

# DFS

## Depth-First Search in Graphs:

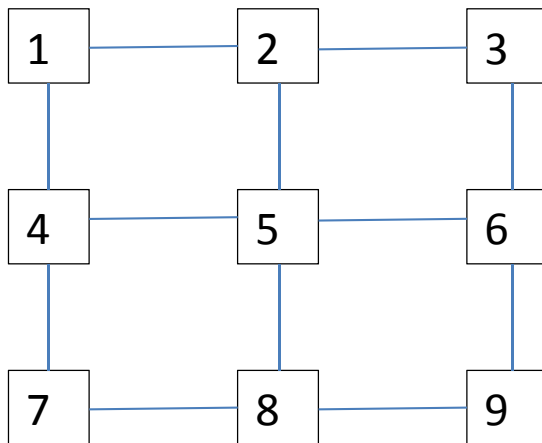
DFS is most similar to pre-order traversal of a tree,  
and can also be similar to post-order traversal

```
boolean seen[1...n];
```

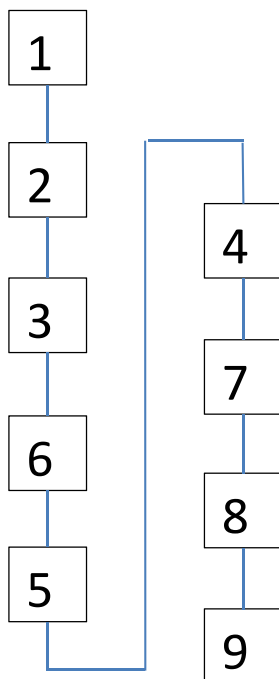
```
DFS (Graph G) {  
    for (k = 1; k<=n; k++)  
        seen[k] = false;  
    choose start vertex;  
    DFS (G, start);  
}
```

```
DFS (G, x) {  
    seen[x] = true;  
    preVisit (x);      // visits nodes in DFS order or pre-order  
    for each vertex y such that (x,y) is an edge  
        if (! seen[y]) {  
            // optionally add edge (x,y) to the DFS tree;  
            DFS (G, y);  
        }  
    postVisit (x);     // optional, visits nodes in post-order  
}
```

## Example: Undirected Graph

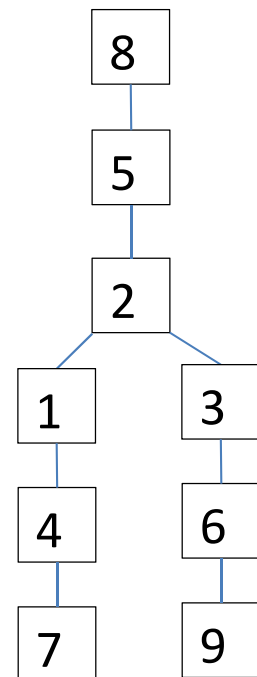


DFS starting at vertex 1:



DFS order: 1,2,3,6,5,4,7,8,9

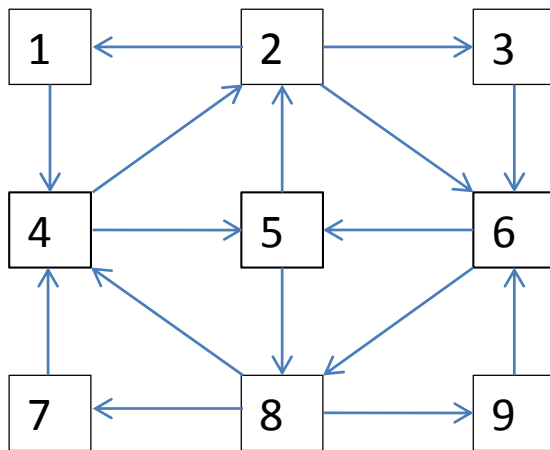
DFS starting at vertex 8:



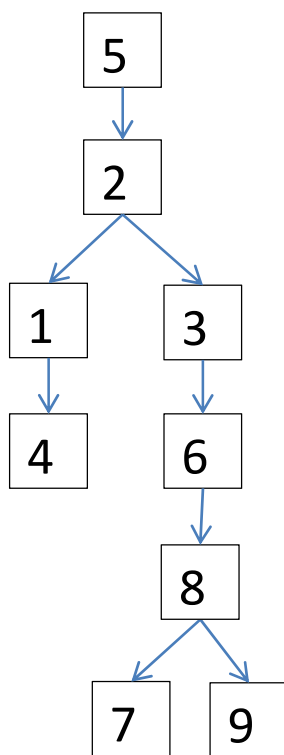
DFS order: 8,5,2,1,4,7,3,6,9

(DFS order is the same as pre-order of the DFS tree)

## Example: Directed Graph

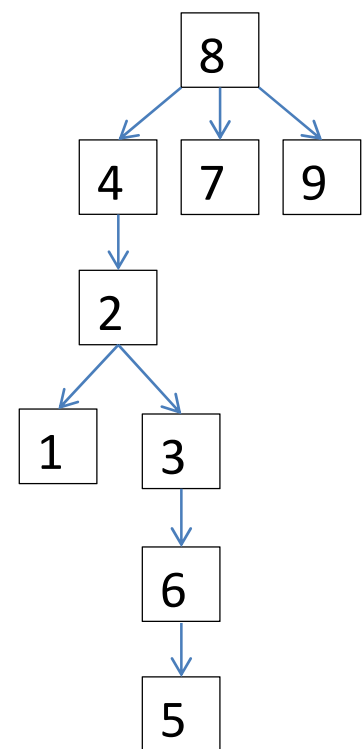


DFS starting at vertex 5:



DFS order: 5,2,1,4,3,6,8,7,9

DFS starting at vertex 8:



DFS order: 8,4,2,1,3,6,5,7,9

(DFS order is the same as pre-order of the DFS tree)

Analysis of DFS (same as for BFS):

If graph represented using adjacency matrix:

$\theta(n^2)$  time

If graph represented using adjacency lists:

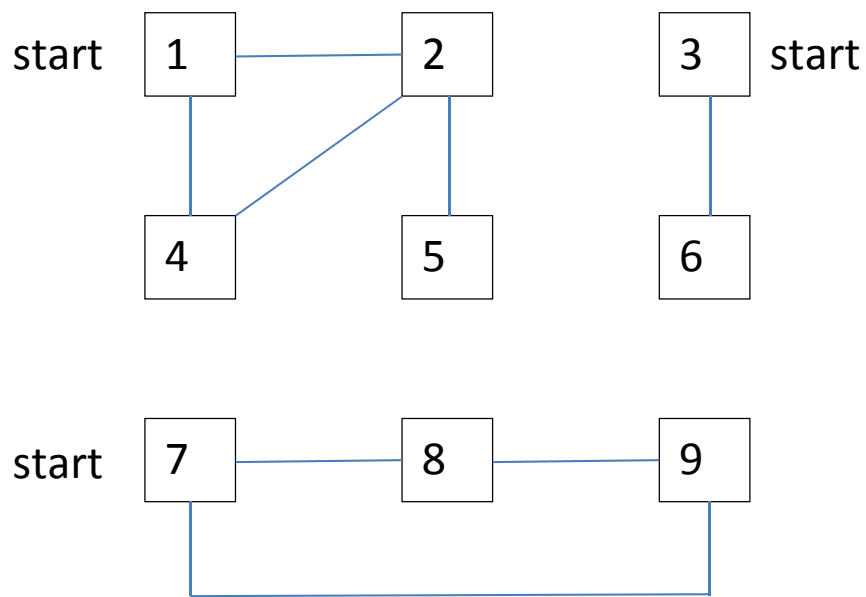
$\theta(n+m)$  time

Simplifies to  $\theta(m)$  time if graph is connected,  
because  $m \geq n-1$  for connected graphs

Note: it is more efficient to use adjacency lists,  
because  $m \leq n^2$  for all graphs

## Applications of DFS:

- Connectedness of undirected graph (same as for BFS):  
Is the graph connected?  
If not, find all the connected components



Components: {1,2,4,5} {3,6} {7,8,9}

More applications of DFS (we'll discuss these next time):

- Strong connectedness of directed graph:  
Is the graph strongly connected?  
If not, find all the strongly connected components
- Detecting a cycle (in undirected or directed graph)
- Topological sort of a Directed Acyclic Graph:  
a linear ordering of the vertices so that whenever  
edge  $x \rightarrow y$  exists, vertex  $x$  must precede vertex  $y$