Bernstein-Vazirani Algorithm

Maria Harrison
Department of Computer Science
School of Engineering
Vanderbilt University
Nashville, TN, USA
maria.harrison@vanderbilt.edu

Abstract—The purpose of this paper is to explore the Bernstein-Vazirani algorithm, and gain a better understanding of its application in cryptography.

Keywords—oracle, quantum states, hidden shift, speed up, key encryption.

I. Introduction

In 1992, Ethan Bernstein and Umesh Vazirani developed the Bernstein-Vazirani algorithm to solve a hidden shift problem. This algorithm has numerous applications in cryptography, error correction, and optimization problems. The goal of this problem was to find a hidden shift.

II. BACKGROUND

The intention of the Bernstein-Vazirani algorithm is to utilize quantum mechanics's superposition and entanglement properties to efficiently determine an unknown string. This algorithm identifies a hidden binary string in one single run. In comparison to classical algorithms, this provides significant speed up times. The Bernstein-Vazirani algorithm provides a quantum algorithm with linear speedup in the amount of queries relative to the best classical algorithm.

A. Hidden Shift Problem

The hidden shift problem, also known as the Bernstein-Vazirani problem, states that when given a function there is a guarantee an undetermined shift exists; An unknown bit string exists. The function is guaranteed to be of type $a \cdot x$ and when given an input, the output will always be of $a \cdot x \pmod{2}$.

When f is a given function of bit strings and a is the phase shift x is a string of n bits n is the length of said bit string $\{x_0, x_1, x_2, ..., x_n\}$ The algorithm defines a function f(x) to determine the

unknown n-bit string a as followed:

$$f(x) = a \cdot x (mod 2)$$

B. The Inner-Product Oracle

This algorithm uses a quantum oracle that applies a given function to a superposition of all possible inputs of a quantum state. Often referred to as the inner-product oracle, the quantum oracle is an inner modulus function that receives queries and applies a single-qubit unitary transformation. The complexity of the Bernstein-Vazirani algorithm is in terms of the number of queries to an oracle.

C. The Quantum Circuit

This algorithm uses n qubits, and outputs each bit of the hidden string after measurements are made. This process

solves the Berstein-Vazirani problem in one query, rather than running multiple passes, at minimum of n runs for a string of length n. This provides significant linear speed up.

III. APPLICATIONS IN CRYPTOGRAPHY

Encryption, in the context of cryptography, is a process in which plain text or a piece of information is converted into cipher text that can only be decoded by the receiver for whom the information is intended. A cryptographic key is a string of randomized characters, often mathematically generated, paired with a cryptographic algorithm to secure data in a message. There are two general categories of cryptographic keys, symmetric and asymmetric encryption. There are four encryption algorithms, triple DES, RSA, Twofish, and AES.

A. Symmetric and Asymmetric Key Encryption

Symmetric key encryption encrypts and decrypts a message using the same key. While this process is very fast, it only provides confidentiality. The length of the key used is either 128 or 256 bits.

Asymmetric key encryption is based on public and private key encryption techniques; one key is used for encryption only and the other for decryption. This process is slower than symmetric, but provides confidentiality, authenticity, and non-repudiation. The length of the key used is 2048 bits or larger.

B. Cracking Cryptographic Key Encryptions

Cryptographic keys vary in length, spanning from 128 to 256 bits, to 2046 bits. The runtime of classical algorithms significantly limits its applications due to its minimum runtime of n runs, for a string of n length. If a classical algorithm was used to crack a cryptographic key of a 128 bit size, the runtime would be a minimum of 128 runs, just to fully traverse the key in its entirety. Bernstein-Vazirani's algorithm is highly applicable to cryptographic key-related attacks because of its significant speedup compared to classical algorithms.

IV. EXPLORATION OF BERNSTEIN-VAZIRANI

To further understand Bernstein-Vazirani's algorithm, I implemented it using quantum circuits, and used it to guess a Fernet encryption key. My coding environment consisted of:

- Visual Studio Code as an IDE
- Python 3.11 language
- Qiskit, Numpy, Pyenv Python submodules

A. Setting Up for a Quantum Circuit

First I generated a random int, in decimal type, whose binary format would be used as my hidden string.

Next, I generated a random int between 1, 39 which represented the number of qubits for my circuit.

Then, I performed a modulus function on my hidden string, checking to ensure that my randomly generated integer would be able to be represented with the randomly generated number of qubits.

When f is my inner product oracle function and r is a randomly generated integer used for my hidden

string n qubits is a randomly generated integer for the number

of qubits my circuit will use secret num is the inner product from a modulus

function The hidden string, number of qubits for my quantum circuit

is as followed:

```
decimal_num = random.randint(1, 1000)
print("Random number HIGH generated in decimal format:", decimal_num)
print("Random number in BINARY format:", bin(decimal_num).replace("00", ""))
               return decimal num
      def getRandomNumberQubits():
              n_qub = random.randint(1, 39)
print("Random number of qubits:", n_qub)
return n_qub
      r = getRandomNumberHigh()
n_qubits = getRandomNumberQubits()
Random number HIGH generated in decimal format: 864
Random number in BINARY format: 1101100000
Random number of qubits: 12
```

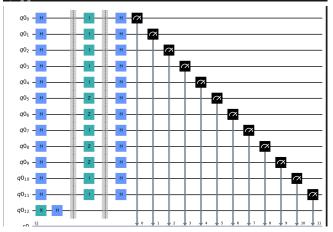
B. Creating Bernstein-Vazirani Algorithm

When bernstein vazirani alg(r, n) is the algorithm and

r is the hidden string in decimal format n is the number of qubits for the quantum circuit qr is a quantum register of size n + 1cr is a classical register of size n bvCirq is a QuantumCircuit . x is the output x gate . *h* is the output *h* gate if (f & 1 << i): is the inner product oracle

The newly created quantum circuit implementing Bernstein-Vazirani's algorithm is as followed:

```
def bernstein_vazirani_alg(r, n):
   qr = QuantumRegister(n+1)
   cr = ClassicalRegister(n)
   bvCirq = QuantumCircuit(qr, cr, name ="BernsteinVazirani")
   # setting output x gate gubits
   bvCirq.x(n)
   # setting output h gate qubits
   bvCirq.h(n)
    for qubit in range(n):
       bvCirq.h(qubit)
   bvCirq.barrier()
   # add return from oracle() func to circuit
    # bvCirq.append(oracle, range(n+1))
    for i in range(n):
       if (r & (1 <<i)):
           bvCirq.z(qr[i])
            bvCirq.id(qr[i])
   bvCirq.barrier()
    for q in range(n):
       bvCirq.h(qr[q])
    for i in range(n):
       bvCirq.measure(i, i)
    return bvCirq
```



C. Creating Quantum Simulator and Executing a Circuit

Next I created a quantum simulator and executed my newly created quantum circuit, implemented with Bernstein-Vazirani's Algorithm.

```
Running Bernstein Vazirani Algorithm
# creating quantum simulator
simulator = Aer.get_backend('qasm_simulator')
shots = 100
shots = 100 # running quantum circuit on simulator and calling .result() to get results of circuit when it runesult = execute(qCirc, backend=simulator, shots=shots).result()
plot_histogram(result.get_counts(qCirc))
```

D. Measuring and Testing Algorithm

Next I measured and tested my algorithm, and compared the hidden string with the algorithm's guess to determine

accuracy.

V. APPLYING BERNSTEIN-VAZIRANI TO GUESS A FERNET ENCRYPTION KEY

To further explore the application of Bernstein-Vazirani in cracking cryptographic keys, I simulated a very elementary level encryption message program.

A. Setting Up

To best replicate and demonstrate the difficulty and randomness of key encryption algorithms, I opted to create a completely random message using a random sentence generator. I utilized Python's Wonderword module to generate 4 complete sentences at random, with the criteria of a noun, verb, adjective, and direct object. I chose a random sentence generator to create my message because I wanted to test my algorithm to determine if it could succinctly behave with messages of varying length, which would result in encryption keys of different sizes.

```
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B. Encrypting a Random Message

Next, I encrypted my randomly generated message using Fernet and then continued the process of decoding the encryption key to use as my hidden string in the algorithm.

```
print("Fernet key encrypted: ", fernet)

v  0.0s

Fernet key encrypted: b'n1QbDGJxjXLzv5od6XZEi0nQN9CrJAqBhPbNqYLe3ys='
```

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```
import binascii
 #fernet key is a base64 encoded 32byte key
 print("Encrypted Key in Base64:")
 print(fernet)
 print("\n")
 b64_ascii = fernet.decode("ascii")
 print("Key of Base64 -> decoded in ASCII")
 print(b64_ascii)
 print("\n")
 b64_binary = binascii.a2b_base64(b64_ascii)
 print("Converting key of B64 -> Binary Array")
 print(b64 binary)
 print("\n")
 print("Converting key of binary array -> bytes")
 bits = ''.join(format(byte, '08b') for byte in b64_binary;
 print(bits)
 print("\n")
 print("Converting key of binary bytes -> decimal")
 converted =[b for b in b64_binary]
 print(converted)
 print("\n")
 print("Sorted number of ints from converted Base64 key")
 converted.sort()
 print(converted)
/ 0.0s
```

The following is the output of the encryption key after it has been decoded and converted into the proper format for my circuit. I opted to convert the final key from Base64 32 bytes to a byte array, and then transformed it into a list of integers. This was the best option for me to be able to randomly select an item from the list to be the hidden string for my algorithm.

```
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```

C. Creating a Simulator and Running Circuit
Next, I created a quantum simulator and executed my
algorithm, using a random part of my decoded Fernet
encryption key to serve as the hidden string. I really wanted
to emphasize the randomness of the hidden string to

demonstrate this algorithm's application capacity and runtime, hence why I opted to randomize a subset of my

decoded encryption key.

D. Measuring and Testing Algorithm

Measuring and testing my algorithm follows the same steps in section IV, part D. The results can be seen as follows:

```
printOutputAndScore(rf, format(fernet_guess), rf_secret)

✓ 0.0s

Random number generated at the beginning and converted to secret string: 84

Secret string: 1818180

Guess: ['0000000010101000']

Yay, the guess was CORRECT!
```

Lastly, after the algorithm finished running, I encoding the guessed string from the algorithm back into the original data encoding and format. I did this to demonstrate the capacity of applying Bernstein-Vazirani's algorithm in cryptography.

VI. CONCLUDING REMARKS

Bernstein-Vazirani's Algorithm poses significant benefits when applied to cryptography. In my example application in this study, I demonstrated the algorithm's ability to identify two independent, randomly created hidden strings. Using a random binary string from a decimal integer as a hidden string was meant to help understand and build the foundations of the Bernstein-Vazirani algorithm. A random binary string, from a subset of a decoded Fernet encryption key demonstrated the capacity of applying this algorithm to crack other encryption keys, such as RSA. The algorithm exploits quantum circuits and querying an oracle, thus surpassing classical algorithms in its runtime alone. The use of querying an inner product oracle poses significant advantages as the algorithm is not required to run n times, in which n is the length of the hidden string. Creating a small scale cryptographic key application in this study barely scrapes the iceberg of cryptography, but provides insight to just one of the applications of the Bernstein-Vazirani algorithm.

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