
- 3 Explain why part (1) and (2) do not contradict the pumping lemma.

## Question 5

Prove  $L = \{1^{n^2} \mid n \geq 0\}$  is not regular.