# Poster: An Algorithm and Tool to Infer Practical Postconditions

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#### **ABSTRACT**

Manually writing pre- and postconditions to document the behavior of a large library is a time-consuming task; what is needed is a way to automatically infer them. Conventional wisdom is that, if one has preconditions, then one can use the strongest postcondition predicate transformer (SP) to infer postconditions. However, we have performed a study using 2,300 methods in 7 popular Java libraries, and found that SP yields postconditions that are exponentially large, which makes them difficult to use, either by humans or by tools.

We solve this problem using a novel algorithm and tool for inferring method postconditions, using the SP, and transmuting the inferred postconditions to make them more concise.

We applied our technique to infer postconditions for over 2,300 methods in seven popular Java libraries. Our technique was able to infer specifications for 75.7% of these methods. Each of these inferred postconditions was verified using an Extended Static Checker. We also found that 84.6% of resulting specifications were less than 1/4 page (20 lines) in length. Our algorithm was able to reduce the length of SMT proofs needed for verifying implementations by 76.7% and reduced prover execution time by 26.7%.

## **KEYWORDS**

Specification inference, JML, predicate transformers

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#### 1 INTRODUCTION

There is a rich body of work on specification mining and inference [1-3]. A major category of these approaches analyze call sites of

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an API method to collect the set of predicates at each of these call sites and then use mining techniques such as frequent items mining to infer preconditions [4, 6]. Another body of work has focused on analyzing call sites to mine temporal patterns over API method calls, e.g. [5, 7]. These works have not focused on inferring postconditions.

The main contribution of this work is a postcondition inference technique to lower the cost of producing specifications. Our technique for method postcondition inference is based on the strongest postcondition predicate transformer: starting from a precondition (e.g., true) this uses the body of a method to produce a logical formula that can be converted into a specification as shown in Figure 1. A key aspect of our work is a novel algorithm for simplifying the often complex specifications that result from the strongest postcondition transformer (SP), which is discussed in Section 2. An illustration of our technique that shows the inferred specification for the method in Figure 1a is shown in Figure 1c.

# 2 TRANSMUTING SP WITH THE FAR ALGORITHM

The main problem with formulas computed by sp is that when the results are converted into specifications, they can become very large, deeply nested, and repetitive. This effect appears in the IF SP rule, which computes the disjunction of two different branch paths according to the equation (sp  $S_1(P \land B)$ )  $\lor$  (sp  $S_2(P \land \neg B)$ ). where P is a precondition, B is the branch conditions, and  $S_1$  and  $S_2$  are the body of the IF and ELSE branches respectively.

At each branch point in the program the branch's precondition is fed back into *sp* as an argument, which causes the number of specification cases in the resulting specification to grow exponentially. In this section we describe the Flattening and Recombination (FAR) algorithm that produces an equivalent specification that is shorter and more practical than the input specification. We show the effect of our algorithm on a small example in Figure 1a, 1b and 1c

### 3 EVALUATION

We designed a series of experiments designed to examine the effectiveness of Strongarm and FAR at inferring practical specifications. We selected a cross section of popular Java libraries: JUnit4 (JU4), JSON-Java (JJA), Commons-CSV (CSV), Commons-CLI (CLI),

```
//@ requires true;
                                              public normal_behavior
                                                                                          public normal_behavior
public int cmp(int a, int b){
                                                                                            requires !(a < b);</pre>
                                                requires true:
    int c = a;
                                                {|
    if (c < b) {
                                                  requires (c < b);
                                                                                               requires (a > b);
        return -1;
                                                                                               ensures \result == 1;
                                                  ensures true:
      else {
                                                  ensures \result == -1;
           (c > b) {
                                                  ensures c == a;
                                                                                               requires !(a > b);
                                                                                               ensures \result == 0;
             return 1;
                                                also
                                                                                             |}
        return 0;
                                                  requires !(c < b);</pre>
                                                                                            also
                                                  requires (c > b);
                                                                                               requires (a < b):
    }
 }
                                                   ensures true;
                                                                                               ensures \result == -1;
                                                  ensures \result == 1;
                    (a)
                                                  ensures c == a:
                                                                                                          (c)
                                                   requires !(c < b);
                                                  requires !(c > b);
                                                  ensures true;
                                                  ensures \result == 0;
                                                  ensures c == a:
                                              |}
                                             |}
                                                               (b)
```

Figure 1: (a) A simple comparison function. Using the standard definition of strongest postconditions, one obtains the postcondition shown in (b) for cmp. This postcondition contains a lot of repetition, tautologies, and is nearly twice the size of the original procedure. (c) The specification for cmp after being processed by our tool. Note that not only is this final specification shorter, it is lexically different than (b) and takes advantage of the fact that !(a < b) is true in multiple branches.

Commons-Codec (COD), Commons-Email (EMA), and Commons-IO (CIO). Code metrics pertaining to these libraries are summarized in Table 1.

Table 1: Code metrics for APIs used in evaluation.

API Name	SLOC	Methods	Files	Version
JUnit4	10,018	1,230	193	4.13
JSON-Java	3,201	200	18	20160212
Commons-CSV	1,501	158	10	1.4
Commons-CLI	2,666	194	22	1.3.1
Commons-Codec	6,607	509	60	1.10
Commons-Email	2,734	192	22	1.4
Commons-IO	9,836	955	115	2.5
Total	36,563	2,331	440	-

**Discussion.** Our evaluation determined that our technique was able to infer specifications for 75.7% of these methods. We also found that 84.6% of resulting specifications were less than 1/4 page (20 lines) in length. Our algorithm was able to reduce the length of SMT proofs needed for verifying implementations by 76.7% and reduced prover execution time by 26.7%. These results indicate the effectiveness of our technique at producing smaller and possibly easier to understand specifications. We plan to further investigate the impact of our technique on usability of inferred specifications in future work.

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