

Design WhatsApp (WA)

FEATURES

- 0) One to one chat
- 1) Group Messin (around 256 mil/grp)
- 2) Sent + Delivered + Read recnts
- 3) Img sharin (covered in Tinder design)
- 4) Online / last seen
- 5) Chats ~~are~~ are temporary / permanent

WORKFLOW

①

Connect to WA on the cloud via a gateway

Reason :- you might be using ~~not~~ a diff protocol (to talk to WA) than the protocol WA uses to talk to its sys. comps. internally

HTTP over TCP ←
websockets (ws) → Reason → you don't need that much security or big headers that HTTP provides you when talking internally

Gateway takes care of most of the security mech.

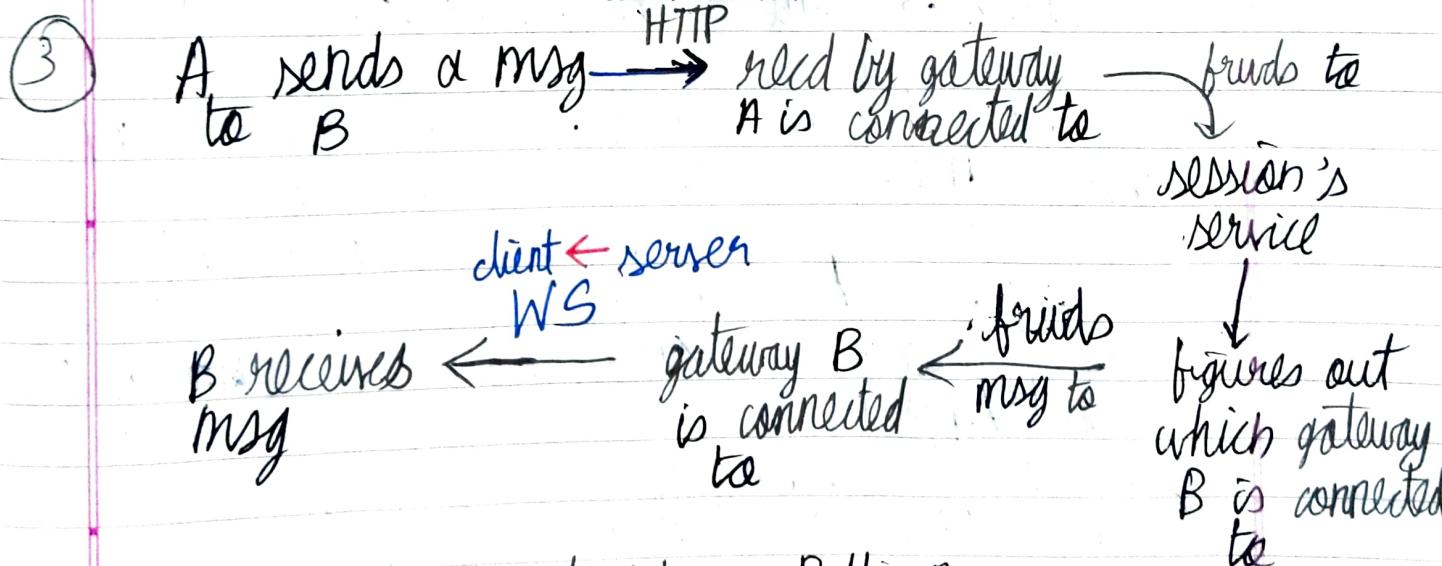
②

Connect user's ~~to~~ via boxes (If A is connected to Box 1, B is connected to Box 1, A can talk to B)

Problem: Currently our gateway is already ~~storing~~ using some mem. to maintain TCP connection's. We shouldn't include data of who's connected to which box here. Plus, if we store this data; we gotta replicate it in other gateways and ∵ this data is transient, we gotta do a lotta syncs → our sys. is too "coupled".

John: Use a separate 'session service' to store New who's connected to which box. sys has been decoupled.

Gateway →
 Stores and maintains more TCP connec's (dumb connec's → dunno what to do with a req → just forward it to sessions service)
Session's Service → Stores which user is connected to which box
 Makes this data accessible to all ~~all~~ gateways.
 Has multiple servers for redundancy



Could also use "Long Polling" instead of WebSockets (WS) but WS is more real time

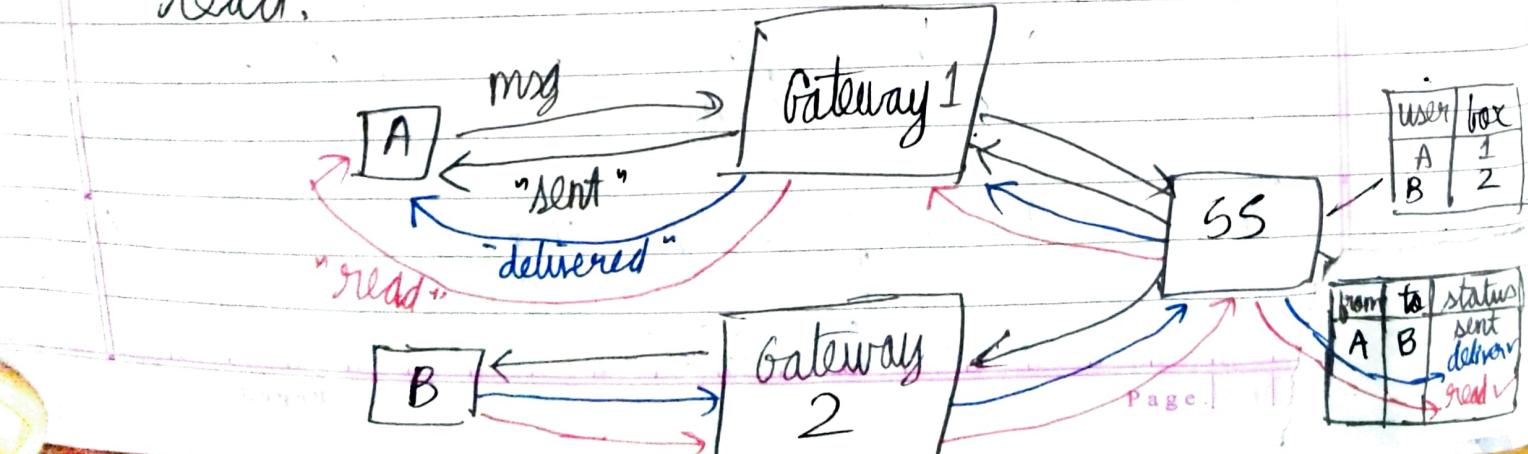
WS → allows P2P comm' i.e. client can send to server and server can also send to client

Read, Sent & Delivered Receipt

Date: M T W T F S S

Parallelly, when session's service (ss) is sending msg from A \rightarrow B to gateway of B, if ~~also~~ also has another dB in which it stores the status of this msg

- ① When the msg ^{from A} is successfully read by ss from A, it sends back a msg to A (sent receipt \rightarrow single tick on WhatsApp) saying it has read the msg and its now gonna try sending it to B. Msg status dB is updated
- ② When B receives the msg forwarded to it by ss, it sends back a res. saying it read the msg from A. Our msg status is updated (msg is now at delivered \rightarrow double tick on WhatsApp stage)
- ③ When B opens and logs into WhatsApp ~~then~~ ~~does~~ automatically and marks a msg as read, then, automatically a res is sent from B to ss saying msg was read \rightarrow double BLUE tick ~~on~~ on WhatsApp. Our msg status dB is updated & ss sends a msg to A that msg sent to B has been read.



Last Seen / Online

Basically a client sends 2 kinda reqs to gateways

- ① User generated → sending, reading texts generate "read" receipts
 - ② Sys. generated → receiving texts from server everytime your ~~mobile~~ internet connection is switched on
- generating "delivered" receipts and sending to SS once msgs are received (not read yet)

- i) :- WhatsApp introduce a new service → Last Seen service which has access to a DB that logs user activity. It stores
 - ① user ② timestamp of latest activity
 a time-series DB
- ii) Anytime a user generates a req to WhatsApp servers (to read, send, delete, star msgs, block ppl, set status etc.), an "activity" is performed and thus, the timestamp in the DB gets updated
- iii) Thus, when another user wants to know if a user is online or not, "last seen services" fetches this data
- iv) IMP : - "activities": only entail user-generated reqs. & not sys-generated reqs

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Grp Msgs

- 1) Now we introduce a separate service called a "Grp Service" (GS) that ~~serves~~ our SS talks to for grp msg functionality. Grp service has a DB that maps 1 grp id \rightarrow many user IDs
- 2) When an SS gets a msg from a user, it asks the GS ~~for~~ that "who are the other members in this grp?" Then, once it receives the members? user ID SS uses its own DB to map which gateway/bar it needs to find msg to.
- 3) Most chat apps tryna limit no. of members in a grp else you gonna be fannin out way too many reqs. to keep track of in an encrypted authorized manner in real-time

J

One soln \rightarrow batch process these reqs.

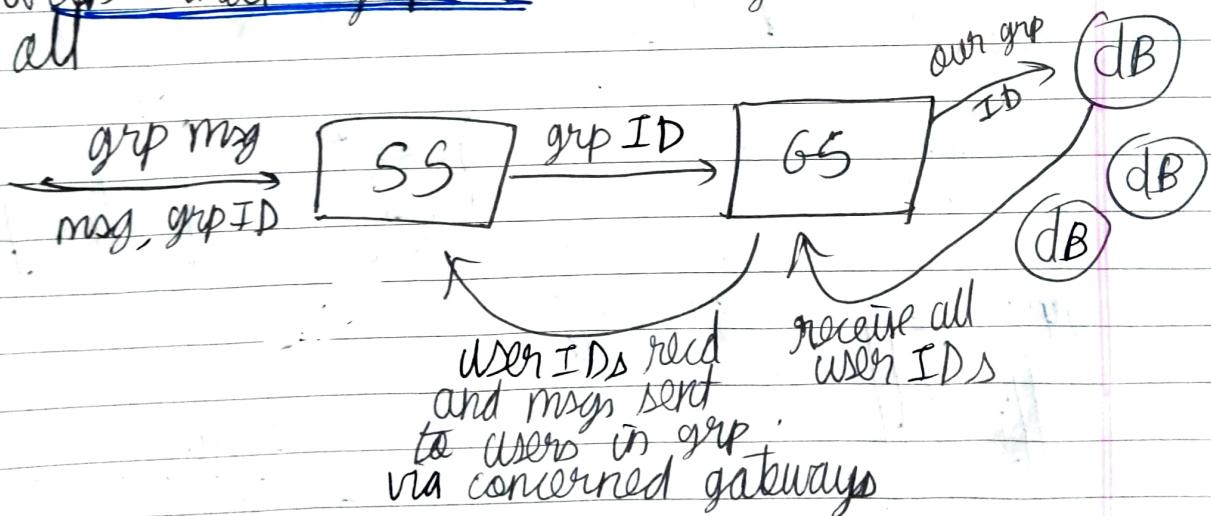
\rightarrow wait for users to pull the msgs (not real-time tho)

- 4) Parser service \rightarrow lies b/w a bar and ss
Parser \rightarrow convert an electronic msg into a sensible msg using lang. obj.
 \rightarrow msg ~~for~~ is unparsed, encrypted, auth. and sent msgs.
 \rightarrow msg ~~for~~ is parsed, decrypted, auth. for ~~other~~ msgs to be delivered

- 5) Persistent Hashin → helps reduce memory footprint across servers by delegating only some info to some servers.
- helps route reqs. to concerned boxes. Routing is based on grp ID to tell users belonging to this grp.

each grp
 Basically, msg has grp IDs attached and, on receiving a msg pertaining to a particular grp ID, the SS routes its req. (to grp service) based on grp ID and when grp service receives this msg, it accesses its DB associated with that grp ID → gets user IDs of all users in grp.

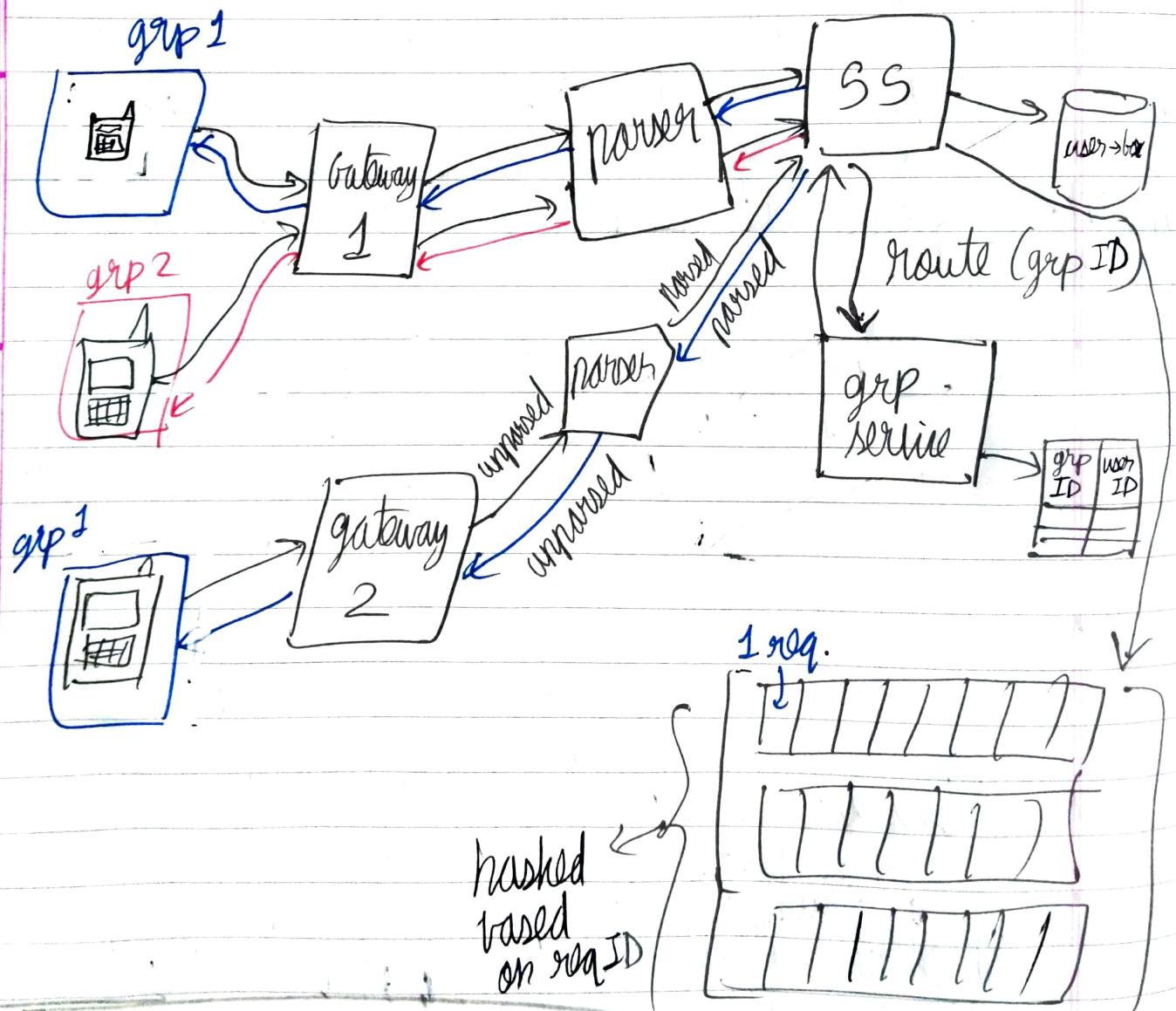
got via consistent hashin. ← with all



- 6) Message Queues → Needed to retry sending a failed msg
- ↳ a part of our SS which ensures that msg will be sent.
- helps us keep track of upto what nt. msgs were successfully sent.
- knows from what msg. it needs to retry sending failed msgs.

NOTE : grp ~~msg~~ receipts (sent, read, delivered) are pretty expensive to manage since every connected grp member needs to +vly respond

NOTE : Deprioritization msgs. → During time periods of excessive load on servers, facebook just serves ~~to~~ msgs pertaining to sender, receiver msgs and all ~~other~~ other op's such as read receipts are suspended. (rate-ltd.)
eg:-festivals



NOTE : Auth. services and Lbs ~~are also introduced in sys.~~

Msg Queue