Algorithm and Problem Solving Quick Guide in C++

Data Structures, Algorithms and Coding Interview Problem Patterns in C++

rfdavid, 2025

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INTRODUCTION

Motivation

The tech industry hiring standard is based on algorithm and data structure.

There are plenty of free resources available around algorithms and data structures. The purpose of this project is to be a quick guide where you can learn and review algorithms and data structures.

Some of the intended **key features:**

- Non-verbose, short-structured, and easy to follow descriptions
- Slide-based, practical for reviewing
- Free and open-source

right in the please add a star at github.com/rfdavid/cpp-algo-cheatsheet

Some Useful Links

Tech Interview Handbook

https://www.techinterviewhandbook.org

A very well-structured resource for interview preparation

Blind 75 Leetcode Questions

https://leetcode.com/discuss/general-discussion/460599/blind-75-leetcode-questions

Blind 75

- Blind 75 is a popular list of algorithm problems that intends to cover the main data structures and patterns.
- It is a curated list of 75 popular coding questions created by an ex-Meta Staff Engineer

Array

- ✓ Two Sum
- ✓ Contains Duplicate
- √ Product of Array Except Self
- ✓ Best Time to Buy and Sell Stock
- ✓ Maximum Subarray
- ✓ Maximum Product Subarray
- ✓ Find Minimum in Rotated Sorted Array

Search in Rotated Sorted Array

3 Sum

✓ Container With Most Water

Binary

- ✓ Sum of Two Integers
- ✓ Number of 1 Bits
- ✓ Counting Bits
- ✓ Missing Number
- ✓ Reverse Bits

Dynamic Programming

- ✓ Climbing Stairs
- Coin Change
- Longest Increasing Subsequence
- ✓ Longest Common Subsequence
- Word Break
- **Combination Sum**
- House Robber
- House Robber II
- Decode Ways
- ✓ Unique Paths

Jump Game

Matrix

- ✓ Set Matrix Zeroes
- ✓ Spiral Matrix
- ✓ Rotate Image

Word Search

Blind 75

Tree

- ✓ Maximum Depth of Binary Tree
- ✓ Same Tree
- ✓ Invert/Flip Binary Tree
- ✓ <u>Binary Tree Maximum Path Sum</u>
- ✓ Binary Tree Level Order Traversal
- ✓ <u>Serialize and Deserialize Binary Tree</u>
- ✓ Subtree of Another Tree

Construct Binary Tree from Preorder and Inorder Traversal

Validate Binary Search Tree

✓ Kth Smallest Element in a Binary Search Tree

Lowest Common Ancestor of Binary Search Tree

<u>Implement Trie (Prefix Tree)</u>

Add and Search Word

Word Search II

Heap

✓ Top K Frequent Elements

Find Median from Data Stream

String

- ✓ Longest Substring Without Repeating Characters
- ✓ Longest Repeating Character Replacement
- ✓ <u>Minimum Window Substring</u>
- √ Valid Anagram
- √ Group Anagrams
- ✓ Valid Parentheses
- ✓ Valid Palindrome

Longest Palindromic Substring

Palindromic Substrings

Encode and Decode Strings &

Linked List

- ✓ Reverse a Linked List
- ✓ <u>Detect Cycle in a Linked List</u>
- ✓ Merge Two Sorted Lists
- ✓ Merge K Sorted Lists
- ✓ Remove Nth Node From End Of List
- ✓ Reorder List

Graph

- ✓ Clone Graph
- ✓ Course Schedule
- ✓ Pacific Atlantic Water Flow
- ✓ Number of Islands
- ✓ Longest Consecutive Sequence

- ✓ <u>Number of Connected Components</u>
 <u>In an Undirected Graph</u>

Interval

- √ Insert Interval
- ✓ <u>Merge Intervals</u>
- ✓ Non-overlapping Intervals
- Meeting Rooms II ☆



Other problems

Tree

- ✓ <u>Maximum Level Sum of a Binary Tree</u>
- ✓ <u>Minimum Number of Increments on Subarrays to Form a Target Array</u>
- ✓ <u>Leaf-Similar Trees</u>
- ✓ Count Good Nodes in Binary Tree

Min Cost Climbing Stairs

Longest Palindromic Subsequence

If the input gets bigger, how many steps does the algorithm take?

- Measure of how much the execution time of an algorithm grows relative to the size of its input (usually called n)
- Expressed in **Big-O notation** (e.g. O(1), O(n), $O(n^2)$, etc) to describe the **upper bound** of how fast the algorithm's runtime grows
- Asymptotic notations
 - **Big-O (O)** Upper Bound (**Worst-case**): Describes the maximum amount of time/memory an algorithm could take.
 - Theta O Tight Bound (Exact): describes both the upper and lower bound (the exact growth rate)
 - **Omega** (Ω) Lower Bound (**Best-case**): describes the minimum time/space the algorithm needs.

Examples

O(1)

Examples	Problems
Accessing an array element (`arr[i]`)	Hash table lookups
Swapping two variables	Checking if a number is even/odd
Stack/Queue `push` or `pop`	Returning first element of a list

O(log n)

Binary Search	Search in sorted array
Balanced BST insert/find (AVL, Red-Black Tree)	Find k-th smallest in BST
Finding floor/ceil in sorted array	Finding square root with binary search

O(n)

Linear Search	Maximum subarray sum (Kadane's algo)
Finding min/max in an array	Counting frequencies with hash map
Traversing linked list or array	One-pass string processing problems

O(n log n)

Merge Sort / Heap Sort	Sorting an array
Efficient algorithms for Closest Pair	Finding inversion count in array
Heapify operations	Kth largest element with heap

O(n²)

Bubble Sort / Insertion Sort	Two Sum (brute force)
Checking all pairs in array	Longest Palindromic Substring (DP)
Floyd-Warshall algorithm	Edit Distance (DP)

O(n³)

Matrix multiplication (naive)	Boolean matrix multiplication
DP on subsequences of length 3	Some DP path-finding problems
Floyd-Warshall for dense graphs	Counting triangles in graph

O(2ⁿ)

Recursive Fibonacci (no memo)	Subset sum (brute force)
Backtracking for combinations/permutations	N-Queens
Traveling Salesman (brute force)	All subsets of array (power set)

O(n!)

Generating all permutations	Traveling Salesman (brute force)
Brute-force anagram check	Word ladder with all transformations
Solving puzzles with all arrangements	Hamiltonian Path

O(V + E)

DFS / BFS (adjacency list)	DFS / BFS (adjacency list)
Dr 37 Dr 3 (dajacerrey list)	Dr or Dr o (adjacerrey net)

O(E log V)

Dijkstra with priority queue Dijkstra with priority queue	Dijkstra with priority queue	Dijkstra with priority queue
---	------------------------------	------------------------------

O(VE)

Bellman-Ford

$O(N_3)$

Floyd-Warshall

O(E log E)

Kruskal's MST algorithm

V = vertices (nodes) E = edges

Space Complexity

As the input size n grows, how much extra memory does the algorithm need to run?

- Measure of how much memory an algorithm uses relative to input size
- Expressed in Big-O notation (e.g. O(1), O(n), O(n²), etc)
- It includes auxiliary space (extra memory used by the algorithm, not counting the input itself) and sometimes considers input space depending on the context
- Count only extra space needed (exclude output)
- The space complexity of a recursive tree traversal is **O(h)**, where h is the height of the tree. This is because each recursive call adds a frame to the call stack, and in the worst case, the maximum stack depth is proportional to the tree's height

Space Complexity

Examples

Algorithm / Operation	Space Complexity	Explanation
Swap two integers	O(1)	Only uses constant space
Iterate through array and sum values	O(1)	No extra memory used besides accumulator
Store array copy	O(n)	Needs space to store the copied array
Recursive factorial (factorial(n))	O(n)	n stack frames in the call stack
Binary search (recursive)	O(log n)	Recursive depth is log(n) for sorted array
Binary search (iterative)	O(1)	No extra space beyond a few variables
Merge sort	O(n)	Needs temp arrays to merge subarrays
Quick sort (in-place)	O(log n)	Call stack for recursive calls
Depth-first search (recursive) in tree	O(h)	h = height of the tree (stack frames)
Breadth-first search (using queue)	O(n)	Stores all nodes at current level in queue
DP with full 2D table (e.g., LCS)	O(m*n)	Stores results of all subproblems
Optimized Fibonacci with two variables	O(1)	Only tracks last two results
Memoized Fibonacci (top-down DP)	O(n)	Memoization table + recursion stack
Using a hash map to count frequencies	O(n)	Stores one count per element
Storing all substrings of a string	O(n^2)	Total number of substrings is ~n²
Adjacency list for graph with V nodes, E edges	O(V + E)	One list per node, total edges stored

DATA STRUCTURES IN C++

Data Structure Decision Diagram



 The following diagram gives you the direction to which data structure to use in C++ according to the problem you are trying to solve

Note: I don't have the source of this diagram. If you know it, please drop me a msg so I can add it here.

Arrays

std::vector is a sequence container that encapsulates dynamic sized arrays*

Linked List

Stack

Queue

Heap

Hash Table

Tree

ARRAY

Arrays

- Memory layout: hold values in a contiguous block of memory.
- **Fixed Size**: the size of an array is defined when it is created and cannot be changed. However, high-level languages have different implementations, making it dynamic.
- Homogeneous elements: all elements are of the same data type (int, float, char...)
- **Efficiency**: accessing elements by index is very efficient *O(1)*, since each index maps directly to a memory location. Also, range scans benefit from CPU cache lines since arrays are stored in contiguous blocks of memory.



leetcode.com/problems/two-sum

Problem

- Given an array of numbers and a target, example: array [2,7,11,15] and target 9
- Return indices of two numbers where they add up to target
- **Output**: [0,1]

```
array[0] + array[1] = 2 + 7 = 9
```



leetcode.com/problems/two-sum

Solution

- Iterative over each number in the array
- Calculate the difference between target and each number, example:

```
array[0] = 2, target 9, then 9 - 2 = 7
```

- Now we know we need the number 7 to sum up to 9
- Check in a hashmap if we have 7 in some part of the array

```
hash[7] exists?
```

- If yes, return the current index and the index of 7
- If not, store the index of the current number in the hashmap for future evaluation

$$hash[2] = 0$$

Code - 1. Two Sum



LeetCode leetcode.com/problems/two-sum

```
Code Time: O(n) Space: O(n)
vector<int> twoSum(vector<int>& nums, int target) {
    std::unordered map<int, int> numMap;
   // n being the size of nums
   for (int i = 0; i < nums.size(); i++) {</pre>
        // current number of the array
        int number = nums[i];
        int diff = target - number;
        // check if the difference is in some part of the array
        // by using a hashmap
        if (numMap.find(diff) != numMap.end()) {
            return { numMap[diff], i};
        // register the current number index
        numMap[number] = i;
   // no matches
   return {};
```

Problem – 217. Contains Duplicate



leetcode.com/problems/contains-duplicate

Problem

- You are given an array of numbers
- Return any value that appears at least twice

Solution

- Loop through the array
- Check if the value is in a hash table
- Return **true** if the value exist
- The problem requires at least twice, but one modification may be having a specific count

Code - 217. Contains Duplicate



LeetCode leetcode.com/problems/contains-duplicate

```
bool containsDuplicate(vector<int>& nums) {
   unordered_map<int, int> seen;
   for (int i = 0; i < nums.size(); ++i) {
      if (seen[nums[i]] == 1) {
        return true;
      }
      seen[nums[i]]++;
   }
   return false;</pre>
```

Another solution (less flexible)

```
bool containsDuplicate(vector<int>& nums) {
    unordered_map<int, bool> seen;
    for (const auto& num : nums) {
        if (seen[num]) {
            return true;
        }
        seen[num] = true;
    }
    return false;
}
```





leetcode.com/problems/product-of-array-except-self

Problem Statement

- You are given an integer array nums
- Return another array where each element is multiplied by all the elements except itself
- Example:

```
nums = [14,2,5,99]
nums[0] = 2 * 5 * 99 (all except 14)
nums[1] = 14 * 5 * 99 (all except 2)
nums[2] = 14 * 2 * 99
nums[3] = 14 * 2 * 5
```



leetcode.com/problems/product-of-array-except-self

Solution

• Go over the array once and calculate the product of the left side. Example:

```
nums = [14,2,5]
left[0] = 1 (think of multiplying all elements before 14, so 1 because there is none)
left[1] = 14 (all elements from the left multiplied, except 2)
left[2] = 14 * 2 = 28 (all elements from the left multiplied, except 5)
left = [1, 14, 28]
```

Using the same logic, do the same calculation but starting from the right

```
right[2] = 1 (no elements after 5)
right[1] = 5 (only 5 after 2)
right[0] = 2 * 5 = 10
right = [10, 5, 1]
```

• Multiply each element from left and right:

```
left \odot right = [1, 14, 28] \odot [10,5,1] = [10, 70, 28]
```

Code - 238. Product of Array Except Self

C LeetCode

leetcode.com/problems/product-of-array-except-self

```
Code Time: O(n) Space: O(n)
vector<int> productExceptSelf(vector<int>& nums) {
    int n = nums.size();
    vector<int> output(n, 1);
    vector<int> right(n, 1);
    // calculate left first
    for (int i = 1; i < n; ++i) {
        output[i] = nums[i - 1] * output[i - 1];
    // calculate right
    for (int i = n - 1; i >= 0; --i) {
        right[i] = nums[i + 1] * right[i + 1];
        output[i] = output[i] * right[i];
    /* or you can save some space using this logic,
      although I don't find it as intuitive as the previous one
    int right = 1;
    for (int i = n - 1; i >= 0; --i) {
        output[i] *= right;
        right *= nums[i];
    return output;
```



leetcode.com/problems/maximum-subarray

Problem

- You are given an array nums
- Find the subarray with the largest sum

• Example:

nums =
$$[-2,1,-3,4,-1,2,1,-5,4]$$

output = 6

The subarray [4,-1,2,1] has the largest sum 6.





leetcode.com/problems/maximum-subarray

Solution

- Use Kadane's algorithm to find the maximum sum of a contiguous subarray in linear time
- Core idea:

at each index, either:

- 1. start a new subarray at **nums[i]** or
- 2. extend the current one by adding nums[i]

Arrays – Kadane's algorithm

- Kadane's algorithm is a dynamic programming algorithm to solve maximum subarray sum
- At every index i: start a new subarray at i extend the previous subarray to include array[i]

Algorithm

1. Initialize:

```
int maxSoFar = array[0];
int maxEndingHere = array[0];
```

2. Loop through the array

```
for (int i = 1; i < array.size(); ++i) {
    maxEndingHere = max(array[i], maxEndingHere + array[i]);
    maxSoFar = max(maxSoFar, maxEndingHere);
}</pre>
```

3. Return maxSoFar;

Problem - 53. Maximum Subarray

```
E LeetCode
```

leetcode.com/problems/maximum-subarray

```
int maxSubArray(vector<int>& nums) {
  int maxSum = nums[0];
  int currentSum = nums[0];
  for (int i = 1; i < nums.size(); ++i) {
     currentSum = max(nums[i], currentSum + nums[i]);
     maxSum = max(maxSum, currentSum);
  }
  return maxSum;
}</pre>
```



leetcode.com/problems/maximum-product-subarray

Problem

- You are given an array nums
- Find a subarray that has the largest product and return the product
- The array may contain negative numbers

• Example:

```
nums = [2, 3, -2, 4]

output = 6

[2,3] has the largest 6 (2 * 3)
```

Problem - 152. Maximum Product Subarray





leetcode.com/problems/maximum-product-subarray

Solution

- Use a modified version of Kadane's algorithm
- Keep track of the minimum and maximum product
- Once the current number is negative, swap minimum product with maximum product
- Check the largest product between maximum product and the final result

Problem - 152. Maximum Product Subarray

C LeetCode

return result;

leetcode.com/problems/maximum-product-subarray

```
int maxProduct(vector<int>& nums) {
  int result = nums[0];
  int maxProd = nums[0];
  int minProd = nums[0];
  for (int i = 1; i < nums.size(); ++i) {
    if (nums[i] < 0) {
       swap(minProd, maxProd);
    }

  minProd = min(nums[i], nums[i] * minProd); // -2
    maxProd = max(nums[i], nums[i] * maxProd); // -30

  result = max(result, maxProd);
}</pre>
```



leetcode.com/problems/find-minimum-in-rotated-sorted-array

Problem

- You are given a sorted array but "rotated"
- Rotated means the elements are displaced in order
- Return the minimum element

• Example:

```
nums = [3,4,5,1,2]
output = 1(minimum element)
```



leetcode.com/problems/find-minimum-in-rotated-sorted-array

Solution

- Perform an adapted binary search
- Example:

```
[3,4,5,1,2]

left = 3, mid = 5, right = 2

You find mid (5), but have to go right, so adjust left:
```

```
if (mid > right)
   left = mid + 1
else
   right = mid
```

Problem – 153. Find Minimum in Rotated Sorted Array

LeetCode

leetcode.com/problems/find-minimum-in-rotated-sorted-array

```
int findMin(vector<int>& nums) {
  int left = 0;
  int right = nums.size() - 1;
  while (left < right) {
    int mid = left + (right - left) / 2;
    if (nums[mid] > nums[right]) {
        left = mid + 1;
        } else {
            right = mid;
        }
    }
    return nums[left];
}
```

Problem - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Problem Statement

- You are given an integer array of stock prices
- Choose a price[i] to buy and price[i] to sell where you achieve maximum profits
- Example:

```
prices = [9, 1, 3, 4]
```

• Output: [1,3]

```
array[3] - array[1] = 4 + 1 = 3
```

Solution - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Solution

- Initialize profit = 0
- Initialize lowestBuyPrice = prices[0]
- Loop through the prices
- Track the lowest buy price → min(lowestBuyPrice, prices[i])
- Check if selling "today" will make the maximum profit and update profit:
 max(prices[i] buy > profit, profit)
- Update profit max(prices[i] - buy

Code - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Code (simplified) Time: O(n) Space: O(1)

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      buy = min(buy, prices[i]);
      profit = max(profit, prices[i] - buy)
   }
   return profit;
}</pre>
```

Code - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Code (optimized) Time: O(n) Space: O(1)

Same logic, but with better branch prediction and less computation

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      if (prices[i] < buy) {
        buy = prices[i];
      } else if (prices[i] - buy > profit) {
           profit = prices[i] - buy;
      }
   }
   return profit;
}
```



leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

Problem

- You are given an integer array of stock prices
- Choose a price[i] to buy and price[i] to sell where you achieve maximum profits
- You can buy/sell multiple times, but only hold at most one transaction at a time
- Output is the maximum profits

• Example:

```
prices = [9, 1, 3, 4]

Output: 2 + 1 = 3

buy (price = 1), sell (price = 3), profit = 2

buy (price = 3), sell (price = 4), profit = 1
```



leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

Solution

- Loop through the array starting from index 1
- If current price[i] is lower than previous price[i 1], buy and sell

• Example:

```
prices = [1, 8, 4] prices[0] = 1, prices[1] = 8, prices[2] = 4
prices[0] < \text{prices}[1] \rightarrow true, profit = 8 - 1 = 7
prices[2] < \text{prices}[1] \rightarrow \text{false}, do nothing
```

Code - Best Time to Buy and Sell Stock II

LeetCode

leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

```
int maxProfit(vector<int>& prices) {
  int profit = 0;
  for (int i = 1; i < prices.size(); ++i) {
    if (prices[i] > prices[i-1]) {
       profit += prices[i] - prices[i - 1];
    }
  }
  return profit;
}
```



leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Problem Statement

• ..

Solution - Best Time to Buy and Sell Stock IV





leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Solution

• ..

Code - Best Time to Buy and Sell Stock IV



leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Code (simplified) Time: O(n) Space: O(n)

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      buy = min(buy, prices[i]);
      profit = max(profit, prices[i] - buy)
   }
   return profit;
}</pre>
```

STRING

Problem – 3. Longest Substring Without Repeating Characters





leetcode.com/problems/longest-substring-without-repeating-characters

Problem Statement

 You are given a string and the goal is to find the longest substring without repeating characters

Example

Input: "abcdb"

Output: 4 (abcd since "b" is repeated)

Solution – 3. Longest Substring Without Repeating Characters





leetcode.com/problems/longest-substring-without-repeating-characters

Solution

- Use sliding window algorithm (left and right)
- Loop through the string
- Try to find if the current character is already added by using unordered set or bitmap
- If added, remove from the set alongside with others using left pointer
- If not, add to the unordered set or bitmap
- Maximum length will be right left + 1



LeetCode leetcode.com/problems/longest-substring-without-repeating-characters

Example

- String: abcbd. Our goal is to return 3 (abcbd)
- Initialize maxLength = 0
- Loop through the string

```
lteration 1: left = 0, right = 0, string[left] = 'a',
          bitmap = ['a'] ('a' is not in bitmap, add), maxLength = max(maxLength, right - left + 1) = 1
lteration 2: left = 0, right = 1, string[right] = 'b'
          bitmap = ['a','b'], maxLength = 2
lteration 3: left = 0, right = 2, string[right] = 'c'
          bitmap = ['a','b', 'c'], maxLength = 3
lteration 4: left = 0, right = 3, string[right] = 'b'
          bitmap = ['a','b','c','b']
          'b' is already in the bitmap. start "clearing" the character using left:
          Iteration 4a: left = 0, string[left] = 'a' is different from 'b', so remove 'a'
                      bitmap = ['b', 'c','b']
          Iteration 4b: left = 1, string[left] = 'b' is the same as the repeated one, remove
                       bitmap = [`c', 'b']
Iteration 5: left = 1, right = 4, string[right] = 'd'
            bitmap = ['c','b','d']
```

Code – 3. Longest Substring Without Repeating Characters

Code (unordered_set)

Use unordered_set when question requires unicode chars

```
int lengthOfLongestSubstring(string s) {
    int maxLength = 0;
    int left = 0, right = 0;
    // track the seen characters
    unordered set<char> seen;
    for (right = 0; right < s.size(); ++right) {</pre>
        char currentChar = s[right];
        // if currentChar is in the set, clean
        // the character and everything from left of it
        // basically, reset the longest substring
        while (seen.count(currentChar)) {
            char c = s[left];
            seen.erase(c);
            left++;
        // insert the current read character
        seen.insert(currentChar);
        // set max length
        maxLength = max(maxLength, right - left + 1);
    return maxLength;
```

Code – 3. Longest Substring Without Repeating Characters

Code (bitmap)

- Using bitset: create a bitmask with 128 bits where each bit represent a character
- Optimal solution for ASCII since ASCII size is 127 characters
- Unicode / UTF-8 can represent over 1.1 million characters, so use unordered_set approach instead

```
int lengthOfLongestSubstring(string s) {
    std::bitset<128> bitmask;
    uint32_t left = 0;
    uint32_t maxLength = 0;

for (uint32_t right = 0; right < s.length(); ++right) {
        uint32_t bitIndex = s[right];
        // if char is already in the bitmask, move left until we reset the bits
        while (bitmask.test(bitIndex)) {
            bitmask.reset(s[left]);
            ++left;
        }

        bitmask.set(bitIndex);
        maxLength = std::max(maxLength, right - left + 1);
    }
    return maxLength;
}</pre>
```

Problem – 424. Longest Repeating Character Replacement



LeetCode

https://leetcode.com/problems/longest-repeating-character-replacement

Problem Statement / Solution / Code Time: O(-) Space: O(-)

•

Problem - 76. Maximum Window Substring



C LeetCode

leetcode.com/problems/minimum-window-substring

Problem Statement / Solution / Code Time: O(-) Space: O(-)

• ...

Problem – 242. Valid Anagram



LeetCode

leetcode.com/problems/valid-anagram

Problem Statement / Solution / Code Time: O(-) Space: O(-)

• ..

Problem – 49. Group Anagrams



E LeetCode

leetcode.com/problems/group-anagrams

Problem Statement / Solution / Code Time: O(-) Space: O(-)

•



LeetCode leetcode.com/problems/valid-parentheses

Problem Statement

- You are given a string containing only the characters '(', ')', '{', '}', '[' and ']'
- A valid input have closed brackets by its own type
- Example

()[] $\{\}$ \rightarrow valid

[]{}(\rightarrow invalid

 $\{()\} \rightarrow \text{valid}$

Solution - Valid Parentheses



leetcode.com/problems/valid-parentheses

Solution

- Loop through the string
- If open brackets ([{ push to a stack
- If closed brackets:
 - pop the last added bracket
 - **check** if the **closed** bracket corresponds to the **popped** bracket
 - if not, return false
- after the loop, return true if the size of the stack is empty (all brackets closed)

Code – Valid Parentheses



LeetCode leetcode.com/problems/valid-parentheses

```
Code
          Time: O(n) Space: O(n)
bool isValid(string s) {
   // stack (LIFO)
   std::stack<char> brackets;
   // O(n)
   for (int i = 0; i < s.size(); ++i) {</pre>
        char bracket = s[i];
        if (bracket == '(' || bracket == '[' || bracket == '{'}) {
           brackets.push(bracket);
       } else {
           if (brackets.size() == 0) return false;
           char lastBracket = brackets.top();
           if (bracket == ')' && lastBracket != '(') return false;
           if (bracket == '}' && lastBracket != '{') return false;
           if (bracket == ']' && lastBracket != '[') return false;
           brackets.pop();
   // all brackets must be closed
   return brackets.size() == 0;
```

Problem - 125. Valid Palindrome



LeetCode

leetcode.com/problems/valid-palindrome

Problem Statement / Solution / Code Time: O(-) Space: O(-)

•

Problem – Minimum Number of Increments on Subarrays



leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Problem Statement

- You are given an array of integers initialized with zeros (e.g. [0,0,0,0])
- The goal is to reach some target (e.g. [1, 2, 2, 3])
- The valid operations is to increment a subarray by one
- The output is the total number of operations In this case:
 - $[1,1,1,1] \rightarrow$ increment the subarray starting from 0 to total size
 - [1,2,2,2] → increment the subarray starting from 1 to total size
 - $[1,2,2,3] \rightarrow$ increment the subarray starting and ending from the last element

Output: 3 (total number of operations)

Solution – Minimum Number of Increments on Subarrays





leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Solution

Explain...

Code (2) - Minimum Number of Increments on Subarrays



leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Code (optimized)

Code – Minimum Number of Increments on Subarrays

LeetCode

leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Code

```
int minNumberOperations(vector<int>& target) {
    int totalOp = target[0];
    for (int i = 1; i < target.size(); ++i) {
        // can't reuse
        if (target[i - 1] < target[i]) {
            totalOp += target[i] - target[i - 1];
        }
    }
    return totalOp;
}</pre>
```

TWO POINTERS



leetcode.com/problems/container-with-most-water

Problem Statement

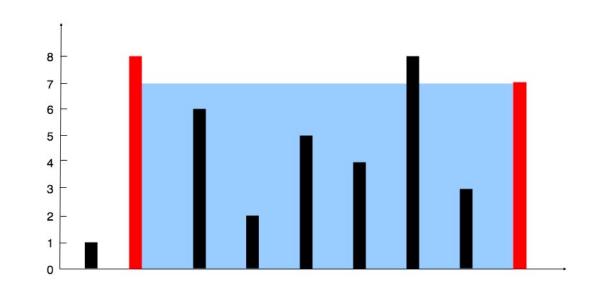
- You are given an integer array height
- Find two lines that together with x-axis form a container with most water
- Example:

Input:

height = [1,8,6,2,5,4,8,3,7]

Output:

49





leetcode.com/problems/container-with-most-water

Solution

- Initialize the maximum area maxArea = 0
- Initialize two pointers, left = 0 and right = height.size 1
- Loop while pointer left < right
- Calculate the area:

area = min(height[left], height[right]) * (right - left)

Update the global maximum area:

maxArea = max(maxArea, area)

- Move the smallest pointer (increment left or decrement right)
- Return maxArea

LeetCode

leetcode.com/problems/container-with-most-water

Code

```
int maxArea(vector<int>& height) {
   // initialize the two pointers (left and right)
   int left = 0;
   int right = height.size() - 1;
   int maxArea = 0;
   while (left < right) {</pre>
       // calculate the area, think about the x-axis and y-axis
        int area = min(height[left], height[right]) * (right - left);
       // update maximum area
       maxArea = max(area, maxArea);
       // is the left pointer (y) smaller than right?
        if (height[left] < height[right]) {</pre>
           // move left pointer to right
           left++;
        } else {
            // otherwise, move right pointer to left
            right--;
   return maxArea;
```

BINARY

Bit Manipulation in C

Operators

```
& AND OR ^ XOR ~ NOT << LEFT SHIFT >> RIGHT SHIFT
```

Common Operations

```
set bit: num |= (1 << pos)

clear bit: num &= ~(1 << pos)

toggle bit: num ^= (1 << pos)

check bit: (num & (1 << pos)) != 0

extract bit: (num >> pos) & 1

extract a range of bits: (num >> pos) & ((1 << length) - 1)</pre>
```

Example

```
void copyBit(int *dst, int src, int srcPos, int dstPos) {
    int bit = (src >> srcPos) & 1; // extract bit
    *dst &= ~(1 << dstPos); // clear destination bit
    *dst |= (bit << dstPos); // set destination bit
}</pre>
```

Binary

- In C++, **std::bitset** represents a fixed-size sequence of N bits
- Example:

```
std::bitset<8> bitmask;
bitmask.reset(1)
bitmask.set(1)
if (bitmask.test(1)) { // true
...
```

- reset : set bit to false
- **set** : set a specific bit
- **test** : check a specific bit
- **count** : return the number of bits set to true
- **flip**: toggle the value of the bits (if true, set to false and vice-versa)

Problem – 371. Sum of Two Integers





leetcode.com/problems/sum-of-two-integers

Problem

Problem – 371. Sum of Two Integers





leetcode.com/problems/sum-of-two-integers

Solution

Problem – 371. Sum of Two Integers



E LeetCode

leetcode.com/problems/sum-of-two-integers

Code Time: O(-) Space: O(-)

• ...



LeetCode

leetcode.com/problems/number-of-1-bits

Problem



LeetCode

leetcode.com/problems/number-of-1-bits

Solution

...



LeetCode

leetcode.com/problems/number-of-1-bits

Code



LeetCode

leetcode.com/problems/number-of-1-bits

Problem Statement / Solution / Code Time: O(-) Space: O(-)

•



LeetCode

leetcode.com/problems/counting-bits

Problem Statement / Solution / Code Time: O(-) Space: O(-)

•



leetcode.com/problems/counting-bits

Problem



leetcode.com/problems/counting-bits

Solution



```
LeetCode
```

leetcode.com/problems/counting-bits



• ...

Problem – 268. Missing Number





https://leetcode.com/problems/missing-number

Problem

Problem – 268. Missing Number



LeetCode

https://leetcode.com/problems/missing-number

Solution

Problem – 268. Missing Number



LeetCode

https://leetcode.com/problems/missing-number

Code Time: O(-) Space: O(-)

• ...

Problem - 190. Reverse Bits





LeetCode leetcode.com/problems/reverse-bits

Problem

Problem - 190. Reverse Bits





leetcode.com/problems/reverse-bits

Solution

Problem - 190. Reverse Bits



LeetCode

leetcode.com/problems/reverse-bits

Code Time: O(-) Space: O(-)

• ...

Negabinary

- Non-standard positional numeral system that uses base of -2
- Allow representing negative numbers in binary
- Example:

$$1101_{-2}$$

$$(-2)^3 + (-2)^2 + 0 + (-2)^0 = -8 + 4 + 0 + 1 = -3$$

Summing Negabinary

Add as a regular binary number, but with negative carry

$$0 + 0 = 0$$

 $1 + 0 = 1$
 $1 + 1 = 0$ with a negative carry 1
 $1 + 1 = 0$ (subtract)
 $1 + 0 = 1$ with a positive carry 1

Negabinary

Example 1

Example 2

$$\begin{array}{r}
 1111 \\
 101010 \\
 + 101100 \\
\hline
 = 11110110
 \end{array}$$

Reference

https://leetcode.com/problems/adding-two-negabinary-numbers

Given two numbers arr1 and arr2 in base -2, return the result of adding them together.

Each number is given in array format: as an array of 0s and 1s, from most significant bit to least significant bit. For example, arr = [1,1,0,1] represents the number $(-2)^3 + (-2)^2 + (-2)^0 = -3$. A number arr in array, format is also guaranteed to have no leading zeros: either arr == [0] or arr[0] == 1.

Return the result of adding arr1 and arr2 in the same format: as an array of 0s and 1s with no leading zeros.

Example 1

```
Input: arr1 = [1,1,1,1,1], arr2 = [1,0,1]
```

Output: [1,0,0,0,0]

Explanation: arr1 represents 11, arr2 represents 5, the output represents 16.

Example 2

```
Input: arr1 = [0], arr2 = [0]
```

Output: [0]

Example 3

```
Input: arr1 = [0], arr2 = [1]
```

Output: [1]

Solution 1073 – Adding Two Negabinary Numbers



https://leetcode.com/problems/adding-two-negabinary-numbers

GRAPH (DFS)

Problem - Keys and Rooms

https://leetcode.com/problems/keys-and-rooms

```
int maxProfit(vector<int>& prices) {
    int profit = 0;
    int buy = prices[0];
    for (auto i = 1; i < prices.size(); i++) {
        if (prices[i] < buy) {
            buy = prices[i];
        } else if (prices[i] - buy > profit) {
                profit = prices[i] - buy;
        }
    }
    return profit;
}
```



https://leetcode.com/problems/clone-graph

Problem Statement

- Given a node reference, create a deep copy of the graph
- The class node has two variables: val and neighbours

```
class Node {
  public int val;
  public List<Node> neighbors;
}
```

Output is the node reference of the copy



https://leetcode.com/problems/clone-graph

Solution

- First check the edge cases (is the node null?)
- Create a hash map to store the nodes that is already created unordered<int, Node*> graph;
- Check if the current node already exists in the graph
- If not, create a new Node object and store in the hashmap
- Visit all the neighbors and add the neighbors to this current node

Code - Clone Graph

```
E LeetCode
```

https://leetcode.com/problems/clone-graph

```
std::unordered_map<int, Node*> graph;
Node* cloneGraph(Node* node) {
   if (node == NULL) {
        return NULL;
    // does this node object exists?
   if (graph.find(node->val) == graph.end()) {
        // node wasn't visited yet, store in the hashmap
        graph[node->val] = new Node(node->val);
        // visit all neighnours
        for (const auto& n : node->neighbors) {
            graph[node->val]->neighbors.push_back(cloneGraph(n));
   return graph[node->val];
```

Problem - 207. Course Schedule





leetcode.com/problems/course-schedule

Problem

Solution – 207. Course Schedule



LeetCode

leetcode.com/problems/course-schedule

Solution

Code - 207. Course Schedule

LeetCode

leetcode.com/problems/course-schedule

Code Time: O(-) Space: O(-)

•

Problem – 417. Pacific Atlantic Water Flow





leetcode.com/problems/pacific-atlantic-water-flow

Problem

Solution - 417. Pacific Atlantic Water Flow



LeetCode

leetcode.com/problems/pacific-atlantic-water-flow

Solution

Code - 417. Pacific Atlantic Water Flow



E LeetCode

leetcode.com/problems/pacific-atlantic-water-flow

Code Time: O(-) Space: O(-)

• ...

Problem - 200. Number of Islands





leetcode.com/problems/number-of-islands

Problem

• ..

Solution – 200. Number of Islands



LeetCode

leetcode.com/problems/number-of-islands

Solution

•

Code - 200. Number of Islands

LeetCode

leetcode.com/problems/number-of-islands

Code Time: O(-) Space: O(-)

• ...

Problem – 128. Longest Consecutive Sequence





leetcode.com/problems/longest-consecutive-sequence

Problem

• ..

Solution - 128. Longest Consecutive Sequence





leetcode.com/problems/longest-consecutive-sequence

Solution

•

Code – 128. **Longest Consecutive Sequence**



E LeetCode

leetcode.com/problems/longest-consecutive-sequence

Code Time: O(-) Space: O(-)

•

Problem – 261. Graph Valid Tree





leetcode.com/problems/graph-valid-tree

Problem

...

Solution – 261. Graph Valid Tree



LeetCode

leetcode.com/problems/graph-valid-tree

Solution

• ..

Code – 261. Graph Valid Tree

LeetCode

leetcode.com/problems/graph-valid-tree

Code Time: O(-) Space: O(-)

• ...

Problem – 323. Number of Connected Components





leetcode.com/problems/number-of-connected-components-in-an-undirected-graph

Problem

•

Problem – 323. Number of Connected Components





leetcode.com/problems/number-of-connected-components-in-an-undirected-graph

Solution

• ..

Problem – 323. Number of Connected Components



E LeetCode

leetcode.com/problems/number-of-connected-components-in-an-undirected-graph

Code Time: O(-) Space: O(-)

• ...

GRAPH (BFS)





https://leetcode.com/problems/maximum-level-sum-of-a-binary-tree

Problem Statement

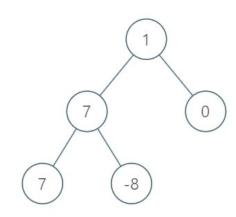
- Given the root of a binary tree, find the smallest level with the maximum sum
- For example, the tree below has the follow sums for each level:

level
$$1 \text{ (root)} = 1$$

level
$$2 = 7 + 0 = 7$$

$$|eve| 3 = 7 - 8 = -1$$

Therefore, level 2 has the maximum sum



Solution – Maximum Level Sum of a Binary Tree





https://leetcode.com/problems/maximum-level-sum-of-a-binary-tree

Solution

- Have a queue with the nodes for the current level
- Sum the values from that level by taking the nodes from the queue
- Example, we know that level 1 has one node. Hence, pop the first node from the queue
 If level 2 has 2 nodes, pop two nodes, sum the values
- In addition, add left and right to the end of the queue to process the next level

Code – Maximum Level Sum of a Binary Tree

```
E LeetCode
```

https://leetcode.com/problems/maximum-level-sum-of-a-binary-tree

```
int maxLevelSum(TreeNode* root) {
    std::queue<TreeNode*> nodes;
    int currentLevel = 0;
    int maxLevel = 1;
    int maxSum = INT MIN;
    nodes.push(root);
    // traverse the graph
    while(!nodes.empty()) {
        int levelSum = 0;
        int levelSize = nodes.size();
        currentLevel++;
        // sum the values in current level
        for (int i = 0; i < levelSize; ++i) {</pre>
            TreeNode* node = nodes.front();
            levelSum += node->val;
            nodes.pop();
            if (node->left) nodes.push(node->left);
            if (node->right) nodes.push(node->right);
        if (levelSum > maxSum) {
            maxLevel = currentLevel;
            maxSum = levelSum;
    return maxLevel;
```

SHORTEST PATH

Shortest Path Algorithms

Algorithms

- BFS
- Dijkstra
- Bellman-Ford
- Floyd-Warshall
- A* search
- Johnson's
- SPFA (Shortest Path Faster)
- Bidirectional Search

TREE

Depth-First Traversals

• **Pre-order**: Root – Left – Right



• In-order: Left - Root - Right



• **Post-order**: Left – Right – Root



Breadth-First Traversal (Level Order Traversal)

Visit every node on a level before moving to a lower level.

Depth-First Traversals

Use a recursive algorithm to traverse according to the order

```
if (!root) return;
• Pre-order: Root – Left – Right
                                                       doSomething();
                                                       visit(node->left);
                                                       visit(node->right);
                                                       if (!root) return;
• In-order: Left – Root – Right
                                                       visit(node->left);
                                                       doSomething();
                                                       visit(node->right);
                                                       if (!root) return;
• Post-order: Left – Right – Root
                                                       visit(node->left);
                                                       visit(node->right);
```

doSomething();

Example of pre-order and in-order

```
struct TreeNode {
    int val;
    TreeNode *left, *right;
    TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
};
// Pre-order traversal
void preorderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    cout << root->val << " ";</pre>
    preorderTraversal(root->left);
    preorderTraversal(root->right);
// In-order traversal
void inorderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    inorderTraversal(root->left);
    cout << root->val << " ";</pre>
    inorderTraversal(root->right);
```

Example of post-order and level-order

```
struct TreeNode {
    int val;
    TreeNode *left, *right;
    TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
};
// Post-order traversal
void postorderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    postorderTraversal(root->left);
    postorderTraversal(root->right);
    cout << root->val << " ";</pre>
// Level-order traversal using a queue
void levelOrderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    queue<TreeNode*> q;
    q.push(root);
    while (!q.empty()) {
        TreeNode* current = q.front();
        q.pop();
        cout << current->val << " ";</pre>
        if (current->left != nullptr) q.push(current->left);
        if (current->right != nullptr) q.push(current->right);
```

BFS Using Stack

BFS with std::stack

 This might be useful for problems when you want to return and resume (for example, <u>872. Leaf-Similar Trees</u>)

```
struct TreeNode {
    int val;
    TreeNode *left, *right;
    TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
};

// Pre-order traversal
void bfs(std::stack<TreeNode*>& tree) {
    while(!tree.empty()) {
        TreeNode* root = tree.top();
        tree.pop();
        // do something ...
        if (root->right) tree.push(root->right);
        if (root->left) tree.push(root->left);
    }
}
```



leetcode.com/problems/same-tree

Problem

- You are given the root of two trees
- Write a function to check if they are the same
- Example:

$$p = [1,2,3], q = [1,2,3]$$

Output: true



leetcode.com/problems/same-tree

Solution

- Traverse both trees (**p** and **q**) recursively and check if the nodes are the same
- Start by the base case:are **p** and **q** null? return true
- One of them are null? return false, because they should be the same
- Finally, check if **p->val** is equal to **q->val** and also for both and left, recursively

Problem - 100. Same Tree



leetcode.com/problems/same-tree

Code

Time: O(n) where n is the number of nodes Space: O(h) where h is the height of the tree. Best case is usually O(log n) for balanced trees, but skewed trees is usually O(n)

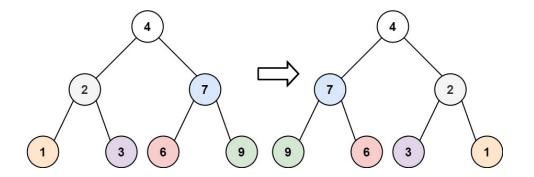
```
bool isSameTree(TreeNode* p, TreeNode* q) {
    // base case: leaf is null. If both are null, then return true
    if (!p && !q) return true;
    // if both are not NULL, then they must have value.
    // If one of them doesn't have value, then they're different, return false
    if (!p || !q) return false;
    // they must have the same value
    // as any other nodes in the tree
    return p->val == q->val &&
        isSameTree(p->left, q->left) &&
        isSameTree(p->right, q->right);
}
```



leetcode.com/problems/invert-binary-tree

Problem

- You are given the root of a binary tree
- Invert the tree and return the root
- Example:







leetcode.com/problems/invert-binary-tree

Solution

- Recursively traverse the tree
- Create a new pointer temp that points to left node
- Set left node to right
- Set right node to temp
- Call the function recursively for left and right
- Return root

Problem - 226. Invert Binary Tree

```
E LeetCode
```

leetcode.com/problems/invert-binary-tree

```
Time: O(n) Space: O(h) where h is the height of the tree
Code
TreeNode* invertTree(TreeNode* root) {
    // base case
    if (!root) return nullptr;
    // create a new pointer to left
    TreeNode* temp = root->left;
    // invert
    root->left = root->right;
    root->right = temp;
    // recursively invert left and right
    invertTree(root->left);
    invertTree(root->right);
    return root;
```





LeetCode https://leetcode.com/problems/maximum-depth-of-binary-tree

Problem Statement

- Given the root of a binary tree, find the <u>maximum depth</u>
- Example:

Output: 4



Solution – Maximum Depth of Binary Tree



LeetCode https://leetcode.com/problems/maximum-depth-of-binary-tree

Solution

- Perform post-order traversal: left right root
- Recursively go left and right to find each value
- Return the max of each one

Code – Maximum Depth of Binary Tree

LeetCode https://leetcode.com/problems/maximum-depth-of-binary-tree

```
int maxDepth(TreeNode* root) {
   if (!root) return 0;
   // find max left
   int maxLeft = maxDepth(root->left);
   // find max right
   int maxRight = maxDepth(root->right);
   // return max +1 (account for root)
   return std::max(maxLeft, maxRight) + 1;
```



https://leetcode.com/problems/path-sum

Problem Statement

- It is given the root of a binary tree and an integer target sum
- Example:



Output: true

Node
$$1 + Node 7 + Node 2 = 10$$





https://leetcode.com/problems/path-sum

Solution

- Start from root node (1)
- Subtract from target number (example 10 1 = 9)
- Continue going down the tree, until the target is 0, return true
- After visiting all nodes, if the target is not zero, return false



Code – Path Sum

```
E LeetCode
```

https://leetcode.com/problems/path-sum

```
bool hasPathSum(TreeNode* root, int targetSum) {
    if (!root) {
        return false;
    // we want targetSum to be zero
   targetSum -= root->val;
   // if there is no left, no right, we've reached the end of the path
    // so if the targetSum is zero, then the nodes summed up to the targetSum
    if (!root->left && !root->right && targetSum == 0) {
        return true;
    // propagate to left and right
    return hasPathSum(root->left, targetSum) || hasPathSum(root->right, targetSum);
```

Also, a small performance tweak can be made by avoiding writing targetSum: targetSum -= root->val

This will avoid a memory write access, making the calculation directly in the CPU, but also at a cost of readability

```
if (!root->left && !root->right && targetSum - root->val == 0) {
    ...
return hasPathSum(root->left, targetSum - root->val) || hasPathSum(root->right, targetSum - root->val);
```

Problem - 297. Serialize and Deserialize Binary Tree



leetcode.com/problems/serialize-and-deserialize-binary-tree

Problem

- Design an algorithm to serialize and deserialize a binary tree
- You have to build two interfaces: serialize that returns a string, and deserialize that returns the whole tree as TreeNode pointer
- The string can be represented at any format (comma-separated, space separated etc)

Solution – 297. Serialize and Deserialize Binary Tree



leetcode.com/problems/serialize-and-deserialize-binary-tree

Solution

• Serialize: traverse the tree pre-order, and append its value to a string

Null value should also be represented

Example: [1,2,null,null,3 ...]

Call "traverse" to do it recursively

Deserialize: split the string into tokens

read each token and re-build the tree by adding a new node

Call "buildTree" to do it recursively

Code - 297. Serialize and Deserialize Binary Tree

```
E LeetCode
```

leetcode.com/problems/serialize-and-deserialize-binary-tree

```
Code Time: O() Space: O()
string serialize(TreeNode* root) {
    // traverse the tree in pre-order: root, left, right
    // generate a string with comma separator,
    // example: 1,2,N,N,3 ...
    string result;
    traverse(root, result);
    return result;
TreeNode* deserialize(string data) {
   // split the input data
   vector<string> tokens = split(data);
   // index to be used to access the elements from tokens recursively.
   // Hence, we need to create it here to pass by reference.
   // Note that index is bounded by the number of tokens, so it won't overflow
   int index = 0:
   TreeNode* root = buildTree(tokens, index);
   return root;
```

continue...

Code - 297. Serialize and Deserialize Binary Tree

```
E LeetCode
```

leetcode.com/problems/serialize-and-deserialize-binary-tree

```
TreeNode* buildTree(vector<string>& tokens, int& index) {
    // read the current token based on the index
    const string& token = tokens[index];
    // increment index before checking for null
    ++index;
    // base case: null node
    if (token == "N") {
        return nullptr;
    }
    // build root
    TreeNode* node = new TreeNode(stoi(token));
    // build left
    node->left = buildTree(tokens, index);
    // build right
    node->right = buildTree(tokens, index);
    return node;
}
```

```
// traverse in pre-order (root, left, right)
// and append the values to the string 's'
// append 'N' if it is NULL
void traverse(TreeNode* root, string& s) {
    if (!s.empty()) s += ",";
    // base case, we need to append null
    if (!root) {
        s += "N";
        return;
    // visit root
    s += to string(root->val);
    // visit left
    traverse(root->left, s);
    // visit right
    traverse(root->right, s);
// helper function in C++ to split string
vector<string> split(const string& s) {
    vector<string> result;
    stringstream ss(s);
    string token;
    while(getline(ss, token, ',')) {
        result.push back(token);
    return result;
```



leetcode.com/problems/serialize-and-deserialize-binary-tree

Some interesting alternative to split

C++ 23 have an interesting way to split using std::views::split

```
vector<string> split(string s) {
    auto result = s |
        views::split(',') |
        views::transform([](auto&& subRange) {
            return string(subRange.start(), subRange.end());
        });
}
```

To understand, this follow a structure similar to unix pipes:

```
echo "123,N,556" | split | transform
```

std::views::split returns ranges, something like:

```
[ range("123"), range("N"), range("556") ]
```

std::views::transform converts each subrange into an actual string



LeetCode leetcode.com/problems/kth-smallest-element-in-a-bst

Problem Statement / Solution

- You are given the root of a binary search tree and an integer k
- Find the kth smallest value

Example

From all values in the tree: 1,2,3,4,5,6

 $\mathbf{k} = \mathbf{3}$ so find the 3^{th} smallest value

Output is 3: 1,2,**3**,4,5,6 (3th)





leetcode.com/problems/kth-smallest-element-in-a-bst

Solution

- Note that the smallest element is in the left leaf.
- Therefore, there is an order from small \rightarrow big values from left \rightarrow root \rightarrow right
- Perform in-order traversal k times and stop in the desired node



```
LeetCode
```

LeetCode leetcode.com/problems/kth-smallest-element-in-a-bst

Time: O(k) Space: O(h) where h is the height of the tree // in-order traversal: left, node: right void traverse(TreeNode* node, int& k, int& result) { // base case if (!node) return; // visit left first traverse(node->left, k, result); // visit node k--; if (k == 0) { result = node->val; return; // visit right traverse(node->right, k, result); int kthSmallest(TreeNode* root, int k) { // perform pre-order traversal int result; traverse(root, k, result); return result;

Problem - 124. Binary Tree Maximum Path Sum





leetcode.com/problems/binary-tree-maximum-path-sum

Problem

• ...

Problem - 124. Binary Tree Maximum Path Sum





leetcode.com/problems/binary-tree-maximum-path-sum

Solution

• ..

```
LeetCode
```

leetcode.com/problems/binary-tree-maximum-path-sum



• ...

Problem – 102. Binary Tree Level Order Traversal



LeetCode

leetcode.com/problems/binary-tree-level-order-traversal

Problem Statement / Solution / Code Time: O(-) Space: O(-)

• ..

Problem – 572. Subtree of Another Tree



LeetCode

leetcode.com/problems/subtree-of-another-tree

Problem Statement / Solution / Code Time: O(-) Space: O(-)

•



leetcode.com/problems/leaf-similar-trees

Problem Statement

- You are given two trees
- The goal is to compare if they have the same leaves
- The leaves should be in the same order
- Example:

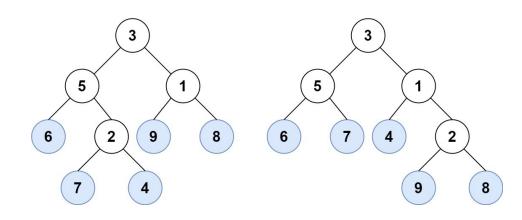
First tree:

leaves = 6,7,4,9,8 (blue nodes)

Second tree:

leaves =
$$6,7,4,9,8$$

Return true if the leaves are the same





leetcode.com/problems/leaf-similar-trees

Solution

- Get the first leaf value from tree 1
- Get the first leaf value from tree 2
- Compare, if they are different, return false immediately
- Otherwise, continue finding the next leaf value for tree 1 and 2

Implementation

- Create two stacks stack<TreeNode*> left and stack<TreeNode*> right
- Add the

Code – Leaf-Similar Trees

LeetCode

leetcode.com/problems/leaf-similar-trees

Code Time: O(n + m) where n and m are the numbers of nodes for trees 1 and 2 Space: O(h1 + h2) where h1 and h2 represents the height of the tree

```
// returns the value of the leaf, or -1 if empty
int getLeaf(stack<TreeNode*>& tree) {
   // tree is a reference, we will always pop an element from it
   while(!tree.empty()) {
       // get the top element from the stack
       TreeNode* node = tree.top();
       // already visited, so remove from stack
       tree.pop();
       // is this a leaf?
       if (!node->left && !node->right) {
           // yes, return the value
           return node->val;
        // push the right FIRST to the stack
       if (node->right) tree.push(node->right);
        // left should be on top of the stack
       if (node->left) tree.push(node->left);
   return -1;
```

```
bool leafSimilar(TreeNode* root1, TreeNode* root2) {
   // initialize the stacks, add root1 and root2
   std::stack<TreeNode*> leftTree, rightTree;
   leftTree.push(root1);
   rightTree.push(root2);
   while(true) {
       // get the leaves to compare
       int leaf1 = getLeaf(leftTree);
       int leaf2 = getLeaf(rightTree);
       // exit immediately if one leaf is different
       if (leaf1 != leaf2) return false;
       // stop when there are no leaves left
       if (leaf1 == -1 | leaf2 == -1) break;
   return true;
```





leetcode.com/problems/count-good-nodes-in-binary-tree

Problem Statement

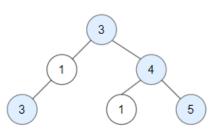
- You are given a binary tree and have to find "good" nodes
- A good node is a node where the values in the path are always than the node
- The root node is always a good node
- Example:
 - root 3 is a good node

left side:

- left leaf 1 is not a good node because 1 < 3
- leaf 3 is a good node because 3 > 1 and 3 == 3

right side:

- leaf 4 is a good node because 4 > 3
- **leaf 1** is not a good node because **1 < 4**
- leaf 5 is a good node because 5 > 4 > 3



Solution - 1448. Count Good Nodes in Binary Tree





leetcode.com/problems/count-good-nodes-in-binary-tree

Solution

..

LeetCode

leetcode.com/problems/count-good-nodes-in-binary-tree

Code Time: O(n + m) where n and m are the numbers of nodes for trees 1 and 2 Space: O(h1 + h2) where h1 and h2 represents the height of the tree

```
// returns the value of the leaf, or -1 if empty
int getLeaf(stack<TreeNode*>& tree) {
   // tree is a reference, we will always pop an element from it
   while(!tree.empty()) {
       // get the top element from the stack
       TreeNode* node = tree.top();
       // already visited, so remove from stack
       tree.pop();
       // is this a leaf?
       if (!node->left && !node->right) {
           // yes, return the value
           return node->val;
        // push the right FIRST to the stack
       if (node->right) tree.push(node->right);
        // left should be on top of the stack
       if (node->left) tree.push(node->left);
   return -1;
```

```
bool leafSimilar(TreeNode* root1, TreeNode* root2) {
   // initialize the stacks, add root1 and root2
   std::stack<TreeNode*> leftTree, rightTree;
   leftTree.push(root1);
   rightTree.push(root2);
   while(true) {
       // get the leaves to compare
       int leaf1 = getLeaf(leftTree);
       int leaf2 = getLeaf(rightTree);
       // exit immediately if one leaf is different
       if (leaf1 != leaf2) return false;
       // stop when there are no leaves left
       if (leaf1 == -1 | leaf2 == -1) break;
   return true;
```

Code – Count Good Nodes in Binary Tree

E LeetCode

leetcode.com/problems/count-good-nodes-in-binary-tree

Code (recursive) Time: O(n + m) where n and m are the numbers of nodes for trees 1 and 2 Space: O(h1 + h2) where h1 and h2 represents the height of the tree

INTERVAL

greedy strategy: sort by the end time

Because ending earlier gives **more room** for future intervals. It's a classic greedy trick: choose the interval that **frees up time** as quickly as possible.

Problem - 57. Insert Interval





leetcode.com/problems/insert-interval

Problem Statement

- You are given an array of intervals, where intervals[i] = [start, endi] and newInterval = [start, end]
- newInterval must be inserted into intervals
- Overlapping intervals must be merged
- Example

intervals = [[1,2],[3,5],[6,7],[8,10],[12,16]] newInterval = [4,8]

Output: [[1,2],[3,10],[12,16]]

Solution - 57. Insert Interval





leetcode.com/problems/insert-interval

Solution

- Sort intervals by the first element (start)
- Initialize result
- Solve in three loops:
 - 1. While there is no overlap with **newInterval**, add to **intervals[i]** to **result**
 - 2. While it overlaps, merge **newInterval**
 - 3. While until the end intervals and add the remaining intervals[i]



leetcode.com/problems/insert-interval

Code

Time: **O(n)** Space: **O(n)** where n is the size of intervals

```
vector<vector<int>> insert(vector<vector<int>>& intervals, vector<int>& newInterval) {
    vector<vector<int>> result;
    int tupleIndex = 0;
    int totalTuples = intervals.size();
    // 1. check if it overlaps
    // 1 ----- 2
    while (tupleIndex < totalTuples && intervals[tupleIndex][1] < newInterval[0]) {</pre>
        result.push_back(intervals[tupleIndex]);
        ++tupleIndex;
    // 2. merge overlap. We already know there is an overlap here,
    // otherwise it should be sorted out in the previous step
    // 3 ---- 5
          4 ---- 8
    while (tupleIndex < totalTuples && intervals[tupleIndex][0] <= newInterval[1]) {</pre>
        newInterval[0] = min(newInterval[0], intervals[tupleIndex][0]);
        newInterval[1] = max(newInterval[1], intervals[tupleIndex][1]);
        ++tupleIndex;
    result.push back(newInterval);
    // 3. add remaining parts
    while (tupleIndex < totalTuples) {</pre>
        result.push back(intervals[tupleIndex]);
        ++tupleIndex;
    return result;
```



leetcode.com/problems/merge-intervals

Problem Statement

•



leetcode.com/problems/merge-intervals

Solution

•

leetcode.com/problems/merge-intervals

Code Time: O(n) Space: O(n)

Problem – 435. Non-overlapping Intervals





leetcode.com/problems/non-overlapping-intervals

Problem Statement

• ...

Solution – 435. Non-overlapping Intervals





leetcode.com/problems/non-overlapping-intervals

Solution

•



leetcode.com/problems/non-overlapping-intervals



Time: O(n) Space: O(n)

• ..

LINKED LIST

Problem – 206. Reverse Linked List



LeetCode

leetcode.com/problems/reverse-linked-list

Problem Statement / Solution / Code Time: O(n) Space: O(n)

• ...

Problem – 141. Linked List Cycle



LeetCode

leetcode.com/problems/linked-list-cycle

Problem Statement / Solution / Code Time: O(n) Space: O(n)

•

Problem – 21. Merge Two Sorted Lists



LeetCode

leetcode.com/problems/merge-two-sorted-lists

Problem Statement / Solution / Code Time: O(n) Space: O(n)

•



leetcode.com/problems/merge-k-sorted-lists

Problem Statement / Solution / Code Time: O(n) Space: O(n)

•

Problem - 19. Remove Nth Node From End of List



E LeetCode

leetcode.com/problems/remove-nth-node-from-end-of-list

Problem Statement / Solution / Code Time: O(n) Space: O(n)

• ...

Problem - 143. Reorder List



LeetCode

leetcode.com/problems/reorder-list

Problem Statement / Solution / Code Time: O(n) Space: O(n)

• ..

https://leetcode.com/problems/swap-nodes-in-pairs

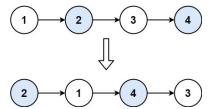
Problem

Given a linked list, swap every two adjacent nodes and return its head. You must solve the problem without modifying the values in the list's nodes (i.e., only nodes themselves may be changed.)

Example 1

Input: head = [1,2,3,4]

Output: [2,1,4,3]



Example 2

Input: head = []

Output: []

Example 3:

Example 3

Input: head = [1]

Output: [1]

Solution – Swap Nodes in Pair

https://leetcode.com/problems/swap-nodes-in-pairs

```
ListNode* swapPairs(ListNode* head) {
   if (head == NULL | head->next == NULL) {
        return head;
    ListNode *node = head;
    ListNode *prev = NULL;
    head = head->next;
    while (node && node->next) {
        ListNode *second = node->next;
        ListNode *next_pair = second->next;
        second->next = node;
       node->next = next_pair;
       if (prev) {
           prev->next = second;
        prev = node;
       node = next_pair;
    return head;
```

Solution (recursive) – Swap Nodes in Pair

https://leetcode.com/problems/swap-nodes-in-pairs

```
ListNode* swapPairs(ListNode* head) {
    if(!head || !head->next)
        return head;
    ListNode* newHead = head->next;
    head->next = swapPairs(head->next->next);
    newHead->next = head;
    return newHead;
}
```

HEAP / PRIORITY QUEUE

Heap

- Heap is a complete binary tree that satisfy the heap property (max or min)
- Min heap: root node contains the minimum value
- Max heap: root node contains the maximum value



Heap in C++

Two main ways to implement:

1. Using std::make_heap from <algorithm>

```
std::make_heap(RandomIt first, RandomIt last)
std::push_heap(RandomIt first, RandomIt last)
std::pop_heap(RandomIt first, RandomIt last)
std::sort_heap(RandomIt first, RandomIt last)
```

2. Using std::priority_queue from <queue> (recommended)

```
std::priority queue<T, Container, Compare>
```

Heap in C++ - std::priority_queue example

Min heap

```
std::priority_queue<int, std::vector<int>, std::greater<int>>
May boan
```

```
Max heap
std::priority_queue<int> or
std::priority queue<int, std::vector<int> std::less<int>>
// Min heap
std::priority queue<int, std::vector<int>, std::greater<int>> minHeap;
minHeap.push(3);
minHeap.push(6);
minHeap.push(4);
// remove top element (3)
minHeap.pop();
// root node (top) is now 4
std::cout << minHeap.top();</pre>
```

https://leetcode.com/problems/kth-largest-element-in-an-array

Problem

Given an integer array nums and an integer k, return the k^{th} largest element in the array. Note that it is the k^{th} largest element in the sorted order, not the k^{th} distinct element.

Example 1

Input: nums = [3,2,1,5,6,4], k = 2
Output: 5

Example 2

Input: nums = [3,2,3,1,2,4,5,5,6], k = 4

Output: 4

Although this problem is classified as "medium", in my opinion it should be classified as "easy"

Solution 1 – Kth Largest Element in an Array

https://leetcode.com/problems/kth-largest-element-in-an-array

```
// SOLUTION 1
int findKthLargest(vector<int>& nums, int k) {
    std::priority_queue<int, std::vector<int>, std::greater<int>> minHeap;
    for (const auto& num : nums) {
        if (minHeap.size() < k) {
            minHeap.push(num);
        } else if (num > minHeap.top()) {
            minHeap.pop();
            minHeap.push(num);
        }
    }
    return minHeap.top();
}
```

Solution 2 – Kth Largest Element in an Array

https://leetcode.com/problems/kth-largest-element-in-an-array

```
// SOLUTION 2 - Simpler approach
int findKthLargest(vector<int>& nums, int k) {
    // min heap: minimum values will be always at the top
    std::priority_queue<int, std::vector<int>, std::greater<int>> minHeap;
    for (const auto& num : nums) {
        // push each num to the heap
        minHeap.push(num);
        // we need the kth largest element only, so once after pushing more than k
        // elements, remove the smallest one (the top)
        if (minHeap.size() > k) {
            minHeap.pop();
        }
    }
    return minHeap.top();
}
```





leetcode.com/problems/top-k-frequent-elements

Problem

- You are given an array of numbers and an integer k
- Return an array with the \mathbf{k} most frequent elements

Example

Input:

```
nums = [1,1,1,2,2,3], k = 2
```

Output:

[1,2]





leetcode.com/problems/top-k-frequent-elements

Solution (1) - hashmap + array sort

Go over the array, count the numbers and store them in an unordered_map

Example:

```
nums = [1,1,1,2,2,3], k = 2
freq[1] = 3
freq[2] = 2
```

- Go over the unordered_map, add to an array and sort descending
- lacktriangle Create another array adding the **k** first elements and return

Code – 347. Top K Frequent Elements

E LeetCode

leetcode.com/problems/top-k-frequent-elements

Code (1) Time: O(n log n) Space: O(n)

```
vector<int> topKFrequent(vector<int>& nums, int k) {
   // 1. Create the number's frequency map
   // O(n)
   unordered map<int, int> freq;
   for (const auto& num : nums) {
        freq[num] += 1;
   // 2. Create an array with the frequencies
   vector<pair<int, int>> freqVec(freq.begin(), freq.end());
   // 3. Sort by the frequency O(n log n)
    sort(freqVec.begin(), freqVec.end(), [](auto& a, auto& b) {
            return a.second > b.second;
            });
   // 4. Create the result with the k first elements
   // 0(k)
   vector<int> result;
   for (int i = 0; i < k; ++i) {
        result.push back(freqVec[i].first);
   return result;
```





leetcode.com/problems/top-k-frequent-elements

Solution (2) - hashmap + min heap

Go over the array, count the numbers and store them in an unordered_map

Example:

```
nums = [1,1,1,2,2,3], k = 2
freq[1] = 3
freq[2] = 2
...
```

- Go over the frequencies, add to a min heap. If the size of the heap exceeds \mathbf{k} , remove the top one (the minimum value)
- Create another array result adding all elements from the heap and return it

Code – 347. Top K Frequent Elements

E LeetCode

leetcode.com/problems/top-k-frequent-elements

Code (2) Time: O(n log k) Space: O(n)

```
vector<int> topKFrequent(vector<int>& nums, int k) {
    // 1. Create the number's frequency map
   // O(n)
    unordered map<int, int> freq;
    for (const auto& num : nums) {
        freq[num] += 1;
    // 2. Create the min heap with priority queue
   // O(n log k)
    priority queue<pair<int, int>, vector<pair<int,int>>, greater<>> minHeap;
    for (const auto& [num, count] : freq) {
        minHeap.push({count, num});
        if (minHeap.size() > k) minHeap.pop();
    // 3. build the result
    vector<int> result;
    while (!minHeap.empty()) {
        auto num = minHeap.top().second;
        minHeap.pop();
        result.push back(num);
    return result;
```





leetcode.com/problems/top-k-frequent-elements

Solution (3) - hashmap + bucket sort

Go over the array, count the numbers and store them in an unordered_map

Example:

```
nums = [1,1,1,2,2,3], k = 2
freq[1] = 3
freq[2] = 2
...
```

Create buckets for each frequency and add the corresponding numbers:

```
bucket[1] = [3] \rightarrow 3 only appears once in nums
bucket[2] = [2] \rightarrow 2 appears twice
bucket[3] = [1] \rightarrow 1 appears three times
```

Go over each bucket, add to the result and return it

Code – 347. Top K Frequent Elements

```
LeetCode
```

leetcode.com/problems/top-k-frequent-elements

Code (3) Time: O(n) Space: O(n)

```
vector<int> topKFrequent(vector<int>& nums, int k) {
    // Create the number's frequency map
    unordered map<int, int> freq;
    for (const auto& num : nums) {
        freq[num]++;
    // create the buckets
    // e.g. [[1,2,3],[4,5,6]] ...
    vector<vector<int>> buckets(nums.size() + 1);
    for (const auto& [num, count] : freq) {
        buckets[count].push back(num);
    // go over each bucket to build the result
    vector<int> result;
    for (int i = buckets.size() - 1; i >= 0; --i) {
        for (const auto& num : buckets[i]) {
            result.push back(num);
            if (result.size() == k) return result;
    return result;
```

Problem – 347. Top K Frequent Elements





leetcode.com/problems/top-k-frequent-elements

Some considerations

- Theoretically, bucket sort should be the fastest solution O(n) < O(n log k)
- In practice, min heap end up being faster:
 - fewer allocations: priority_queue stores flat pairs rather than inner vectors
 - better cache locality: heap is built over a single array (binary heap)
 - if **k** is small, heap touches fewer elements

MATRIX

Problem - 73. Set Matrix Zeroes



LeetCode

leetcode.com/problems/set-matrix-zeroes

Problem Statement / Solution / Code Time: O(-) Space: O(-)

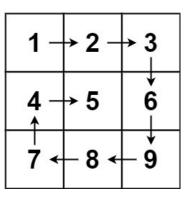
• ..



leetcode.com/problems/spiral-matrix

Problem Statement / Solution / Code Time: O(-) Space: O(-)

• ..



Problem – 48. Rotate Image



LeetCode

leetcode.com/problems/rotate-image

Problem Statement / Solution / Code Time: O(-) Space: O(-)

• ...

DYNAMIC PROGRAMMING

Dynamic Programming

Dynamic Programming (DP) is an algorithm technique used to solve problems that can be broken down into **simpler, overlapping subproblems.**

Key Concepts of Dynamic Programming

- Overlapping subproblems: a problem has overlapping subproblems if it can be broken down into subproblems.
- **Memoization (Top-Down Approach)**: store the results in a cache (typically a dictionary or array) to avoid recalculation recursion and caching approach.
- **Tabulation (Bottom-Up Approach)**: first solve all possible subproblems iteratively, and store them in a table.

Common Patterns in Dynamic Programming

- Toy example (Fibonacci): Climbing Stairs, N-th Tribonacci Number, Perfect Squares
- Constant Transition: Min Cost Climbing Stairs, House Robber, Decode Ways, Minimum Cost For Tickets, Solving Questions With Brainpower
- Grid: Unique Paths, Unique Paths II, Minimum Path Sum, Count Square Submatrices with All Ones, Maximal Square,
 Dungeon Game
- Dual-Sequence: Longest Common Subsequence, Uncrossed Lines, Minimum ASCII Delete Sum for Two Strings, Edit
 Distance, Distinct Subsequences, Shortest Common Supersequence
- Interval: Longest Palindromic Subsequence, Stone Game VII, Palindromic Substrings, Minimum Cost Tree From Leaf Values, Burst Balloons, Strange Printer
- Longest Increasing Subsequence: Count Number of Teams, Longest Increasing Subsequence, Partition Array for Maximum Sum, Largest Sum of Averages, Filling Bookcase Shelves
- Knapsack: Partition Equal Subset Sum, Number of Dice Rolls With Target Sum, Combination Sum IV, Ones and Zeroes,
 Coin Change, Coin Change II, Target Sum, Last Stone Weight II, Profitable Schemes
- Topological Sort on Graphs: Longest Increasing Path in a Matrix, Longest String Chain, Course Schedule III
- DP on Trees: House Robber III, Binary Tree Cameras
- Other problems: 2 Keys Keyboard, Word Break, Minimum Number of Removals to Make Mountain Array, Out of Boundary Paths

Credits

Dynamic Programming – Example – Fibonacci Sequence

```
Naive Recursive Approach

int fib(int n) {
   if (n <= 1) {
      return n;
   }
   return fib(n - 1) + fib(n - 2);
}</pre>
```

```
Memoization (Top-Down DP)

std::unordered_map<int, int> memo;

int fib(int n) {
    if (n <= 1) {
        return n;
    }
    if (memo.find(n) != memo.end()) {
        return memo[n];
    }
    memo[n] = fib(n - 1) + fib(n - 2);
    return memo[n];
}</pre>
```

```
Tabulation (Bottom-up DP)

int fib(int n) {
    if (n <= 1) {
        return n;
    }
    int dp[n + 1];
    dp[0] = 0;
    dp[1] = 1;
    for (int i = 2; i <= n; i++) {
        dp[i] = dp[i - 1] + dp[i - 2];
    }
    return dp[n];
}</pre>
```

Problem – Climbing Stairs



leetcode.com/problems/climbing-stairs

Problem Statement

You need to climb a staircase with n steps to get to the top. Each time you can choose to climb either 1 step or 2 steps at a time. Find out how many different ways you can climb to the top of the staircase.

Example 1

Input: n = 2

Output: 2

Explanation: There are two ways to get to the top

- 1. Climb 1 step at a time, twice
- 2. Climb 2 steps in one go

Example 2:

Input: n = 3

Output: 3

Explanation: There are three ways to get to the top:

- 1. Climb 1 step at a time, three times
- 2. Climb 1 step, then 2 steps
- 3. Climb 2 steps, then 1 ste.

Solution – Climbing Stairs

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LeetCode
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leetcode.com/problems/climbing-stairs

```
std::unordered map<int, int> memo;
int climbStairs(int n) {
   // Identify the sequence, when:
   // n = 0 (0 way), there is no way to get up
   // n = 1 (1 way): only one way : 1-step
   // n = 2 (2 ways): 1s + 1s | 2s
   // n = 3 (3 ways): 1s + 1s + 1s | 1s + 2s | 2s + 1s
   // n = 4 (5 ways): 1s + 1s + 1s + 1s | 1s + 1s + 2s | 1s + 2s + 1s | 2s + 1s + 1s | 2s + 2s |
   if (n <= 2) {
       return n;
   if (memo.find(n) != memo.end()) {
       return memo[n];
   memo[n] = climbStairs(n - 1) + climbStairs(n - 2);
   return memo[n];
```

Problem – 1143. Longest Common Subsequence



LeetCode

https://leetcode.com/problems/longest-common-subsequence

Problem Statement / Solution / Code Time: O(-) Space: O(-)

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Problem – 62. Unique Paths



LeetCode

https://leetcode.com/problems/unique-paths

Problem Statement / Solution / Code Time: O(-) Space: O(-)

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Tips

Problem Statement / Solution / Code Time: O(n) Space: O(n)

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Problem – number. name









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Problem Statement / Solution / Code Time: O(-) Space: O(-)

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