

# Network Security

## Project 1 - Hacking the Cipher

### Tools

Language: C/C++

Used Library : OpenSSL

### Program flow (P.S. 咖啡色字為使用到的 functions)

#### 1. 從.pub 檔讀入各個 Public Keys

```
RSA *PEM_read_RSA_PUBKEY(FILE *fp, RSA **x,  
                           pem_password_cb *cb, void *u)
```

將檔案中各個 Public keys 讀入, 並且存成 RSA 這個 Data structure, 透過 RSA->n 即可取得 BIGNUM 形態的 n value .

#### 2. 找尋兩個不互值的 n, 並求出其 GCD

```
int BN_gcd(BIGNUM *r, BIGNUM *a, BIGNUM *b, BN_CTX *ctx)
```

取 12 個 Public Key 的 n, 兩兩互相做 GCD .

#### 3. 令他們的 GCD 為 p, 利用 $n/p$ 求出 q

```
int BN_div(BIGNUM *dv, BIGNUM *rem, const BIGNUM *a,  
           const BIGNUM *d, BN_CTX *ctx)
```

## 4. 求出各別的 $\Phi(n)$

```
int BN_sub(BIGNUM *r, const BIGNUM *a, const BIGNUM *b)
```

```
int BN_mul(BIGNUM *r, BIGNUM *a, BIGNUM *b, BN_CTX *ctx)
```

利用  $\Phi(n) = (q - 1)(p - 1)$  這個公式求出小於  $n$  與  $n$  互值的數的個數。

## 5. 有了 $\Phi(n)$ , Publics Key, 即可求出 Private Key

```
BIGNUM *BN_mod_inverse(BIGNUM *r, BIGNUM *a,  
                        const BIGNUM *n, BN_CTX *ctx)
```

Private Key 為 Public Key 在模  $\Phi(n)$  下的乘法反元素, 因此有了 Public Key 與  $\Phi(n)$ , 即可推出 Private Key。

## 6. 將 Private Key 輸出成 PEM format (.pem 檔)

```
int PEM_write_RSAPrivateKey(FILE *fp, RSA *x,  
                             const EVP_CIPHER *enc, unsigned char *kstr,  
                             int klen, pem_password_cb *cb, void *u)
```

## Result

```
Ricky publicKeys ./main
Target Public Key Number : 3, 8

Common Factor :
D0B7C45918753A4C0A819AA2E7797124C24ACFF9A0F093F4A86641CFB73CC63BF564988ED9A86B0E6F716616C1042A7277EB4F81BA079526609AE74CBBB1C900B8D29E2A0B568E5881516704D6EF5
AD98DD23A86E47C3B727CC5BD078E022EF2F656CC240B8E8EB1F6602F52C74D880550A3AC6F2444FAA5759686B0AD7C69

Public Key 3's q :
C18E34C5788F13D16F45946EF99689A3EC54361F3E8FC0AEFA3BA110A618CF46DD67C93FFE550FAEC58D78377AF75D15C603F71DCCFE5F4B934FF59F708E57D63234C2FA486C384B528D185CEBE
97610262E059F182CDD0DED802B03709264D8D3F4F68E0F1276DBAF81398D3FBF4F8FF9A7337C7B9D80B0260FD2F88884D

Public Key 8's q :
D87CCB4BC31EDE5EE7EA65480DE4CE42C6828D281FEC8126C517A922E28BF00E2E5BF7F0662D95AC69AA8C06DC7C7D1B6D0EDB2E5BD8E8D8DB939D55EC28A0AB6FBF352481E32F17C14F915E89A022
D17872C45F894B93BE6F333320E20E56A9332D48C93157A36163F7BD4FAC4CC133A880F9A523F4A298BE6336570428B86D

Private Key 3 :
C47D0D820B578A3512292B069A1875140EDEEC55C7040674FA35F5EB80CF87E590D118590385F38A5666980210063D8C9425632E28403EBB503A89CE080DE5CF040819BDC802BB733572F8F8E992
B27DEAC2E5E133E80657E0E485077C5465B090DEAE2D4B455C388C0A0DE88888B4BA8B208A0412D4D22E63BA2C063F59F72B08F13754E12625B87D7F7E1E88ED5B5236C0F59D89A2E4193D982B8F333
28B117F71D0912ED178A66581D6FB8837E53A79D622D533DEB386321DB2A42F118384AC30B3A9965C303C68357099EC2BE249AFD715CC2ECCCE3E45832BCBD1FA1A162E56F130F22AECDD077D57B7C
EAC84298F5693310DD862C4A5764B39A659BD

Private Key 8 :
AA1DC41706D587D555D688B75FFDA71FCCD064E87B8C4C3512373523AC58C382C554E9B184101CCD50C169DCAE1150188AF342C720FD3A39DED2CDA2A5DA325FF45D7DB47C7230568824ACD1CAF2
88B32A96CF9F0705E6DB7CDDAF452BF47DD97FA4DC60B40630297B9999C68ADF1112BEC1043EC450988AD334A8988BB76B35E3808937FE7A12BF7A4DC9FD8B76FC365ECB238DC0BF15DA3D65F54CD
7CB33797D2CA7B8F257AFAC79BD4332B96688B2A9102D61F6584BD1ECF048C488D73B1BF8D3FECF2A76BAD13C221296260115879649BE0042EB3061F81D6F0E0E6EADD225CC417712D867465AE488B
1A84C47700C21860B50D16FFC00DFBEFE0195D
```

## Reference

- **Wiki OpenSSL – Manual: Bn**  
[https://wiki.openssl.org/index.php/Manual:Bn\(3\)](https://wiki.openssl.org/index.php/Manual:Bn(3))
- **Wiki OpenSSL – Manual: Rsa**  
[https://wiki.openssl.org/index.php/Manual:Rsa\(3\)](https://wiki.openssl.org/index.php/Manual:Rsa(3))
- **Wiki OpenSSL – Manual: Pem**  
[https://wiki.openssl.org/index.php/Manual:Pem\(3\)](https://wiki.openssl.org/index.php/Manual:Pem(3))
- **Understanding Common Factor Attacks:  
An RSA-Cracking Puzzle**  
<http://www.loyalty.org/~schoen/rsa/>