



BITS Pilani

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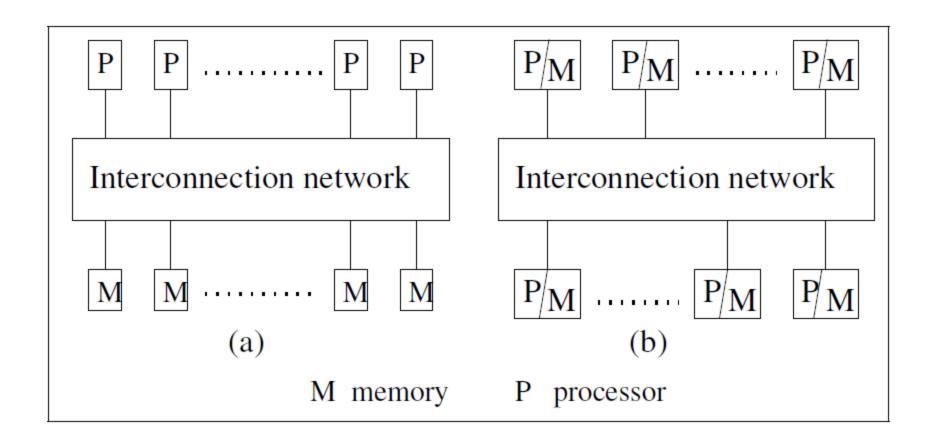
SS ZG 526: Distributed Computing

Lect 2: Communication models, and Design Issues

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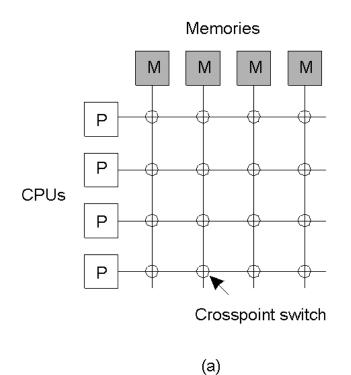


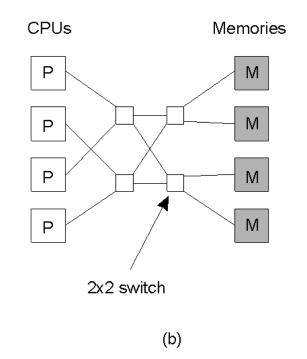
Multiprocessor Vs Multi-computer systems



Multiprocessors







- a) A crossbar switch
- b) An omega switching network



Distributed Computing models

Minicomputer model (e.g. VAXs)

- Each computer supports many users
- Local processing but can fetch remote data (files, databases)

Workstation model (e.g. Athena, and Andrew)

- Most of the work is locally done
- Using a dfs, a user can access remote data

Processor pool model (e.g. Amoeba is a combination of workstation and processor pool model)

- Terminals are Xterms or diskless terminals
- Pool of backend processors handle processing



Distributed Computing Systems: Pros and Cons

Advantages

- Communication and resource sharing possible
- Economy : Price-performance
- Reliability & scalability
- Potential for incremental growth

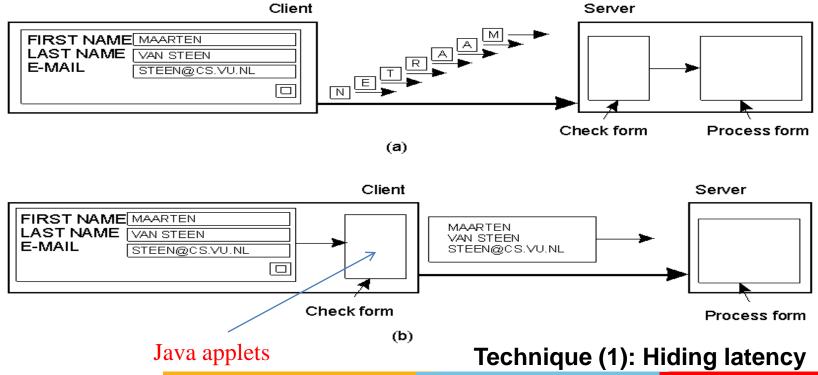
Disadvantages

- Distribution-aware OSs and applications
- High network connectivity is essential
- Security and privacy



Design Issues: Scaling

- If possible, do asynchronous communication
 - Not always possible if the client has nothing to do
- Alternatively, by moving part of the computation





Scaling Technique (2): Distribution aware

Examples: DNS resolutions



More Design Issues

- Lack of Global Knowledge
- Naming
- Compatibility
- Process Synchronization
- Resource Management
- Security



Img. Source: www.bbc.co.uk

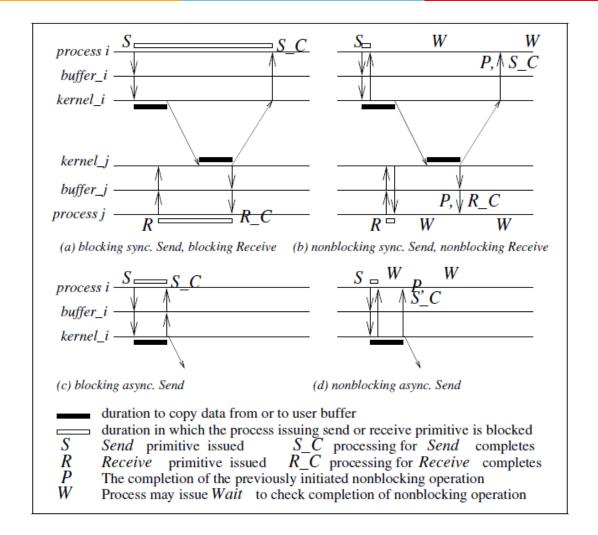


Communication primitives

- Synchronous (send/receive)
 - Handshake between sender and receiver
 - Send completes when Receive completes
 - Receive completes when data copied into buer
- Asynchronous (send)
 - Control returns to process when data copied out of user-specified buffer
- Blocking (send/receive)
 - Control returns to invoking process after processing of primitive (whether sync / async) completes
- Non-blocking (send/receive)
 - Control returns to process immediately after invocation
 - Send: even before data copied out of user buffer
 - Receive: even before data may have arrived from sender

Distributed Communication: Blocking /non-blocking; Synchronous/Asynchronous; send/receive primitives







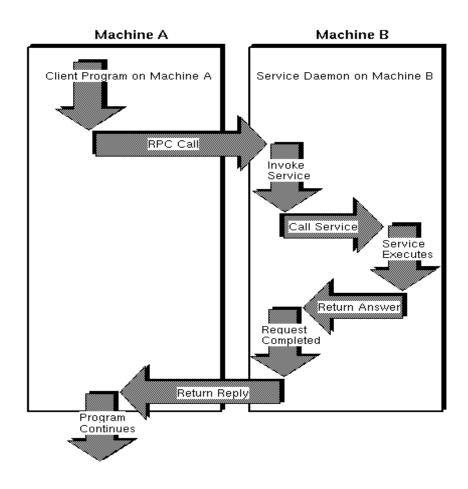
RPC: Remote Procedure Call

- Issues:
 - identifying and accessing the remote procedure
 - parameters
 - return value
- Sun RPC
- Microsoft's DCOM
- OMG's CORBA
- Java RMI
- XML/RPC
- SOAP/.NET
- AJAX (Asynchronous Javascript and XML)

Many types



Remote Procedure Calls (RPC)



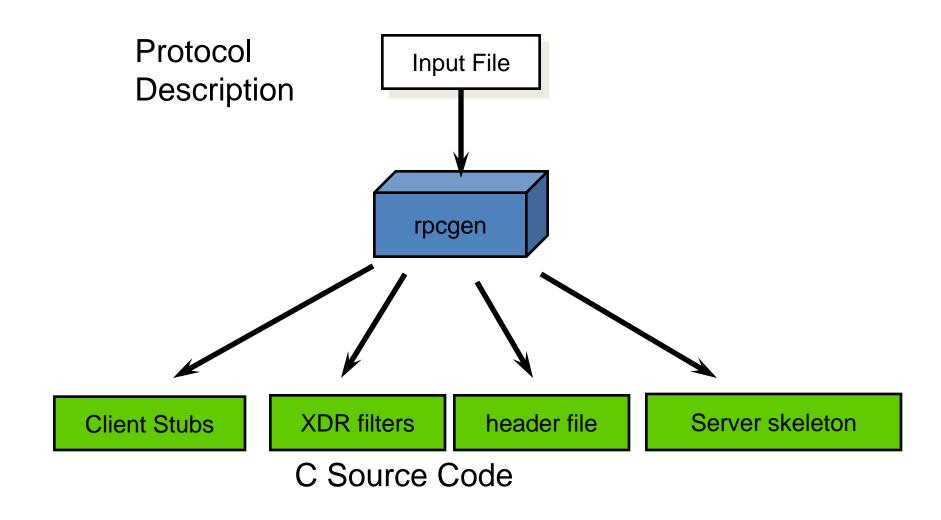


SUN RPC

```
struct square_in {
     long arg1;
struct square_out {
     long res1;
program SQUARE_PROG {
     version SQUARE_VERS {
         square_out SQUAREPROC(square_in) = 1;
          \} = 1;
} = 0x13451111;
```

Rpcgen







Rpcgen continued...

bash\$ rpcgen square.x

produces:

- square.hheader
- square_svc.c server stub
- square_clnt.c client stub
- square_xdr.c XDR conversion routines

Function names derived from IDL function names and version numbers



Square Client: Client.c

```
#include "square.h"
#include <stdio.h>
int main (int argc, char **argv)
{ CLIENT *cl;
 square_in in;
 square_out *out;
 if (argc != 3) { printf("client <localhost> <integer>"); exit (1); }
 cl = clnt_create (argv[1], SQUARE_PROG, SQUARE_VERS, "tcp");
 in.arg1 = atol (argv [2]);
  if ((out = squareproc_1(&in, cl)) == NULL)
      { printf ("Error"); exit(1); }
  printf ("Result %Id\n", out -> res1);
 exit(0);
```



Square server: server.c

```
#include "square.h"
#include <stdio.h>
square_out *squareproc_1_svc (square_in *inp, struct
svc_req *rqstp)
 static square_out *outp;
 outp.res1 = inp -> arg1 * inp -> arg1;
 return (&outp);
```

Exe creation



gcc –o client client.c square_clnt.c square_xdr.c –lnsl

gcc –o server server.c square_svc.c square_xdr.c –
 lrpcsvc -lnsl