

# **SHORTEST PATH**

# Shortest Path Algorithms

Algorithm	Use Case	Graph Type	Time Complexity	Notes
BFS	When all edges have <b>equal weight</b>	Unweighted / same-weighted edges	$O(V + E)$	Simplest and fastest when all weights are equal.
Dijkstra	When edge weights are <b>non-negative</b>	Weighted, no negative weights	$O((V + E) \log V)$ with heap	Greedy, efficient for SSSP (Single Source Shortest Path).
Bellman-Ford	When edge weights can be <b>negative</b>	Weighted, allows negative weights	$O(V * E)$	Slower, but handles negative weights and detects negative cycles.
Floyd-Warshall	For <b>all-pairs shortest paths</b>	Dense graphs, small number of nodes	$O(V^3)$	Easy to implement; handles negative weights (but not negative cycles).
A* Search	For shortest path with a <b>goal node</b> and <b>heuristic</b>	Weighted, heuristic needed	Depends on heuristic quality	Often used in pathfinding (e.g. maps, games); faster than Dijkstra if heuristic is good.
Johnson's Algorithm	<b>All-pairs shortest paths</b> in sparse graphs with <b>negative weights</b>	Weighted, allows negative weights	$O(V^2 \log V + V * E)$	Reweights graph with Bellman-Ford, then runs Dijkstra from each node.
SPFA	Practical variant of Bellman-Ford, often faster	Weighted, allows negative weights	Avg: $O(E)$ , Worst: $O(VE)$	Queue-based; faster in practice, not guaranteed. Handles negative weights.
Bidirectional Search	When you know <b>start and target</b> , speeds up search in undirected graphs	Unweighted or uniformly weighted	$O(b^{(d/2)})$ in best case	Runs two simultaneous searches (from start and end); fast when goal is known.

# Problem – 743. Network Delay Time

Medium



LeetCode

[leetcode.com/problems/network-delay-time](https://leetcode.com/problems/network-delay-time)

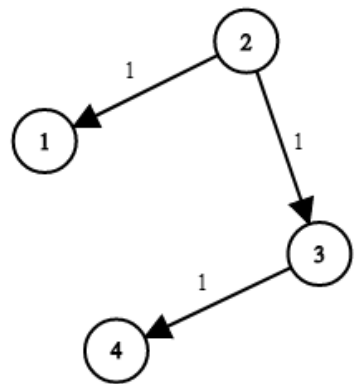
## Problem

- You are given a network of nodes  $n$  with destination and time to reach that node
- You are given a starting node  $k$  and the number of nodes in the network  $n$
- A signal is sent from node  $k$  to all nodes in the network
- Find the minimum time required for all nodes to receive the signal from  $k$

- **Example:**

**$k = 2$     $n = 4$**

**Output: 3**



# Solution – 743. Network Delay Time

Medium



LeetCode

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## Solution

- This is solved using Dijkstra algorithm
- Build the graph by storing in the destination node and the time from a source node:  
`graph[node] = [[node, distance]]`  
`graph[2] = [[1,1], [2,3]]`
- Set up a min-heap for Dijkstra (priority\_queue) with distance and node
- Perform Dijkstra algorithm and store the shortest paths
- Check if all nodes were reached
- Return the longest distance among shortest paths

# Code – 743. Network Delay Time

Medium



LeetCode

leetcode.com/problems/network-delay-time

**Code** Time:  $O((n + e) \log n)$  Space:  $O(n + e)$  where  $n$  is the number of nodes and  $e$  the number of edges

```
int networkDelayTime(vector<vector<int>>& times, int n, int k) {
    // node => (destination, distance)
    unordered_map<int, vector<pair<int, int>>> graph;
    for (const auto& time : times) {
        // time[0] = source node, time[1] = dest node, time[2] = time
        graph[time[0]].emplace_back(time[1], time[2]);
    }

    // min heap: distance from the origin 'k' to 'node'
    priority_queue<pair<int, int>, vector<pair<int, int>>, greater<>> minHeap;
    minHeap.emplace(0, k); // distance, starting node

    // shortest path from each node to the origin 'k' (node, distance)
    unordered_map<int, int> dist;

    // we start exploring the nodes from the minimum
    // distance to the origin 'k'
    while (!minHeap.empty()) {
        auto [distance, node] = minHeap.top();
        minHeap.pop();
        // already visited, skip
        if (dist.count(node)) continue;
        // set the distance
        dist[node] = distance;
        // look at the connections
        for (const auto& [n, d] : graph[node]) {
            // n[node, distance]
            // quick optimization, not necessary
            if (dist.count(n)) continue;
            // add the distance since we want
            // the distance from the origin
            minHeap.emplace(distance + d, n);
        }
    }

    // check if all nodes were visited
    if (dist.size() != n) return -1;

    // all the minimum distances are calculated
    // find the max one since we want to reach
    // all nodes
    int minTime = 0;
    for (const auto& [n, d] : dist) {
        minTime = max(minTime, d);
    }

    return minTime;
}
```