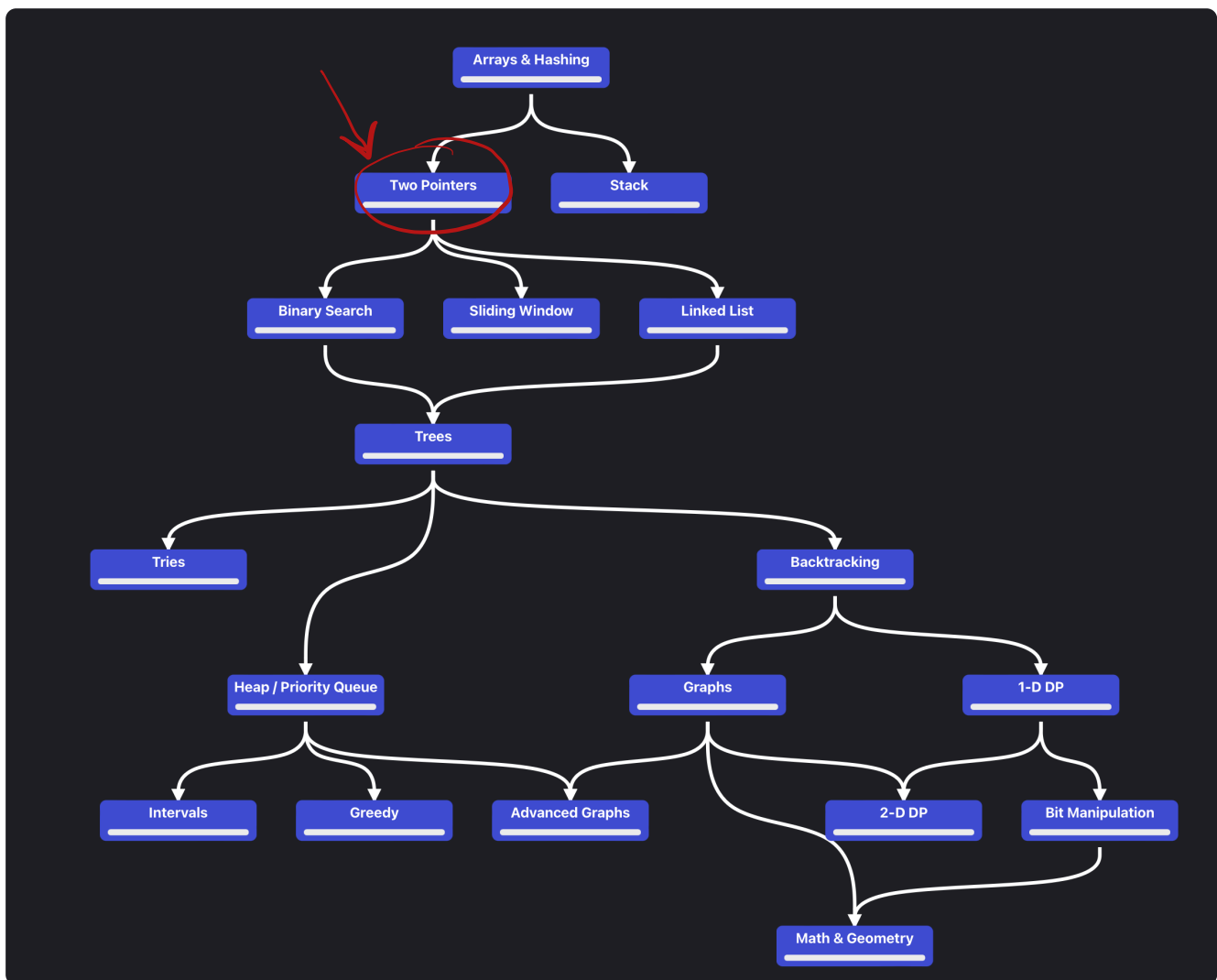


Two-Pointers



125. Valid Palindrome

Easy

8.7K

8.1K

☆

🔒

Companies

A phrase is a **palindrome** if, after converting all uppercase letters into lowercase letters and removing all non-alphanumeric characters, it reads the same forward and backward. Alphanumeric characters include letters and numbers.

Given a string `s`, return `true` if it is a **palindrome**, or `false` otherwise.

Example 1:

Input: `s = "A man, a plan, a canal: Panama"`
Output: `true`
Explanation: "amanaplanacanalpanama" is a palindrome.

Example 2:

Input: `s = "race a car"`
Output: `false`
Explanation: "raceacar" is not a palindrome.

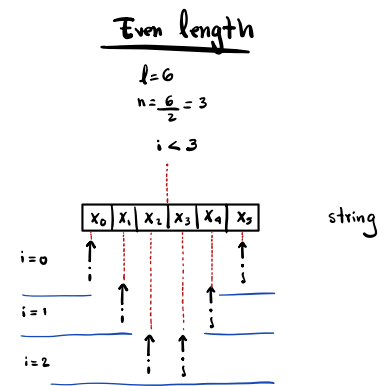
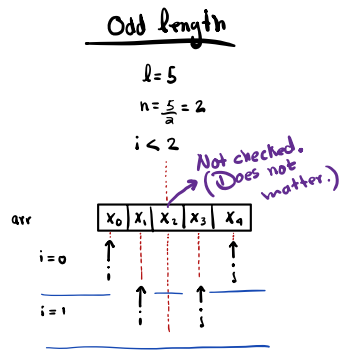
Example 3:

Input: `s = ""`
Output: `true`
Explanation: `s` is an empty string "" after removing non-alphanumeric characters. Since an empty string reads the same forward and backward, it is a palindrome.

125. Valid Palindrome

Clean up String (Remove non-alpha numeric characters)

Traverse half the string.
 Using two pointers to compare characters in a symmetric fashion.



Note: For:

$$l = 0 \Rightarrow n = \frac{0}{2} = 0$$

$$l = 1 \Rightarrow n = \frac{1}{2} = 0$$

For loop is Not executed
 ⇒ return TRUE

```

for ( i < l/2 ) {
  if ( arr[i] != arr[i] ) {
    return false;
  }
}

```

```

bool isPalindrome(string s) {
  string cleaned;
  for (char c : s) { // Create an entirely new string by
    // filtering out alpha numeric characters
    // and transforming them lower case
    if (isalnum(c)) {
      cleaned += tolower(c);
    }
  }

  long left = 0, right = cleaned.size() - 1;

  while (left < right) {
    if (cleaned[left] != cleaned[right]) {
      return false;
    }
    left++;
    right--;
  }

  return true;
}

```

Return the indices of the two numbers, `index1` and `index2`, added by one as an integer array `[index1, index2]` of length 2.

The tests are generated such that there is **exactly one solution**. You **may not** use the same element twice.

Your solution must use only constant extra space.

Example 1:

```
Input: numbers = [2,7,11,15], target = 9
```

Output: [1,2]

Explanation: The sum of 2 and 7 is 9. Therefore, $\text{index}_1 = 1$, $\text{index}_2 = 2$. We return [1, 2].

Example 2:

Input: numbers = [2,3,4], target = 6

Output: [1,3]

Explanation: The sum of 2 and 4 is 6. Therefore $index_1 = 1$, $index_2 = 3$. We return $[1, 3]$.

Example 3:

```
Input: numbers = [-1,0], target = -1
```

Output: [1,2]

Explanation: The sum of -1 and 0 is -1. Therefore $index_1 = 1$, $index_2 = 2$. We return [1, 2].

Target = 15

1	3	4	5	8	10	16
---	---	---	---	---	----	----

$$a[i] + a[j]$$

i	j	$1 + 16 = 17 > 15 \Rightarrow$ Decrease j
i	j	$1 + 10 = 11 < 15 \Rightarrow$ Increase i
i	j	$3 + 10 = 13 < 15 \Rightarrow$ Increase i
i	j	$4 + 10 = 14 < 15 \Rightarrow$ Increase i
i	j	$5 + 10 = 15 = 15$

Note: This might be an example of a greedy algorithm.

```
vector<int> twoSum(vector<int>& numbers, int target) {
    vector<int> ans;

    int i = 0; // Pointer to the first element
    long j = numbers.size() - 1; // Pointer to the last element

    while (i < j) {

        int num1 = numbers[i];
        int num2 = numbers[j];

        int dif = target - (num1 + num2); // compute difference

        if (dif > 0) { // We need a larger number -> increase left pointer
            i++;
        } else if (dif < 0) { // We need a smaller number -> decrease right pointer
            j--;
        } else { // dif == 0, then we've found our answer.
            ans.push_back(i + 1);
            ans.push_back(j + 1);
            break;
        }
    }

    return ans;
}
```

15. 3Sum

Hint

Medium

29.5K

2.7K



Companies

Given an integer array `nums`, return all the triplets `[nums[i], nums[j], nums[k]]` such that `i != j`, `i != k`, and `j != k`, and `nums[i] + nums[j] + nums[k] == 0`.

Notice that the solution set must not contain duplicate triplets.

Example 1:

Input: `nums = [-1,0,1,2,-1,-4]`**Output:** `[[-1,-1,2],[-1,0,1]]`**Explanation:**`nums[0] + nums[1] + nums[2] = (-1) + 0 + 1 = 0.``nums[1] + nums[2] + nums[4] = 0 + 1 + (-1) = 0.``nums[0] + nums[3] + nums[4] = (-1) + 2 + (-1) = 0.`The distinct triplets are `[-1,0,1]` and `[-1,-1,2]`.

Notice that the order of the output and the order of the triplets does not matter.

Example 2:

Input: `nums = [0,1,1]`**Output:** `[]`**Explanation:** The only possible triplet does not sum up to 0.

Example 3:

Input: `nums = [0,0,0]`**Output:** `[[0,0,0]]`**Explanation:** The only possible triplet sums up to 0.

```
vector<vector<int>> threeSum(vector<int>& nums) {  
  
    vector<vector<int>> ans;  
    sort(nums.begin(), nums.end()); // First sort  
  
    long size = nums.size();  
  
    for (int i = 0; i < size - 2; i++) { // i only gets up to third to last element  
  
        int j = i + 1; // middle pointer  
        long k = size - 1; // right pointer  
  
        if (i > 0 && (nums[i] == nums[i - 1])) { // Checked only after first iteration  
            continue; // We already found all the triples containing nums[i - 1], we must skip,  
        } // otherwise, this leads to repeated triple. Example: [-4,-1,-1,0,1,2]  
  
        while (j < k) {  
  
            int first = nums[i];  
            int second = nums[j];  
            int third = nums[k];  
  
            int sum = first + second + third;  
  
            if (sum < 0) { // We need a larger number, so increase left pointer.  
                j++;  
            } else if (sum > 0) { // We need a smaller number, so decrease right pointer.  
                k--;  
            } else { // sum == 0, we've found a triple  
  
                ans.push_back({first, second, third});  
  
                // Must update values right away.  
                j++;  
                k--;  
  
                // Check if current pointer values are the same as before. Skip all repeated values.  
                // We also have to check that j does not surpass k.  
                // Example [0,0,0,0]. Here, j goes out of bounds if we don't check j < k.  
                while(j < k && (nums[j] == nums[j - 1])) j++;  
                while(j < k && (nums[k] == nums[k + 1])) k--;  
  
            }  
        }  
    }  
    return ans;  
}
```

15. 3Sum

11. Container With Most Water

Medium 27.7K 1.6K

Companies

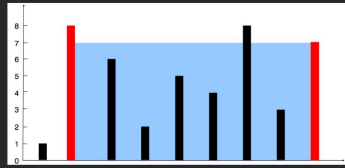
You are given an integer array `height` of length `n`. There are `n` vertical lines drawn such that the two endpoints of the i^{th} line are $(i, 0)$ and $(i, \text{height}[i])$.

Find two lines that together with the x-axis form a container, such that the container contains the most water.

Return the maximum amount of water a container can store.

Notice that you may not slant the container.

Example 1:



Input: `height = [1,8,6,2,5,4,8,3,7]`

Output: 49

Explanation: The above vertical lines are represented by array `[1,8,6,2,5,4,8,3,7]`. In this case, the max area of water (blue section) the container can contain is 49.

Example 2:

Input: `height = [1,1]`

Output: 1

```
long maxArea(vector<int>& height) {  
    long size = height.size();  
    long maxArea = 0;  
    long i = 0, j = size - 1; // Prepare pointers  
    while (i < j) {  
        long left = height[i];  
        long right = height[j];  
        long A = 0;  
        if (left < right) { // Calculate area depending on the smallest 'wall'  
            A = (j - i) * left;  
            i++; // Move pointer corresponding to the lower 'wall'  
        } else { // left >= right  
            A = (j - i) * right;  
            j--; // Move pointer corresponding to the lower 'wall'  
        }  
        if (maxArea < A) { // Update maxArea if necessary  
            maxArea = A;  
        }  
    }  
    return maxArea;  
}
```

11. Container With Most Water

42. Trapping Rain Water

Hard 30.2K 449

Companies

Given n non-negative integers representing an elevation map where the width of each bar is 1, compute how much water it can trap after raining.

Example 1:



Input: height = [0,1,0,2,1,0,1,3,2,1,2,1]
Output: 6
Explanation: The above elevation map (black section) is represented by array [0,1,0,2,1,0,1,3,2,1,2,1]. In this case, 6 units of rain water (blue section) are being trapped.

Example 2:

Input: height = [4,2,0,3,2,5]
Output: 9

```
int trap(vector<int>& height) {
    int n = height.size();
    if (n == 0) return 0; // handle empty vector

    int left = 0, right = n - 1; // These are pointers
    int leftMax = 0, rightMax = 0; // These are max heights found so far, going leftwards and rightwards

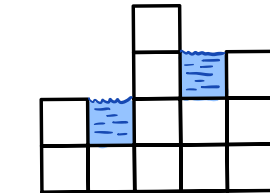
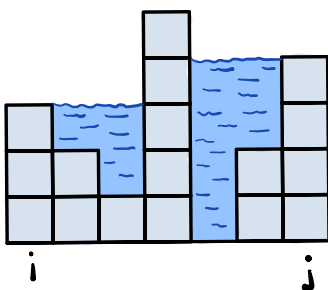
    int trappedWater = 0;

    while (left < right) {
        if (height[left] < height[right]) {
            // Left height is smaller than right height.
            // Then, there is the potential for the left height
            // to contain the water

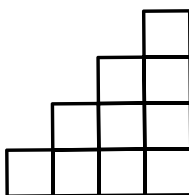
            if (height[left] >= leftMax)
                leftMax = height[left];
            else
                // height[left] < leftMax -> Then we are 100% sure there is water above this height.
                // Why? We know that to the right there is a height higher than the maxLeft.
                // And we know that the current height is lower than the maxLeft
                trappedWater += leftMax - height[left];
            left++;
        } else {
            // height[right] >= height[left]
            if (height[right] >= rightMax)
                rightMax = height[right];
            else
                trappedWater += rightMax - height[right];
            right--;
        }
    }

    return trappedWater;
}
```

42. Trapping Rain Water



i	j	$h[i] > \max L$ \downarrow $h[i] < h[j]$ $\left. \begin{matrix} \max L: 2 \\ \max R: 0 \end{matrix} \right\}$ water: 0
i	j	$h[i] < \max L$ \downarrow $h[i] < h[j]$ $\left. \begin{matrix} \max L: 2 \\ \max R: 0 \end{matrix} \right\}$ water: 1
i	j	$h[i] > \max R$ \downarrow $h[i] > h[j]$ $\left. \begin{matrix} \max L: 2 \\ \max R: 3 \end{matrix} \right\}$ water: 1
i	j	$h[i] < \max R$ \downarrow $h[i] > h[j]$ $\left. \begin{matrix} \max L: 2 \\ \max R: 3 \end{matrix} \right\}$ water: 2



i	j	$h[i] > \max L$ \downarrow $h[i] < h[j]$ $\left. \begin{matrix} \max L: 1 \\ \max R: 0 \end{matrix} \right\}$ water: 0
i	j	$h[i] > \max L$ \downarrow $h[i] < h[j]$ $\left. \begin{matrix} \max L: 2 \\ \max R: 0 \end{matrix} \right\}$ water: 0
i	j	$h[i] > \max L$ \downarrow $h[i] < h[j]$ $\left. \begin{matrix} \max L: 3 \\ \max R: 0 \end{matrix} \right\}$ water: 0