

# Problem – Climbing Stairs

Easy



LeetCode

[leetcode.com/problems/climbing-stairs](https://leetcode.com/problems/climbing-stairs)

## Problem

- You need to climb a staircase with **n steps**
- Each time, you can only climb either **1 step** or **2 steps**
- Find out how **many different ways** you can **climb to the top**
- **Example:**

## Input

$n = 2$

**Output:** 2

## Explanation:

1. Climb 1 step, then climb 1 step again
2. Climb 2 steps in one go

# Solution – Climbing Stairs

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## Solution

- **You need to climb a staircase with  $n$  steps** and want to find the total number of distinct ways to reach the top
- **At each step, you can either take 1 step or 2 steps** - this gives you two choices at most positions
- **This follows the Fibonacci sequence pattern** - the number of ways to reach step  $n$  equals ways to reach  $(n-1)$  plus ways to reach  $(n-2)$
- **Use dynamic programming to avoid recalculating subproblems** - either bottom-up tabulation or memoized recursion works well
- **Base cases are crucial**: typically  $f(1) = 1$  and  $f(2) = 2$ , representing the ways to reach the first and second steps

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## Code

```
std::unordered_map<int, int> memo;
```

```
int climbStairs(int n) {  
    // Identify the sequence, when:  
    // n = 0 (0 way), there is no way to get up  
    // n = 1 (1 way): only one way : 1-step  
    // n = 2 (2 ways): 1s + 1s | 2s  
    // n = 3 (3 ways): 1s + 1s + 1s | 1s + 2s | 2s + 1s  
    // n = 4 (5 ways): 1s + 1s + 1s + 1s | 1s + 1s + 2s | 1s + 2s + 1s | 2s + 1s + 1s | 2s + 2s |  
  
    if (n <= 2) {  
        return n;  
    }  
  
    if (memo.find(n) != memo.end()) {  
        return memo[n];  
    }  
  
    memo[n] = climbStairs(n - 1) + climbStairs(n - 2);  
    return memo[n];  
}
```

# Problem – 1143. Longest Common Subsequence

Medium



LeetCode

<https://leetcode.com/problems/longest-common-subsequence>

## Problem

- You are given two strings, example:  
text1 = "abcd"  
text2 = "ace"
- Find the longest subsequence between them

# Problem – 62. Unique Paths

Medium



LeetCode

<https://leetcode.com/problems/unique-paths>

## Problem

- The robot is placed in a  $m \times n$  grid
- It starts at the top-left cell  $(0,0)$  and must reach the bottom-right  $(m - 1, n - 1)$
- The robot can only move right or down at any point
- Return the number of unique paths the robot can take to reach the destination



# Solution – 62. Unique Paths

Medium



LeetCode

<https://leetcode.com/problems/unique-paths>

## Solution 1 (recursive)

- Define a recursive function `countPaths(m, n)`

- **Base case**

If  $m == 1$  or  $n == 1$ , there's only one way to reach that cell (either all downs or all rights).

- **Recursive case**

To reach cell  $(m, n)$  the robot must come from:

Cell  $(m - 1, n) \rightarrow$  from above

Cell  $(m, n - 1) \rightarrow$  from left

So the number of of paths to  $(m, n)$  is the **sum** of the paths to those two cells

- **Memoization**

Use a 2D vector  $[m + 1][n + 1]$

# Code – 62. Unique Paths

Medium



LeetCode

<https://leetcode.com/problems/unique-paths>

**Code** Time:  $O(m * n)$  Space:  $O(m * n)$

```
int countPaths(int m, int n, vector<vector<int>>& memo) {
    if (m == 1 || n == 1) return 1;
    if (memo[m][n] != -1) return memo[m][n];
    memo[m][n] = countPaths(m, n - 1, memo) + countPaths(m - 1, n, memo);
    return memo[m][n];
}
int uniquePaths(int m, int n) {
    /*
        count(m, n) = count(m, n + 1) + count(m + 1, n)
        same as (imagine robot going from m,n to 0,0 up and left)
        count(m, n) = count(m, n - 1) + count(m - 1, n)
    */
    vector<vector<int>> memo(m + 1, vector<int>(n + 1, -1));
    return countPaths(m, n, memo);
}
```

# Solution – 62. Unique Paths

Medium



LeetCode

<https://leetcode.com/problems/unique-paths>

## Solution 2 (iterative)

- Create a 2D vector `dp` of size  $(m+1) \times (n+1)$  to store intermediate results
  - Set `dp[1][1] = 1` because there is exactly one way to stand on the starting cell
  - Iterate through each cell `(row, col)` from `(1, 1)` to `(m, n)`:
    - Skip `(1, 1)` since it's already initialized
    - For every other cell, the number of unique paths to it is the sum of:
      - Paths from the cell above: `dp[row-1][col]`
      - Paths from the cell to the left: `dp[row][col-1]`
- `dp[row][col] = dp[row - 1][col] + dp[row][col - 1]`**