#### **Design Document**

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## **Description:**

Our design uses a single register accumulator to store data and compare it to inputs. At all times we only use one register and use an allocated space in Memory for data. For our addresses we will use sign extensions to target specific places in memory to receive either data for destination. The input must have the correct first bit to target the proper place in memory.

We are going to use 2 registers, the accumulator(\$acc) and stack pointer(\$sp). The accumulator is the only register available by the programmer.

<u>I:</u>					•
Opcode		Immediate		Unused	
5		8		3	
AI:					
Opcode	Address		Immediate	;	
5	8		3		
PC relative for	bne and beq				
<u>A</u> :					
Opcode	Address		Unused		
5		8	3		

We left shift by 1 bit then we sign extent (the most significant bit will be 1 if it is a data and 0 if it is a instruction)

## **Instructions**

	<u> </u>	<u> </u>							
Name		Туре	Operation Opcode						
load	а	А	acc = rt	00001					
			es an 8 bit address a and loads the value at memory address a to the umulator, using the address rule.						
save	а	А	Mem[getAddr(rt)] = acc 00						
		Take an 8 bit address a and save the value in the accumulator into the memory with address a, using the address rule.							
loadui	imm	I	acc = {imm, 8b'0} 00011						
		Takes an 8 bit immediate and load it to the upper 8 bits of the accumulator							
bne	a, imm	AI	if(acc != Mem[getAddr(rt)]) 00100						

			PC = PC + 2 + getAddr(imm)				
		Takes an 8 bit address and a 3 bit immediate. If the value stored at address a is not equal to the value of the accumulator, then jump to the address calculated from the immediate using the branch address rule.					
beq	a, imm	AI if(acc == Mem[getAddr(rt)]) 00101 PC = PC + 2 + getAddr(imm)					
		a is equal to	oit address and a 3 bit immediate. If the value the value of the accumulator, then jump to the com the immediate using the branch address	ne address			
slt	а	А	acc = acc < Mem[getAddr(rt)] ? 1:0	00110			
		the accumul	e value in the accumulator with the value stor ator is less than a then we set the accumulat mulator to 0.				
slti	imm	I	acc = acc < SignExtent(imm) ? 1:0	00111			
		Compare the value in the accumulator with the immediate, if the accumulator is less than the immediate then we set the accumulator to 1, else we set the accumulator to 0.					
j	а	А	PC = getAddr(rt)	01000			
		Jump to the instruction with address a, calculated using the address rule.					
jal	а	I	Men[ra] = PC + 2 PC = getAddr(imm)  01001				
			instruction with address a, calculated using trent PC + 2 to a fix memory location.	he address rule.			
sw	imm	1	sp + SignExtent(imm) = acc 01010				
		Stored the value in the accumulator onto the stack where it is offset imm to the stack pointer.					
lw	imm	I	acc = sp + SignExtent(imm) 01011				
		Stored the value from the stack where it is off offset imm to the stack pointer to the accumulator.					
ms	imm	I	sp = sp + SignExtent(imm)	01100			
		Move the stack pointer with the sign extended immediate.					

sub	а	А	acc = acc - Mem[getAddr(rt)]	01101			
		Subtract the value stored at address a from the accumulator and store the result in the accumulator					
add	а	А	A acc = acc + Mem[getAddr(rt)] 01110				
		Add the valu	ue stored at address a to the accumulator and ator	d store the result in			
addi	imm	I	acc = acc + SignExtent(imm)	01111			
		Add the sigr	n extended immediate to the accumulator and ator	I store the result in			
and	а	А	acc = acc & Mem[getAddr(rt)]	10000			
		And the value stored at address a to the accumulator and store the result in the accumulator					
or	а	A acc = acc   Mem[getAddr(rt)] 10001		10001			
		Or the value stored at address a to the accumulator and store the result in the accumulator					
ori	imm	I	acc = acc   ZeroExtent(imm) 10010				
		Or the zero extended immediate to the accumulator and store the result in the accumulator					
loadi	imm	I	I acc = SignExtent(imm) 10011				
		Load the sign extended immediate to the accumulator.					
ZeroE SignE ra = 0	Extent = {8	dress[7]}, add o'0, imm} [address[7]},ir					

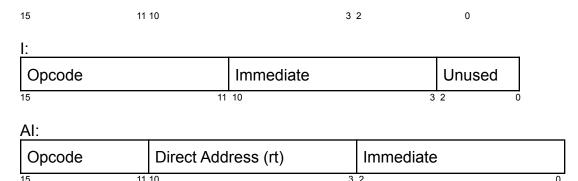
Address rule: We left shift by 1 bit then we sign extent (the most significant bit will be 1 if it is a data and 0 if it is a instruction)

Branch address: Left shift the immediate by 1, sign extend it to 16 bits then add it to the value of the current PC plus 2.

#### **Types**

A:

<u> </u>		
Opcode	Direct Address (rt)	Unused



# Call procedure

For the callers, they are responsible to store the \$acc register value, and put the return address on the \$acc. For callees, they are responsible to restore the value in the Data memory, and callees will move the stack to store the original value of the data in the stack memory, and restore them back before return. Also, it's callee's responsibility to store the return address in the stack memory and use them for return.

Example program(s)

High Level Code	<del>1 - 1</del>	A	ssembly	<u>N</u>	lachine Code	<u>)</u>	Addresses
High Level Code  int relPrime(int n) {     int m;     m = 2;     while (gcd(n, m) != 1) {         m = m + 1;     }     return m; }	loop:	loadi save ms load sw save load sw jal save lw save lw save loadi bne load add save	ssembly  2 m -12 m 0 a n 4 b ra 8 gcd 0 0 m 4 n 8 ra 1 o, end m 1 m	10011 00010 01100 00001 01010 00001 01010 00001 01010 01001 01011 00010 01011 00010 01011 00010 01011 00010 01011 00010	00000010 10000011 11110110 1000001 000000	000 000 000 000 000 000 000 000 000 00	Addresses  0x 0030 0x 0032 0x 0034 0x 0036 0x 0038 0x 003A 0x 003C 0x 003E 0x 0040 0x 0042 0x 0044 0x 0046 0x 0048 0x 004A 0x 004A 0x 004C 0x 004E 0x 004C 0x 0050 0x 0052 0x 0054 0x 0056 0x 0058 0x 005C 0x 005E
	end:	j ms	loop 12	01000 01100	00011011 00001100	000 000	0x 0060 0x 0062

		j	ra	01000	11111111	000	0x 0064
<pre>int gcd(int a, int b) {     if (a == 0) {         return b;     }     while (b!= 0) {         if (a &gt; b) {             a = a - b;         } else {             b = b - a;         }     return a; }</pre>	gcd: loop: go: else:	loadi bne load i bne load slt save load sub save j load sub save j load j	0 a, loop b ra  0 b, go end b a i 1 i, else a b a loop b a bo loop a ra	10011 00100 00001 01000 10010 01000 00001 00110 00010 10011 00010 01001 01001 01101 00010 00001 01101 00010	00000000 10000000 10000001 111111111 000000	000 010 000 000 000 001 000 000 000 000	0x 0002 0x 0004 0x 0006 0x 0008 0x 000A 0x 000C 0x 000E 0x 0010 0x 0012 0x 0014 0x 0016 0x 0018 0x 001A 0x 001C 0x 001E 0x 0020 0x 0022 0x 0022 0x 0024 0x 0028 0x 002A 0x 002C 0x 002E
if (n == 0) {     n++; } else {     n = 2; }  while (n != 0) {     n = n - m }	else: done: loop:	loadi beq load sub save j	0 n, else n 1 n done 2 n  0 n, done n m n loop	10011 00100 00001 01110 00010 01000 10011 00010 10011 00101 01010 01000	00000000 10000100 10000100 00000001 10000100 00000100 10000100 10000100 10000100 10000100 10000100 000000	000 100 000 000 000 000 000 000 100 000 000 000	0x 0002 0x 0004 0x 0006 0x 0008 0x 000A 0x 000C 0x 000E 0x 0010 0x 0012 0x 0002 0x 0004 0x 0006 0x 0008 0x 000A 0x 000C 0x 000C 0x 000C

<pre>int count = 0; for (int i = 0; i &lt; n; i++) {           count++; }</pre>	loadi 0 save count save i loop: beq n, done add 1 save count save i j loop done:	10011         00000000         000         0x 0002           00010         10000110         000         0x 0004           00010         10000010         000         0x 0006           00101         10000100         100         0x 0008           01110         00000001         000         0x 000A           00010         10000110         000         0x 000C           00010         10000010         000         0x 0010           01000         00000100         0x 0010           0x 0012         0x 0012
Data: 0xFF00 a(value = m) 0xFF02 b(value = n) 0xFF04 i 0xFF06 m 0xFF08 n 0xFF0A o 0xFF0C count		Stack: 0x1FFF

Team repo: set upped, goto link in M1 and join yellow