



Data Structures & Algorithms
Design- SS ZG519
Lecture - 10

Dr. Padma Murali



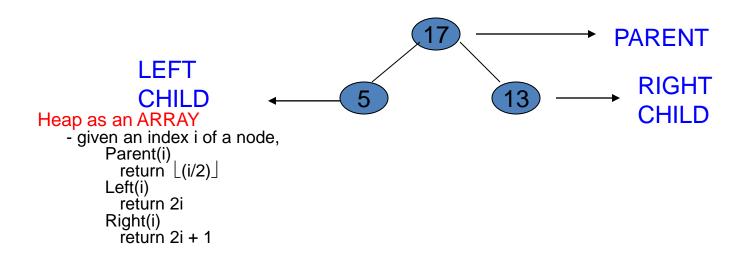
Lecture 10 Topics

Heap Sort

HEAPSORT

Heap

- The heap data structure is an array object which can be viewed as a nearly complete binary tree.
- the two attributes of heap are length[A] and heap-size[A]

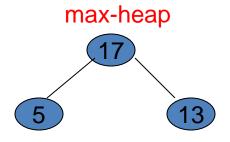


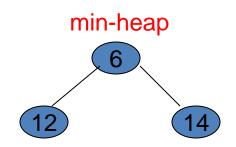
Max-heap

- max-heap property is A[parent(i)] >= A[i]

Min-heap

- min-heap property is A[parent(i)] <= A[i]







HEAPSORT

Basic Procedures used in Heapsort algorithm

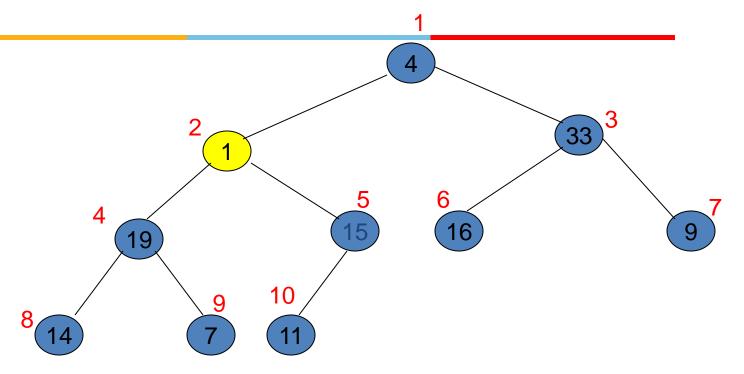
- 1. MAX-HEAPIFY maintains the heap property.
- BUILD-MAX-HEAP produces a max-heap from an unordered input array.
- 3. HEAPSORT sorts an array in place



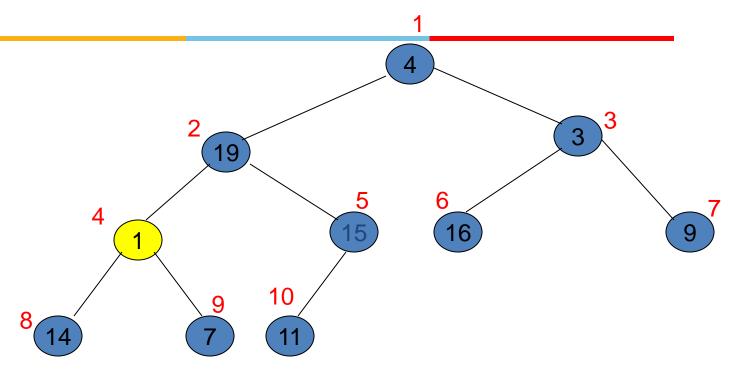
HEAPSORT

```
MAX-HEAPIFY (A,i)
      I ← left(i)
      r \leftarrow right(i)
   if I <= heap-size[A] and A[I] > A[i]
        then largest 

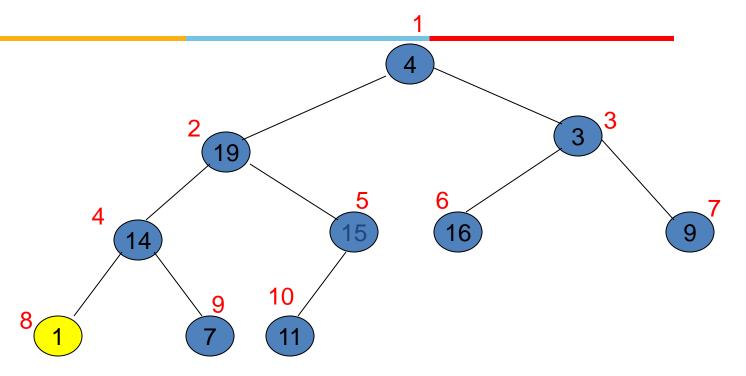
        else largest ← i
   if r <= heap-size[A] and A[r] > A[largest]
        then largest ← r
   <mark>if</mark> largest ≠ i
        then exchange A[i] <-> A[largest]
                MAX-HEAPIFY (A, largest)
```



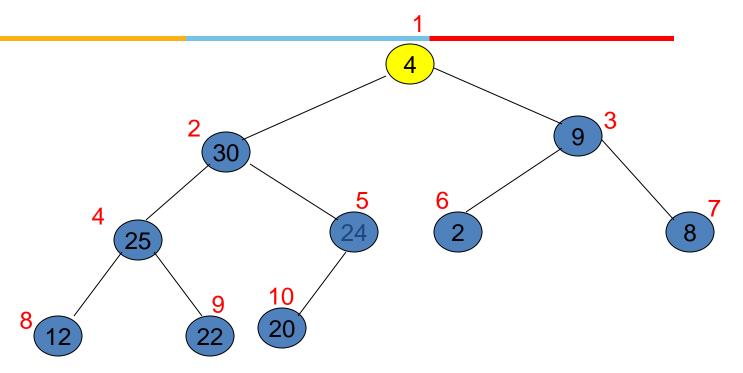
MAX-HEAPIFY(A,2)



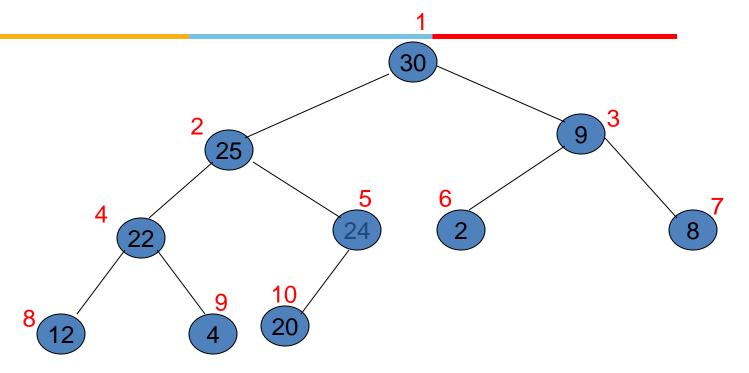
MAX-HEAPIFY(A,4)



MAX-HEAPIFY(A,8)



MAX-HEAPIFY(A,1)



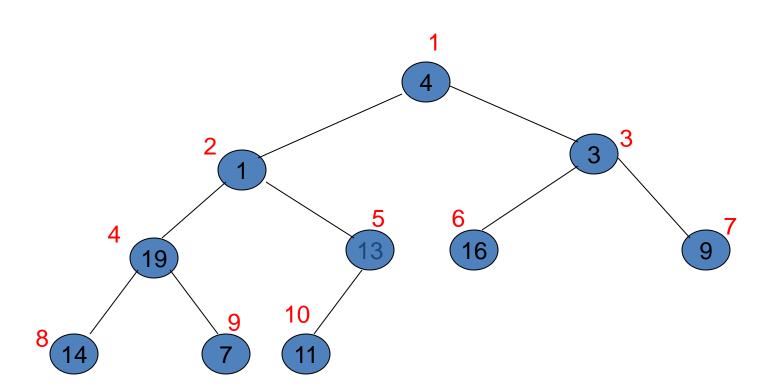
MAX-HEAPIFY(A,1)



BUILD-MAX-HEAP (A) Heap-size[A] ← length[A] for i ← length[A]/2 downto 1 do MAX-HEAPIFY (A,i)

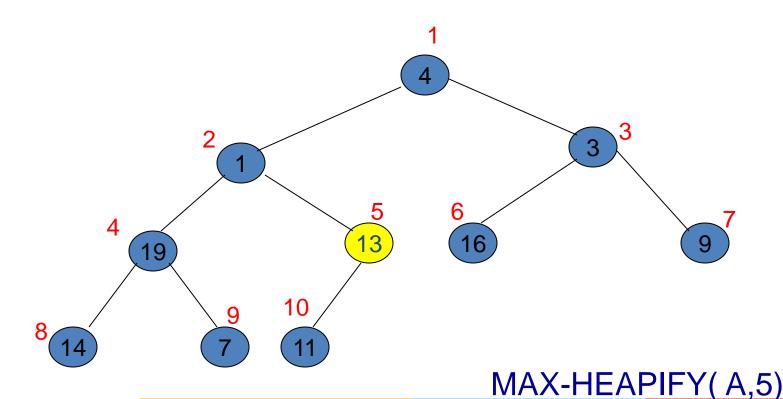


4 1 3 19 13 16 9 14 7 11



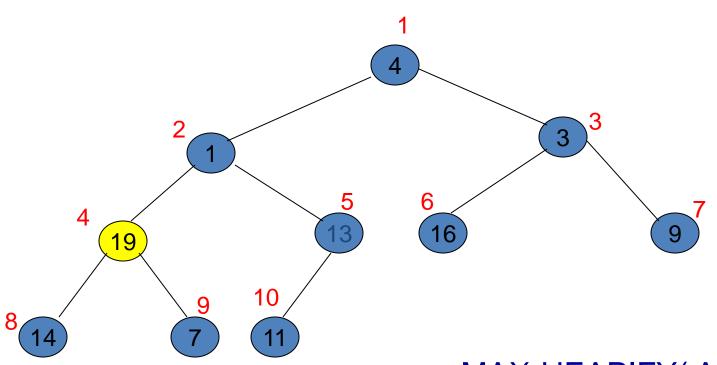


4 1 3 19 13 16 9 14 7 11	4	1	3	19	13	16	9	14	7	11
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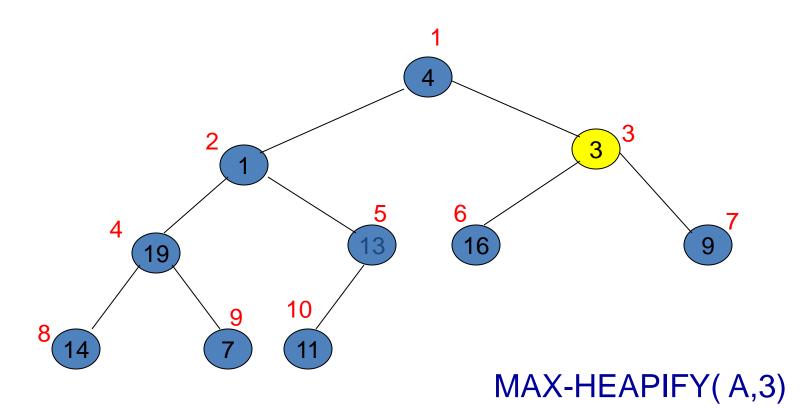


4 1 3 19 13 16 9 14 7 11
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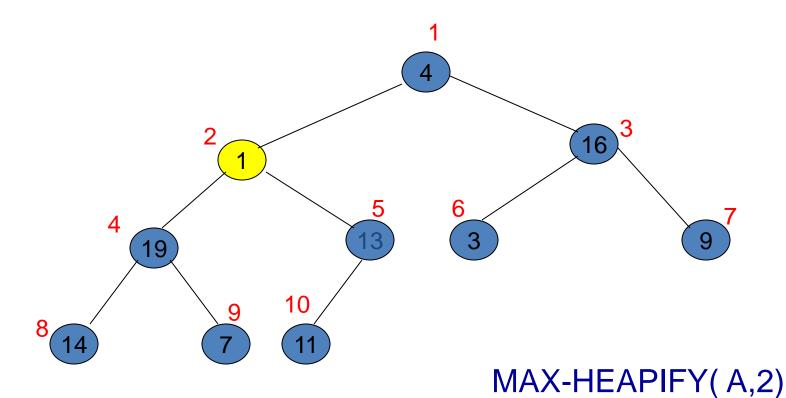


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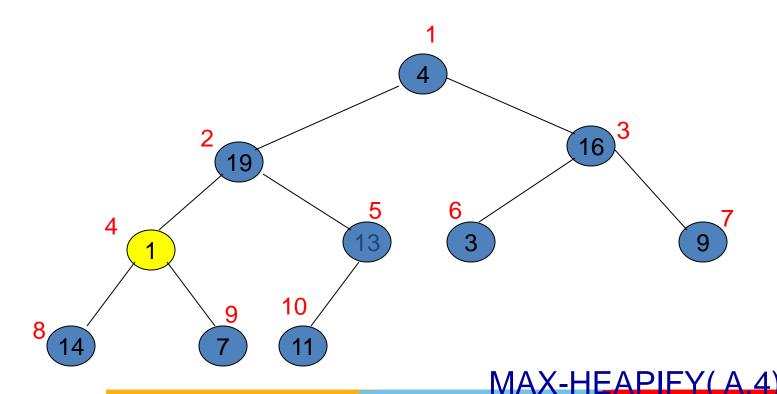


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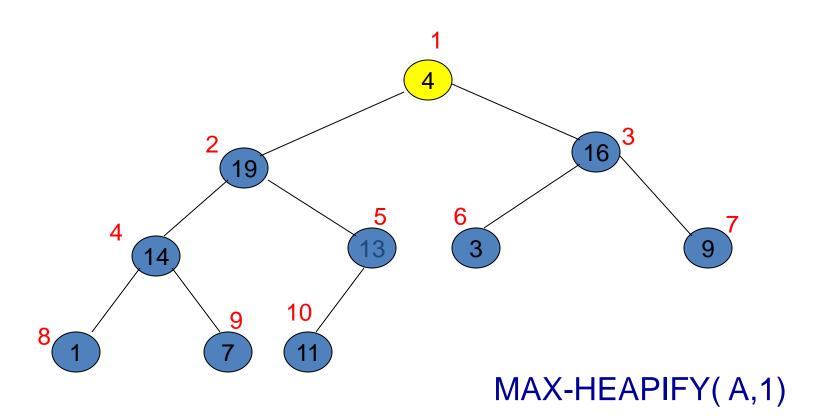


4 1 3 19 13 16 9 14 7 11



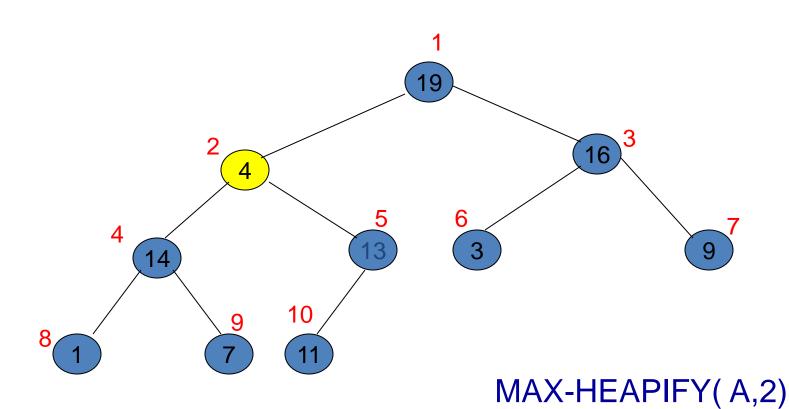


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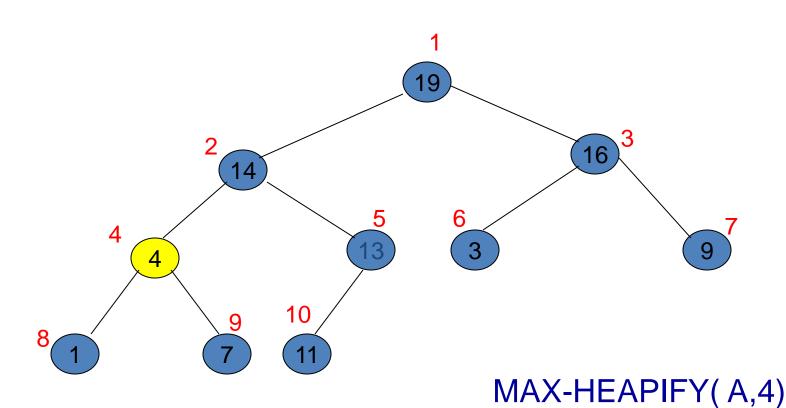


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	4	1	3	19	13	16	9	14	7	11





4 1 3 19 13 16 9 14 7 11

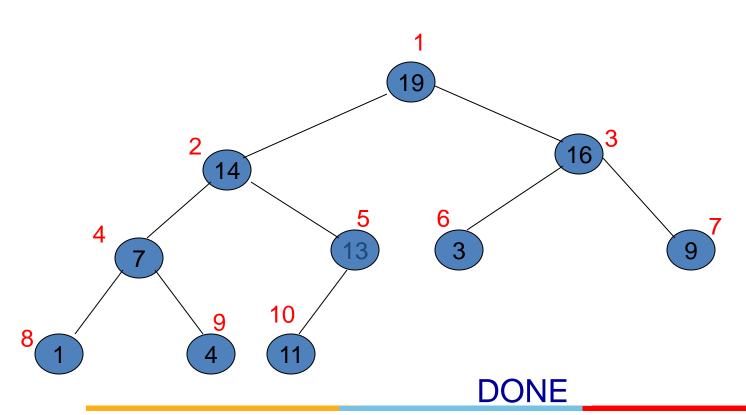












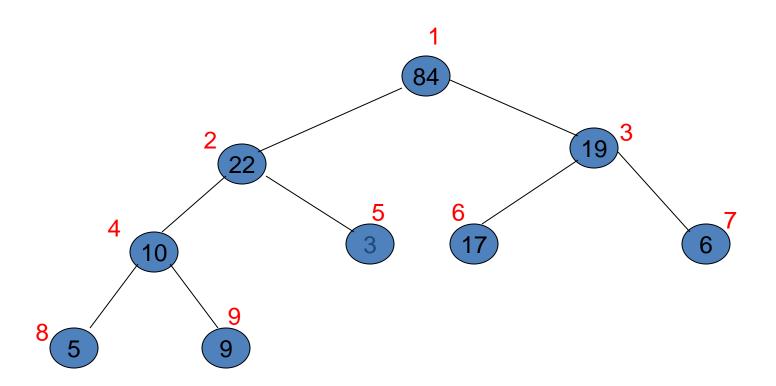


5 3 17 10 84 19 6 22 9

BUILD a heap on the above array



5 3 17 10 84 19 6 22 9





```
HEAPSORT(A)

BUILD-MAX-HEAP(A)

for i ← length[A] downto 2

do exchange A[1] <-> A[i]

heap-size [A] ← heap-size[A] – 1

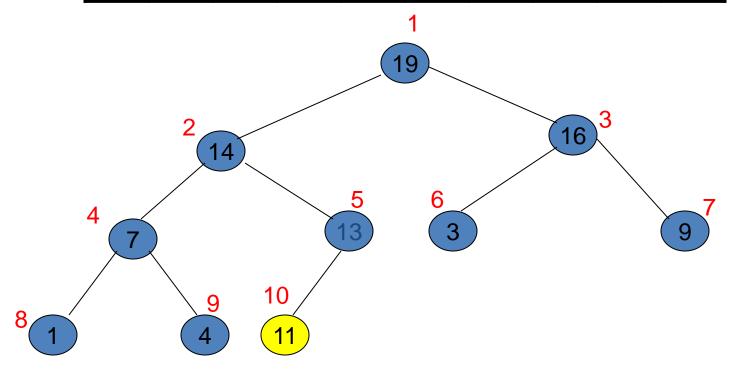
MAX-HEAPIFY(A,1)
```

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HeapSort algorithm





A _____

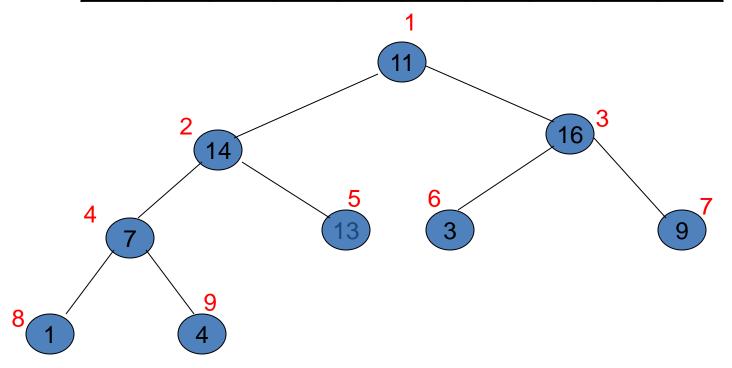
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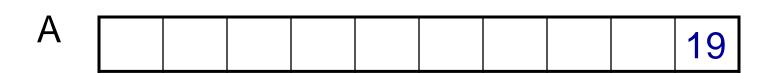
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HeapSort algorithm







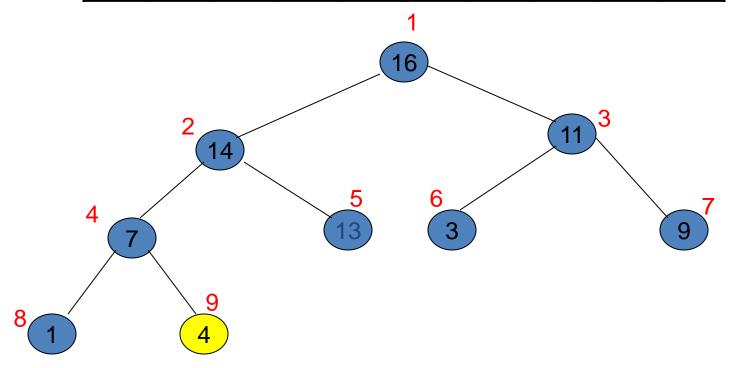
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HeapSort algorithm





A 19

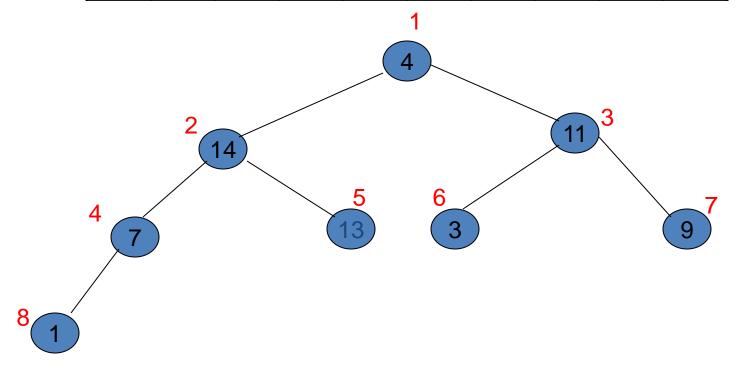
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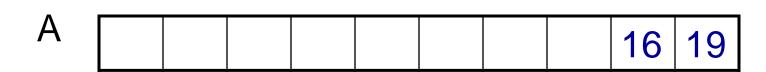
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HeapSort algorithm







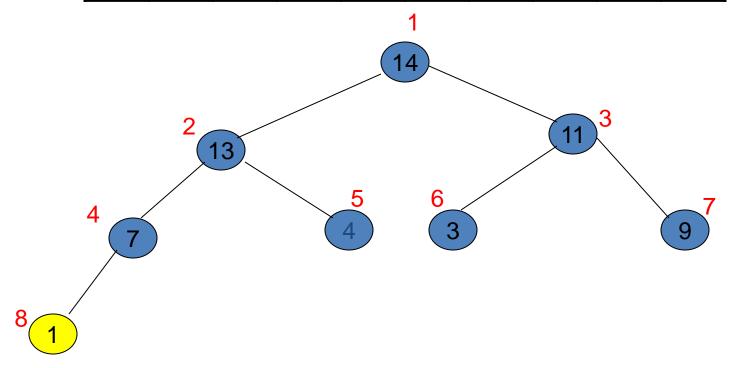
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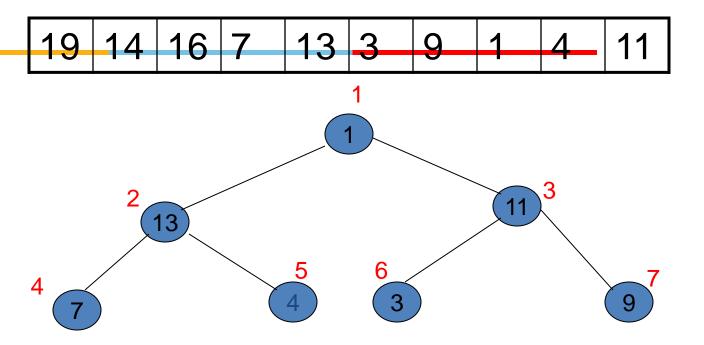
HeapSort algorithm

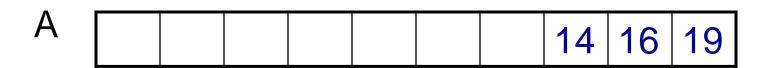




A 16 19

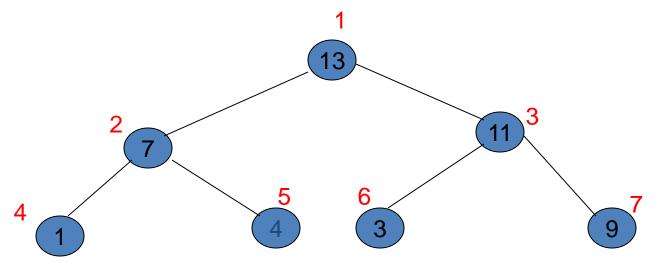
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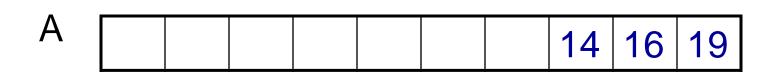




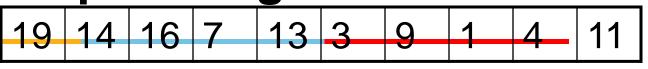


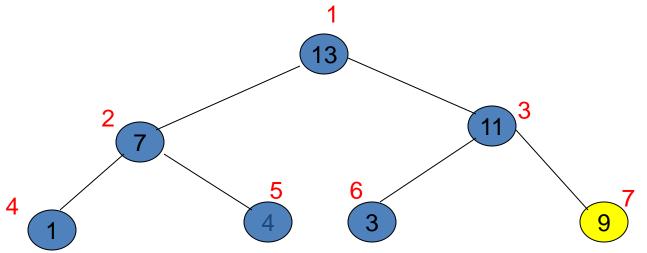


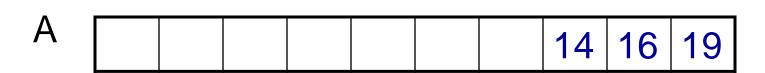






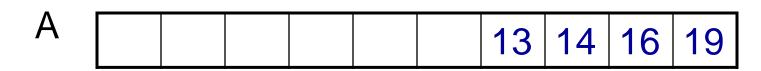








19 14 16 7 13 3 9 1 4 11

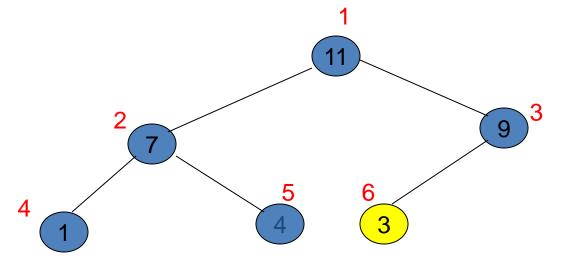


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HeapSort algorithm





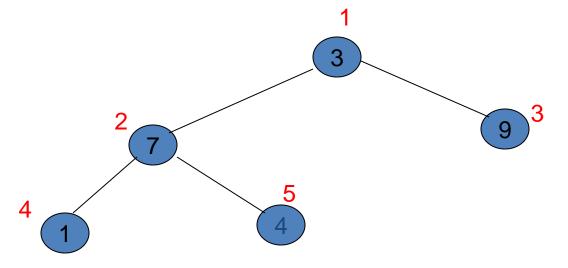
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HeapSort algorithm





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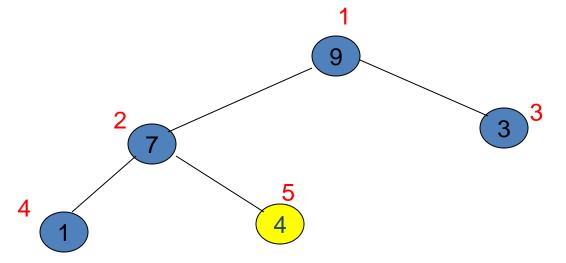
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HeapSort algorithm





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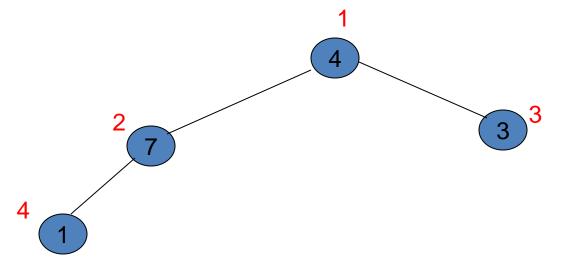
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HeapSort algorithm







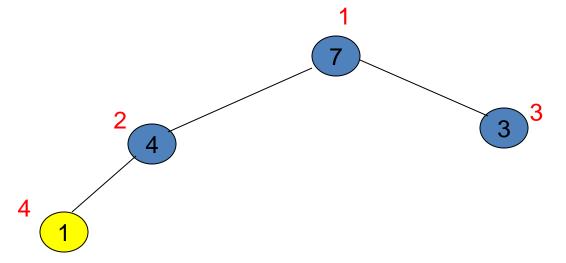
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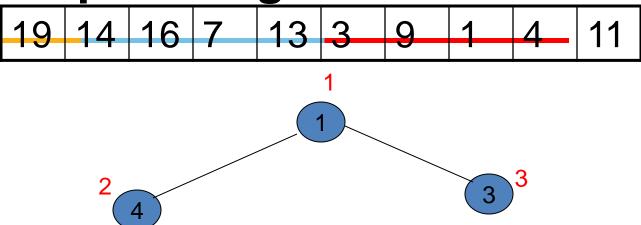
HeapSort algorithm





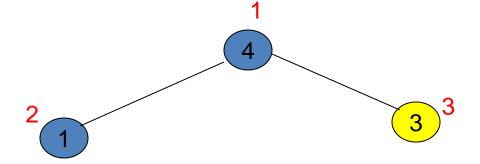
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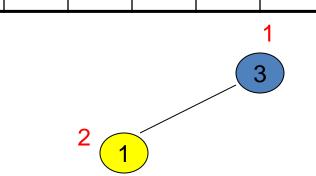


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A



19 14 16 7 13 3 9 1 4 11

A 01 03 04 07 09 11 13 14 16 19

RUNNING TIME OF HEAPSORT

RUNNING TIME OF MAX-HEAPIFY

Running time of MAX-HEAPIFY depends on the height h of a heap i.e O(h)

Height of a heap = lg n

$$T(n) = O(\lg n)$$

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```
HEAPSORT(A)

BUILD-MAX-HEAP(A)

for i ← length[A] downto 2

do exchange A[1] <-> A[i]

heap-size [A] ← heap-size[A] – 1

MAX-HEAPIFY(A,1)
```

- Running time of BUILD-MAX-HEAP is O(n)
- MAX-HEAPIFY is executed n-1 times with running time O(log n)
- Total Running time of Heap Sort is
 O(n) + n-1(O(log n) = O(n log n)

RUNNING TIME OF HEAPSORT INTOVALED

Running time of BUILD-MAX-HEAP

$$\sum_{h=0}^{\infty} h/2^{h} = \frac{1}{2} = 2.$$

$$(1 - \frac{1}{2})^{2}$$

Running time of BUILD-MAX-HEAP can be bounded as