

CS653: Functional Programming

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G Machine

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Agenda

- ▶ G Machine (continued)

E is a letrec-expression

$$\mathbb{C}[\text{letrec } \mathcal{D}s \text{ in } E_b] \rho d = \mathbb{C}[\text{letrec}[\mathcal{D}s]] \rho' d' \\ \mathbb{C}[E_b] \rho' d' \\ \text{slide } (d' - d)$$

where

$$\mathcal{D}s \equiv \begin{bmatrix} x_1 = E_1 \\ \dots \\ x_n = E_n \end{bmatrix}$$

$$(\rho', d') = Xr[\mathcal{D}s] \rho d$$

$$Xr[\mathcal{D}s] \rho d = (\rho[x_1 = d + 1, \dots, x_n = d + n], d + n)$$

$$\begin{aligned} \mathbb{C}[\text{letrec}[\mathcal{D}s]] \rho d &= \text{alloc } n \\ &\quad \mathbb{C}[E_1] \rho d; \text{update } n \\ &\quad \dots \\ &\quad \mathbb{C}[E_n] \rho d; \text{update } 1 \end{aligned}$$

E is a letrec-expression

- ▶ $Xr[\]$ creates an environment in which $x_1 \dots x_n$ refer to the n locations above the top of the stack and lengthens the current context by n
- ▶ $\mathcal{C}letrec[\]$ actually allocates n locations and creates pointers pointing to them from the top of stack.
- ▶ Then it constructs the graphs corresponding to E_1 to E_n , and overwrites the allocated locations with these graphs.
- ▶ Finally E_b is evaluated in the letrec environment and slid into position

Code for Built-in Functions

E is cons

$\mathbb{C}[\text{cons}] \rho d = \text{pushglobal } \cons

- ▶ $\$cons$ code is

```
cons  
update 1  
unwind / return
```

- ▶ `cons` is the g-machine instruction that pops the top two elements of the stack and makes a cons out of them.
- ▶ The root of the redex is updated
- ▶ Since `unwind` will find `:` which is in WHNF, we can use a `return` instead.

E is +

$\mathbb{C}[\![+]\!] \rho \ d = \text{pushglobal } \$+$

► $\$+$ code is

```
push 1
eval
push 1
eval
add
update 3
pop 2
unwind / return
```

E is head

$\mathbb{C}[\text{head}] \rho d = \text{pushglobal } \head

► $\$head$ code is

```
eval      -- expose the cons cell
head      -- pick the head
eval      -- evaluate the head before updation
update 1  -- else there will be duplicate
           -- evaluation (see SPJ - section 12.4)
unwind    -- a return is not correct here (why?)
```


E is UNPACK_SUM_1_2

C[[UNPACK_SUM_1_2]] ρ d = pushglobal \$UNPACK_SUM_1_2

► \$UNPACK_SUM_1_2 code is

```

                                -- remember that $UNPACK_SUM_... is called with
                                -- a function f and a value v as arguments
push 1                          -- push the value
eval                            -- evaluate
testcons                        -- test.1.2 actually, tests whether the WHNF
                                -- matches the first alternative
jfail L                         -- jump on fail to L
push 1                          -- push the value once again
SEL12                          -- push the second component
push 2                          -- push the value yet again
SEL11                          -- push the first component
push 2                          -- push the function
mkap
mkap
update 3
pop 2
unwind
L: pushfail                    -- push the value fail, the previous fail has
                                -- been consumed by jfail and the surrounding
                                -- FATBAR needs to see a fail.

update 3
pop 2
return
```

E is if

$\mathbb{C}[\text{if}] \rho d = \text{pushglobal } \if

► \$if code is

```
push 0
eval
jfalse L1
push 1
jump L2
L1:  push 2
L2:  update 4
pop 3
unwind
```

E is [] (FATBAR)

$\mathbb{C}[\text{[]}] \rho d = \text{pushglobal } \FATBAR

► **\$FATBAR code is**

```
push 0
eval
jfail L1
push 0
jump L2
L1:  push 1
L2:  update 3
pop 2
unwind
```

A large example

Supercombinators:

```
$xxs f x xs = f x : $map f xs
$map f l     = if l == [] then []
              else unpack.1.2 ($xxs f) l
              [] error
```

\$xxs code

```
globstart $xxs, 3:
push 2      -- xs
push 1      -- f
pushglobal $map
mkap
mkap        -- ($map f x)
push 2      -- x
push 2      -- f
mkap        -- (f x)
pushglobal PACK_SUM_1_2
mkap
mkap        -- f x : $map f xs
update 4
pop 3
unwind
```

\$map code

```
globstart $map, 2
push error
push 2 -- l
push 2 -- f
pushglobal $xxs
mkap -- ($xxs f)
pushglobal $UNPACK_SUM_1_2
mkap
mkap
pushglobal $FATBAR
mkap
mkap -- $UNPACK_SUM_1_2 ($xxs f) 1
pushglobal $PACK_SUM_2_0 -- Nil
pushglobal $PACK_SUM_2_0
-- continued on next slide
```

\$map code (continued)

```
push 4
pushglobal $==
mkap
mkap    -- 1 == []
pushglobal $if
mkap
mkap
mkap    -- if ... [] error
update 3
pop 2
unwind
```

That's all, folks!