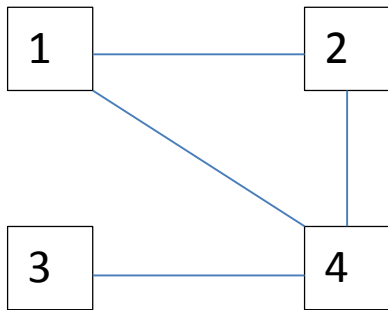


Graphs

Graph $G = (V, E)$ with $n = |V|$, $m = |E|$

- Adjacency matrix data structure
- Adjacency lists data structure

Example: Unweighted Undirected Graph



Adjacency matrix

$M[1...n][1...n]$ of booleans

$M[j][k] = 1$ iff (j,k) is edge

$M[j][k] = M[k][j]$

M	1	2	3	4
1	0	1	0	1
2	1	0	0	1
3	0	0	0	1
4	1	1	1	0

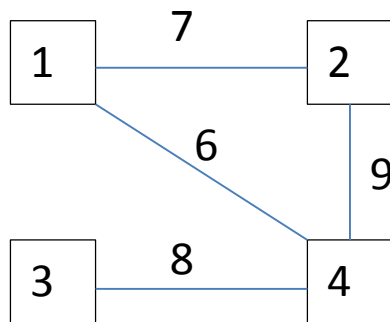
Adjacency lists

$L[1...n]$ of lists of vertices

Each $L[k]$ has neighbors of k
(usually in ascending order)

	L	
1	<div></div>	→ 2 → 4
2	<div></div>	→ 1 → 4
3	<div></div>	→ 4
4	<div></div>	→ 1 → 2 → 3

Example: Weighted Undirected Graph



Weighted adjacency matrix
 $M[1\dots n][1\dots n]$ of ints/floats
 $M[j][k] = \text{weight}(j, k)$
 $M[j][k] = M[k][j]$

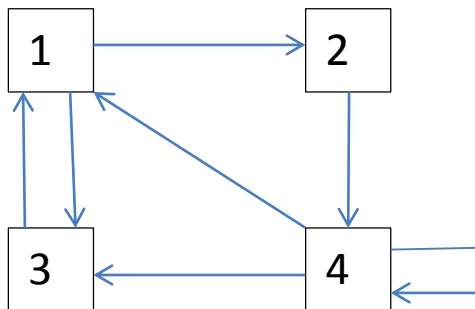
M	1	2	3	4
1	0	7	∞	6
2	7	0	∞	9
3	∞	∞	0	8
4	6	9	8	0

Weighted adjacency lists
 $L[1\dots n]$ of lists of pairs
[neighbor, edge weight]

L	
1	$\rightarrow [2,7] \rightarrow [4,6]$
2	$\rightarrow [1,7] \rightarrow [4,9]$
3	$\rightarrow [4,8]$
4	$\rightarrow [1,6] \rightarrow [2,9] \rightarrow [3,8]$

Typically use 0 or ∞ for
non-existent edges

Example: Unweighted Directed Graph



Adjacency matrix

$M[1\dots n][1\dots n]$ of booleans

$M[j][k] = 1$ iff (j,k) is edge

$M[j][k]$ can be $\neq M[k][j]$

M	1	2	3	4
1	0	1	1	0
2	0	0	0	1
3	1	0	0	0
4	1	0	1	1

Adjacency lists

$\text{Out}[1\dots n], \text{In}[1\dots n]$

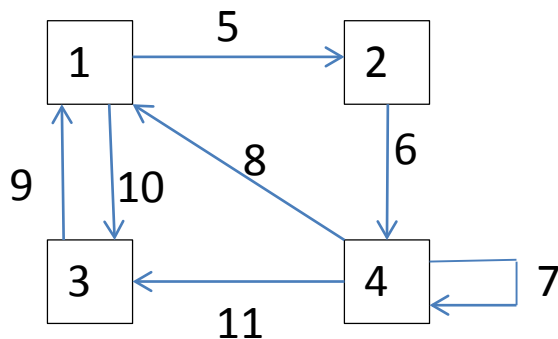
$\text{Out}[k]$ has outgoing neighbors

$\text{In}[k]$ has incoming neighbors

	Out
1	<div></div> $\rightarrow 2 \rightarrow 3$
2	<div></div> $\rightarrow 4$
3	<div></div> $\rightarrow 1$
4	<div></div> $\rightarrow 1 \rightarrow 3 \rightarrow 4$

	In
1	<div></div> $\rightarrow 3 \rightarrow 4$
2	<div></div> $\rightarrow 1$
3	<div></div> $\rightarrow 1 \rightarrow 4$
4	<div></div> $\rightarrow 2 \rightarrow 4$

Example: Weighted Directed Graph



Weighted adjacency matrix
 $M[1\dots n][1\dots n]$ of ints/floats
 $M[j][k] = \text{weight}(j, k)$
 $M[j][k]$ can be $\neq M[k][j]$

	1	2	3	4
1	0	5	10	∞
2	∞	0	∞	6
3	9	∞	0	∞
4	8	∞	11	7

Weighted adjacency lists

$\text{Out}[1\dots n], \text{In}[1\dots n]$

$\text{Out}[k]$ and $\text{In}[k]$ are lists of pairs
 [neighbor, edge weight]

	Out	
1		→[2,5]→[3,10]
2		→[4,6]
3		→[1,9]
4		→[1,8]→[3,11]→[4,7]

	In	
1		→ [3,9] → [4,8]
2		→ [1,5]
3		→ [1,10] → [4,11]
4		→ [2,6] → [4,7]

Analysis:

Each kind of adjacency matrix:

$\theta(n^2)$ space/memory used

Each kind of adjacency lists:

$\theta(n+m)$ space/memory used

Recall m = number of edges

Upcoming we will look at several algorithms on graphs.

We'll analyze the running time of each algorithm using the adjacency matrix and/or adjacency lists data structures.