Algorithm and Problem Solving Cheatsheet in C++

Data Structures, Algorithms and Coding Interview Problem Patterns in C++

rfdavid, 2025

rfdavid.com

MOTIVATION

Motivation

The tech industry hiring standard is based on algorithm and data structure.

There are plenty of free resources available around algorithms and data structures. The purpose of this project is to be a quick guide where you can learn and review learned algorithms and data structures.

Some of the intended **key features:**

- Non-verbose, short-structured, and easy to follow descriptions
- Slide-based, practical for reviewing
- Free and open-source

right in the please add a star at github.com/rfdavid/cpp-algo-cheatsheet

Some Useful Links

Tech Interview Handbook

https://www.techinterviewhandbook.org

A very well-structured resource for interview preparation

TUF

https://takeuforward.org/interviews/blind-75-leetcode-problems-detailed-video-solutions

Contains explanation and some videos for the problems from blind 75 list

Blind 75 Leetcode Questions

https://leetcode.com/discuss/general-discussion/460599/blind-75-leetcode-guestions

Blind 75

- Blind 75 is a popular list of algorithm problems that intends to cover the main data structures and patterns.
- It is a curated list of 75 popular coding questions created by an ex-Meta Staff Engineer

Array

✓ Two Sum

✓ Best Time to Buy and Sell Stock

Contains Duplicate

Product of Array Except Self

Maximum Subarray

Maximum Product Subarray

Find Minimum in Rotated Sorted Array

Search in Rotated Sorted Array

3 Sum

✓ Container With Most Water

Binary

<u>Sum of Two Integers</u>

Number of 1 Bits

Counting Bits

Missing Number

Reverse Bits

Dynamic Programming

✓ Climbing Stairs

Coin Change

Longest Increasing Subsequence

Longest Common Subsequence

Word Break Problem

Combination Sum

House Robber

House Robber II

Decode Ways

Unique Paths

Jump Game

Matrix

Set Matrix Zeroes

Spiral Matrix

Rotate Image

Word Search

Blind 75

Tree

✓ Maximum Depth of Binary Tree

Same Tree

Invert/Flip Binary Tree

Binary Tree Maximum Path Sum

Binary Tree Level Order Traversal

Serialize and Deserialize Binary Tree

Subtree of Another Tree

Construct Binary Tree from Preorder and Inorder Traversal

Validate Binary Search Tree

Kth Smallest Element in a BST

Lowest Common Ancestor of BST

Implement Trie (Prefix Tree)

Add and Search Word

Word Search II

Heap

Merge K Sorted Lists

Top K Frequent Elements

Find Median from Data Stream

String

✓ Longest Substring Without Repeating Characters

Longest Repeating Character Replacement

Minimum Window Substring

Valid Anagram

Group Anagrams

✓ Valid Parentheses

Valid Palindrome

Longest Palindromic Substring

Palindromic Substrings

Encode and Decode Strings ☆

Linked List

Reverse a Linked List

Detect Cycle in a Linked List

Merge Two Sorted Lists

Merge K Sorted Lists

Remove Nth Node From End Of List

Reorder List

Graph

✓ Clone Graph

Course Schedule

Pacific Atlantic Water Flow

Number of Islands

Longest Consecutive Sequence

Graph Valid Tree ☆

Number of Connected Components

In an Undirected Graph 🖈

Interval

√ Insert Interval

✓ Merge Intervals

✓ Non-overlapping Intervals

Meeting Rooms II ☆



Other problems

Tree

- ✓ <u>Maximum Level Sum of a Binary Tree</u>
- ✓ <u>Minimum Number of Increments on Subarrays to Form a Target Array</u>
- ✓ <u>Leaf-Similar Trees</u>
- ✓ Count Good Nodes in Binary Tree

DATA STRUCTURES IN C++

Data Structure Decision Diagram



 The following diagram gives you the direction to which data structure to use in C++ according to the problem you are trying to solve

Note: I don't have the source of this diagram. If you know it, please drop me a msg so I can add it here.

Vectors

std::vector is a sequence container that encapsulates dynamic sized arrays*

Linked List

Stack

Queue

Heap

Hash Table

Tree

ARRAY

Arrays

Characteristics

- Memory layout: hold values in a contiguous block of memory.
- **Fixed Size**: the size of an array is defined when it is created and cannot be changed. However, high-level languages have different implementations, making it dynamic.
- Homogeneous elements: all elements are of the same data type (int, float, char...)
- **Efficiency**: accessing elements by index is very efficient *O(1)*, since each index maps directly to a memory location. Also, range scans benefit from CPU cache lines since arrays are stored in contiguous blocks of memory.

Arrays – Kadane's algorithm

- Kadane's algorithm is a dynamic programming algorithm to solve maximum subarray sum
- At every index i: start a new subarray at i extend the previous subarray to include array[i]

Algorithm

1. Initialize:

```
int maxSoFar = array[0];
int maxEndingHere = array[0];
```

2. Loop through the array

```
for (int i = 1; i < array.size(); ++i) {
    maxEndingHere = max(array[i], maxEndingHere + array[i]);
    maxSoFar = max(maxSoFar, maxEndingHere);
}</pre>
```

3. Return maxSoFar;



Problem Statement

- Given an array of numbers and a target, example: array [2,7,11,15] and target 9
- Return indices of two numbers where they add up to target
- **Output**: [0,1]

```
array[0] + array[1] = 2 + 7 = 9
```



Solution

- Iterative over each number in the array
- Calculate the difference between target and each number, example:

```
array[0] = 2, target 9, then 9 - 2 = 7
```

- Now we know we need the number 7 to sum up to 9
- Check in a hashmap if we have 7 in some part of the array

```
hash[7] exists?
```

- If yes, return the current index and the index of 7
- If not, store the index of the current number in the hashmap for future evaluation

$$hash[2] = 0$$

Code - Two Sum



LeetCode leetcode.com/problems/two-sum

```
Code Time: O(n) Space: O(n)
vector<int> twoSum(vector<int>& nums, int target) {
    std::unordered map<int, int> numMap;
   // n being the size of nums
   for (int i = 0; i < nums.size(); i++) {</pre>
        // current number of the array
        int number = nums[i];
        int diff = target - number;
        // check if the difference is in some part of the array
        // by using a hashmap
        if (numMap.find(diff) != numMap.end()) {
            return { numMap[diff], i};
        // register the current number index
        numMap[number] = i;
   // no matches
   return {};
```

Problem - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Problem Statement

- You are given an integer array of stock prices
- Choose a price[i] to buy and price[i] to sell where you achieve maximum profits
- Example:

```
prices = [9, 1, 3, 4]
```

• **Output**: [1,3]

```
array[3] - array[1] = 4 + 1 = 3
```

Solution - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Solution

- Initialize profit = 0
- Initialize lowestBuyPrice = prices[0]
- Loop through the prices
- Track the lowest buy price → min(lowestBuyPrice, prices[i])
- Check if selling "today" will make the maximum profit and update profit:
 max(prices[i] buy > profit, profit)
- Update profit max(prices[i] - buy

Code - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Code (simplified) Time: O(n) Space: O(n)

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      buy = min(buy, prices[i]);
      profit = max(profit, prices[i] - buy)
   }
   return profit;
}</pre>
```

Code - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Code (optimized) Time: O(n) Space: O(n)

Same logic, but with better branch prediction and less computation

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      if (prices[i] < buy) {
        buy = prices[i];
      } else if (prices[i] - buy > profit) {
           profit = prices[i] - buy;
      }
   }
   return profit;
}
```





leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

Problem Statement

- You are given an integer array of stock prices
- Choose a price[i] to buy and price[i] to sell where you achieve maximum profits
- You can buy/sell multiple times, but only hold at most one transaction at a time
- Output is the maximum profits

• Example:

```
prices = [9, 1, 3, 4]

Output: 2 + 1 = 3

buy (price = 1), sell (price = 3), profit = 2

buy (price = 3), sell (price = 4), profit = 1
```



leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

Solution

- Loop through the array starting from index 1
- If current price[i] is lower than previous price[i 1], buy and sell

• Example:

```
prices = [1, 8, 4] prices[0] = 1, prices[1] = 8, prices[2] = 4
prices[0] < \text{prices}[1] \rightarrow true, profit = 8 - 1 = 7
prices[2] < \text{prices}[1] \rightarrow \text{false}, do nothing
```

Code - Best Time to Buy and Sell Stock II

LeetCode

leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

```
int maxProfit(vector<int>& prices) {
  int profit = 0;
  for (int i = 1; i < prices.size(); ++i) {
    if (prices[i] > prices[i-1]) {
       profit += prices[i] - prices[i -1];
    }
  }
  return profit;
}
```



leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Problem Statement

• ..

Solution - Best Time to Buy and Sell Stock IV





leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Solution

• ..

Code - Best Time to Buy and Sell Stock IV



leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Code (simplified) Time: O(n) Space: O(n)

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      buy = min(buy, prices[i]);
      profit = max(profit, prices[i] - buy)
   }
   return profit;
}</pre>
```

STRING

Problem – 3. Longest Substring Without Repeating Characters





leetcode.com/problems/longest-substring-without-repeating-characters

Problem Statement

 You are given a string and the goal is to find the longest substring without repeating characters

Example

Input: "abcdb"

Output: 4 (abcd since "b" is repeated)

Solution – 3. Longest Substring Without Repeating Characters





leetcode.com/problems/longest-substring-without-repeating-characters

Solution

- Use sliding window algorithm (left and right)
- Loop through the string
- Try to find if the current character is already added by using unordered set or bitmap
- If added, remove from the set alongside with others using left pointer
- If not, add to the unordered set or bitmap
- Maximum length will be right left + 1



LeetCode leetcode.com/problems/longest-substring-without-repeating-characters

Example

- String: abcbd. Our goal is to return 3 (abcbd)
- Initialize maxLength = 0
- Loop through the string

```
lteration 1: left = 0, right = 0, string[left] = 'a',
          bitmap = ['a'] ('a' is not in bitmap, add), maxLength = max(maxLength, right - left + 1) = 1
lteration 2: left = 0, right = 1, string[right] = 'b'
          bitmap = ['a','b'], maxLength = 2
lteration 3: left = 0, right = 2, string[right] = 'c'
          bitmap = ['a','b', 'c'], maxLength = 3
lteration 4: left = 0, right = 3, string[right] = 'b'
          bitmap = ['a','b','c','b']
          'b' is already in the bitmap. start "clearing" the character using left:
          Iteration 4a: left = 0, string[left] = 'a' is different from 'b', so remove 'a'
                      bitmap = ['b', 'c','b']
          Iteration 4b: left = 1, string[left] = 'b' is the same as the repeated one, remove
                       bitmap = [`c', 'b']
Iteration 5: left = 1, right = 4, string[right] = 'd'
            bitmap = ['c','b','d']
```

Code – 3. Longest Substring Without Repeating Characters

Code (unordered_set)

Use unordered_set when question requires unicode chars

```
int lengthOfLongestSubstring(string s) {
    int maxLength = 0;
    int left = 0, right = 0;
    // track the seen characters
    unordered set<char> seen;
    for (right = 0; right < s.size(); ++right) {</pre>
        char currentChar = s[right];
        // if currentChar is in the set, clean
        // the character and everything from left of it
        // basically, reset the longest substring
        while (seen.count(currentChar)) {
            char c = s[left];
            seen.erase(c);
            left++;
        // insert the current read character
        seen.insert(currentChar);
        // set max length
        maxLength = max(maxLength, right - left + 1);
    return maxLength;
```

Code – 3. Longest Substring Without Repeating Characters

Code (bitmap)

- Using bitset: create a bitmask with 128 bits where each bit represent a character
- Optimal solution for ASCII since ASCII size is 127 characters
- Unicode / UTF-8 can represent over 1.1 million characters, so use unordered_set approach instead

```
int lengthOfLongestSubstring(string s) {
    std::bitset<128> bitmask;
    uint32_t left = 0;
    uint32_t maxLength = 0;

for (uint32_t right = 0; right < s.length(); ++right) {
        uint32_t bitIndex = s[right];
        // if char is already in the bitmask, move left until we reset the bits
        while (bitmask.test(bitIndex)) {
            bitmask.reset(s[left]);
            ++left;
        }

        bitmask.set(bitIndex);
        maxLength = std::max(maxLength, right - left + 1);
    }
    return maxLength;
}</pre>
```



LeetCode leetcode.com/problems/valid-parentheses

Problem Statement

- You are given a string containing only the characters '(', ')', '{', '}', '[' and ']'
- A valid input have closed brackets by its own type
- Example

()[] $\{\}$ \rightarrow valid

[]{}(\rightarrow invalid

 $\{()\} \rightarrow \text{valid}$

Solution - Valid Parentheses



leetcode.com/problems/valid-parentheses

Solution

- Loop through the string
- If open brackets ([{ push to a stack
- If closed brackets:
 - pop the last added bracket
 - **check** if the **closed** bracket corresponds to the **popped** bracket
 - if not, return false
- after the loop, return true if the size of the stack is empty (all brackets closed)

Code – Valid Parentheses



LeetCode leetcode.com/problems/valid-parentheses

```
Code
          Time: O(n) Space: O(n)
bool isValid(string s) {
   // stack (LIFO)
   std::stack<char> brackets;
   // O(n)
   for (int i = 0; i < s.size(); ++i) {</pre>
        char bracket = s[i];
        if (bracket == '(' || bracket == '[' || bracket == '{'}) {
           brackets.push(bracket);
       } else {
           if (brackets.size() == 0) return false;
           char lastBracket = brackets.top();
           if (bracket == ')' && lastBracket != '(') return false;
           if (bracket == '}' && lastBracket != '{') return false;
           if (bracket == ']' && lastBracket != '[') return false;
           brackets.pop();
   // all brackets must be closed
   return brackets.size() == 0;
```

Problem – Minimum Number of Increments on Subarrays



leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Problem Statement

- You are given an array of integers initialized with zeros (e.g. [0,0,0,0])
- The goal is to reach some target (e.g. [1, 2, 2, 3])
- The valid operations is to increment a subarray by one
- The output is the total number of operations In this case:
 - $[1,1,1,1] \rightarrow$ increment the subarray starting from 0 to total size
 - [1,2,2,2] → increment the subarray starting from 1 to total size
 - $[1,2,2,3] \rightarrow$ increment the subarray starting and ending from the last element

Output: 3 (total number of operations)

Solution – Minimum Number of Increments on Subarrays





leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Solution

Explain...

Code (2) - Minimum Number of Increments on Subarrays



leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Code (optimized)

Code – Minimum Number of Increments on Subarrays

LeetCode

leetcode.com/problems/minimum-number-of-increments-on-subarrays-to-form-a-target-array

Code

```
int minNumberOperations(vector<int>& target) {
    int totalOp = target[0];
    for (int i = 1; i < target.size(); ++i) {
        // can't reuse
        if (target[i - 1] < target[i]) {
            totalOp += target[i] - target[i - 1];
        }
    }
    return totalOp;
}</pre>
```

TWO POINTERS



leetcode.com/problems/container-with-most-water

Problem Statement

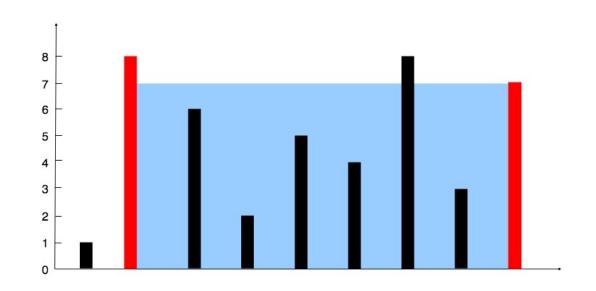
- You are given an integer array height
- Find two lines that together with x-axis form a container with most water
- Example:

Input:

height = [1,8,6,2,5,4,8,3,7]

Output:

49





leetcode.com/problems/container-with-most-water

Solution

- Initialize the maximum area maxArea = 0
- Initialize two pointers, left = 0 and right = height.size 1
- Loop while pointer left < right
- Calculate the area:

area = min(height[left], height[right]) * (right - left)

Update the global maximum area:

maxArea = max(maxArea, area)

- Move the smallest pointer (increment left or decrement right)
- Return maxArea

LeetCode

leetcode.com/problems/container-with-most-water

Code

```
int maxArea(vector<int>& height) {
   // initialize the two pointers (left and right)
   int left = 0;
   int right = height.size() - 1;
   int maxArea = 0;
   while (left < right) {</pre>
       // calculate the area, think about the x-axis and y-axis
        int area = min(height[left], height[right]) * (right - left);
       // update maximum area
       maxArea = max(area, maxArea);
       // is the left pointer (y) smaller than right?
        if (height[left] < height[right]) {</pre>
           // move left pointer to right
           left++;
        } else {
            // otherwise, move right pointer to left
            right--;
   return maxArea;
```

BINARY

Bit Manipulation in C

Operators

```
& AND OR ^ XOR ~ NOT << LEFT SHIFT >> RIGHT SHIFT
```

Common Operations

```
set bit: num |= (1 << pos)

clear bit: num &= ~(1 << pos)

toggle bit: num ^= (1 << pos)

check bit: (num & (1 << pos)) != 0

extract bit: (num >> pos) & 1

extract a range of bits: (num >> pos) & ((1 << length) - 1)</pre>
```

Example

```
void copyBit(int *dst, int src, int srcPos, int dstPos) {
    int bit = (src >> srcPos) & 1; // extract bit
    *dst &= ~(1 << dstPos); // clear destination bit
    *dst |= (bit << dstPos); // set destination bit
}</pre>
```

Binary

- In C++, **std::bitset** represents a fixed-size sequence of N bits
- Example:

```
std::bitset<8> bitmask;
bitmask.reset(1)
bitmask.set(1)
if (bitmask.test(1)) { // true
...
```

- reset : set bit to false
- **set** : set a specific bit
- **test** : check a specific bit
- **count** : return the number of bits set to true
- **flip**: toggle the value of the bits (if true, set to false and vice-versa)

Negabinary

- Non-standard positional numeral system that uses base of -2
- Allow representing negative numbers in binary
- Example:

$$1101_{-2}$$

$$(-2)^3 + (-2)^2 + 0 + (-2)^0 = -8 + 4 + 0 + 1 = -3$$

Summing Negabinary

Add as a regular binary number, but with negative carry

$$0 + 0 = 0$$

 $1 + 0 = 1$
 $1 + 1 = 0$ with a negative carry 1
 $1 + 1 = 0$ (subtract)
 $1 + 0 = 1$ with a positive carry 1

Negabinary

Example 1

Example 2

$$\begin{array}{r}
 1111 \\
 101010 \\
 + 101100 \\
\hline
 = 11110110
 \end{array}$$

Reference

https://leetcode.com/problems/adding-two-negabinary-numbers

Given two numbers arr1 and arr2 in base -2, return the result of adding them together.

Each number is given in array format: as an array of 0s and 1s, from most significant bit to least significant bit. For example, arr = [1,1,0,1] represents the number $(-2)^3 + (-2)^2 + (-2)^0 = -3$. A number arr in array, format is also guaranteed to have no leading zeros: either arr == [0] or arr[0] == 1.

Return the result of adding arr1 and arr2 in the same format: as an array of 0s and 1s with no leading zeros.

Example 1

```
Input: arr1 = [1,1,1,1,1], arr2 = [1,0,1]
```

Output: [1,0,0,0,0]

Explanation: arr1 represents 11, arr2 represents 5, the output represents 16.

Example 2

```
Input: arr1 = [0], arr2 = [0]
```

Output: [0]

Example 3

```
Input: arr1 = [0], arr2 = [1]
```

Output: [1]

Solution 1073 – Adding Two Negabinary Numbers



https://leetcode.com/problems/adding-two-negabinary-numbers

GRAPH (DFS)

Problem - Keys and Rooms

https://leetcode.com/problems/keys-and-rooms

```
int maxProfit(vector<int>& prices) {
    int profit = 0;
    int buy = prices[0];
    for (auto i = 1; i < prices.size(); i++) {
        if (prices[i] < buy) {
            buy = prices[i];
        } else if (prices[i] - buy > profit) {
                profit = prices[i] - buy;
        }
    }
    return profit;
}
```



https://leetcode.com/problems/clone-graph

Problem Statement

- Given a node reference, create a deep copy of the graph
- The class node has two variables: val and neighbours

```
class Node {
  public int val;
  public List<Node> neighbors;
}
```

Output is the node reference of the copy



https://leetcode.com/problems/clone-graph

Solution

- First check the edge cases (is the node null?)
- Create a hash map to store the nodes that is already created unordered<int, Node*> graph;
- Check if the current node already exists in the graph
- If not, create a new Node object and store in the hashmap
- Visit all the neighbors and add the neighbors to this current node

Code - Clone Graph

```
E LeetCode
```

https://leetcode.com/problems/clone-graph

```
std::unordered_map<int, Node*> graph;
Node* cloneGraph(Node* node) {
   if (node == NULL) {
        return NULL;
    // does this node object exists?
   if (graph.find(node->val) == graph.end()) {
        // node wasn't visited yet, store in the hashmap
        graph[node->val] = new Node(node->val);
        // visit all neighnours
        for (const auto& n : node->neighbors) {
            graph[node->val]->neighbors.push_back(cloneGraph(n));
   return graph[node->val];
```

GRAPH (BFS)





https://leetcode.com/problems/maximum-level-sum-of-a-binary-tree

Problem Statement

- Given the root of a binary tree, find the smallest level with the maximum sum
- For example, the tree below has the follow sums for each level:

level
$$1 \text{ (root)} = 1$$

level
$$2 = 7 + 0 = 7$$

$$|eve| 3 = 7 - 8 = -1$$

Therefore, level 2 has the maximum sum



Solution – Maximum Level Sum of a Binary Tree





https://leetcode.com/problems/maximum-level-sum-of-a-binary-tree

Solution

- Have a queue with the nodes for the current level
- Sum the values from that level by taking the nodes from the queue
- Example, we know that level 1 has one node. Hence, pop the first node from the queue
 If level 2 has 2 nodes, pop two nodes, sum the values
- In addition, add left and right to the end of the queue to process the next level

Code – Maximum Level Sum of a Binary Tree

```
E LeetCode
```

https://leetcode.com/problems/maximum-level-sum-of-a-binary-tree

```
int maxLevelSum(TreeNode* root) {
    std::queue<TreeNode*> nodes;
    int currentLevel = 0;
    int maxLevel = 1;
    int maxSum = INT MIN;
    nodes.push(root);
    // traverse the graph
    while(!nodes.empty()) {
        int levelSum = 0;
        int levelSize = nodes.size();
        currentLevel++;
        // sum the values in current level
        for (int i = 0; i < levelSize; ++i) {</pre>
            TreeNode* node = nodes.front();
            levelSum += node->val;
            nodes.pop();
            if (node->left) nodes.push(node->left);
            if (node->right) nodes.push(node->right);
        if (levelSum > maxSum) {
            maxLevel = currentLevel;
            maxSum = levelSum;
    return maxLevel;
```

SHORTEST PATH

Shortest Path Algorithms

Algorithms

- BFS
- Dijkstra
- Bellman-Ford
- Floyd-Warshall
- A* search
- Johnson's
- SPFA (Shortest Path Faster)
- Bidirectional Search

TREE

Depth-First Traversals

• **Pre-order**: Root – Left – Right



• In-order: Left - Root - Right



• **Post-order**: Left – Right – Root



Breadth-First Traversal (Level Order Traversal)

Visit every node on a level before moving to a lower level.

Depth-First Traversals

Use a recursive algorithm to traverse according to the order

if (!root) return; • **Pre-order**: Root – Left – Right doSomething(); visit(node->left); visit(node->right); if (!root) return; • In-order: Left – Root – Right visit(node->left); doSomething(); visit(node->right); if (!root) return; • **Post-order**: Left – Right – Root visit(node->left); visit(node->right);

doSomething();

Example of pre-order and in-order

```
struct TreeNode {
    int val;
    TreeNode *left, *right;
    TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
};
// Pre-order traversal
void preorderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    cout << root->val << " ";</pre>
    preorderTraversal(root->left);
    preorderTraversal(root->right);
// In-order traversal
void inorderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    inorderTraversal(root->left);
    cout << root->val << " ";</pre>
    inorderTraversal(root->right);
```

Example of post-order and level-order

```
struct TreeNode {
    int val;
    TreeNode *left, *right;
    TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
};
// Post-order traversal
void postorderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    postorderTraversal(root->left);
    postorderTraversal(root->right);
    cout << root->val << " ";</pre>
// Level-order traversal using a queue
void levelOrderTraversal(TreeNode* root) {
    if (root == nullptr) return;
    queue<TreeNode*> q;
    q.push(root);
    while (!q.empty()) {
        TreeNode* current = q.front();
        q.pop();
        cout << current->val << " ";</pre>
        if (current->left != nullptr) q.push(current->left);
        if (current->right != nullptr) q.push(current->right);
```

BFS Using Stack

BFS with std::stack

 This might be useful for problems when you want to return and resume (for example, <u>872. Leaf-Similar Trees</u>)

```
struct TreeNode {
    int val;
    TreeNode *left, *right;
    TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
};

// Pre-order traversal
void bfs(std::stack<TreeNode*>& tree) {
    while(!tree.empty()) {
        TreeNode* root = tree.top();
        tree.pop();
        // do something ...
        if (root->right) tree.push(root->right);
        if (root->left) tree.push(root->left);
    }
}
```





LeetCode https://leetcode.com/problems/maximum-depth-of-binary-tree

Problem Statement

- Given the root of a binary tree, find the <u>maximum depth</u>
- Example:

Output: 4



Solution – Maximum Depth of Binary Tree



LeetCode https://leetcode.com/problems/maximum-depth-of-binary-tree

Solution

- Perform post-order traversal: left right root
- Recursively go left and right to find each value
- Return the max of each one

Code – Maximum Depth of Binary Tree

LeetCode https://leetcode.com/problems/maximum-depth-of-binary-tree

```
int maxDepth(TreeNode* root) {
   if (!root) return 0;
   // find max left
   int maxLeft = maxDepth(root->left);
   // find max right
   int maxRight = maxDepth(root->right);
   // return max +1 (account for root)
   return std::max(maxLeft, maxRight) + 1;
```



https://leetcode.com/problems/path-sum

Problem Statement

- It is given the root of a binary tree and an integer target sum
- Example:



Output: true

Node
$$1 + Node 7 + Node 2 = 10$$





https://leetcode.com/problems/path-sum

Solution

- Start from root node (1)
- Subtract from target number (example 10 1 = 9)
- Continue going down the tree, until the target is 0, return true
- After visiting all nodes, if the target is not zero, return false



Code – Path Sum

```
E LeetCode
```

https://leetcode.com/problems/path-sum

```
bool hasPathSum(TreeNode* root, int targetSum) {
    if (!root) {
        return false;
    // we want targetSum to be zero
   targetSum -= root->val;
   // if there is no left, no right, we've reached the end of the path
    // so if the targetSum is zero, then the nodes summed up to the targetSum
    if (!root->left && !root->right && targetSum == 0) {
        return true;
    // propagate to left and right
    return hasPathSum(root->left, targetSum) || hasPathSum(root->right, targetSum);
```

Also, a small performance tweak can be made by avoiding writing targetSum: targetSum -= root->val

This will avoid a memory write access, making the calculation directly in the CPU, but also at a cost of readability

```
if (!root->left && !root->right && targetSum - root->val == 0) {
    ...
return hasPathSum(root->left, targetSum - root->val) || hasPathSum(root->right, targetSum - root->val);
```



LeetCode leetcode.com/problems/kth-smallest-element-in-a-bst

Problem Statement / Solution

- You are given the root of a binary search tree and an integer k
- Find the kth smallest value

Example

From all values in the tree: 1,2,3,4,5,6

 $\mathbf{k} = \mathbf{3}$ so find the 3^{th} smallest value

Output is 3: 1,2,**3**,4,5,6 (3th)





leetcode.com/problems/kth-smallest-element-in-a-bst

Solution

- Note that the smallest element is in the left leaf.
- Therefore, there is an order from small \rightarrow big values from left \rightarrow root \rightarrow right
- Perform in-order traversal k times and stop in the desired node



```
LeetCode
```

LeetCode leetcode.com/problems/kth-smallest-element-in-a-bst

Time: O(k) Space: O(h) where h is the height of the tree // in-order traversal: left, node: right void traverse(TreeNode* node, int& k, int& result) { // base case if (!node) return; // visit left first traverse(node->left, k, result); // visit node k--; if (k == 0) { result = node->val; return; // visit right traverse(node->right, k, result); int kthSmallest(TreeNode* root, int k) { // perform pre-order traversal int result; traverse(root, k, result); return result;

BINARY TREE (DFS)



leetcode.com/problems/leaf-similar-trees

Problem Statement

- You are given two trees
- The goal is to compare if they have the same leaves
- The leaves should be in the same order
- Example:

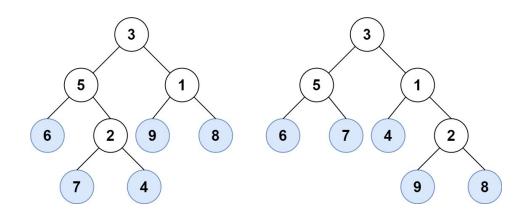
First tree:

leaves = 6,7,4,9,8 (blue nodes)

Second tree:

leaves = 6,7,4,9,8

Return true if the leaves are the same





leetcode.com/problems/leaf-similar-trees

Solution

- Get the first leaf value from tree 1
- Get the first leaf value from tree 2
- Compare, if they are different, return false immediately
- Otherwise, continue finding the next leaf value for tree 1 and 2

Implementation

- Create two stacks stack<TreeNode*> left and stack<TreeNode*> right
- Add the

Code – Leaf-Similar Trees

LeetCode

leetcode.com/problems/leaf-similar-trees

Code Time: O(n + m) where n and m are the numbers of nodes for trees 1 and 2 Space: O(h1 + h2) where h1 and h2 represents the height of the tree

```
// returns the value of the leaf, or -1 if empty
int getLeaf(stack<TreeNode*>& tree) {
   // tree is a reference, we will always pop an element from it
   while(!tree.empty()) {
       // get the top element from the stack
       TreeNode* node = tree.top();
       // already visited, so remove from stack
       tree.pop();
       // is this a leaf?
       if (!node->left && !node->right) {
           // yes, return the value
           return node->val;
        // push the right FIRST to the stack
       if (node->right) tree.push(node->right);
        // left should be on top of the stack
       if (node->left) tree.push(node->left);
   return -1;
```

```
bool leafSimilar(TreeNode* root1, TreeNode* root2) {
   // initialize the stacks, add root1 and root2
   std::stack<TreeNode*> leftTree, rightTree;
   leftTree.push(root1);
   rightTree.push(root2);
   while(true) {
       // get the leaves to compare
       int leaf1 = getLeaf(leftTree);
       int leaf2 = getLeaf(rightTree);
       // exit immediately if one leaf is different
       if (leaf1 != leaf2) return false;
       // stop when there are no leaves left
       if (leaf1 == -1 | leaf2 == -1) break;
   return true;
```



leetcode.com/problems/count-good-nodes-in-binary-tree

Problem Statement

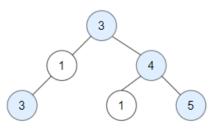
- You are given a binary tree and have to find "good" nodes
- A good node is a node where the values in the path are always than the node
- The root node is always a good node
- Example:
 - root 3 is a good node

left side:

- left leaf 1 is not a good node because 1 < 3
- leaf 3 is a good node because 3 > 1 and 3 == 3

right side:

- leaf 4 is a good node because 4 > 3
- leaf 1 is not a good node because 1 < 4
- leaf 5 is a good node because 5 > 4 > 3



Solution – Count Good Nodes in Binary Tree



E LeetCode

leetcode.com/problems/count-good-nodes-in-binary-tree

Solution

•

Code – Count Good Nodes in Binary Tree

LeetCode

leetcode.com/problems/count-good-nodes-in-binary-tree

Code Time: O(n + m) where n and m are the numbers of nodes for trees 1 and 2 Space: O(h1 + h2) where h1 and h2 represents the height of the tree

```
// returns the value of the leaf, or -1 if empty
int getLeaf(stack<TreeNode*>& tree) {
   // tree is a reference, we will always pop an element from it
   while(!tree.empty()) {
       // get the top element from the stack
       TreeNode* node = tree.top();
       // already visited, so remove from stack
       tree.pop();
       // is this a leaf?
       if (!node->left && !node->right) {
           // yes, return the value
           return node->val;
        // push the right FIRST to the stack
       if (node->right) tree.push(node->right);
        // left should be on top of the stack
        if (node->left) tree.push(node->left);
   return -1;
```

```
bool leafSimilar(TreeNode* root1, TreeNode* root2) {
   // initialize the stacks, add root1 and root2
   std::stack<TreeNode*> leftTree, rightTree;
   leftTree.push(root1);
   rightTree.push(root2);
   while(true) {
       // get the leaves to compare
       int leaf1 = getLeaf(leftTree);
       int leaf2 = getLeaf(rightTree);
       // exit immediately if one leaf is different
       if (leaf1 != leaf2) return false;
       // stop when there are no leaves left
       if (leaf1 == -1 | leaf2 == -1) break;
   return true;
```

INTERVAL

greedy strategy: sort by the end time

Because ending earlier gives **more room** for future intervals. It's a classic greedy trick: choose the interval that **frees up time** as quickly as possible.

Problem - 57. Insert Interval





leetcode.com/problems/insert-interval

Problem Statement

- You are given an array of intervals, where intervals[i] = [start, endi] and newInterval = [start, end]
- newInterval must be inserted into intervals
- Overlapping intervals must be merged
- Example

intervals = [[1,2],[3,5],[6,7],[8,10],[12,16]] newInterval = [4,8]

Output: [[1,2],[3,10],[12,16]]

Solution - 57. Insert Interval





leetcode.com/problems/insert-interval

Solution

- Sort intervals by the first element (start)
- Initialize result
- Solve in three loops:
 - 1. While there is no overlap with **newInterval**, add to **intervals[i]** to **result**
 - 2. While it overlaps, merge **newInterval**
 - 3. While until the end intervals and add the remaining intervals[i]



leetcode.com/problems/insert-interval

Code

Time: **O(n)** Space: **O(n)** where n is the size of intervals

```
vector<vector<int>> insert(vector<vector<int>>& intervals, vector<int>& newInterval) {
    vector<vector<int>> result;
    int tupleIndex = 0;
    int totalTuples = intervals.size();
    // 1. check if it overlaps
    // 1 ----- 2
    while (tupleIndex < totalTuples && intervals[tupleIndex][1] < newInterval[0]) {</pre>
        result.push_back(intervals[tupleIndex]);
        ++tupleIndex;
    // 2. merge overlap. We already know there is an overlap here,
    // otherwise it should be sorted out in the previous step
    // 3 ---- 5
          4 ---- 8
    while (tupleIndex < totalTuples && intervals[tupleIndex][0] <= newInterval[1]) {</pre>
        newInterval[0] = min(newInterval[0], intervals[tupleIndex][0]);
        newInterval[1] = max(newInterval[1], intervals[tupleIndex][1]);
        ++tupleIndex;
    result.push back(newInterval);
    // 3. add remaining parts
    while (tupleIndex < totalTuples) {</pre>
        result.push back(intervals[tupleIndex]);
        ++tupleIndex;
    return result;
```



leetcode.com/problems/merge-intervals

Problem Statement

•



leetcode.com/problems/merge-intervals

Solution

•

leetcode.com/problems/merge-intervals

Code Time: O(n) Space: O(n)

Problem – 435. Non-overlapping Intervals





leetcode.com/problems/non-overlapping-intervals

Problem Statement

• ...

Solution – 435. Non-overlapping Intervals





leetcode.com/problems/non-overlapping-intervals

Solution

•



leetcode.com/problems/non-overlapping-intervals



Time: O(n) Space: O(n)

• ..



leetcode.com/problems/...

Problem Statement / Solution / Code Time: O(n) Space: O(n)

• ..

LINKED LIST

https://leetcode.com/problems/swap-nodes-in-pairs

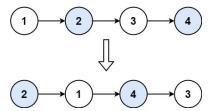
Problem

Given a linked list, swap every two adjacent nodes and return its head. You must solve the problem without modifying the values in the list's nodes (i.e., only nodes themselves may be changed.)

Example 1

Input: head = [1,2,3,4]

Output: [2,1,4,3]



Example 2

Input: head = []

Output: []

Example 3:

Example 3

Input: head = [1]

Output: [1]

Solution – Swap Nodes in Pair

https://leetcode.com/problems/swap-nodes-in-pairs

```
ListNode* swapPairs(ListNode* head) {
   if (head == NULL | head->next == NULL) {
        return head;
    ListNode *node = head;
    ListNode *prev = NULL;
    head = head->next;
    while (node && node->next) {
        ListNode *second = node->next;
        ListNode *next_pair = second->next;
        second->next = node;
       node->next = next_pair;
       if (prev) {
           prev->next = second;
        prev = node;
       node = next_pair;
    return head;
```

Solution (recursive) – Swap Nodes in Pair

https://leetcode.com/problems/swap-nodes-in-pairs

```
ListNode* swapPairs(ListNode* head) {
    if(!head || !head->next)
        return head;
    ListNode* newHead = head->next;
    head->next = swapPairs(head->next->next);
    newHead->next = head;
    return newHead;
}
```

HEAP / PRIORITY QUEUE

Heap

- Heap is a complete binary tree that satisfy the heap property (max or min)
- Min heap: root node contains the minimum value
- Max heap: root node contains the maximum value



Heap in C++

Two main ways to implement:

1. Using std::make_heap from <algorithm>

```
std::make_heap(RandomIt first, RandomIt last)
std::push_heap(RandomIt first, RandomIt last)
std::pop_heap(RandomIt first, RandomIt last)
std::sort_heap(RandomIt first, RandomIt last)
```

2. Using std::priority_queue from <queue> (recommended)

```
std::priority queue<T, Container, Compare>
```

Heap in C++ - std::priority_queue example

Min heap

```
std::priority_queue<int, std::vector<int>, std::greater<int>>
May boan
```

```
Max heap
std::priority_queue<int> or
std::priority queue<int, std::vector<int> std::less<int>>
// Min heap
std::priority queue<int, std::vector<int>, std::greater<int>> minHeap;
minHeap.push(3);
minHeap.push(6);
minHeap.push(4);
// remove top element (3)
minHeap.pop();
// root node (top) is now 4
std::cout << minHeap.top();</pre>
```

https://leetcode.com/problems/kth-largest-element-in-an-array

Problem

Given an integer array nums and an integer k, return the k^{th} largest element in the array. Note that it is the k^{th} largest element in the sorted order, not the k^{th} distinct element.

Example 1

Input: nums = [3,2,1,5,6,4], k = 2
Output: 5

Example 2

Input: nums = [3,2,3,1,2,4,5,5,6], k = 4

Output: 4

Although this problem is classified as "medium", in my opinion it should be classified as "easy"

Solution 1 – Kth Largest Element in an Array

https://leetcode.com/problems/kth-largest-element-in-an-array

```
// SOLUTION 1
int findKthLargest(vector<int>& nums, int k) {
    std::priority_queue<int, std::vector<int>, std::greater<int>> minHeap;
    for (const auto& num : nums) {
        if (minHeap.size() < k) {
            minHeap.push(num);
        } else if (num > minHeap.top()) {
            minHeap.pop();
            minHeap.push(num);
        }
    }
    return minHeap.top();
}
```

Solution 2 – Kth Largest Element in an Array

https://leetcode.com/problems/kth-largest-element-in-an-array

```
// SOLUTION 2 - Simpler approach
int findKthLargest(vector<int>& nums, int k) {
    // min heap: minimum values will be always at the top
    std::priority_queue<int, std::vector<int>, std::greater<int>> minHeap;
    for (const auto& num : nums) {
        // push each num to the heap
        minHeap.push(num);
        // we need the kth largest element only, so once after pushing more than k
        // elements, remove the smallest one (the top)
        if (minHeap.size() > k) {
            minHeap.pop();
        }
    }
    return minHeap.top();
}
```

DYNAMIC PROGRAMMING

Dynamic Programming

Dynamic Programming (DP) is an algorithm technique used to solve problems that can be broken down into **simpler, overlapping subproblems.**

Key Concepts of Dynamic Programming

- Overlapping subproblems: a problem has overlapping subproblems if it can be broken down into subproblems.
- **Memoization (Top-Down Approach)**: store the results in a cache (typically a dictionary or array) to avoid recalculation recursion and caching approach.
- **Tabulation (Bottom-Up Approach)**: first solve all possible subproblems iteratively, and store them in a table.

Dynamic Programming – Example – Fibonacci Sequence

```
Naive Recursive Approach

int fib(int n) {
   if (n <= 1) {
      return n;
   }
   return fib(n - 1) + fib(n - 2);
}</pre>
```

```
Memoization (Top-Down DP)

std::unordered_map<int, int> memo;

int fib(int n) {
    if (n <= 1) {
        return n;
    }
    if (memo.find(n) != memo.end()) {
        return memo[n];
    }
    memo[n] = fib(n - 1) + fib(n - 2);
    return memo[n];
}</pre>
```

```
Tabulation (Bottom-up DP)

int fib(int n) {
    if (n <= 1) {
        return n;
    }
    int dp[n + 1];
    dp[0] = 0;
    dp[1] = 1;
    for (int i = 2; i <= n; i++) {
        dp[i] = dp[i - 1] + dp[i - 2];
    }
    return dp[n];
}</pre>
```

Problem – Climbing Stairs



leetcode.com/problems/climbing-stairs

Problem Statement

You need to climb a staircase with n steps to get to the top. Each time you can choose to climb either 1 step or 2 steps at a time. Find out how many different ways you can climb to the top of the staircase.

Example 1

Input: n = 2

Output: 2

Explanation: There are two ways to get to the top

- 1. Climb 1 step at a time, twice
- 2. Climb 2 steps in one go

Example 2:

Input: n = 3

Output: 3

Explanation: There are three ways to get to the top:

- 1. Climb 1 step at a time, three times
- 2. Climb 1 step, then 2 steps
- 3. Climb 2 steps, then 1 ste.



leetcode.com/problems/climbing-stairs

```
std::unordered map<int, int> memo;
int climbStairs(int n) {
   // Identify the sequence, when:
   // n = 0 (0 way), there is no way to get up
   // n = 1 (1 way): only one way : 1-step
   // n = 2 (2 ways): 1s + 1s | 2s
   // n = 3 (3 ways): 1s + 1s + 1s | 1s + 2s | 2s + 1s
   // n = 4 (5 ways): 1s + 1s + 1s + 1s | 1s + 1s + 2s | 1s + 2s + 1s | 2s + 1s + 1s | 2s + 2s |
   if (n <= 2) {
       return n;
   if (memo.find(n) != memo.end()) {
       return memo[n];
   memo[n] = climbStairs(n - 1) + climbStairs(n - 2);
   return memo[n];
```

EOF



leetcode.com/problems/...

Problem Statement / Solution / Code Time: O(n) Space: O(n)

•



leetcode.com/problems/...

Problem Statement / Solution / Code Time: O(n) Space: O(n)

• ..