Tutorial 3

COMP 5361: Discrete Structures and Formal Languages

Mohammad Reza Davari

Concordia University

Outline

Contents of the section

Nested Quantifiers

Definition

The case where one quantifier is within the scope of another quantifier.

Examples

Assuming the domain of discourse is $\mathbb{R}$ translate the followings to plain English:

1. $\forall x \forall y (x + y = y + x)$

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3. $\forall x \forall y \forall z (x + (y + z) = (x + y) + z)$

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1. $\forall x \forall y (x + y = y + x)$
2. $\forall x \exists y (x + y = 0)$
3. $\forall x \forall y \forall z (x + (y + z) = (x + y) + z)$
4. $\forall x \forall y ((x > 0) \wedge (y < 0) \rightarrow (xy < 0))$

Caution

Order matters...

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- Be careful with the order of existential and universal quantifiers!
- Quantifiers of the same kind can interchange places.

Example

- $\forall x \forall y \exists z (x + y = z)$

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    - True

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- $\forall x \forall y \exists z (x + y = z)$
  - True
- $\exists z \forall x \forall y (x + y = z)$

Example

- $\forall x \forall y \exists z (x + y = z)$
  - True
- $\exists z \forall x \forall y (x + y = z)$
  - False

Translation to English

- Translate the statement, $\forall x(C(x) \vee \exists y(C(y) \wedge F(x, y)))$ where:
  - $C(x)$: x has a computer
  - $F(x, y)$: x and y are friends

Translation to English

- Translate the statement, $\forall x(C(x) \vee \exists y(C(y) \wedge F(x, y)))$ where:
  - $C(x)$: x has a computer
  - $F(x, y)$: x and y are friends
- Translate the statement, $\exists x \forall y \forall z((F(x, y) \wedge F(x, z) \wedge (y \neq z)) \rightarrow \neg F(y, z))$ where:
  - $F(x, y)$: x and y are friends
  - Domain: All student in the class

Translate from English

Translate the followings into logical expressions:

- If a person is female and is a parent, then this person is someone's mother.

Translate from English

Translate the followings into logical expressions:

- If a person is female and is a parent, then this person is someone's mother.
- Everyone has exactly one best friend.

Negating Nested Quantifiers

Recursive Negation

Statements involving nested quantifiers can be negated by successively applying the rules for negating statements involving a single quantifier.

Example

Negate the following statement:

$$\exists w \forall a \exists f (P(w, f) \wedge Q(f, a))$$

for some f, P, and Q.

Contents of the section

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Definition

- **Argument:** An argument is a sequence of statements that end with a conclusion.
- **Valid Argument:** A valid argument is an argument that the conclusion, or final statement of the argument, follows from the **truth** of the preceding statements[1].
- **Fallacy:** An incorrect way of reasoning which lead to invalid arguments.

[1] Also called premise

Example

Check weather the following argument is valid or not.

$$\begin{array}{r} p \rightarrow q \\ p \\ \hline \therefore q \end{array}$$

The most important rule of all time

The tautology $(p \wedge (p \to q)) \to q$ is the basis of the rule of inference called modus ponens, or the law of detachment.

Rules of Inference for Propositional Logic

- **Modus tollens:** $(\neg q \land (p \to q)) \to \neg p$
- **Hypothetical syllogism:** $((p \to q) \land (q \to r)) \to (p \to r)$
- **Disjunctive syllogism:** $((p \lor q) \land \neg p) \to q$
- **Addition:** $p \to (p \lor q)$
- **Simplification:** $(p \land q) \to p$
- **Conjunction:** $((p) \land (q)) \to (p \land q)$
- **Resolution:** $((p \lor q) \land (\neg p \lor r)) \to (q \lor r)$