ARRAY

Arrays

- Memory layout: hold values in a contiguous block of memory.
- **Fixed Size**: the size of an array is defined when it is created and cannot be changed. However, high-level languages have different implementations, making it dynamic.
- Homogeneous elements: all elements are of the same data type (int, float, char...)
- **Efficiency**: accessing elements by index is very efficient *O(1)*, since each index maps directly to a memory location. Also, range scans benefit from CPU cache lines since arrays are stored in contiguous blocks of memory.



Problem

- Given an array of numbers and a target, example: array [2,7,11,15] and target 9
- Return indices of two numbers where they add up to target
- **Output**: [0,1]

```
array[0] + array[1] = 2 + 7 = 9
```



leetcode.com/problems/two-sum

Solution

- Iterative over each number in the array
- Calculate the difference between target and each number, example:

```
array[0] = 2, target 9, then 9 - 2 = 7
```

- Now we know we need the number 7 to sum up to 9
- Check in a hashmap if we have 7 in some part of the array

```
hash[7] exists?
```

- If yes, return the current index and the index of 7
- If not, store the index of the current number in the hashmap for future evaluation

$$hash[2] = 0$$

Code - 1. Two Sum



LeetCode leetcode.com/problems/two-sum

```
Code Time: O(n) Space: O(n)
vector<int> twoSum(vector<int>& nums, int target) {
    std::unordered map<int, int> numMap;
   // n being the size of nums
   for (int i = 0; i < nums.size(); i++) {</pre>
        // current number of the array
        int number = nums[i];
        int diff = target - number;
        // check if the difference is in some part of the array
        // by using a hashmap
        if (numMap.find(diff) != numMap.end()) {
            return { numMap[diff], i};
        // register the current number index
        numMap[number] = i;
   // no matches
   return {};
```

Problem – 217. Contains Duplicate



leetcode.com/problems/contains-duplicate

Problem

- You are given an array of numbers
- Return any value that appears at least twice

Solution

- Loop through the array
- Check if the value is in a hash table
- Return **true** if the value exist
- The problem requires at least twice, but one modification may be having a specific count

Code - 217. Contains Duplicate



LeetCode leetcode.com/problems/contains-duplicate

```
bool containsDuplicate(vector<int>& nums) {
   unordered_map<int, int> seen;
   for (int i = 0; i < nums.size(); ++i) {
      if (seen[nums[i]] == 1) {
        return true;
      }
      seen[nums[i]]++;
   }
   return false;</pre>
```

Another solution (less flexible)

```
bool containsDuplicate(vector<int>& nums) {
    unordered_map<int, bool> seen;
    for (const auto& num : nums) {
        if (seen[num]) {
            return true;
        }
        seen[num] = true;
    }
    return false;
}
```





leetcode.com/problems/product-of-array-except-self

Problem Statement

- You are given an integer array nums
- Return another array where each element is multiplied by all the elements except itself
- Example:

```
nums = [14,2,5,99]
nums[0] = 2 * 5 * 99 (all except 14)
nums[1] = 14 * 5 * 99 (all except 2)
nums[2] = 14 * 2 * 99
nums[3] = 14 * 2 * 5
```



leetcode.com/problems/product-of-array-except-self

Solution

• Go over the array once and calculate the product of the left side. Example:

```
nums = [14,2,5]
left[0] = 1 (think of multiplying all elements before 14, so 1 because there is none)
left[1] = 14 (all elements from the left multiplied, except 2)
left[2] = 14 * 2 = 28 (all elements from the left multiplied, except 5)
left = [1, 14, 28]
```

Using the same logic, do the same calculation but starting from the right

```
right[2] = 1 (no elements after 5)
right[1] = 5 (only 5 after 2)
right[0] = 2 * 5 = 10
right = [10, 5, 1]
```

• Multiply each element from left and right:

```
left \odot right = [1, 14, 28] \odot [10,5,1] = [10, 70, 28]
```

Code - 238. Product of Array Except Self

C LeetCode

leetcode.com/problems/product-of-array-except-self

```
Code Time: O(n) Space: O(n)
vector<int> productExceptSelf(vector<int>& nums) {
    int n = nums.size();
    vector<int> output(n, 1);
    vector<int> right(n, 1);
    // calculate left first
    for (int i = 1; i < n; ++i) {
        output[i] = nums[i - 1] * output[i - 1];
    // calculate right
    for (int i = n - 1; i >= 0; --i) {
        right[i] = nums[i + 1] * right[i + 1];
        output[i] = output[i] * right[i];
    /* or you can save some space using this logic,
      although I don't find it as intuitive as the previous one
    int right = 1;
    for (int i = n - 1; i >= 0; --i) {
        output[i] *= right;
        right *= nums[i];
    return output;
```



leetcode.com/problems/maximum-subarray

Problem

- You are given an array nums
- Find the subarray with the largest sum

• Example:

nums =
$$[-2,1,-3,4,-1,2,1,-5,4]$$

output = 6

The subarray [4,-1,2,1] has the largest sum 6.





leetcode.com/problems/maximum-subarray

Solution

- Use Kadane's algorithm to find the maximum sum of a contiguous subarray in linear time
- Core idea:

at each index, either:

- 1. start a new subarray at **nums[i]** or
- 2. extend the current one by adding nums[i]

Arrays – Kadane's algorithm

- Kadane's algorithm is a dynamic programming algorithm to solve maximum subarray sum
- At every index i: start a new subarray at i extend the previous subarray to include array[i]

Algorithm

1. Initialize:

```
int maxSoFar = array[0];
int maxEndingHere = array[0];
```

2. Loop through the array

```
for (int i = 1; i < array.size(); ++i) {
    maxEndingHere = max(array[i], maxEndingHere + array[i]);
    maxSoFar = max(maxSoFar, maxEndingHere);
}</pre>
```

3. Return maxSoFar;

Problem - 53. Maximum Subarray

```
LeetCode
```

leetcode.com/problems/maximum-subarray

```
int maxSubArray(vector<int>& nums) {
  int maxSum = nums[0];
  int currentSum = nums[0];
  for (int i = 1; i < nums.size(); ++i) {
     currentSum = max(nums[i], currentSum + nums[i]);
     maxSum = max(maxSum, currentSum);
  }
  return maxSum;
}</pre>
```



leetcode.com/problems/maximum-product-subarray

Problem

- You are given an array nums
- Find a subarray that has the largest product and return the product
- The array may contain negative numbers

• Example:

```
nums = [2, 3, -2, 4]

output = 6

[2,3] has the largest 6 (2 * 3)
```

Problem - 152. Maximum Product Subarray





leetcode.com/problems/maximum-product-subarray

Solution

- Use a modified version of Kadane's algorithm
- Keep track of the minimum and maximum product
- Once the current number is negative, swap minimum product with maximum product
- Check the largest product between maximum product and the final result

Problem - 152. Maximum Product Subarray

```
C LeetCode
```

leetcode.com/problems/maximum-product-subarray

```
int maxProduct(vector<int>& nums) {
  int result = nums[0];
  int maxProd = nums[0];
  int minProd = nums[0];
  for (int i = 1; i < nums.size(); ++i) {
    if (nums[i] < 0) {
       swap(minProd, maxProd);
    }

  minProd = min(nums[i], nums[i] * minProd); // -2
    maxProd = max(nums[i], nums[i] * maxProd); // -30

  result = max(result, maxProd);
}
return result;</pre>
```



leetcode.com/problems/find-minimum-in-rotated-sorted-array

Problem

- You are given a sorted array but "rotated"
- Rotated means the elements are displaced in order
- Return the minimum element

• Example:

```
nums = [3,4,5,1,2]
output = 1(minimum element)
```



leetcode.com/problems/find-minimum-in-rotated-sorted-array

Solution

- Perform an adapted binary search
- Example:

```
[3,4,5,1,2]
left = 3, mid = 5, right =2
```

You find mid (5), but have to go right, so adjust left:

```
if (mid > right)
   left = mid + 1
else
   right = mid
```

Code - 153. Find Minimum in Rotated Sorted Array

C LeetCode

leetcode.com/problems/find-minimum-in-rotated-sorted-array

```
Code Time: O(log n) Space: O(1)

int findMin(vector<int>& nums) {
   int left = 0;
   int right = nums.size() - 1;
   while (left < right) {
      int mid = left + (right - left) / 2;
      if (nums[mid] > nums[right]) {
         left = mid + 1;
      } else {
         right = mid;
      }
   }
   return nums[left];
```

Problem - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Problem Statement

- You are given an integer array of stock prices
- Choose a price[i] to buy and price[i] to sell where you achieve maximum profits
- Example:

```
prices = [9, 1, 3, 4]
```

• **Output**: [1,3]

```
array[3] - array[1] = 4 + 1 = 3
```

Solution - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Solution

- Initialize profit = 0
- Initialize lowestBuyPrice = prices[0]
- Loop through the prices
- Track the lowest buy price → min(lowestBuyPrice, prices[i])
- Check if selling "today" will make the maximum profit and update profit:
 max(prices[i] buy > profit, profit)
- Update profit max(prices[i] - buy

Code - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Code (simplified) Time: O(n) Space: O(1)

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      buy = min(buy, prices[i]);
      profit = max(profit, prices[i] - buy)
   }
   return profit;
}</pre>
```

Code - Best Time to Buy and Sell Stock



leetcode.com/problems/best-time-to-buy-and-sell-stock

Code (optimized) Time: O(n) Space: O(1)

Same logic, but with better branch prediction and less computation

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      if (prices[i] < buy) {
        buy = prices[i];
      } else if (prices[i] - buy > profit) {
           profit = prices[i] - buy;
      }
   }
   return profit;
}
```





leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

Problem

- You are given an integer array of stock prices
- Choose a price[i] to buy and price[i] to sell where you achieve maximum profits
- You can buy/sell multiple times, but only hold at most one transaction at a time
- Output is the maximum profits

• Example:

```
prices = [9, 1, 3, 4]

Output: 2 + 1 = 3

buy (price = 1), sell (price = 3), profit = 2

buy (price = 3), sell (price = 4), profit = 1
```



leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

Solution

- Loop through the array starting from index 1
- If current price[i] is lower than previous price[i 1], buy and sell

• Example:

```
prices = [1, 8, 4] prices[0] = 1, prices[1] = 8, prices[2] = 4 prices[0] < \text{prices}[1] \rightarrow true, profit = 8 - 1 = 7 prices[2] < \text{prices}[1] \rightarrow \text{false}, do nothing
```

Code - Best Time to Buy and Sell Stock II

LeetCode

leetcode.com/problems/best-time-to-buy-and-sell-stock-ii

```
int maxProfit(vector<int>& prices) {
  int profit = 0;
  for (int i = 1; i < prices.size(); ++i) {
    if (prices[i] > prices[i-1]) {
       profit += prices[i] - prices[i - 1];
    }
  }
  return profit;
}
```



leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Problem Statement

•

Solution - Best Time to Buy and Sell Stock IV





leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Solution

• ..

Code - Best Time to Buy and Sell Stock IV



leetcode.com/problems/best-time-to-buy-and-sell-stock-iv

Code (simplified) Time: O(n) Space: O(n)

```
int maxProfit(vector<int>& prices) {
   int profit = 0;
   int buy = prices[0];
   for (auto i = 1; i < prices.size(); i++) {
      buy = min(buy, prices[i]);
      profit = max(profit, prices[i] - buy)
   }
   return profit;
}</pre>
```



leetcode.com/problems/search-in-rotated-sorted-array

Problem

- Variation of Find Minimum in Sorted Rotated Array problem
- You are given a sorted array but "rotated" and a target number n
- Rotated means the elements are displaced in order
- Search the number **n** and return its index

Example

nums =
$$[4,5,6,7,0,1,2]$$
, target = 0

Output: 4

4 is the index where the target number **0** is located





leetcode.com/problems/search-in-rotated-sorted-array

Solution

- Perform a binary search with some modification
- One side is always sorted, so find which side (left or right)
- Check if the target is in the range of the sorted side and adjust mid Example:

$$[2,4,5,6,7,0,1]$$
 target = 0

- 1. Find mid (6)
- 2. Find the sorted side (left) = [2,4,5]
- 3. Check if your target is in this side. Is **target** between 2 and 5?
- 4. Adjust mid to search at the other side if not, otherwise continue searching at the same side

Problem – 33. Search in Rotated Sorted Array

```
LeetCode
```

leetcode.com/problems/search-in-rotated-sorted-array

Code Time: O(log n) Space: O(1)

```
int search(vector<int>& nums, int target) {
    int left = 0;
    int right = nums.size() - 1;
    while (left <= right) {</pre>
        int mid = left + (right - left) / 2;
        if (nums[mid] == target) return mid;
        // figure it out the sorted side
        if (nums[mid] >= nums[0]) {
            // left side is sorted
            // is target within this range?
            if (target >= nums[0] && target < nums[mid]) {</pre>
                right = mid - 1;
            } else {
                left = mid + 1;
        } else {
            // right side is sorted
            // is target within this range?
            if (target <= nums[right] && target > nums[mid]) {
                left = mid + 1;
            } else {
                right = mid - 1;
    return -1;
```



leetcode.com/problems/3sum

Problem

- You are given an array of integer nums
- Find distinct triples that the final sum is equal to zero

• Example:

nums =
$$[-1,0,1,2,-1,-4]$$

Output:

$$[[-1,-1,2],[-1,0,1]]$$

Explanation:

$$nums[0] + nums[1] + nums[2] = (-1) + 0 + 1 = 0.$$

$$nums[1] + nums[2] + nums[4] = 0 + 1 + (-1) = 0.$$

$$nums[0] + nums[3] + nums[4] = (-1) + 2 + (-1) = 0.$$

The distinct triplets are [-1,0,1] and [-1,-1,2].

Solution - 15. 3Sum





leetcode.com/problems/3sum

Solution

- Use three pointers: i, j and k
- Sort the array. This is necessary to move the pointers j and k
- Pointer i starts at the beginning the array
- Pointer j starts at i + 1 (second position)
- Pointer k starts at the end of the array
- Pointer i always move forward until the end of the array
- For each value of **i,** j and k will move either forward or backward, depending on the results of the sum
- Once find a sum == 0, add to a set to guarantee no duplicates



leetcode.com/problems/3sum

Code Time: O(n² log n) Space: O(n²)

```
vector<vector<int>> threeSum(vector<int>& nums) {
    set<vector<int>> triplets;
    sort(nums.begin(), nums.end());
    for (int i = 0; i < nums.size() - 2; ++i) {
        int j = i + 1;
        int k = nums.size() - 1;
       // it is a solution
       while (j < k) {
            int sum = nums[i] + nums[j] + nums[k];
            if (sum == 0) {
               triplets.insert({nums[i], nums[j], nums[k]});
                j++;
                k--;
            if (sum < 0) {
                j++;
            } else {
                k--;
    vector<vector<int>> result;
    // convert the solutions to the expected return
    for (const auto& t : triplets) {
        result.push_back(t);
    return result;
```

Note: why not a unordered_set which is hash-based?



leetcode.com/problems/container-with-most-water

Problem Statement

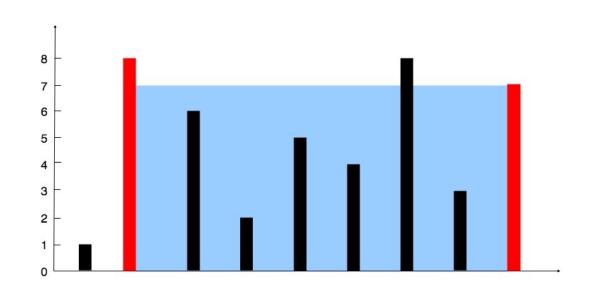
- You are given an integer array height
- Find two lines that together with x-axis form a container with most water
- Example:

Input:

height = [1,8,6,2,5,4,8,3,7]

Output:

49





leetcode.com/problems/container-with-most-water

Solution

- Initialize the maximum area maxArea = 0
- Initialize two pointers, left = 0 and right = height.size 1
- Loop while pointer left < right
- Calculate the area:

area = min(height[left], height[right]) * (right - left)

Update the global maximum area:

maxArea = max(maxArea, area)

- Move the smallest pointer (increment left or decrement right)
- Return maxArea

LeetCode

leetcode.com/problems/container-with-most-water

```
Code Time: O(n) Space: O(1)
int maxArea(vector<int>& height) {
    // left and right = positions
   int left = 0;
    int right = height.size() - 1;
    int maxArea = 0;
    while (left < right) {</pre>
        int area = min(height[left], height[right]) * (right - left);
        maxArea = max(maxArea, area);
        // adjust left and right based on 'greedy' algorithm
        // move from the lowest height
        if (height[left] < height[right]) {</pre>
            left++;
        } else {
            right--;
    return maxArea;
```