

Objectives To experiment with procedural abstraction, using it to produce all of the recursion for the example problems.

Same rules as the previous assignments, including the prohibition of mutation.

#1 (60 points) `snlist-recur`. I have slightly adapted the book's definition of `s-lists` (EoPL, p 8) to allow numbers as well as symbols in the lists.

```
<sn-list>          ::= ( {<sn-expression>}* )
<sn-expression> ::= <number> | <symbol> | <sn-list>
```

Define a procedure `snlist-recur` that is similar to `list-recur` (written in class), but it returns procedures that work on `sn-lists`. When we call `snlist-recur` with arguments of the correct types, it returns a procedure that

- takes an `sn-list` as its only argument, and
- traverses the entire `sn-list` and its sub-lists, and does the intended computation.

After you write `snlist-recur`, test it by using it to define the functions in (a)-(f). Each of them will have an `sn-list` as one of its arguments.

When you use `snlist-recur` to produce a procedure `f`, in some cases (namely parts b and e), you may need to first write a curried version of `f` (as we did in class when we wrote curried versions of `member?` and `map`) in order to get a procedure that recurs on only one argument (the `sn-list`). Procedures that you pass as arguments to `snlist-recur` **must not be explicitly recursive**, nor may they call `map`, which is a substitute for recursion. All of the recursion on the `sn-lists` in your solutions should be produced by `snlist-recur` itself. To reiterate, none of your procedures in parts (a) - (f) should contain any explicit calls to recursive procedures that you write.

There are ways to write all of the required procedures without properly using `snlist-recur`. But that would miss the point of this problem. Thus, if you do not use `snlist-recur` as prescribed above, **you will not earn any points**, even if the grading program says that your code works for all of the test cases.

Hint for this problem: `snlist-recur` will probably need to take three arguments (`snlist` refers to an argument passed to the procedure returned by a call to `snlist-recur`):

A base-value to be returned when `snlist` is the empty list.

A procedure to be applied when the `car` of `snlist` is a list (i.e., it is a pair or the empty list).

A procedure to be applied when the `car` of `snlist` is not a list (i.e., it is a symbol or number).

Note: However you design it, when your `snlist-recur` procedure is applied to arguments of the proper types, it must return a procedure that expects exactly one argument (an `sn-list`).

Notes:

1. As usual, your code for these and for any other procedures you write may assume that all of their arguments are the correct type(s). You do not have to do any checking of arguments for validity.
2. Your code for each of the following parts can be fairly short. My longest procedure definition is only 8 lines long.

(a) (`sn-list-sum snlst`) finds the sum of all of the numbers within `snlst` (which contains no symbols).

```
(sn-list-sum '((2 (3) 4) 5 ((1)) ())) → 15
(sn-list-sum '()) → 0
```

(b) `(sn-list-map proc snlst)` applies `proc` to each element of `snlst` and returns the results in an sn-list that has the same “shape” as `snlst`.

```
(sn-list-map (lambda (x) (+ 1 x))  
  '((2 (3) 4) 5 ((1)) () 5)) → ((3 (4) 5) 6 ((2)) () 6)
```

(c) `(sn-list-paren-count snlst)` counts the number of parentheses required to produce the printed representation of `snlst`. (You can get this count by looking at `cars` and `cdrs` of `snlst`).

```
(sn-list-paren-count '()) → 2  
(sn-list-paren-count '(2 (3 4) 5)) → 4  
(sn-list-paren-count '(2 (3) (4 () ((5)))))) → 12
```

Note: sn-lists
are always
proper lists.

(d) `(sn-list-reverse snlst)` reverses `snlst` and all of its sublists.

```
(sn-list-reverse '(a (b c) ( ) (d (e f)))) → (((f e) d) ( ) (c b) a)
```

(e) `(sn-list-occur s snlst)` counts how many times the symbol `s` occurs in the sn-list `snlst`.

```
(sn-list-occur 'a '(( ) a ((a)) a (a a b a) (a a))) → 8
```

(f) `(sn-list-depth snlst)` finds the maximum nesting-level of parentheses in the printed representation of `snlst`.

```
(sn-list-depth '()) → 1  
(sn-list-depth '(1 2 3)) → 1  
(sn-list-depth '(1 (2 3) 4)) → 2  
(sn-list-depth '(1 (2 (3)) (2 3))) → 3  
(sn-list-depth '(((3) (( ) 2) (2 3) 1))) → 4
```

#2 (20 points) Recall the following syntax definition from page 9 of EOPL:

```
<bintree> ::= <number> | (<symbol> <bintree> <bintree> )
```

Write a `bt-recur` procedure, similar to the `list-recur` and `snlist-recur` procedures from class and this homework. Calling `bt-recur` produces a procedure that recurs over all of the elements of a bintree.

Then use `bt-recur` to create the following two procedures:

- **(bt-sum T)** finds the sum of all of the numbers in the leaves of the bintree `T`.
- **(bt-inorder T)** creates a list of the symbols from the *interior* nodes of `T`, in the order that they would be visited in an inorder traversal of the binary tree.

The following transcript should help your understand what `bt-sum` and `bt-inorder` do.

I do not show the code that was used to construct `t1`.

```
> t1  
(a (b 1 4) (c (d 2 5) 3))  
> (bt-sum t1)  
15  
> (bt-inorder t1)  
(b a d c)  
> (define t2 (list 'e 6 t1))
```

```
> t2  
(e 6 (a (b 1 4) (c (d 2 5) 3)))  
> (bt-sum t2)  
21  
> (bt-inorder t2)  
(e b a d c)
```

Note: As in the `snlist-recur` problems from the previous problem, the definitions of `bt-sum` and `bt-inorder` should not contain any explicit recursive calls. All recursion must be produced by `bt-recur`.

Previous terms' Piazza Q&A

Assignment 9: Efficiency for bt-inorder

Does bt-inorder needs to be in $O(n)$ time? I'm currently using an append, which makes it $O(n \log n)$.

Instructor answer: You are allowed to do it the simple way, using append. The best case is $\Theta(n \log n)$. Worst case is $\Theta(n^2)$.