

max moves

We will get the max number of moves by moving the extreme nodes to fill the internal gaps only one step at a time i.e. after one move one gap is reduced

So max no. of moves will be equal to total no. of gaps

$$\text{no. of gaps} = \underset{\downarrow}{\text{range}} - \underset{\downarrow}{\text{occupied}}$$
$$\left(\text{stones}[n-1] - \text{stones}[0] + 1 \right) - (n)$$

but as soon as we move an extreme
(smallest / largest) stone to

into center all the next gaps (for
smallest) & all the prev gaps (for
largest) will get invalidated
and hence don't need to be filled

X \rightarrow stone

O \rightarrow gap.

X	O	O	X	X	X	O	O	O	O	X
1	2	3	4	5	6	7	8	9	10	11

total gaps = 6 (count)

or

$(11 - 1 + 1) - (5)$
range - occupied

$$11 - 5 = 6$$

now if we move the smallest to reduce the gap.

	X	0	0		X	X	X	0	0	0	0	X
↳	1	2	3	4	5	6	7	8	9	10	11	
	0	0	0	X	X	X	0	0	0	X	X	

↓
you can notice here these gaps
no longer need to be filled.

↳ X X X 0 0 X X

↳ X X X 0 X X

↳ X X X X X

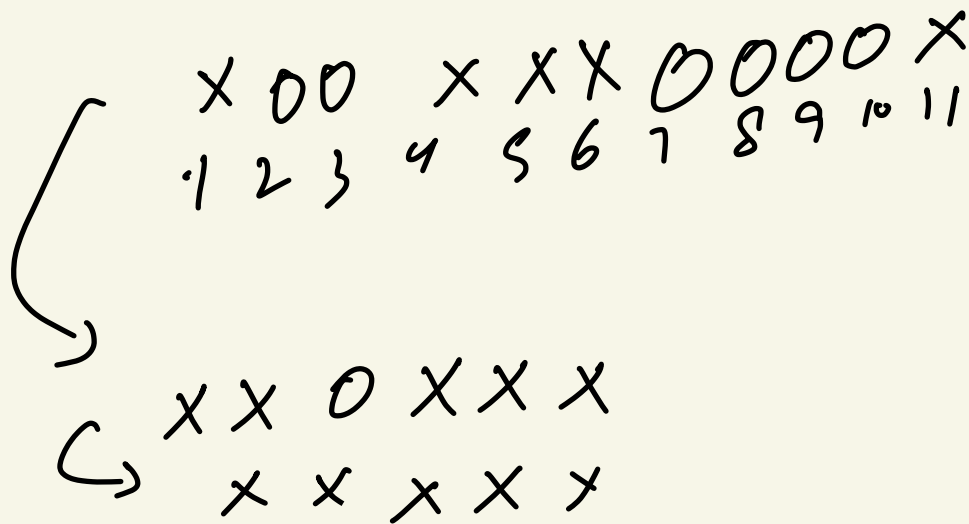
Total steps = 4

so, if we use the smallest first to fill the internal gap

we would take $\text{stones}[1] - \text{stones}[0]$
steps less
-1

so the gap b/w 2nd & 1st stone
needs to be reduced from
no. of total gaps from 1st to nth stone

If we move the largest stone first,
the gap b/w $\text{stones}[n-1] - \text{stones}[n-2]$
would be reduced in 1 move



4 gaps reduced in first move.

$$\text{total moves} = 2$$

i.e.
$$\text{total no. of gaps} = \left[\begin{array}{l} \text{stones } [n-1] - \\ \text{stones } [n-2] - 1 \end{array} \right]$$

$$11 - 6 - 1$$

$$= 6 - 4$$

$$= 2$$

to maximize the no. of moves
we have to move smallest / largest
depending upon which reduces
the gaps by less no.

$$\text{total-gaps} = (\text{stones}[n-1] - \text{stones}[0] + 1) - n$$

$$\text{total-gaps} = \min(\text{stones}[1] - \text{stones}[0] - 1, \text{stones}[n-1] - \text{stones}[n-2] - 1)$$

$$\text{stones}[n-1] - \text{stones}[n-2] - 1)$$

there's another way to calculate this.

that we can calculate the no. of gaps b/w

smallest & second largest

OR

second smallest & largest

& take the maximum

$$\text{max.} \begin{cases} \text{stones}[n-2] - \text{stones}[0] + 1 \\ \text{OR} - (n-1) \end{cases} \rightarrow \text{because we don't consider largest stone}$$

$$\text{stones}[n-1] - \text{stones}[1] + 1 - (n-1)$$

what this means is.

$$\text{stones}[n-2] - \text{stones}[0] + 1 - (n-1)$$

gaps b/w smallest & second largest stones is the max no. of moves because the internal gaps will be filled one step at a time

$\begin{array}{cccccccccc} X & 0 & 0 & 0 & X & X & 0 & 0 & X \\ \hline 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \end{array}$

3 steps

How?

$$\max \left\{ \begin{array}{l} 6-1+1 \\ - (4-1) \end{array} \right\} \text{ OR } \left\{ \begin{array}{l} 9-5+1 \\ - (4-1) \end{array} \right\}$$

$\begin{array}{ccccccc} X & 0 & 0 & 0 & X & X & 0 & 0 & X \\ & & \swarrow & & \searrow & & & & \end{array}$

↳

$\begin{array}{cccccc} X & X & 0 & 0 & X & X \end{array}$

$$\max \left\{ \begin{array}{l} 3 \\ 2 \end{array} \right\}$$

↳

$\begin{array}{ccccc} X & X & 0 & X & X \end{array}$

↳

$\begin{array}{cccc} X & X & X & X \end{array}$

Only consider the example

$\begin{array}{ccccccc} X & 0 & 0 & X & X & X & 0 & 0 & 0 & 0 & X \end{array}$

$$\max \left\{ \begin{array}{l} 2 \\ 4 \end{array} \right\}$$

moving the smallest to right would lead to more steps.