Sum lst 1

n (n)

$n-1$ (n)

$n-2$ (n)

2 (n)

1 (n)
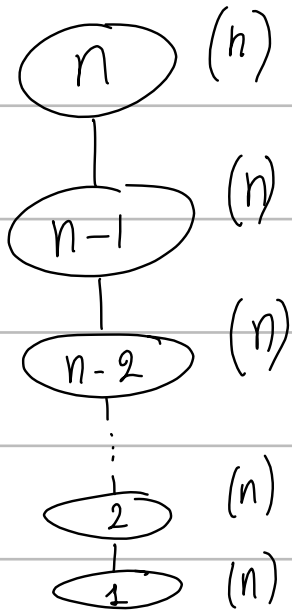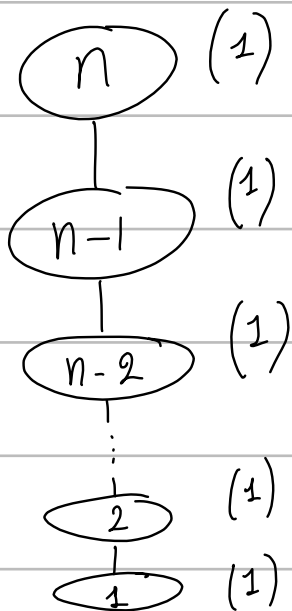
Each call costs $\theta(n)$ asymptotically because

of list slicing being performed at each call

Total cost $= n + n + n + \ldots + n = \theta(n^2)$

Sum lst 2

n (1)

$n-1$ (1)

$n-2$ (1)

2 (1)

1 (1)

✻ In this case n is high-low+1 and not the lenght of the list

Each call costs $\theta(1)$ asymptotically because

only primitive operations are performed

Total cost $= 1 + 1 + 1 + \ldots + 1 = \theta(n)$