

- 1 With the pseudo-code below, determine the final value of sum when $n=23$.

Algorithm 1

```

1: procedure
2:   sum = 0
3:   for (i=2, i<n, i=i+3) do
4:     if i mod 2 == 0
5:       sum = sum + i

```

$$2 \rightarrow 0 + 2 = 2$$

5

$$8 \rightarrow 2 + 8 = 10$$

11

$$14 \rightarrow 10 + 14 = 24$$

17

$$20 \rightarrow 24 + 20 = 44$$

23

44

- 2 Use the pseudo-code below to answer the following question.

- (a) Without working through the pseudo-code, how many times with val be updated in Line 7?
 (b) What is the final value of val.

Algorithm 2

```

1: procedure
2:   dFour = [1, 2, 3, 4]
3:   dSix = [1, 2, 3, 4, 5, 6]
4:   val = 0
5:   for (x in dFour) do
6:     for (y in dSix) do
7:       val = val + x.y

```

$$0 + 1 \cdot 1 + 1 \cdot 2 + 1 \cdot 3 + 1 \cdot 4 + 1 \cdot 5 + 1 \cdot 6$$

$$+ 2 \cdot 1 + 2 \cdot 2 + 2 \cdot 3 + 2 \cdot 4 + 2 \cdot 5 + 2 \cdot 6$$

$$+ 3 \cdot 1 + 3 \cdot 2 + 3 \cdot 3 + 3 \cdot 4 + 3 \cdot 5 + 3 \cdot 6$$

$$+ 4 \cdot 1 + 4 \cdot 2 + 4 \cdot 3 + 4 \cdot 4 + 4 \cdot 5 + 4 \cdot 6 = \underline{210}$$

$$a) 4 \text{ loops} \times 6 \text{ loops} = 24$$

$$b) 210$$

- 3 Using the pseudo-code below, what is the ending value for count when $n = 100$? When $n = 10,000$? When $n = 10^{42}$?

Algorithm 3

```

1: procedure
2:   count=0
3:   while (n>1) do
4:     count++
5:     n = n/10

```

$$n_1 = 100 = 10^2 \rightarrow \text{count}_1 = 2$$

$$n_2 = 10,000 = 10^4 \rightarrow \text{count}_2 = 4$$

$$n_3 = 10^{42} \rightarrow \text{count}_3 = 42$$

- 4 This question assumes familiarity with modular arithmetic discussed in Unit 2. In the pseudo-code below, Line 7 performs string addition (concatenation), for example: "a" + "b" = "ab". Even if the addends look like numbers, it will still do string addition: "2" + "3" = "23". Answer the questions below.

- What is the output of the algorithm when $x = 99$ and $b = 2$?
- What is the output of the algorithm when $x = 1642$ and $b = 8$?

Algorithm 4

```

1: procedure EXPAND(x, b)
2:   Input: integers x and b
3:   Output: outVal, x written in base b
4:
5:   outVal = an empty string
6:   while (x>0) do
7:     outVal = string(x mod b) + outVal
8:     x = x div b
9:   return(outVal)

```

a) $x = 99$ $b = 2$

Blank space

• $(99 \text{ div } 2 = 49, 99 \bmod 2 = 1) \rightarrow (49 \text{ div } 2 = 24, 49 \bmod 2 = 1) \rightarrow (24 \text{ div } 2 = 12, 24 \bmod 2 = 0) \rightarrow (12 \text{ div } 2 = 6, 12 \bmod 2 = 0) \rightarrow (6 \text{ div } 2 = 3, 6 \bmod 2 = 0) \rightarrow (3 \text{ div } 2 = 1, 3 \bmod 2 = 1) \rightarrow (1 \text{ div } 2 = 0, 1 \bmod 2 = 1)$

• $\text{outVal} = (1+1) \rightarrow (0+11) \rightarrow (0+011) \rightarrow (0+0011) \rightarrow (1+00011) \rightarrow (1+00011)$

$= 1100011$

b) $x = 1642$ $b = 8$

• $(1642 \text{ div } 8 = 205, 1642 \bmod 8 = 2) \rightarrow (205 \text{ div } 8 = 25, 205 \bmod 8 = 5) \rightarrow (25 \text{ div } 8 = 3, 25 \bmod 8 = 1) \rightarrow (3 \text{ div } 8 = 0, 3 \bmod 8 = 3)$

• $\text{outVal} = 3152$

- 5 Order each function according to their growth from slowest to fastest (This is the same as execution speed from fastest to slowest for an algorithm).

- (a) $n \cdot \log n$
- (b) $\log_{10} n$
- (c) 10^n
- (d) n^{10}
- (e) $10n$
- (f) 100
- (g) n^2

f) 100 b) $\log_{10} n$ e) $10n$ a) $n \log n$ g) n^2 d) n^{10} c) 10^n

- 6 The provided expressions are the processing time of an algorithm for problems of size n . For each expression, find the *lowest possible* Big- \mathcal{O} complexity stated in simplest form.

- (a) $100n^2 + 0.0002n^3 + 10,000$
- (b) $20n + 2n \log_{10} n + 30$
- (c) $\log_{10} 5 + \log_2 7$
- (d) $5n + \sqrt{n}$
- (e) $0.01(2^n) + n^5 + \ln n$
- (f) $100 \cdot 4^n + 200 \cdot 3^n$

a) $\mathcal{O}(n^3)$ b) $\mathcal{O}(n \log n)$ c) $\mathcal{O}(1)$ d) $\mathcal{O}(n)$ e) $\mathcal{O}(2^n)$ f) $\mathcal{O}(4^n)$

- 7 $f(x) = 6x \log_4 x$ is which of the following? **Mark all that apply.**

☐ $\mathcal{O}(6)$ ☐ $\mathcal{O}(\log x)$ ☐ $\mathcal{O}(6 \log x)$ ☐ $\mathcal{O}(x \log x)$ ☐ $\mathcal{O}(x^2)$ ☐ $\mathcal{O}(2^x)$

Since Big- \mathcal{O} is upper bound, $f(x) \in \mathcal{O}(x \log x) \subset \mathcal{O}(x^2) \subset \mathcal{O}(2^x)$, so

$\mathcal{O}(x \log x)$, $\mathcal{O}(x^2)$, $\mathcal{O}(2^x)$

- 8 Let $h(x) = 6x^4 + 2x^3 + x + 100$. Which of the following is true? **Mark all that apply.**

☐ $h(x) \in \mathcal{O}(x^4)$ ☐ $h(x) \in \Theta(x^4)$ ☐ $h(x) \in \Omega(x^4)$

/ / /

$\mathcal{O}(x^4)$ upper bound $\Omega(x^4)$ lower bound $\Theta(x^4)$ tight bound

9 Let $h(x) = 6 \log_2 x$. Which of the following is true? **Mark all that apply.**

- ☐ $h(x) \in \mathcal{O}(4x^2)$ ☐ $h(x) \in \Theta(4x^2)$ ☐ $h(x) \in \Omega(4x^2)$

✓
6 log x is faster than 4x

10 Let $f(x) = x^2 \log_3 x$. Which of the following is true? **Mark all that apply.**

- ☐ $f(x) \in \mathcal{O}(2x \log_{10} x)$ ☐ $f(x) \in \Theta(2x \log_{10} x)$ ☐ $f(x) \in \Omega(2x \log_{10} x)$

✓ $x^2 > 2x$

11 Let $g(x) = 2x \log_{10} x$. Which of the following is true? **Mark all that apply.**

- ☐ $g(x) \in \mathcal{O}(x^2 \log_3 x)$ ☐ $g(x) \in \Theta(x^2 \log_3 x)$ ☐ $g(x) \in \Omega(x^2 \log_3 x)$

✓ $2x < x^2$

Note: $\Theta \rightarrow$ upper
bigger
slower $\Omega \rightarrow$ lower
smaller
faster

12 Using the provided pseudo-code, find the worst case performance in Big- \mathcal{O} notation.

Algorithm 5 Some more messy pseudocode

```
1: procedure SOMEPROCEDURE2
2:   j=0
3:   for i=0; i<n; i++ do
4:     j=i+j
```

- A. $\mathcal{O}(\log n)$
B. $\mathcal{O}(n \log n)$
C. $\mathcal{O}(n^2)$
D. $\mathcal{O}(n)$

D ✓

- 13 Using the provided pseudo-code, find the worst case performance in Big- \mathcal{O} notation.

Algorithm 6 Some messy pseudo-code

```

1: procedure SOMEPROCEDURE
2:   for i=1 and i<=n do
3:     j=1
4:     while j<n do
5:       j=j+2

```

- A. $\mathcal{O}(\log n)$
 B. $\mathcal{O}(n \log n)$
 C. $\mathcal{O}(n^2)$
 D. $\mathcal{O}(n)$

C

- 14 Using the provided pseudo-code, find the worst case performance in Big- \mathcal{O} notation. Assume we know that **someMethod(n)** is $\mathcal{O}(\log n)$.

Algorithm 7 Some pseudocode with a method in it

```

1: procedure SOMEPROCEDURE3
2:   j=0
3:   for i=0; i<n; i++ do
4:     j=someMethod(n)

```

- A. $\mathcal{O}(\log n)$
 B. $\mathcal{O}(n \log n)$
 C. $\mathcal{O}(n^2)$
 D. $\mathcal{O}(n)$

B

- 15 Using the provided pseudo-code, find the worst case performance in Big- \mathcal{O} notation.

Algorithm 8 Some more messy pseudocode

```

1: procedure SOMEPROCEDURE4
2:   while n>1 do
3:     n=n/2

```

- A. $\mathcal{O}(\log n)$
 B. $\mathcal{O}(n \log n)$
 C. $\mathcal{O}(n^2)$
 D. $\mathcal{O}(n)$

A

- 16 Using the provided pseudo-code, find the worst case performance in Big- \mathcal{O} notation.

Algorithm 9

```

1: procedure
2:   a = 1
3:   c = 0
4:   while a < n do
5:     for i = 0; i < n; i++ do
6:       c = c + 1
7:     a = a * 3

```

- A. $\mathcal{O}(\log n)$
 B. $\mathcal{O}(n \log n)$
 C. $\mathcal{O}(n^2)$
 D. $\mathcal{O}(n)$

B

While loop runs $\log_3 n$ times
 for loop runs n times
 $\rightarrow n \log n$