

# CS653: Functional Programming

## 2017-18 *II<sup>nd</sup>* Semester

### G Machine

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# Agenda

- ▶ Full Laziness
- ▶ G Machine

# Maximal Free Expressions

- ▶ **Free Subexpression:** A subexpression  $E$  of a  $\lambda$  abstraction  $L$  is free in  $L$  if all variables in  $E$  are free in  $L$ .
- ▶ **Maximal Free Subexpression (MFE):** A MFE of  $L$  is a free expression which is not a *proper* subexpression of another free subexpression of  $L$ .

# $\lambda$ -lifting using Maximal Free Expressions

Instead of  $\lambda$ -lifting a free variable, abstract out free subexpressions as parameters

```
f = g 4
g x y = y+(sqrt x)
(f 1) + (f 2)
```

```
$g1 sqrtx y = + y sqrtx
$g x = $g1 (sqrt x)
$f = $g 4
$Prog = + ($f 1) ($f 2)
```

When evaluating the expression, we get

```
$Prog
→ + ( 1) ( 2)
      |      |
      |      |→ ($g 4)
      |      |
      |      |→ ($g1 (sqrt 4))
      |      |
      |      |→ ($g1 (sqrt 4))
      |      |
      |      |→ ($g1 2)
      |      |
      |      |→ ($g1 2)
      |      |
      |      |→ 2
      |      |
      |      |→ 2
      |      |
      |      |→ 3
      |      |
      |      |→ 7
```

# Fully Lazy $\lambda$ -lifting and Letrecs

```
let
  f =  $\lambda x$ .letrec fac =  $\lambda n$ .(...)
      in x + (fac 1000)
in (f 3) + (f 4)
```

Without fully lazy  $\lambda$ -lifting:

```
$fac fac n = (...)
f x =letrec fac = $fac fac
    in + x (fac 1000)
$Prog = + (f 3) (f 4)
```

# Fully Lazy $\lambda$ -lifting and Letrecs

'Float' free letrec-s outward

```
letrec
    fac =  $\lambda n. (...)$ 
in let
    f =  $\lambda x. x + (fac\ 1000)$ 
in (f 3) + (f 4)
```

Now, we can have fully lazy  $\lambda$ -lifting:

```
$fac n = (...)  
$fac1000 = $fac 1000  
$f x = + x $fac1000  
$Prog = + ($f 3) ($f 4)
```

# G - Machine



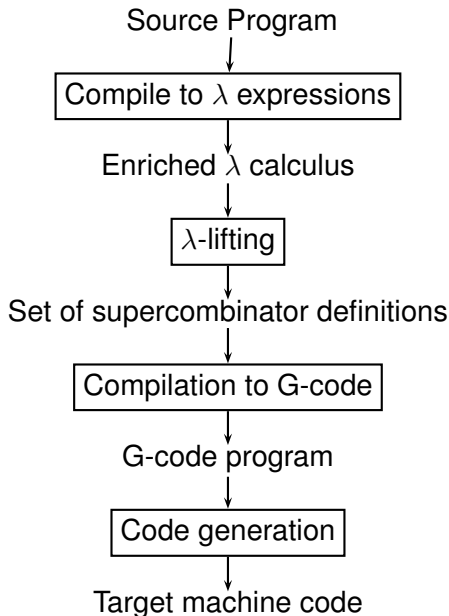
# G-machine

- ▶ Heart of functional programming paradigm is function application
- ▶ Need an efficient way to apply function body to arguments
- ▶ i.e., efficient graph reduction
- ▶ G-machine is an extremely fast implementation of graph reduction based on supercombinator compilation.
  - ▶ developed at Chalmers Institute of Technology, Sweden
  - ▶ by Johnsson and Augustsson

# G-code

- ▶ An *intermediate* code
- ▶ code for an *abstract sequential* machine
- ▶ Machine code for a specific target is generated from G-code

# Structure of G-machine Compiler



# Example of G-machine Execution

- ▶ Consider the following program in supercombinator form:

```
$from n = n : $from ($succ n)
```

```
$succ n = n + 1
```

```
$Prog = $from ($succ 0)
```

- ▶ We shall look at the code along with snapshots of execution.

# Startup code

```
begin  
pushglobal $Prog  
eval  
print  
end
```

◀ Beginning of program  
Push \$Prog onto stack  
Evaluate it  
Print the result  
End of program



# Startup code

```
begin  
pushglobal $Prog  
eval  
print  
end
```

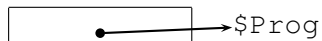
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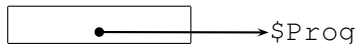


# \$Prog code

```
$Prog = $from ($succ 0)
```

```
pushint 0  
pushglobal $succ  
mkapp  
pushglobal $from  
mkapp  
update 1  
unwind
```

◀ Push constant 0  
Push function \$succ  
Construct (\$succ 0)  
Push function \$from  
Construct (\$from (\$succ 0))  
Update the root of the redex  
Initiate next reduction



# \$Prog code

```
$Prog = $from ($succ 0)
```

```
pushint 0
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```
pushglobal $succ
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```
mkapp
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```
pushglobal $from
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```
mkapp
```

```
update 1
```

```
unwind
```

Push constant 0

◀ Push function \$succ

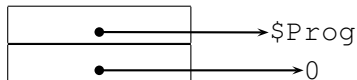
Construct (\$succ 0)

Push function \$from

Construct (\$from (\$succ 0))

Update the root of the redex

Initiate next reduction



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$Prog = $from ($succ 0)
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update 1
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```
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Push function \$succ

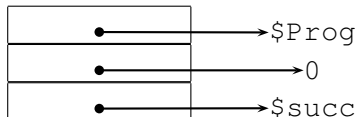
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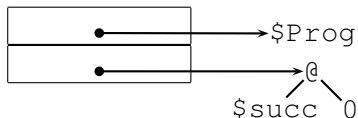
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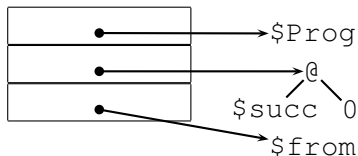
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Push constant 0

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pushglobal $succ
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Push function \$succ

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mkapp
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Construct (\$succ 0)

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pushglobal $from
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Push function \$from

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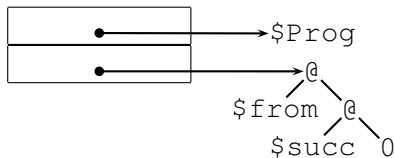
Construct (\$from (\$succ 0))

```
update 1
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◀ Update the root of the redex

```
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Initiate next reduction



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Push function \$succ

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Construct (\$succ 0)

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pushglobal $from
```

Push function \$from

```
mkapp
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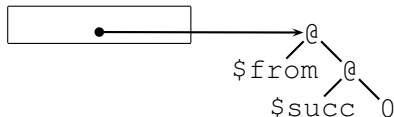
Construct (\$from (\$succ 0))

```
update 1
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Update the root of the redex

```
unwind
```

◀ Initiate next reduction





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```

```
update 1
```

```
unwind
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Push constant 0

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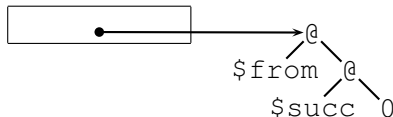
Construct (\$succ 0)

Push function \$from

Construct (\$from (\$succ 0))

Update the root of the redex

Initiate next reduction



travels down the spine, identifies the tip, manipulates pointers so that they start pointing to the arguments, jumps to the code of \$from.

## \$Prog code

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$Prog = $from ($succ 0)
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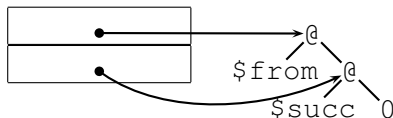
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## \$from code

```
$from = n:$from ($succ 0)
```

push 0	◀ Push n
pushglobal \$succ	Push function \$succ
mkapp	Construct (\$succ n)
pushglobal \$from	Push function \$from
mkapp	Construct (\$from (\$succ n))
push 1	Push n
cons	Construct (n:(\$from(\$succ n)))
update 2	Update the root of the redex
pop 1	Pop the parameter n
unwind	Initiate next reduction

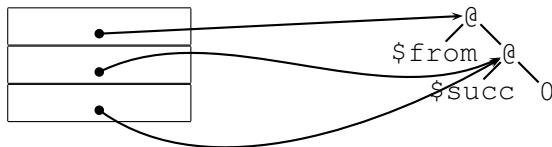


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```
push 0  
pushglobal $succ  
mkapp  
pushglobal $from  
mkapp  
push 1  
cons  
update 2  
pop 1  
unwind
```

Push n  
◀ Push function \$succ  
Construct (\$succ n)  
Push function \$from  
Construct (\$from (\$succ n))  
Push n  
Construct (n:(\$from(\$succ n)))  
Update the root of the redex  
Pop the parameter n  
Initiate next reduction



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$from = n:$from ($succ 0)
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```
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update 2
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```
pop 1
```

```
unwind
```

Push n

Push function \$succ

◀ Construct (\$succ n)

Push function \$from

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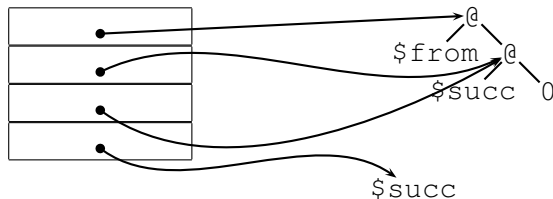
Push n

Construct (n:(\$from(\$succ n)))

Update the root of the redex

Pop the parameter n

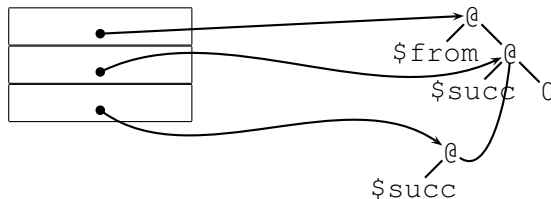
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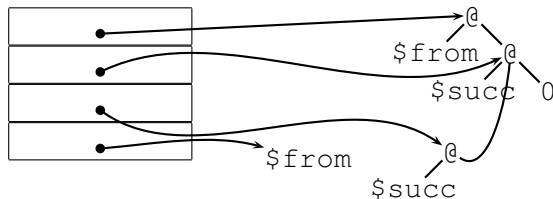
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Update the root of the redex

Pop the parameter n

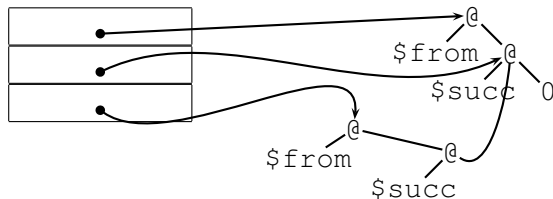
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Push n

Push function \$succ

Construct (\$succ n)

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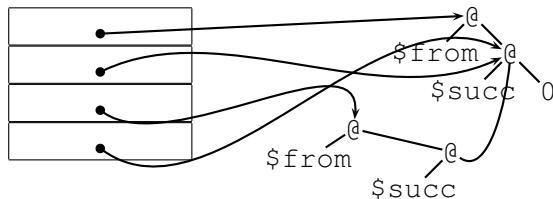
Push n

◀ Construct (n:(\$from(\$succ n)))

Update the root of the redex

Pop the parameter n

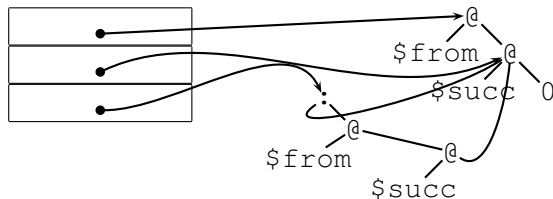
Initiate next reduction



## \$from code

```
$from = n:$from ($succ 0)
```

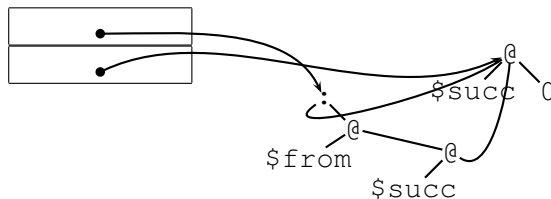
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```
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Push n

```
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## Push function \$succ

mkapp

Construct (\$succ n)

```
pushglobal $from
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## Push function \$from

mkapp

Construct (\$from (\$succ n))

```
push 1
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## Push n

cons

Construct (n:(\$from(\$succ n)))

update 2

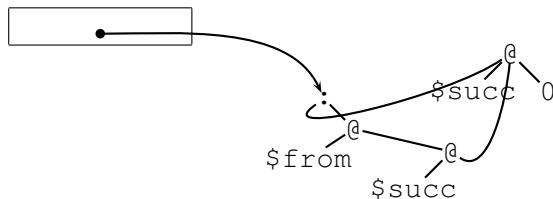
## Update the root of the redex

pop 1

## Pop the parameter n

unwind

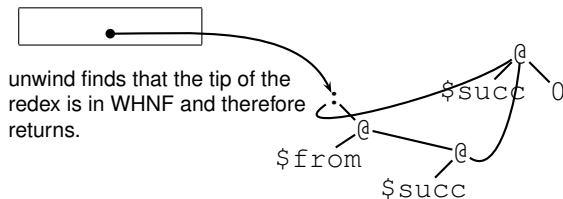
◀ Initiate next reduction



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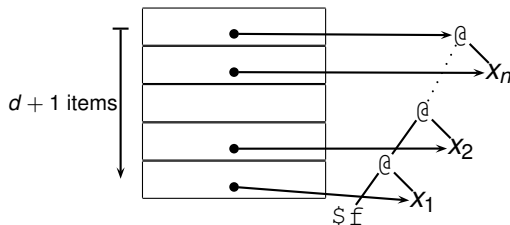
# The General Compilation Scheme

- ▶ The startup code is always

```
begin  
pushglobal $Prog  
eval  
print  
end
```

# Code for Supercombinator

- ▶ The code for supercombinator is given by  $\mathbb{F}$  compilation scheme  
 $\mathbb{F}[\$F\ x_1\ x_2 \dots x_n] = \dots \text{G-code for } \$F \dots$
- ▶ To define  $\mathbb{F}$ , note the situation just before entering a supercombinator body



# Context

- ▶ The distance from the top of the stack to the root of the redex is called the *context*
- ▶ Its length is denoted as  $d + 1$
- ▶ The initial value of  $d$  is  $n$   
and changes as the execution of the supercombinator body proceeds.



# Environment Function $\rho$

- ▶ The distance between  $i^{\text{th}}$  arg and the root of the redex is given by an environment function  $\rho$
- ▶  $\rho x_i = n + 1 - i$
- ▶ Thus, the distance of an argument  $i$  from the top of the stack is  $d - (\rho x_i)$ .
- ▶ This is how we keep track of the position of an argument from the top of the stack.
- ▶ Note that  $\rho$  is constant for a supercombinator, but  $d$  has to be tracked as compilation proceeds.

# FF Compilation Scheme

$$\mathbb{F}[\$F \ x_1 \ x_2 \ \dots \ x_n] = \text{globstart } \$F, n \\ \mathbb{R}[E] [x_1 = n, x_2 = n - 1, \dots x_n = 1] \ n$$

# $\mathbb{R}$ Compilation Scheme

$$\mathbb{R}[E] \rho d = \mathbb{C}[E] \rho d$$

update ( $d + 1$ )  
pop  $d$   
unwind

# $\mathbb{C}$ Compilation Scheme

- ▶  $\mathbb{C}[E] \rho d$
- ▶ The definition of  $\mathbb{C}[\ ]$  is given through case by case basis.

## $E$ is a constant

- ▶  $E$  is an integer  $i$  (or Boolean etc.)

$$\mathbb{C}[\![ i ]\!] \rho d = \text{pushint } i$$

- ▶  $E$  is a supercombinator or built-in function,  $f$

$$\mathbb{C}[\![ f ]\!] \rho d = \text{pushglobal } f$$

## $E$ is a variable

- ▶ The value of a variable  $x$  is in stack
- ▶ At offset  $(d - \rho x)$  from the top
- ▶ We would like to bring it to the top
- ▶ The G-code is

$$\mathbb{C}[\![ x ]\!] \rho d = \text{push } (d - \rho x)$$

## $E$ is an application

- ▶ Application ( $E_1 E_2$ )
- ▶ First construct an instance of  $E_2$
- ▶ Leave a pointer to the  $E_2$  instance at the top of the stack
- ▶ Construct an instance of  $E_1$  - but adjust for the  $E_2$ 's instance in the context ( $d$ )
- ▶ Make an application Cell on the top of the stack

$$\begin{aligned} \mathbb{C}[\![ E_1 E_2 ]\!] \rho d = & \mathbb{C}[\![ E_2 ]\!] \rho d \\ & \mathbb{C}[\![ E_1 ]\!] \rho (d + 1) \\ & \text{mkap} \end{aligned}$$

## $E$ is a let-expression

- ▶  $E$  is  $\text{let } x = E_x \text{ in } E_b$
- ▶ Construct instance of  $E_x$ , leave pointer to it at the stack top
- ▶ Augment  $\rho$  to say that  $x$  is at offset  $d + 1$  from the base of the context
- ▶ Construct instance of  $E_b$ , using updated  $\rho$  and  $d$
- ▶ Now  $E_b$  instance is at the top,  $E_x$  instance is below it. Since we do not want  $E_x$  any more, clean up.

$$\begin{aligned}\mathbb{C}[\![ \text{let } x = E_x \text{ in } E_b ]\!] \rho \ d &= \mathbb{C}[\![ E_x ]\!] \rho \ d \\ &\quad \mathbb{C}[\![ E_b ]\!] \rho[x = d + 1] \ (d + 1) \\ &\quad \text{slide 1}\end{aligned}$$