



Counterfactual Regret Minimization and Domination in Extensive-Form Games

Richard Gibson

University of Alberta

Edmonton, Alberta, Canada

Overview

Counterfactual
Regret Minimization
(CFR)

Overview

Counterfactual
Regret Minimization
(CFR)

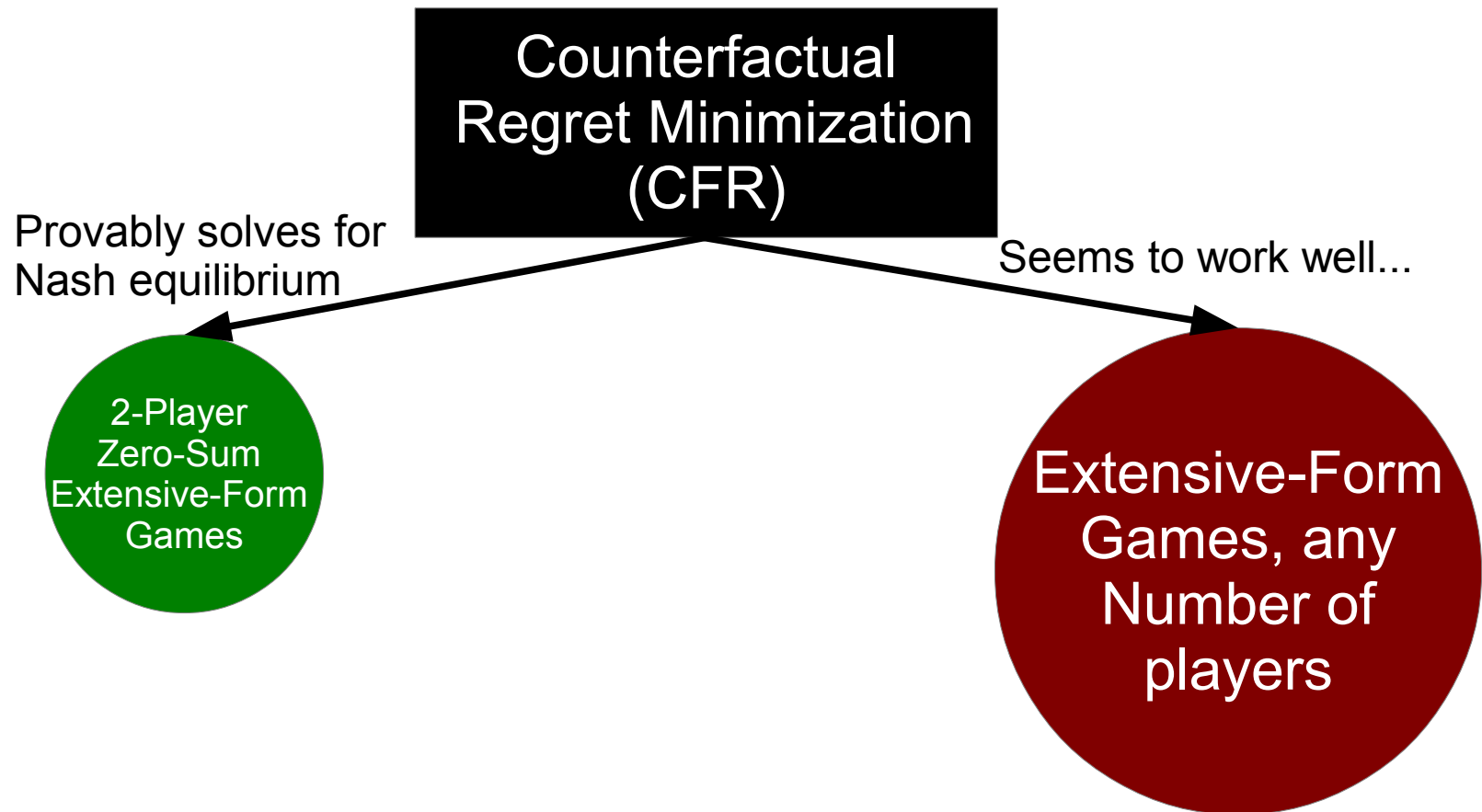
Provably solves for
Nash equilibrium

2-Player
Zero-Sum
Extensive-Form
Games

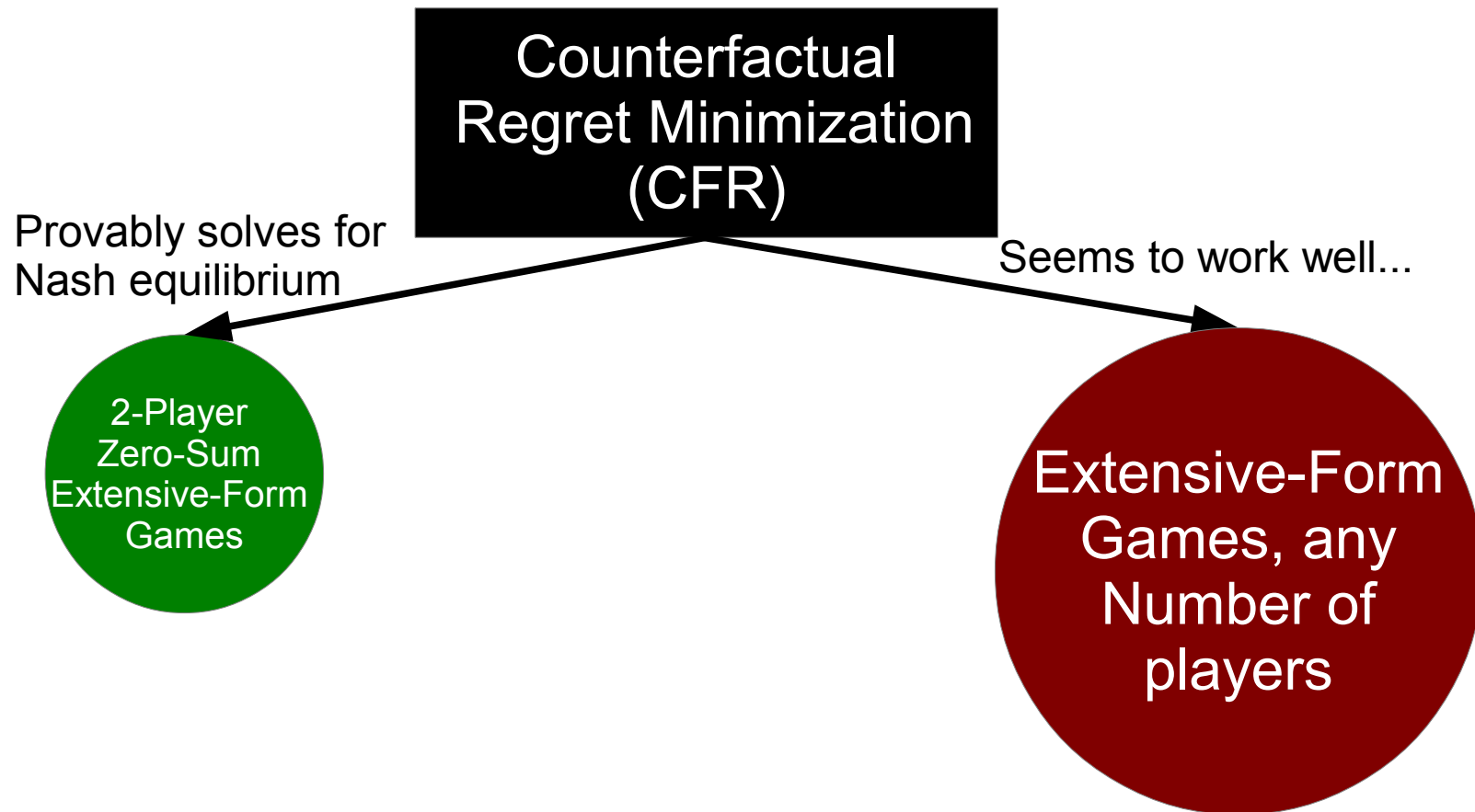
```
graph TD; CFR[Counterfactual Regret Minimization (CFR)] -- "Provably solves for Nash equilibrium" --> G((2-Player Zero-Sum Extensive-Form Games));
```

The diagram illustrates the application of Counterfactual Regret Minimization (CFR) to finding Nash equilibria. A black rectangular box at the top right contains the text 'Counterfactual Regret Minimization (CFR)'. A black arrow points from the bottom-left corner of this box to a green circular node at the bottom left. The green node contains the text '2-Player Zero-Sum Extensive-Form Games'. To the left of the arrow, the text 'Provably solves for Nash equilibrium' is written, indicating the theoretical guarantee of the CFR algorithm for this class of games.

Overview

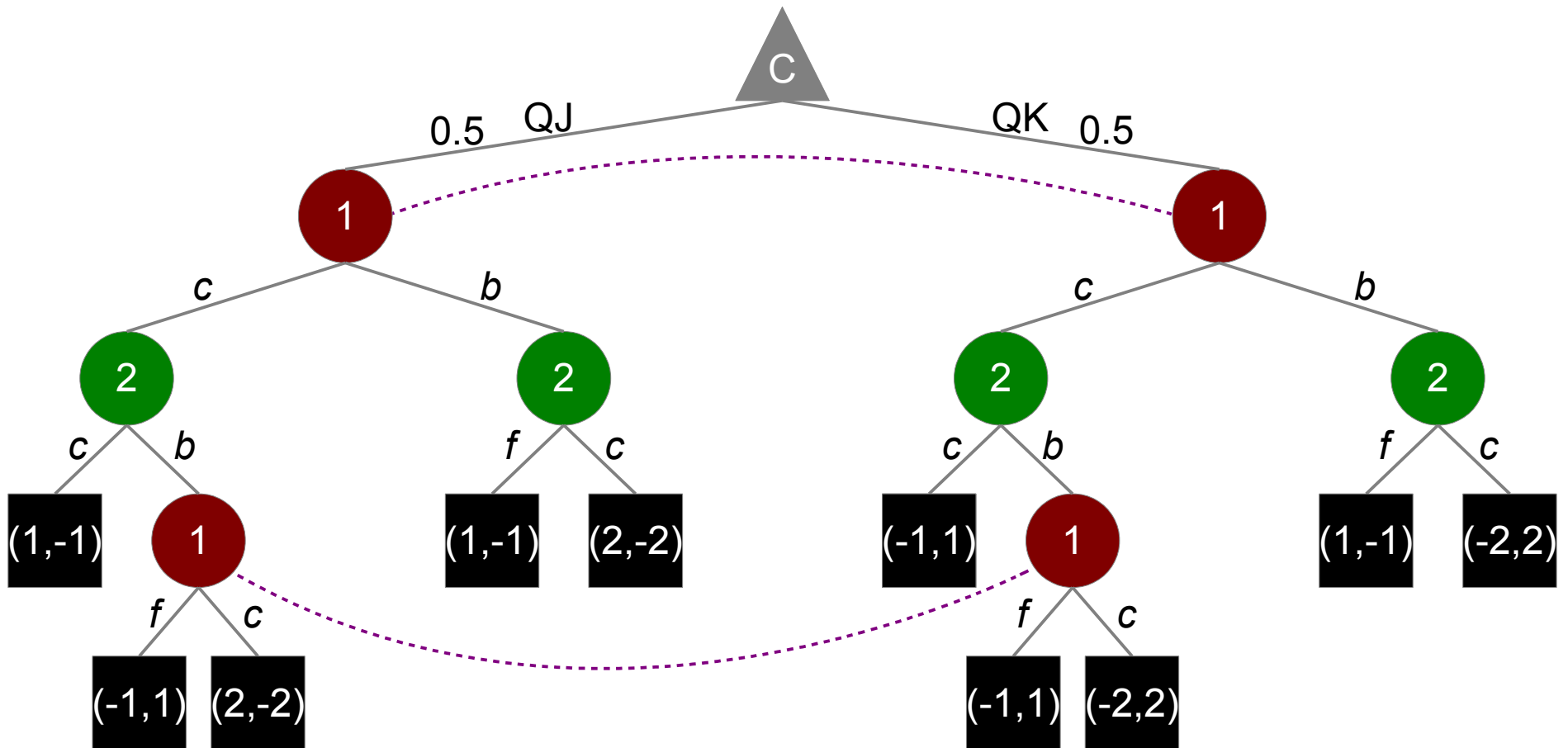


Overview

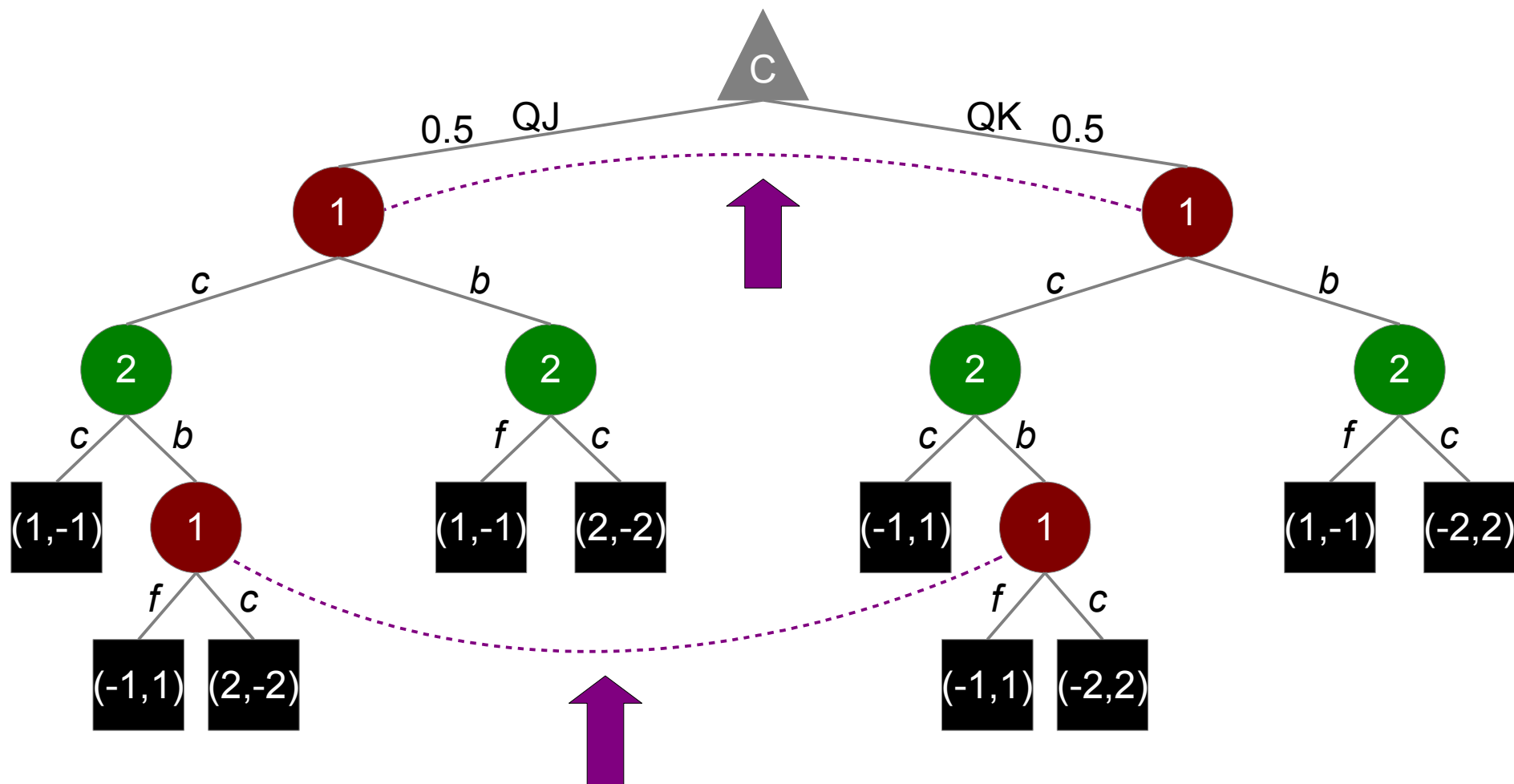


Question: Why do CFR strategies work well in extensive-form games outside of the 2-player zero-sum case?

Extensive-Form Games

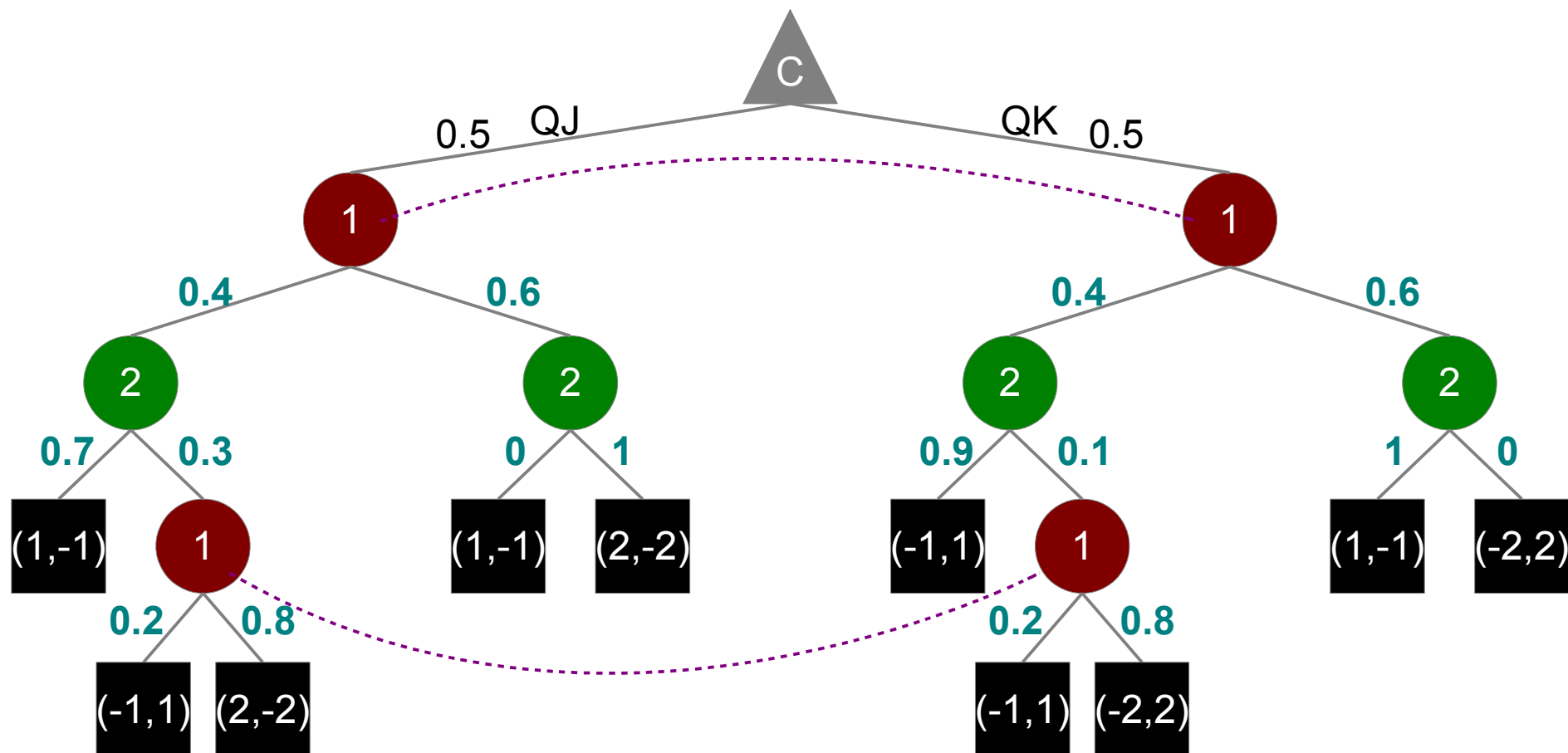


Extensive-Form Games



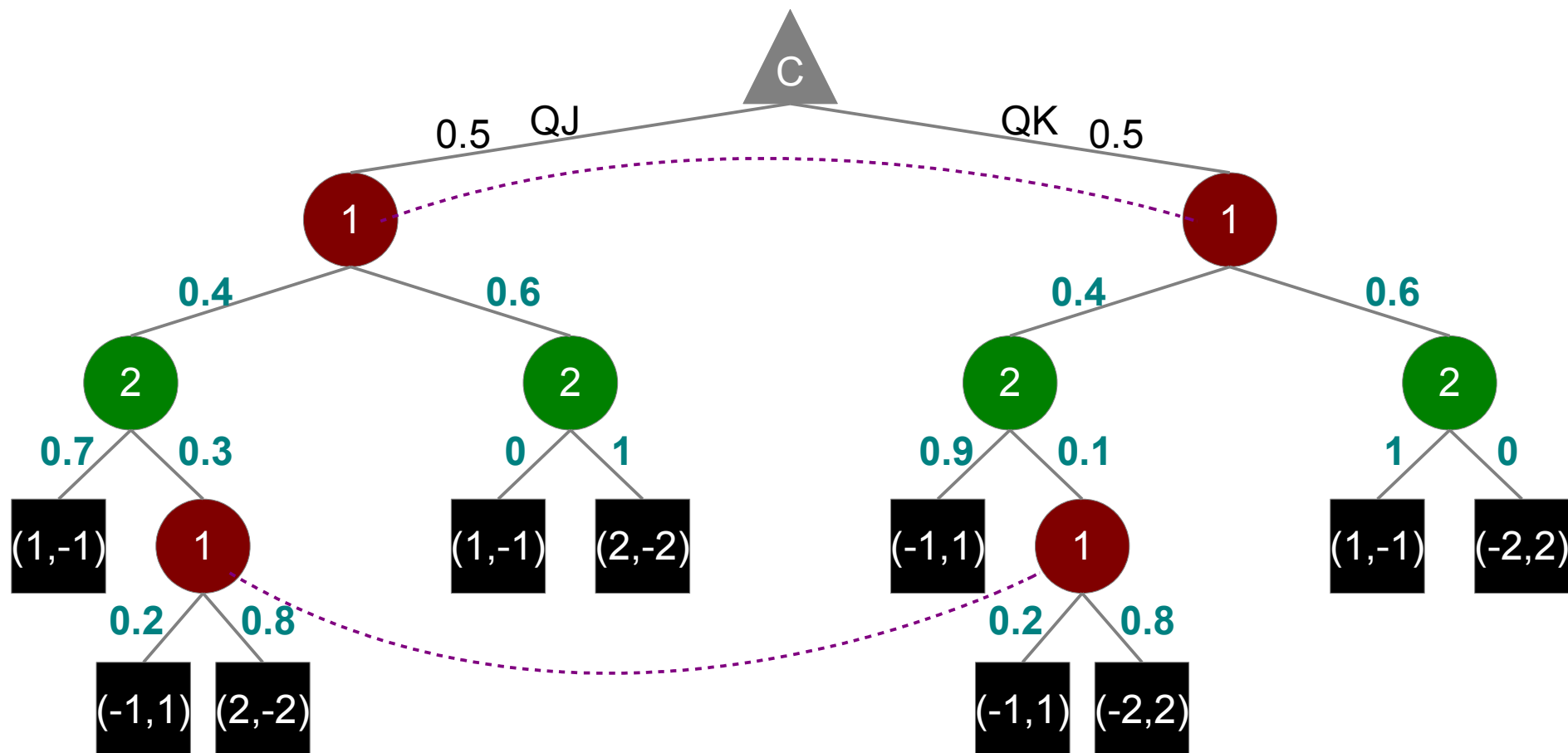
Information sets group states that are indistinguishable to the player.

Extensive-Form Games



A **strategy profile** $\sigma = (\sigma_1, \sigma_2)$ assigns a probability distribution over actions at each information set. Example: Probability player 1 checks is $\sigma_1(Q?, c) = 0.4$.

Extensive-Form Games



$u_i(\sigma)$ is the expected utility for player i , assuming players play according to σ .

Counterfactual Regret Minimization (CFR)

[Zinkevich *et al.*, NIPS 2007]

- CFR is an iterative algorithm that generates strategy profiles $(\sigma^1, \sigma^2, \dots, \sigma^T)$ over many iterations T .
- Final output of CFR: $\sigma^{\text{AVG}} = \text{Average}(\sigma^1, \sigma^2, \dots, \sigma^T)$.
- For 2-player zero-sum games, σ^{AVG} is an ϵ -Nash equilibrium, with $\epsilon \rightarrow 0$ as $T \rightarrow \infty$:

$$u_1(\sigma_1^{\text{AVG}}, \sigma_2^{\text{AVG}}) \geq \max_{\sigma_1^*} u_1(\sigma_1^*, \sigma_2^{\text{AVG}}) - \epsilon$$

$$u_2(\sigma_1^{\text{AVG}}, \sigma_2^{\text{AVG}}) \geq \max_{\sigma_2^*} u_2(\sigma_1^{\text{AVG}}, \sigma_2^*) - \epsilon$$

Counterfactual Regret Minimization (CFR)

- Outside of 2-player zero-sum games, σ^{AVG} is not necessarily an approximate Nash equilibrium [Abou Risk and Szafron, AAMAS 2010].
 - A player may gain by deviating from σ^{AVG} .
- In these games, a Nash equilibrium might not be the most appropriate solution concept anyways.
- On the other hand, σ^{AVG} performs very well in practice...

Annual Computer Poker Competition

3-Player Limit Hold'em - 2009

Agent	Instant Run-off: Round 0
Hyperborean-Eqm	319 \pm 2
Hyperborean-BR	299 \pm 2
akuma	151 \pm 2
dpp	171 \pm 2
CMURingLimit	-37 \pm 2
dcu3pl	-63 \pm 2
Bluechip	-548 \pm 2

3-Player Limit Hold'em - 2010

Agent	Instant Run-off: Round 0
Hyperborean.iro	144 \pm 32
dcu3pl.tbr	98 \pm 30
LittleRock	65 \pm 35
Arnold3	-135 \pm 39
Bender	-172 \pm 16

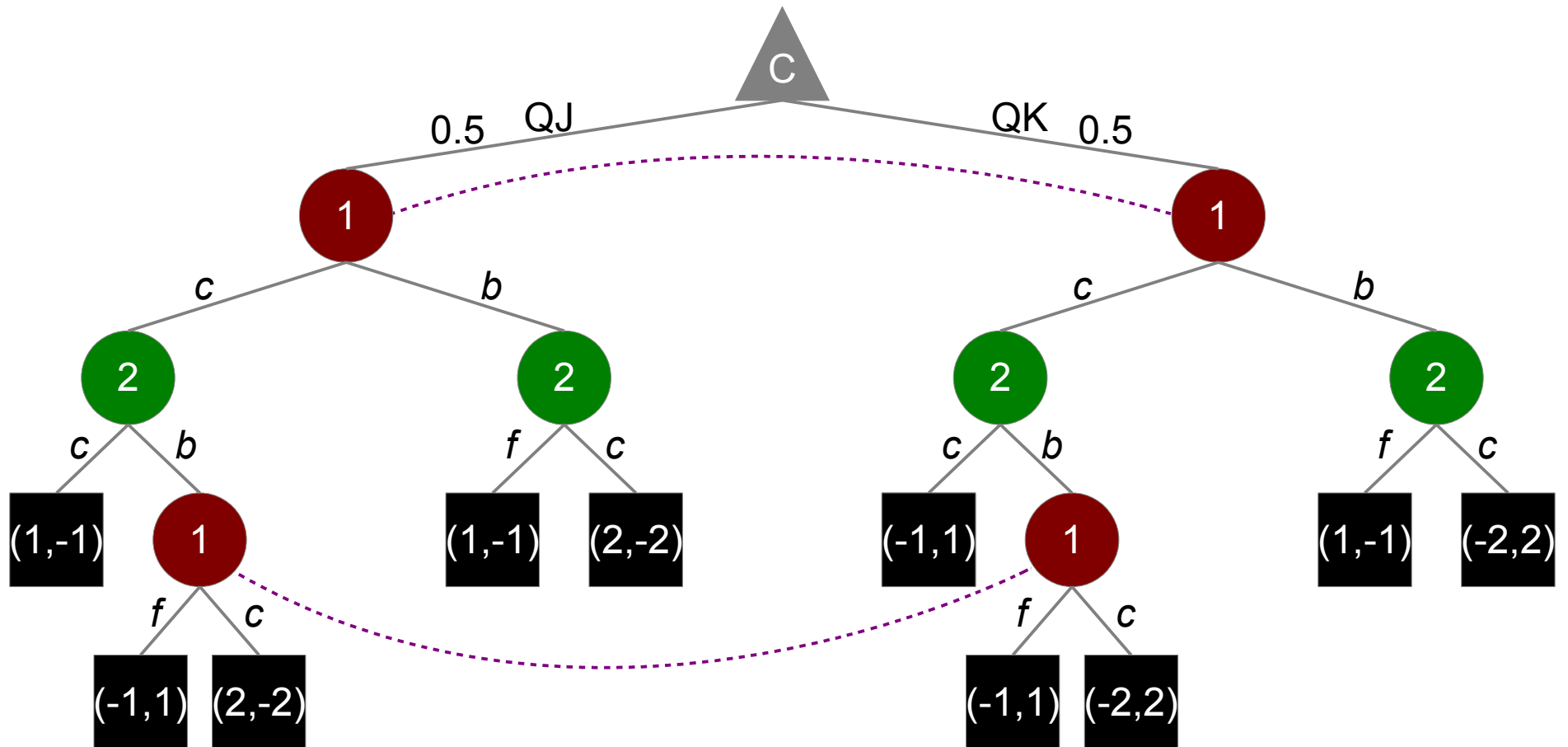
3-Player Limit Hold'em - 2011

Agent	Instant Run-off: Round 0
Sartre3p	243 \pm 20
Hyperborean-3p-limit-iro	204 \pm 20
LittleRock	113 \pm 19
AAIMontybot	96 \pm 44
dcubot3plr	77 \pm 19
OwnBot	-4 \pm 30
Bnold3	-91 \pm 22
Entropy	-108 \pm 36
player.zeta.3p	-530 \pm 33

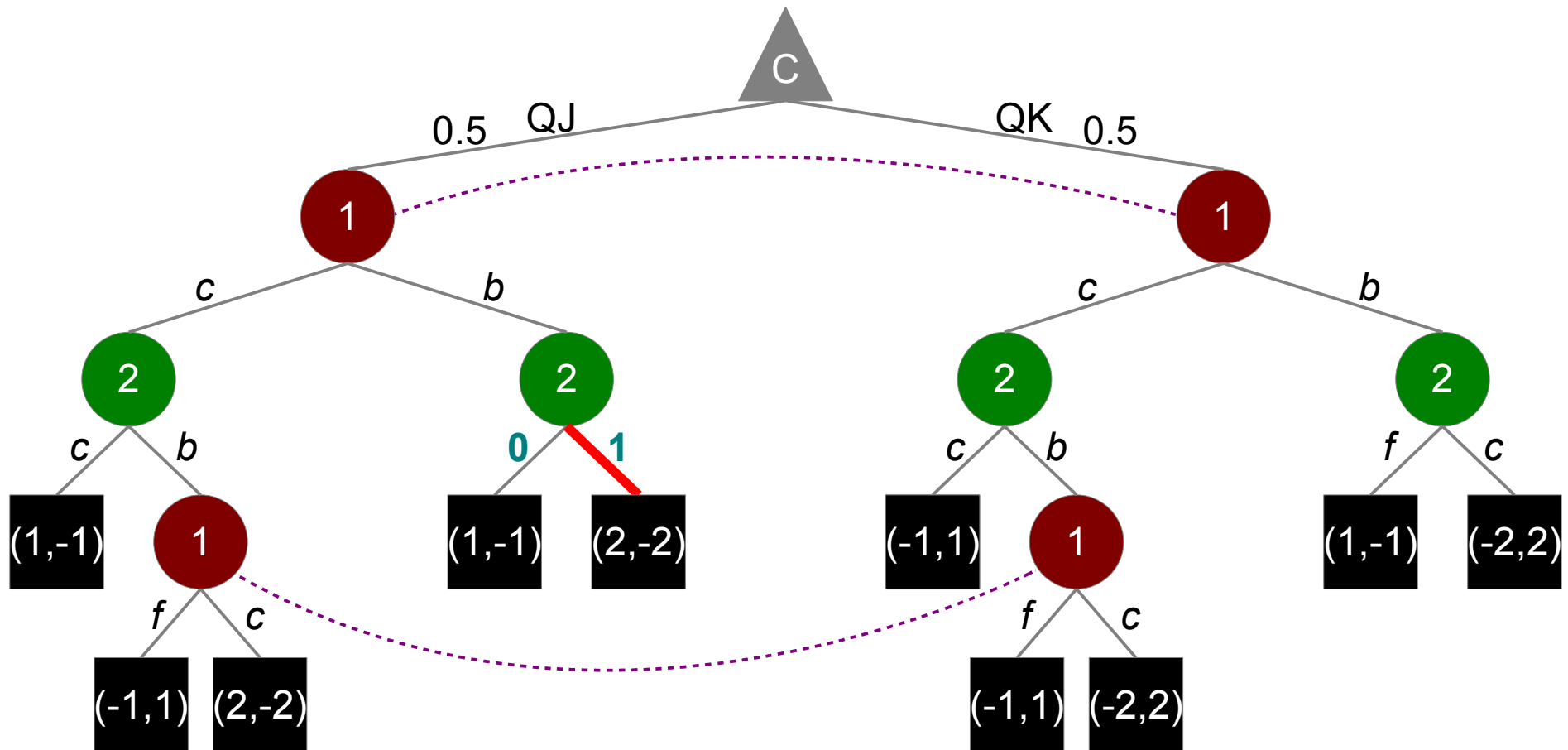
Counterfactual Regret Minimization (CFR)

- In games with more than 2-players, σ^{AVG} is a “good” strategy. **Why?**
- What properties make a strategy good in games with more than 2-players?
- We know what a bad strategy is...

Domination

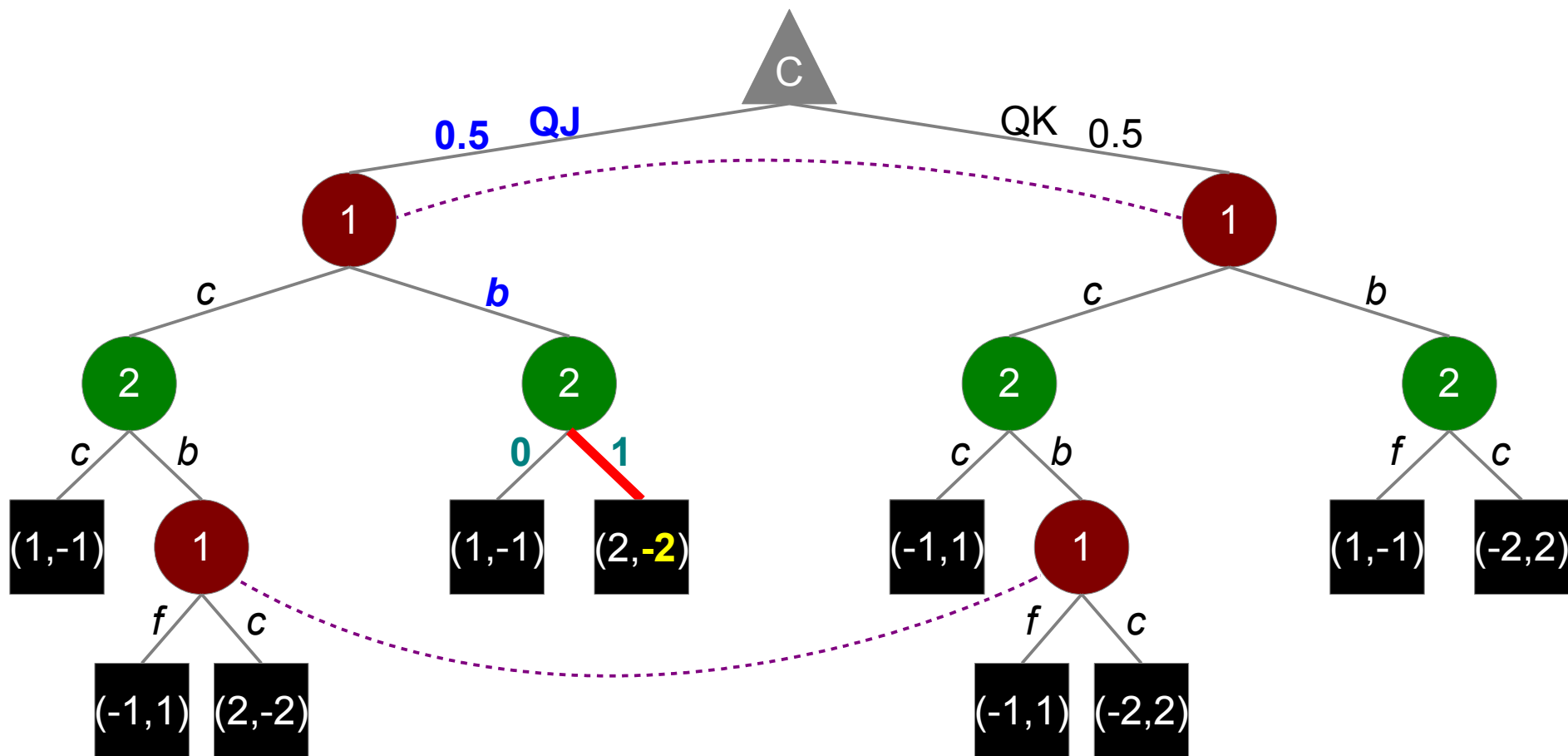


Domination



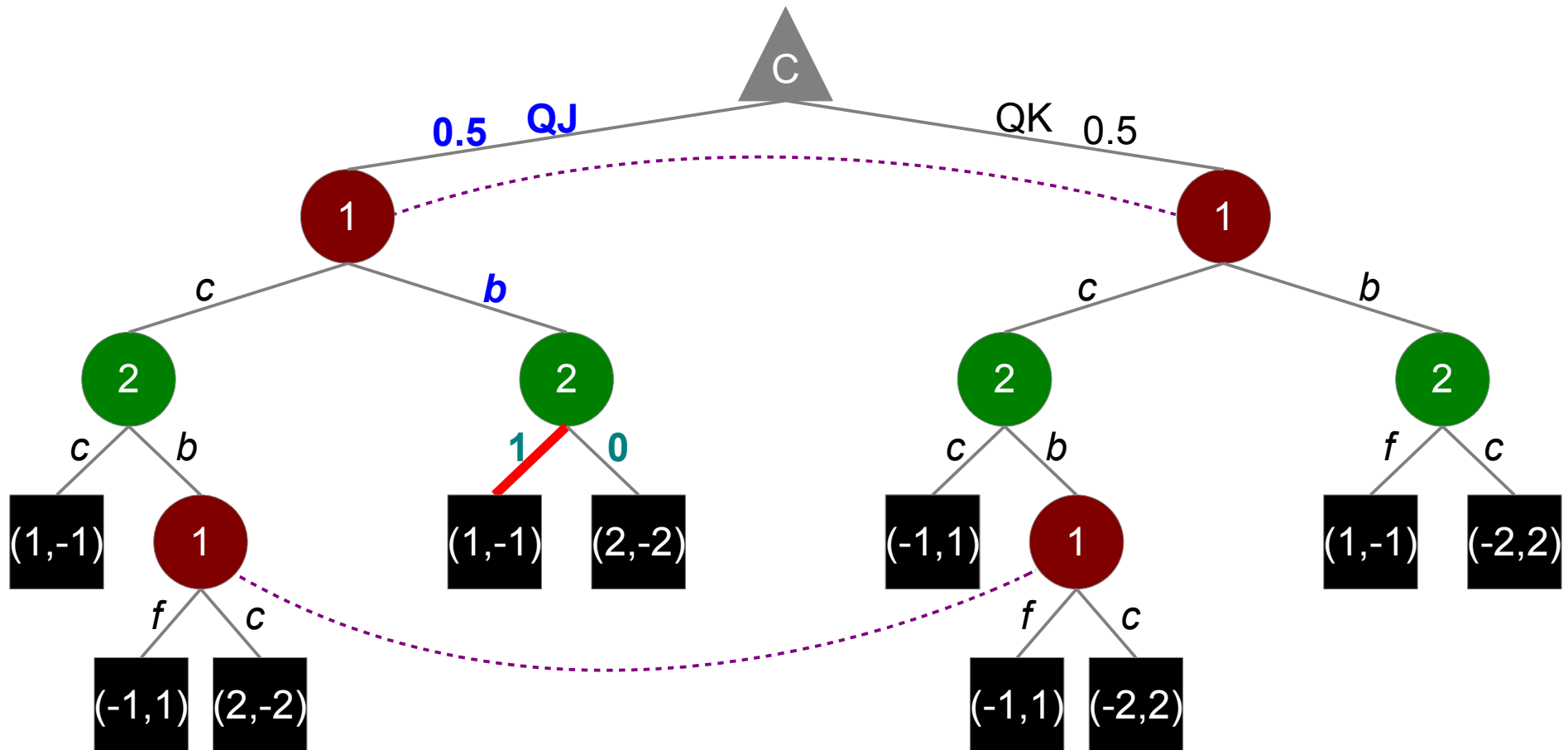
Consider any player 2 strategy $\sigma_2^{J,c}$ that always calls with the Jack when faced with a bet.

Domination



$$u_2(\sigma_1, \sigma_2^{J,c}) = \dots + 0.5\sigma_1(Q?, b)1(-2) + \dots$$

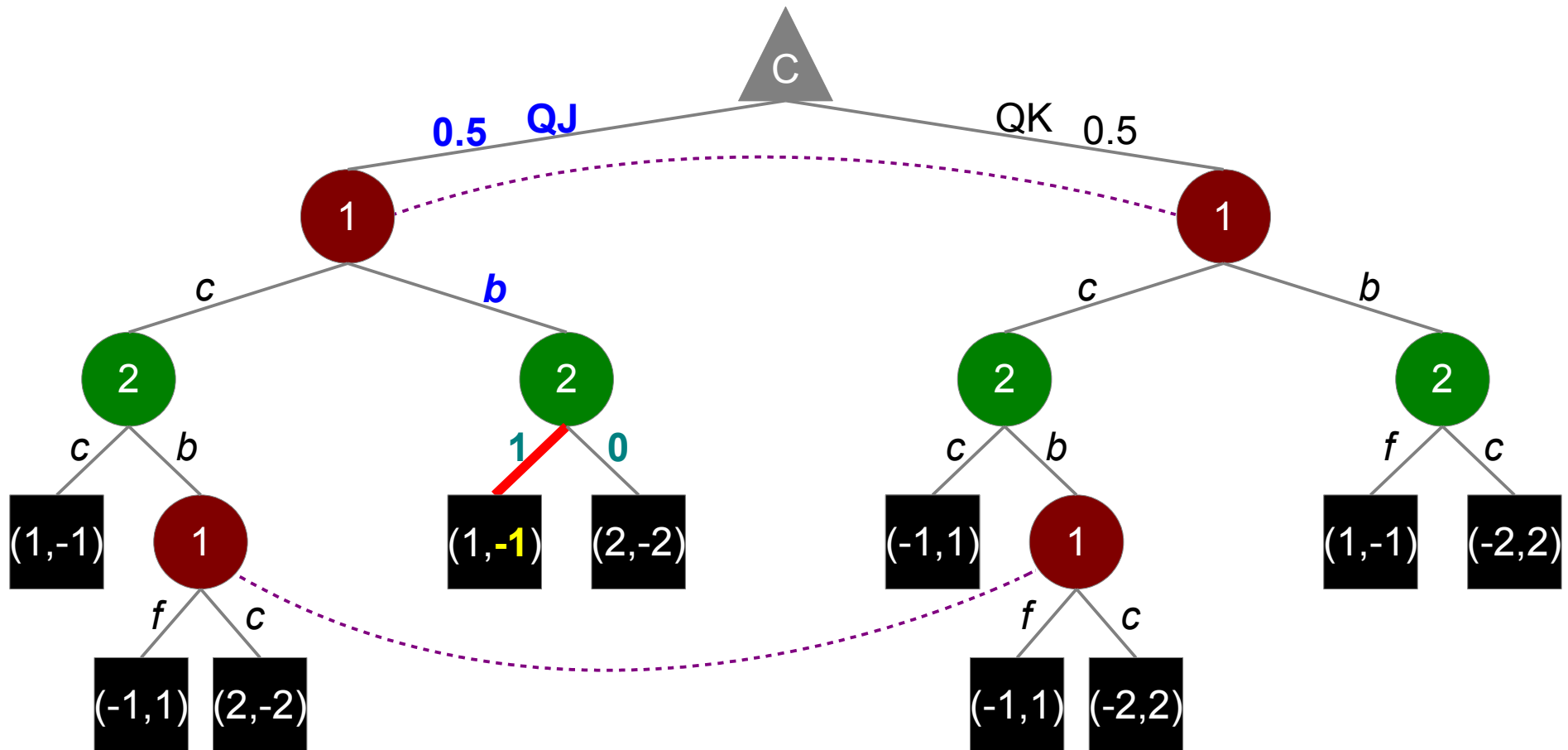
Domination



$$u_2(\sigma_1, \sigma_2^{J,c}) = \dots + 0.5\sigma_1(Q?, b)1(-2) + \dots$$

Now consider the same player 2 strategy, except always folds the J. Call it $\sigma_2^{J,f}$.

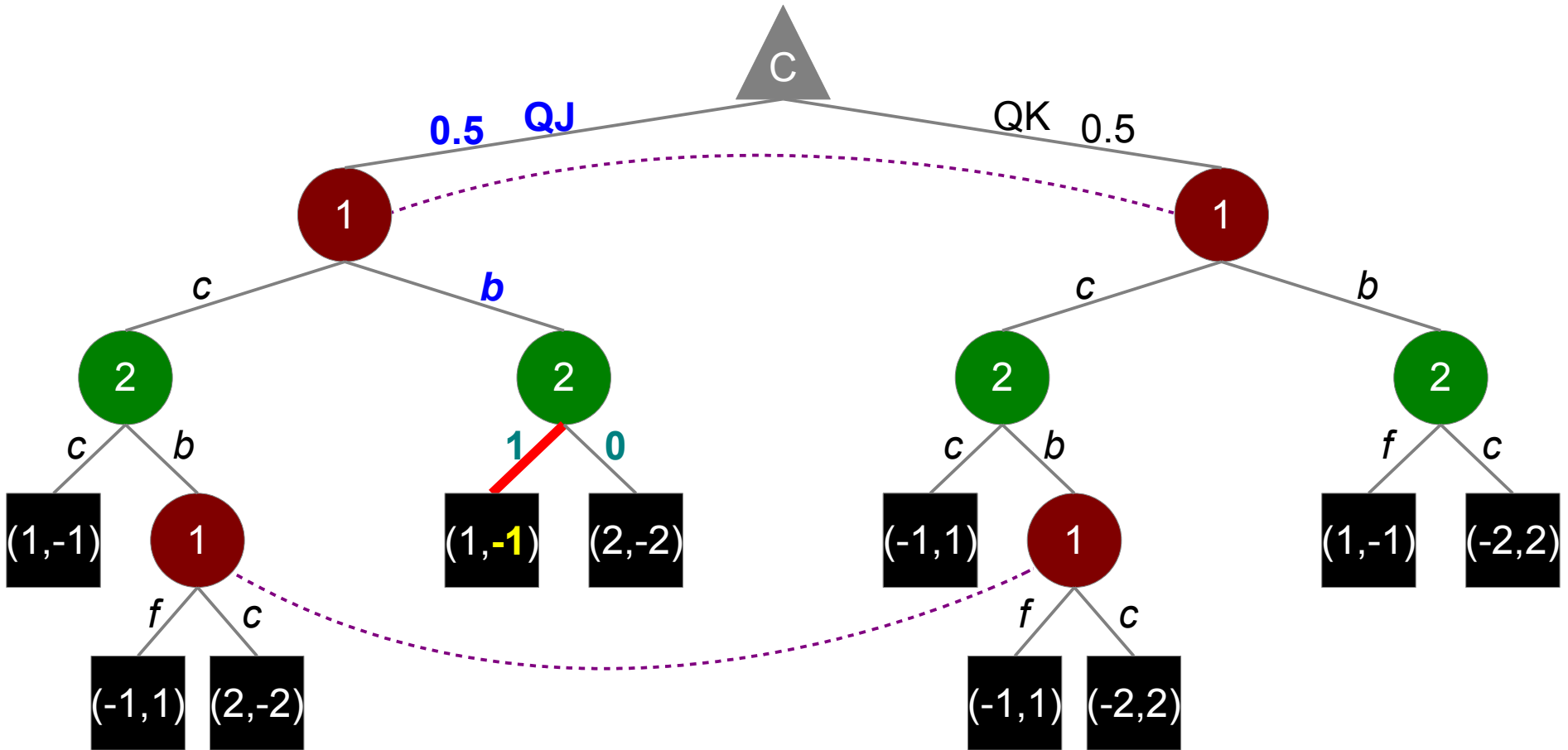
Domination



$$u_2(\sigma_1, \sigma_2^{J,c}) = \dots + 0.5\sigma_1(Q?, b)1(-2) + \dots$$

$$u_2(\sigma_1, \sigma_2^{J,f}) = \dots + 0.5\sigma_1(Q?, b)1(-1) + \dots$$

Domination



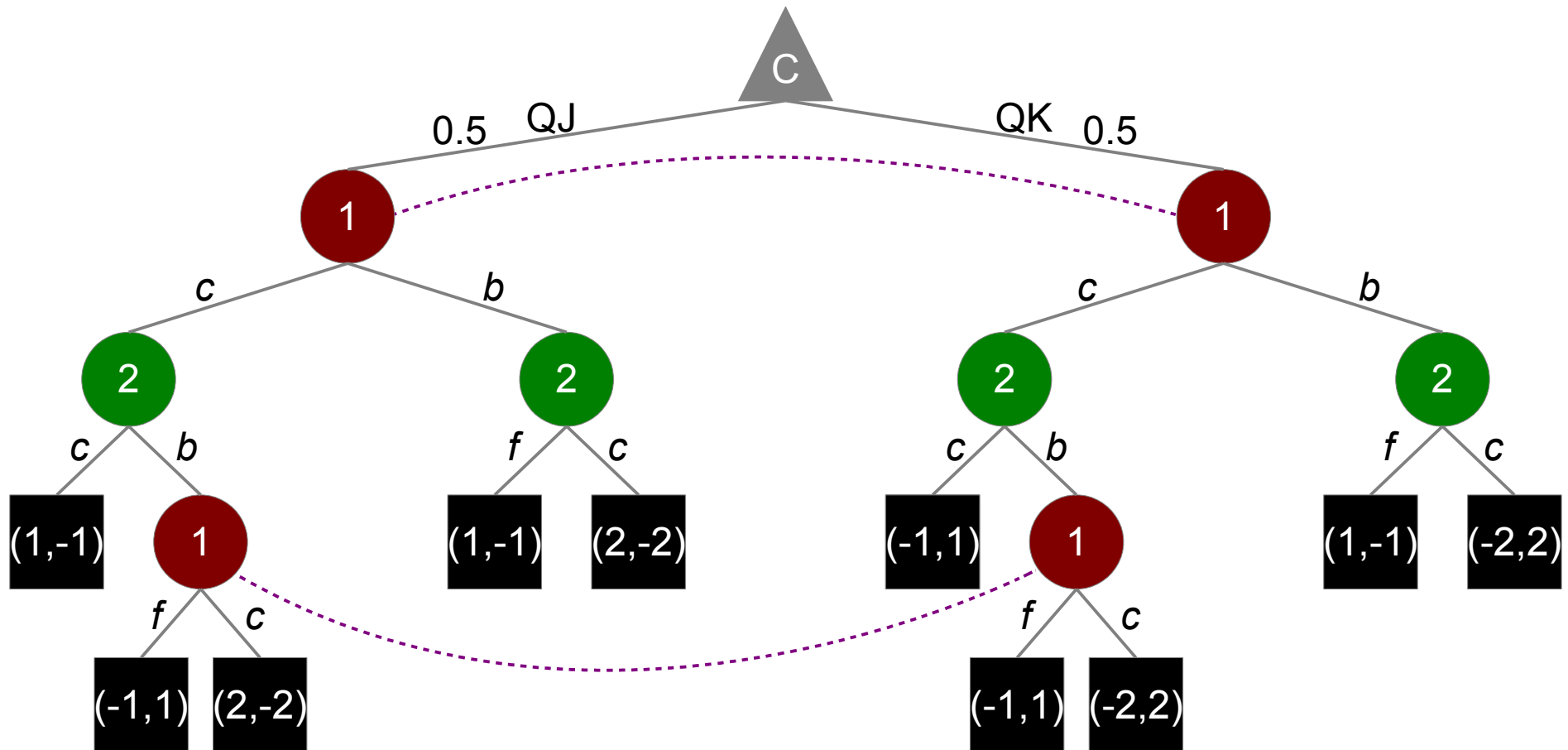
$$u_2(\sigma_1, \sigma_2^{J,c}) = \dots + 0.5\sigma_1(Q?, b)1(-2) + \dots$$

$$u_2(\sigma_1, \sigma_2^{J,f}) = " + 0.5\sigma_1(Q?, b)1(-1) + "$$

$$u_2(\sigma_1, \sigma_2^{J,c}) \leq u_2(\sigma_1, \sigma_2^{J,f}) \text{ for all } \sigma_1.$$

$$u_2(\sigma_1, \sigma_2^{J,c}) < u_2(\sigma_1, \sigma_2^{J,f}) \text{ if } \sigma_1(Q?, b) > 0$$

Domination

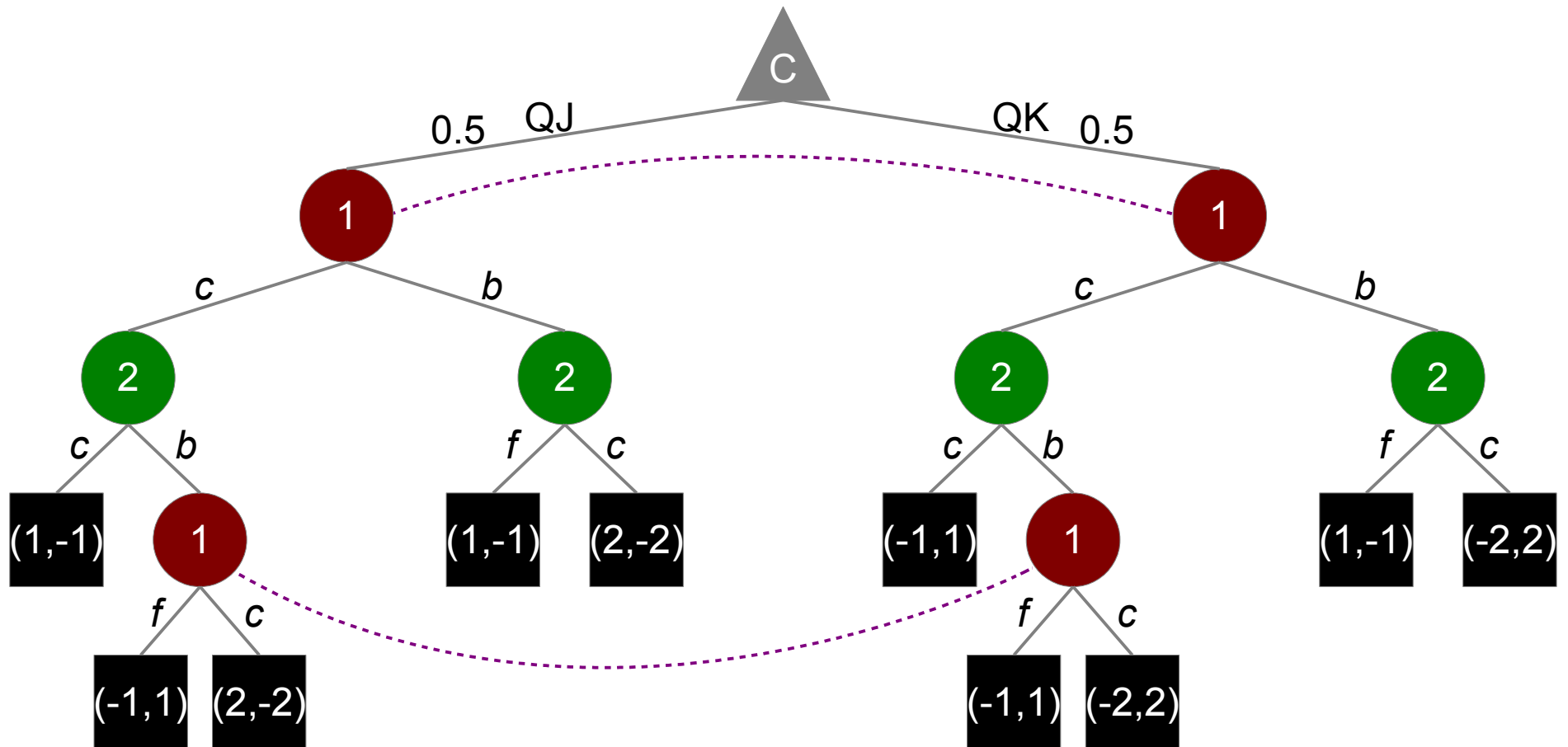


σ_2 is a **dominated strategy** if there exists σ_2' such that

$$u_2(\sigma_1, \sigma_2, \sigma_3, \dots) \leq u_2(\sigma_1, \sigma_2', \sigma_3, \dots) \text{ for all } \sigma_1, \sigma_3, \dots$$

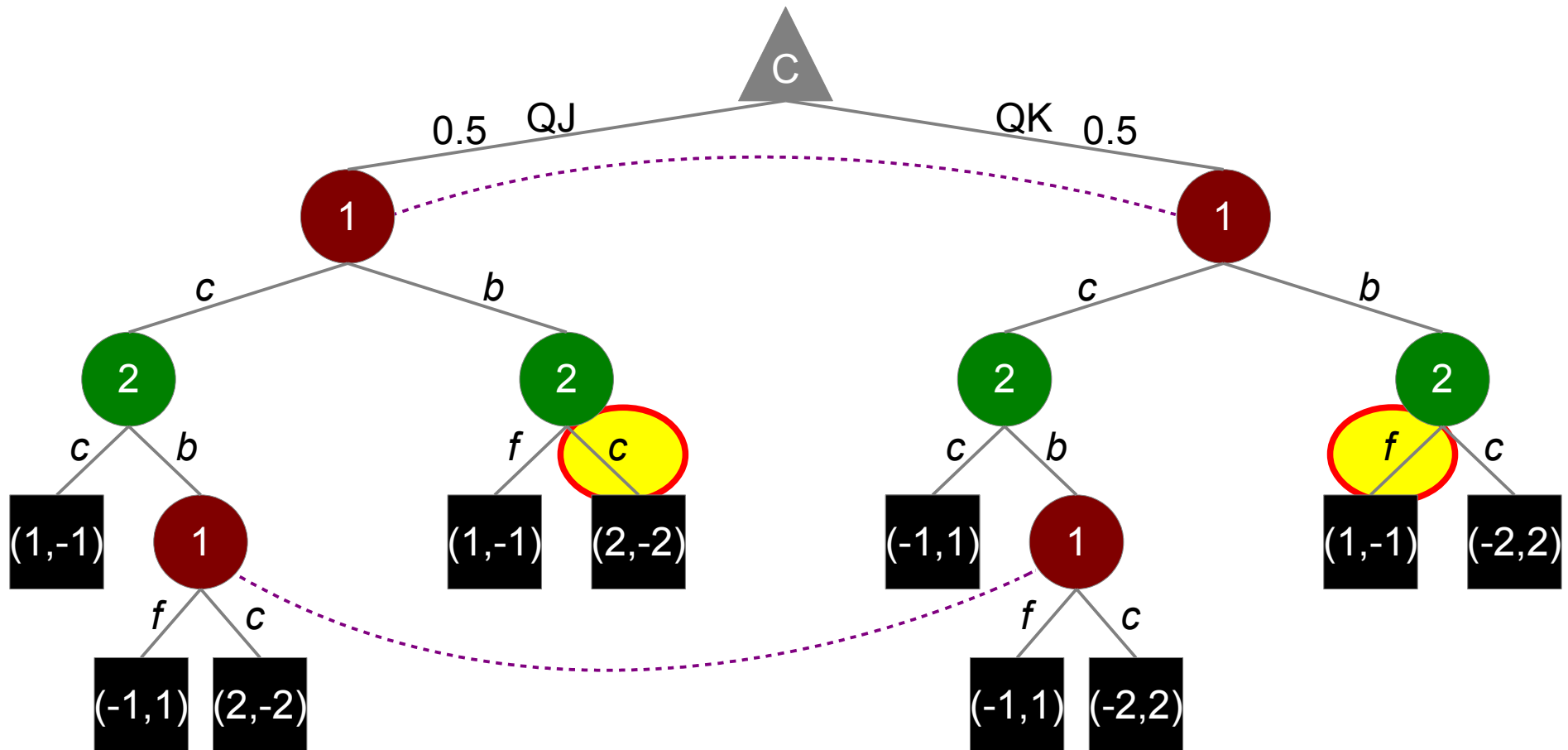
$$u_2(\sigma_1, \sigma_2, \sigma_3, \dots) < u_2(\sigma_1, \sigma_2', \sigma_3, \dots) \text{ for some } \sigma_1, \sigma_3, \dots$$

Domination



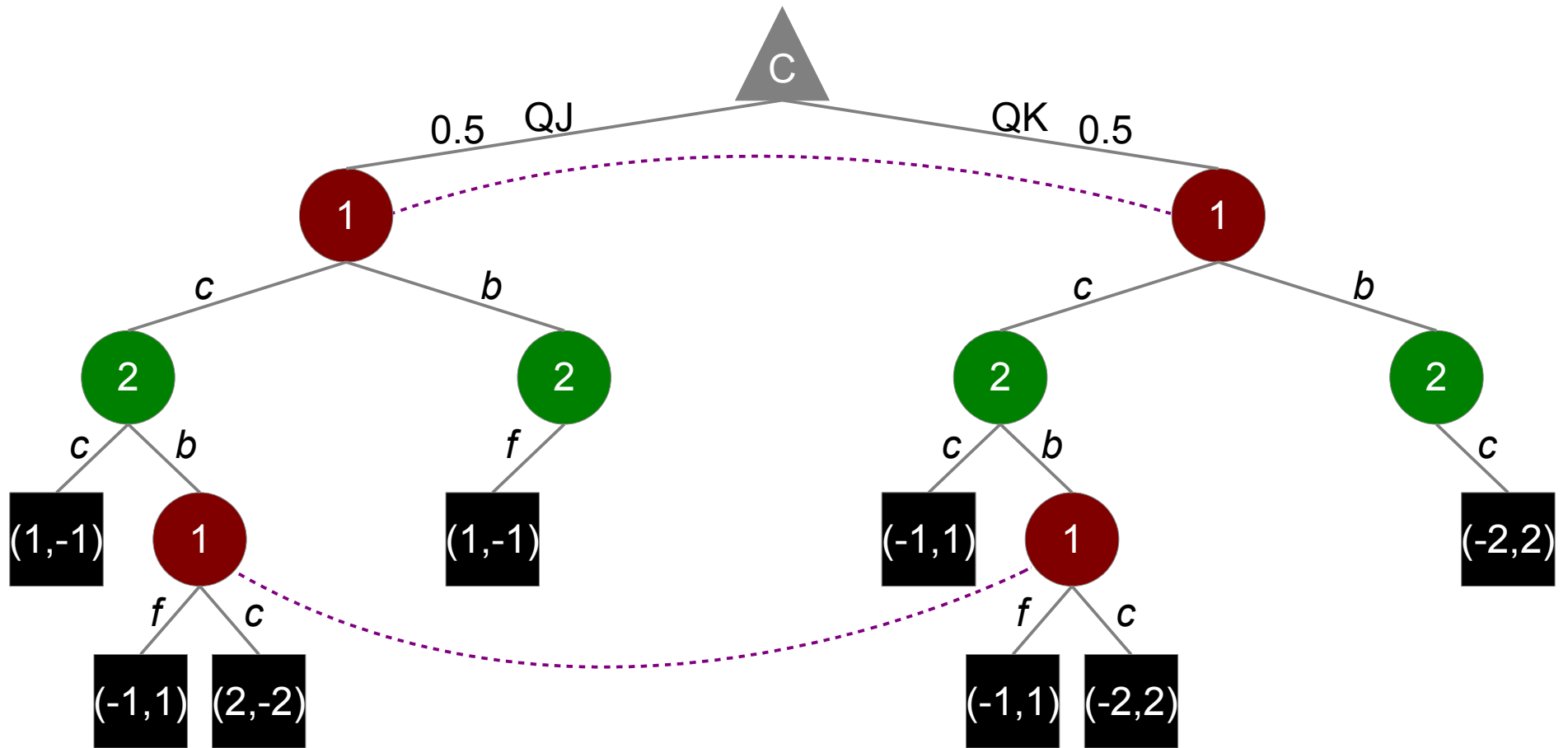
$\sigma_2^{J,c}$ is **dominated** by $\sigma_2^{J,f}$
 $\sigma_2^{K,f}$ is **dominated** by $\sigma_2^{K,c}$

Domination

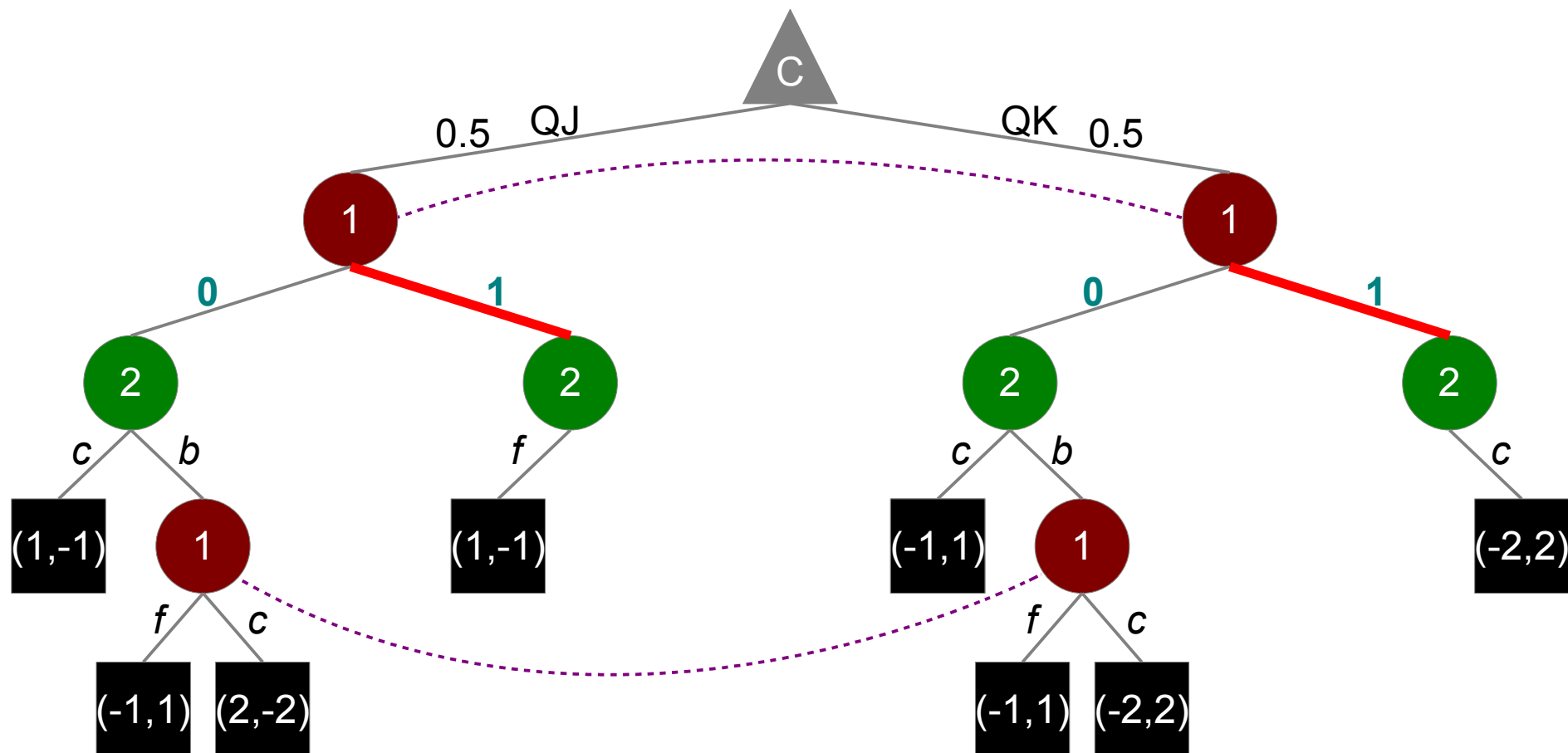


Define a **dominated action** to be an action such that any strategy that always plays that action is dominated (assuming that player plays to reach that action).

Domination

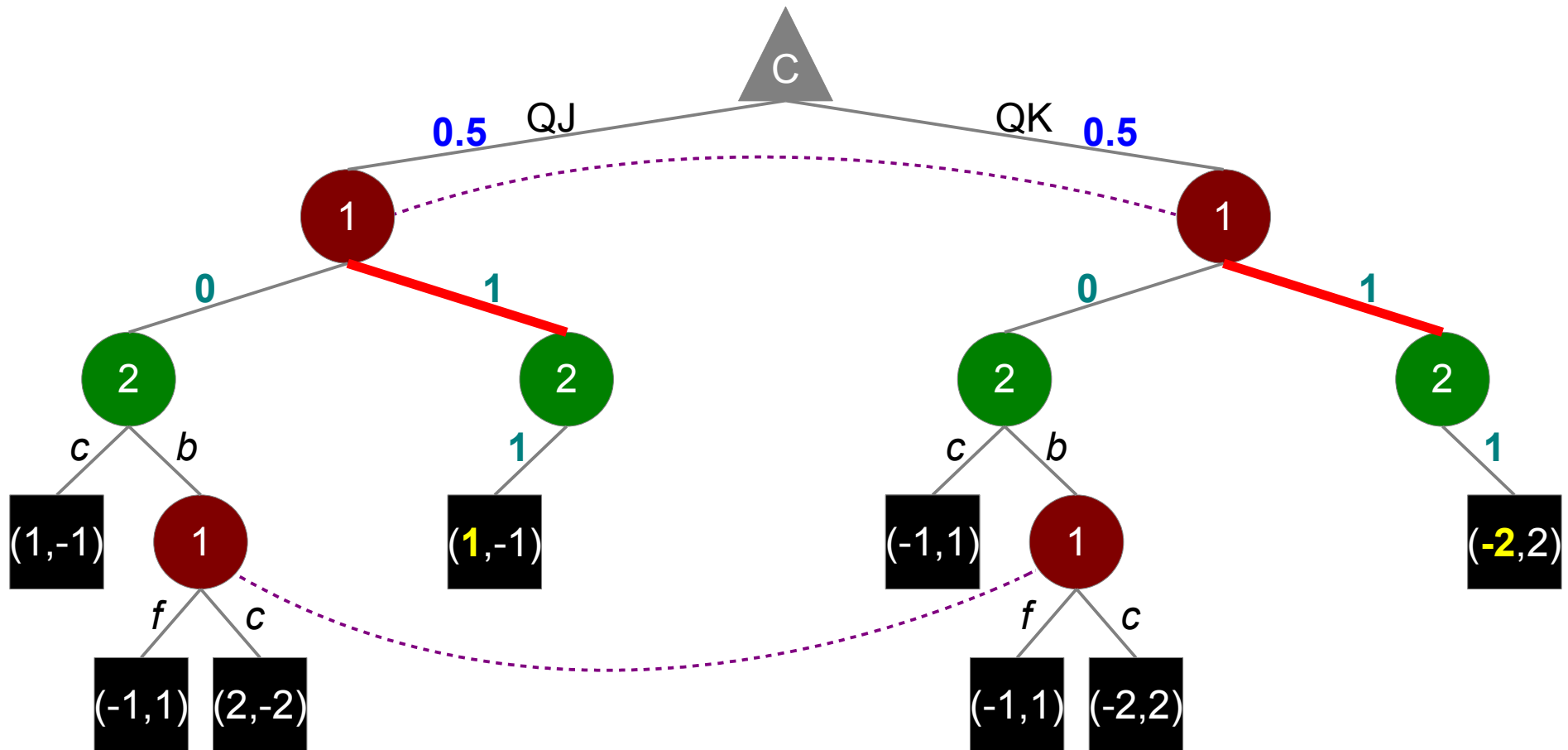


Domination



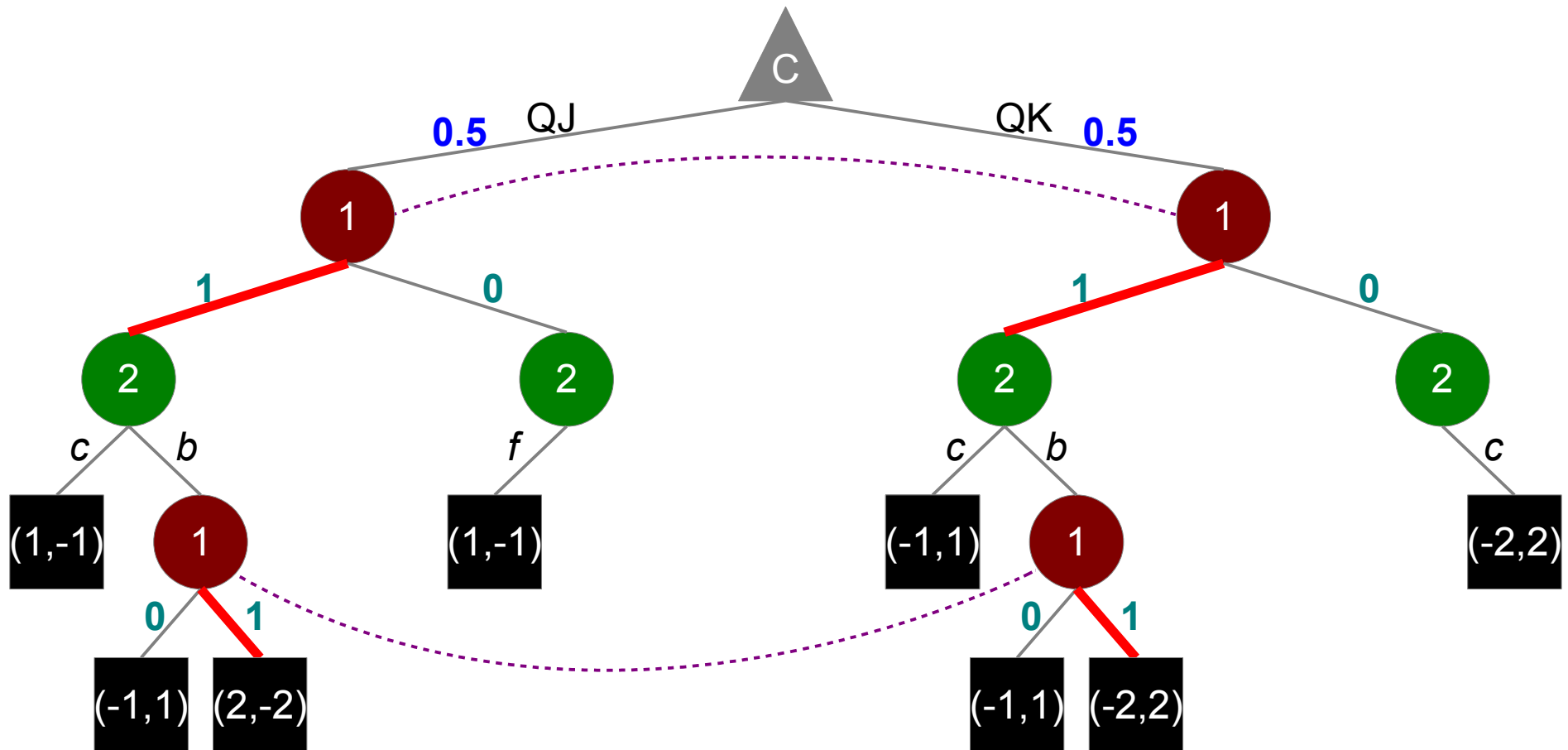
Consider the player 1 strategy σ_1^b that always bets.

Domination



$$u_1(\sigma_1^b, \sigma_2) = 0.5(1)(1)(1) + 0.5(1)(1)(-2) = -0.5$$

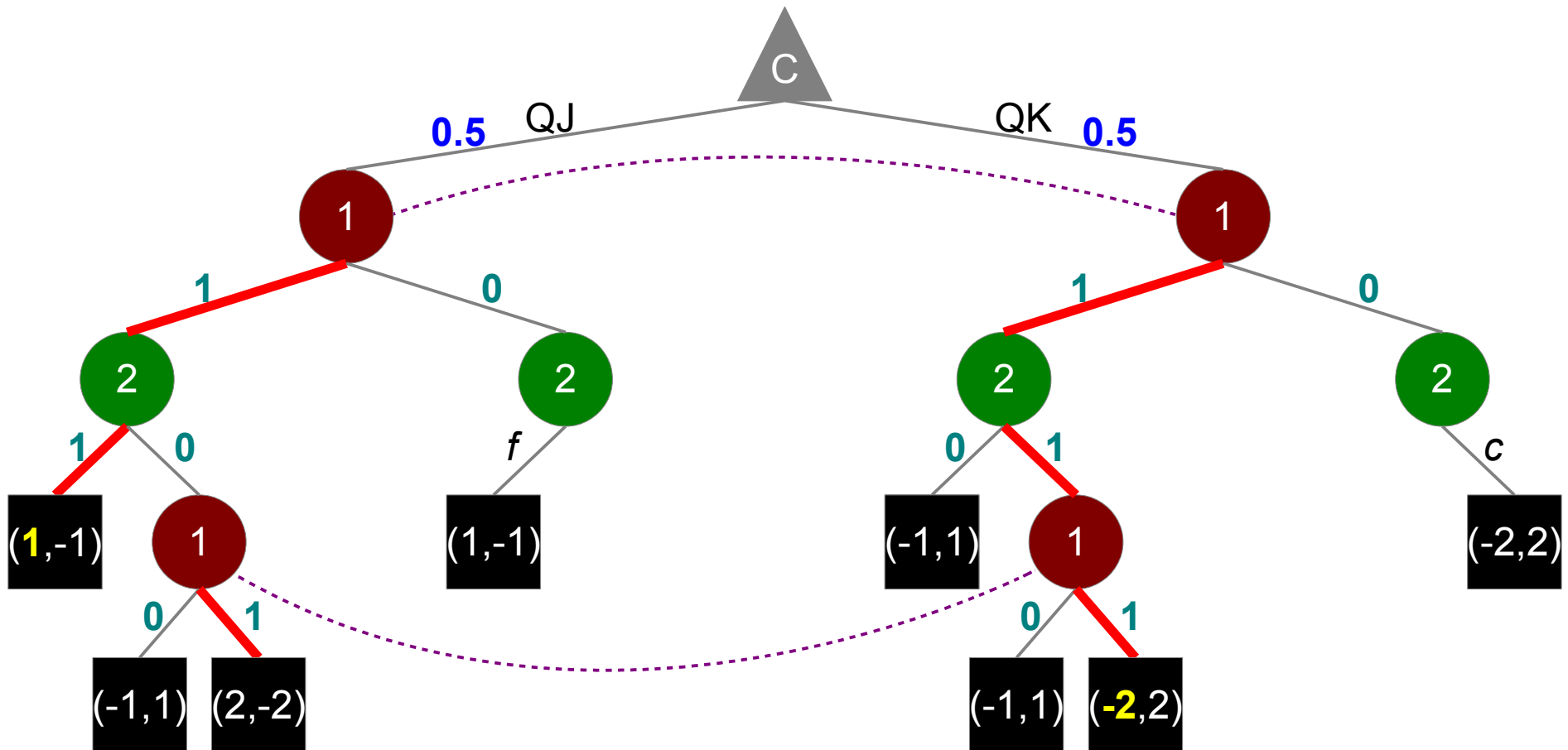
Domination



$$u_1(\sigma_1^b, \sigma_2) = 0.5(1)(1)(1) + 0.5(1)(1)(-2) = -0.5$$

Now consider the player 1 strategy σ_1^{cc} that checks then calls.

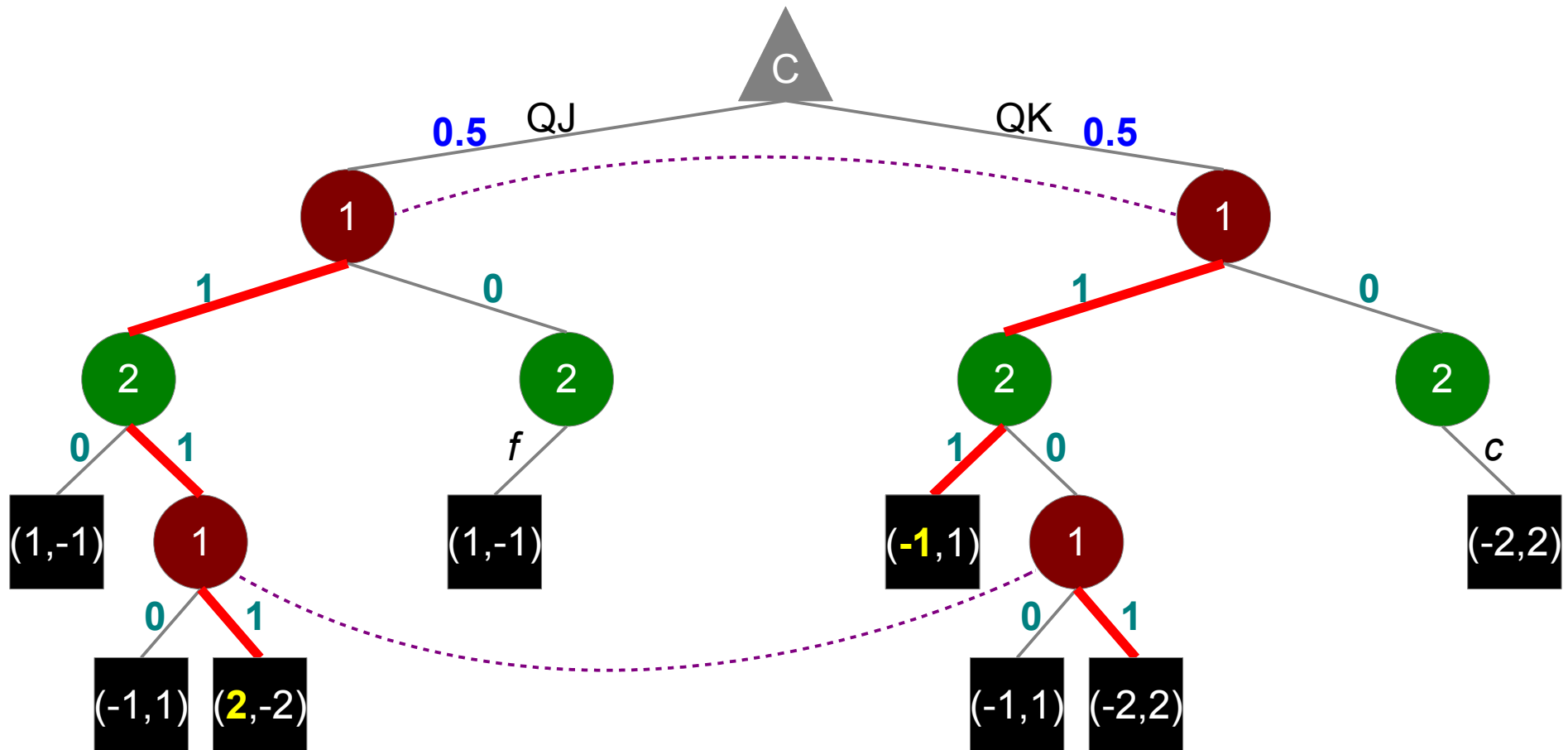
Domination



$$u_1(\sigma_1^b, \sigma_2) = 0.5(1)(1)(1) + 0.5(1)(1)(-2) = -0.5$$

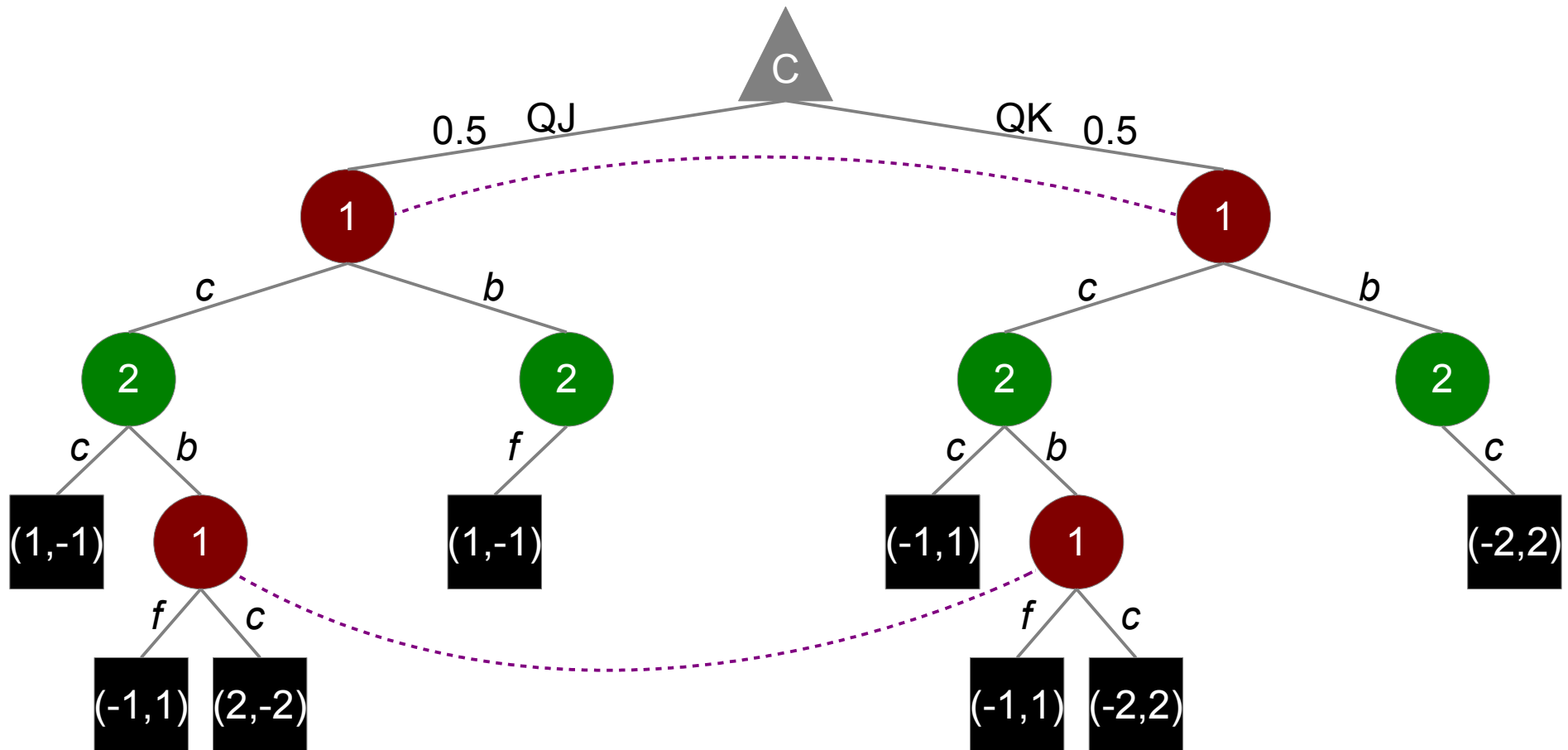
$$u_1(\sigma_1^{cc}, \sigma_2^{Jc, Kb}) = 0.5(1)(1)(1) + 0.5(1)(1)(1)(-2) = -0.5$$

Domination



$$\begin{aligned}
 u_1(\sigma_1^b, \sigma_2) &= 0.5(1)(1)(1) + 0.5(1)(1)(-2) = -0.5 \\
 u_1(\sigma_1^{cc}, \sigma_2^{Jc,Kb}) &= 0.5(1)(1)(1) + 0.5(1)(1)(1)(-2) = -0.5 \\
 u_1(\sigma_1^{cc}, \sigma_2^{Jb,Kc}) &= 0.5(1)(1)(1)(2) + 0.5(1)(1)(-1) = +0.5
 \end{aligned}$$

Domination

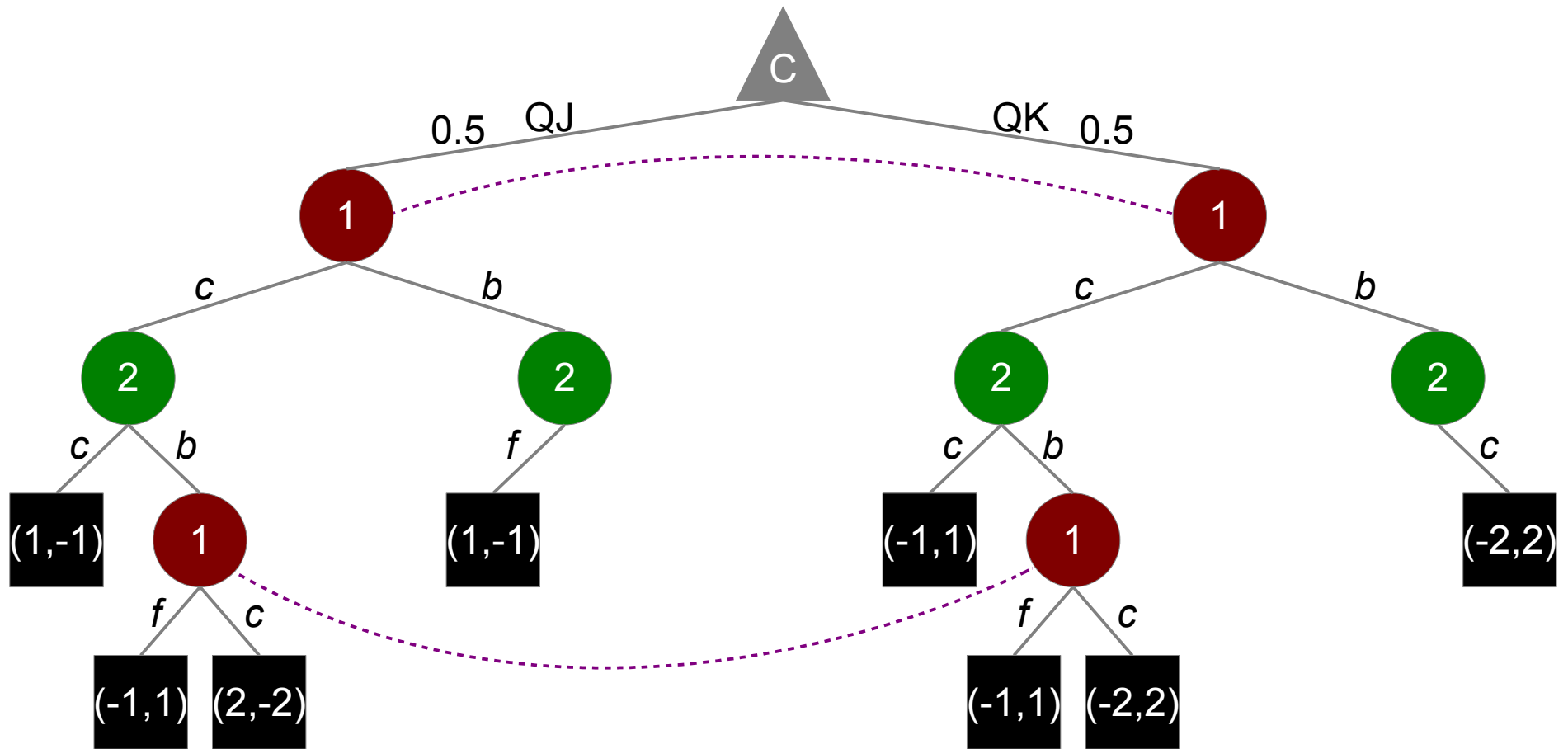


σ_1 is an **iteratively dominated strategy** if there exists σ_1' such that

$$u_1(\sigma_1, \sigma_2, \sigma_3, \dots) \leq u_1(\sigma_1', \sigma_2, \sigma_3, \dots) \text{ for all non-iteratively dominated } \sigma_2, \sigma_3, \dots$$

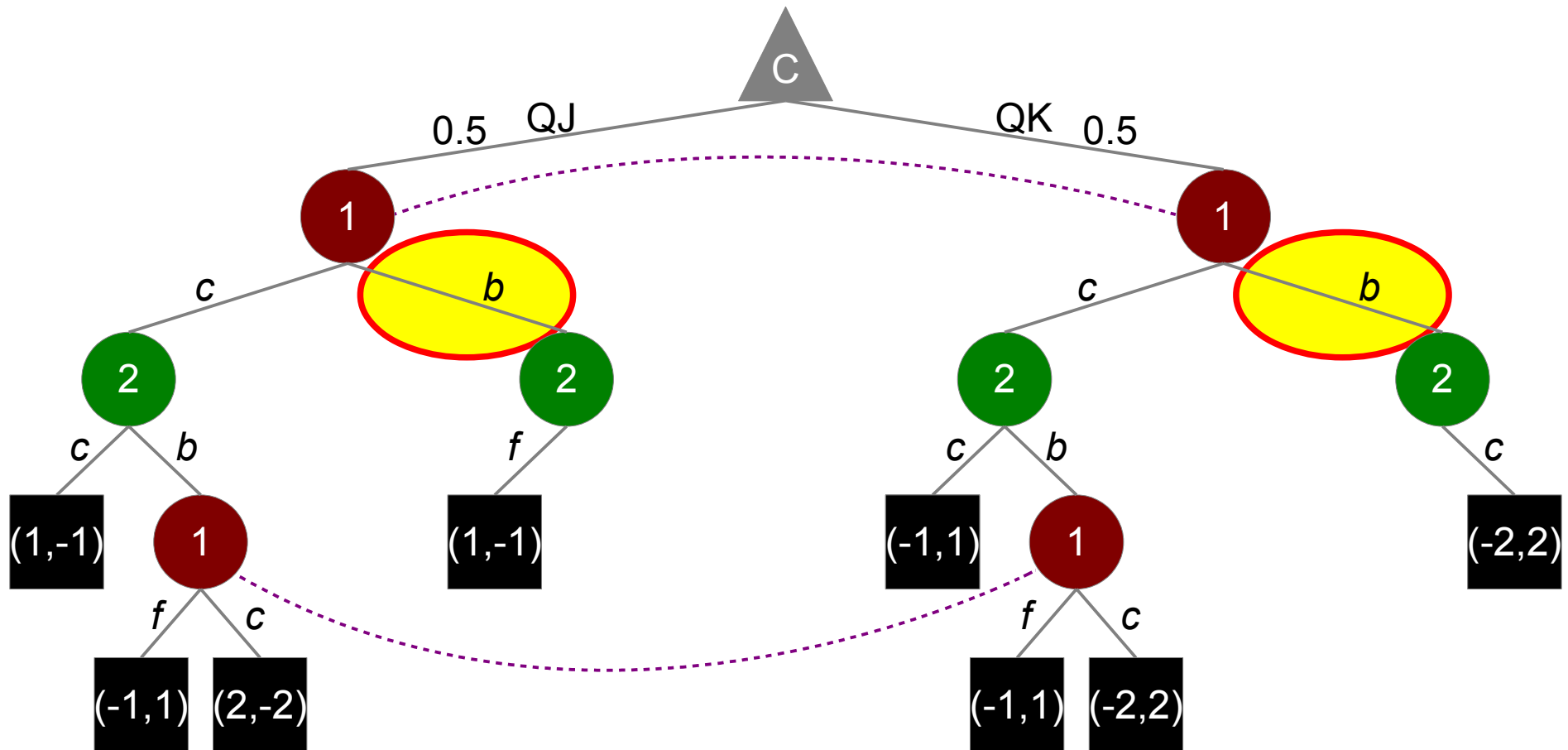
$$u_1(\sigma_1, \sigma_2, \sigma_3, \dots) < u_1(\sigma_1', \sigma_2, \sigma_3, \dots) \text{ for some non-iteratively dominated } \sigma_2, \sigma_3, \dots$$

Domination



σ_1^b is **iteratively dominated** by σ_1^{cc}

Domination



Define an **iteratively dominated action** to be an action such that any strategy that always plays that action is iteratively dominated (assuming that player plays to reach that action).

Domination and CFR

- Clearly, one should not play a **dominated action**.
- If we assume our opponents are rational, then we should also not play an **iteratively dominated action**.
- **Theorem:** If a is an iteratively **strictly** dominated action, and the players play to reach a “often enough,” then when running CFR,

$$\sigma^{\text{AVG}}(a) \rightarrow 0 \text{ as } T \rightarrow \infty.$$

- Can also prove a weaker result regarding CFR avoiding strictly dominated strategies.

Discussion

- We can show that CFR avoids dominated actions and strategies, but how important is it to avoid such actions and strategies?
 - Need to measure correlation between playing dominated actions or strategies and performance.
 - Hard to identify all dominated actions in large games, but may be computationally possible in smaller games.

Discussion

- Recall that CFR generates a sequence of strategy profiles, $(\sigma^1, \sigma^2, \dots, \sigma^T)$ over many iterations T .
- Can show that for an iteratively strictly dominated action a , after a finite number of iterations T_0 , the profiles generated play a with probability 0.
 - If avoiding iteratively dominated actions is enough to perform well, then perhaps there is no need to use the average profile σ^{AVG} as is needed in 2-player zero-sum games.

Conclusion

- CFR can generate strong strategies outside of 2-player zero-sum games, but we do not have a good understanding of why this is so.
- Iteratively dominated actions and strategies should typically be avoided in any game.
- We have shown that the strategies produced by CFR tend to avoid playing iteratively strictly dominated actions.
 - More work is required to conclude that this really does help explain the strong performance of CFR-generated strategies.



Thanks for listening!

Richard Gibson

Twitter: @RichardGGibson

Email: rggibson@cs.ualberta.ca

Website: <http://cs.ualberta.ca/~rggibson>

CPRG Website: <http://cs.ualberta.ca/~poker>