## Chapter III Simulating random variables

## Sampling from U(0,1)

Need to simulate independent random variables uniformly distributed on [0, 1].

**Definition:** A sequence of pseudo-random numbers  $\{u_i\}$  is a deterministic sequence of numbers in [0,1] having the same statistical properties as a similar sequence of random numbers. Ripley 1987.

The sequence  $\{u_i\}$  is reproducible provided  $u_1$  is known.

A good sequence would be "unpredictable to the uninitiated".

#### 3.1 Generating Random Variables

- ► Inference techniques used so far have been based on simulation
- We now consider how to simulate X from  $f_X(x)$ .
- ▶ In semester 1 used MCMC but simpler methods needed in order to do MCMC.
- ightharpoonup Starting point: generate U from U[0,1] distribution
- ▶ Then consider transformation g(U) to obtain a random draw from  $f_X(x)$ .

How could we generate U[0,1] r.v.s with coin tosses?

## Congruential generators (D.H. Lehmer, 1949)

The general form of a congruential generator is

$$N_i = (aN_{i-1} + c) \mod M,$$
  
 $U_i = N_i/M$ , where integers  $a, c \in [0, M-1]$ 

If c = 0, it is called a multiplicative congruential generator (otherwise, mixed).

These numbers are restricted to the M possible values

$$0, \quad \frac{1}{M}, \quad \frac{2}{M}, \quad \dots, \quad \frac{M-1}{M}.$$

Clearly, they are rational numbers, but if M is large they will practically cover the reals in [0,1].

 $N_1$ : the **seed**. Can be re-set so you can reproduce same set of uniform random numbers. In R, use **set.seed(i)**, where i an integer.

As soon as some  $N_i$  repeats, say,  $N_i = N_{i+T}$ , then the whole subsequence repeats, i.e.  $N_{i+t} = N_{i+T+t}$ , t = 1, 2, ...

The least such T is called the *period*.

A good generator will have a long period.

The period cannot be longer than M and also depends on a and c.

Several useful Theorems exist concerning periods of congruential generators. For example, for c > 0, T = M if and only if

- 1. c and M have no common factors (except 1),
- 2.  $1 = a \pmod{p}$  for every prime number that divides M,
- 3.  $1 = a \pmod{4}$  if 4 divides M.

#### Lattice structure

Notice that for a congruential generator

$$N_i - aN_{i-1} = c - bM,$$

where b > 0 is an integer. Therefore,

$$U_i - aU_{i-1} = \frac{c}{M} - b.$$

The LHS lies in (-a, 1) since  $U_i \in [0, 1)$ .

Therefore, b can take at most a+1 distinct values.

If we plot points  $(U_{i-1}, U_i)$ , all the points will lie on at most a+1 parallel lines.

Usually M is chosen to make the modulus operation efficient, and then a and c are chosen to make the period as long as possible. Ripley suggests c = 0 or c = 1 is usually a good choice.

The NAG Fortran Library G05CAF

$$M = 2^{59}$$
  $a = 13^{13}$   $c = 0$ 

Another recommended one is

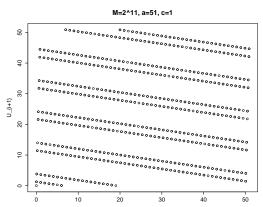
$$M = 2^{32}$$
  $a = 69069$   $c = 1$ .

so that

$$N_i = (69069N_{i-1} + 1) \mod 2^{32}$$

and

$$U_i = 2^{-32} N_i$$



All linear congruential generators exhibit this kind of lattice structure, not just for pairs  $(U_{i-1}, U_i)$ , but also for triples  $(U_{i-2}, U_{i-1}, U_i)$ , and in higher dimensions.

A good generator is expected to have *fine lattice structure*, that is, points  $(U_{i-k+1}, \ldots, U_{i-1}, U_i) \in [0,1)^k$  must lie on many hyperplanes in  $\mathbb{R}^k$  for all small k  $(k \ll M)$ .

#### RANDU - lattice structure

 $M = 2^{31}$ ,  $a = 2^{16} + 3 = 65539$ , and c = 0.

Once very popular, RANDU has eventually been found out to be a rather poor generator.

## Generation from non-U(0,1)

We have a sequence  $U_1, U_2, U_3, \ldots$  of independent uniform random numbers in [0, 1].

We want  $X_1, X_2, \ldots$  distributed independently and identically from some specified distribution.

The answer is to transform the  $U_1, U_2, \ldots$  sequence into  $X_1, X_2, \ldots$  sequence.

The idea is to find a function  $g(U_1, U_2, U_3, ...)$  that has the required distribution.

There are always many ways of doing this. A good algorithm should be quick because millions of random numbers may be required.

#### RANDU - lattice structure II

Let  $U_i = N_i/m$  then for this generator

$$U_{i+2} - 6U_{i+1} + 9U_i = k$$
 an integer.

Since  $0 \le U_i < 1$ 

$$-6 < U_{i+2} - 6U_{i+1} + 9U_i < 10.$$

Therefore  $k = -5, -4, \dots, -1, 0, +1, \dots, 9$ .

Hence k can take on 15 integer values only, and subsequently  $(U_{i-2}, U_{i-1}, U_i)$  must lie on at most 15 parallel planes.

This is an example of *coarse lattice structure*, unsatisfactory coverage of  $[0,1)^3$ .

#### 3.2 The inversion method

Let X be any continuous random variable and define  $Y = F_X(X)$ , where  $F_X$  is the distribution function of X:  $F_X(x) = P(X \le x)$ .

**Claim:**  $Y \sim U[0, 1]$ .

**Proof**  $Y \in [0,1]$  and the distribution function of Y is

$$F_Y(y) = P(Y \le y) = P(F_X(X) \le y)$$
  
=  $P(X \le F_X^{-1}(y)) = F_X(F_X^{-1}(y)) = y$ 

which is the distribution function of a uniform random variable on [0, 1].

So whatever the distribution of X,  $Y = F_X(X)$  is uniformly distributed on [0,1]. The inversion method turns this backwards. Let  $U = F_X(X)$ , then  $X = F_X^{-1}(U)$ .

▶ So to generate  $X \sim F_X$  take a single uniform variable U, and set  $X = F_X^{-1}(U)$ .

## Example: exponential distribution

Let  $X \sim Exp(1/\lambda)$  (mean  $\lambda$ ), i.e.

$$f(x) = \lambda^{-1} e^{-x/\lambda} \quad (x \ge 0)$$

$$F(x) = \int_0^x \lambda^{-1} e^{-z/\lambda} dz = [-e^{-z/\lambda}]_0^x = 1 - e^{-x/\lambda}.$$

Set  $U = 1 - e^{-X/\lambda}$  and solve for X

$$X = -\lambda \ln(1 - U).$$

Note that 1-U is uniformly distributed on [0,1], so we might as well use

$$X = -\lambda \ln U$$
.

**Question:** What are the limitations of the inversion method?

#### Discrete distributions - example

Let  $X \sim \text{Bin}(4, 0.3)$ . The probabilities are

$$P(X = 0) = .2401, \quad P(X = 1) = .4116, \quad P(X = 2) = .2646$$
  
 $P(X = 3) = .0756, \quad P(X = 4) = .0081.$ 

The algorithm says X=0 if  $0 \le U \le .2401$ , X=1 if  $.2401 < U \le .6517$ , X=2 if  $.6517 < U \le .9163$ , X=3 if  $.9163 < U \le .9919$ , X=4 if .9919 < U < 1.

Carrying out the binomial algorithm means the following. Let  $U \sim U(0,1)$ .

- 1. Test  $U \leq .2401$ . If true, return X = 0.
- 2. If false, test  $U \leq .6517$ . If true, return X = 1.
- 3. If false, test  $U \leq .9163$ . If true, return X = 2.
- 4. If false, test  $U \leq .9919$ . If true, return X = 3.
- 5. If false, return X = 4.

#### Discrete distributions

The inversion method works for discrete random variables in the following sense.

Let X be discretely distributed with possible values  $x_i$  having probabilities  $p_i$ . So

$$P(X = x_i) = p_i, \qquad \sum_{i=1}^{k} p_i = 1.$$

Then  $F_X(x) = \sum_{x_i \leq x} p_i$  is a step function.

Inversion gives  $X = x_i$  if  $\sum_{x_j < x_i} p_j < U \le \sum_{x_j \le x_i} p_j$  which clearly gives the right probability values.

▶ Think of this as splitting [0,1] into intervals of length  $p_i$ . The interval in which U falls is the value of X.

**Question:** What problems might we face using this method? Eg Consider a Poisson(100) distribution.

#### Discrete distributions - example

Consider the speed of this. The expected number of steps (which roughly equates to speed) is

$$1 \times .2401 + 2 \times .4116 + 3 \times .2646 + 4 \times .0756 + 4 \times .0081$$
  
=  $1 + E(X) - 0.0081 = 2.1919$ 

To speed things up we can rearrange the order so that the later steps are less likely.

- 1. Test  $U \leq .4116$ . If true return X = 1.
- 2. If false, test  $U \leq .6762$ . If true return X = 2.
- 3. If false, test U < .9163. If true return X = 0.
- 4. and 5. as before.

Expected number of steps:

 $1 \times .4116 + 2 \times .2646 + 3 \times .2401 + 4 \times (0.0956 + 0.0081) = 1.9959$ . Approximate 10% speed increase.

#### 3.3 Transformations

- (a) If  $U \sim U(0,1)$  set V = (b-a)U + a then  $V \sim U(a,b)$  where a < b.
- (b) If  $Y_i$  are iid exponential with parameter  $\lambda$  then

$$X = \sum_{i=1}^{n} Y_i = -\frac{1}{\lambda} \sum_{i=1}^{n} \log U_i = -\frac{1}{\lambda} \log \left( \prod_{i=1}^{n} U_i \right)$$

has a  $Ga(n, \lambda)$  distribution.

- (c) If  $X_1 \sim Ga(p, 1)$ ,  $X_2 \sim Ga(q, 1)$ ,  $X_1$  and  $X_2$  independent then  $Y = X_1/(X_1 + X_2) \sim Be(p, q)$ .
- (d) Composition: if

$$f = \sum_{i=1}^{r} p_i f_i$$

where  $\sum p_i = 1$  and each  $f_i$  is a density, then we can sample from f by first sampling I from the discrete distribution  $p = \{p_1, \dots, p_r\}$  and then taking a sample from  $f_I$ .

## 3.4 Rejection Algorithm

#### Fundamental Theorem of Simulation:

Simulating

$$X \sim f(x)$$

is equivalent to simulating

$$(X, U) \sim U\{(x, u) : 0 < u < f(x)\}.$$

Note that  $f(x, u) = \mathbb{I}_{0 < u < f(x)}$  so that

$$\int f(x, u) du = \int_0^{f(x)} du = f(x)$$

as required.

Hence, f is the marginal density of the joint distribution  $(X, U) \sim U\{(x, u) : 0 < u < f(x)\}.$ 

## The Box-Müller algorithm for the normal distribution

We cannot generate a normal random variable by inversion, because  $F_X$  is not known in closed form (nor its inverse). The Box–Müller method (1958). Let  $U_1, U_2 \sim U[0, 1]$ . Calculate

$$X_1 = \sqrt{-2 \ln U_1} \cos(2\pi U_2),$$
  
$$X_2 = \sqrt{-2 \ln U_1} \sin(2\pi U_2).$$

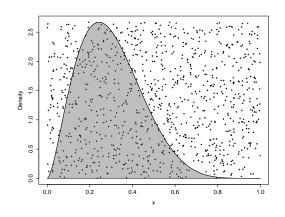
Then  $X_1$  and  $X_2$  are independent N(0,1) variables. The method is not particularly fast, but is easy to program and quite memorable.

## Rejection Algorithm Explained

The problem with this result is that simulating uniformly from the set

$$\{(x, u) : 0 < u < f(x)\}$$

may not be possible. A solution is to simulate the pair (X, U) in a bigger set, where simulation is easier, and then take the pair if the constraint is satisfied.



## Rejection: Uniform bounding box

Suppose that f(x) is zero outside the interval [a, b] (so that  $\int_a^b f(x) dx = 1$ ) and that f is bounded above by m.

- Simulate the pair  $(Y, U) \sim U[a, b] \times [0, m]$   $(Y \sim U[a, b], U \sim U[0, m]$  independently).
- Accept the pair if the constraint 0 < U < f(Y) is satisfied.

This results in the correct distribution for the accepted Y value, call it X.

$$\mathbb{P}(X \le x) = \mathbb{P}(Y \le x | U < f(Y))$$

$$= \frac{\int_a^x \int_0^{f(y)} dudy}{\int_a^b \int_0^{f(y)} dudy}$$

$$= \int_a^x f(y)dy.$$

Note: we can use the rejection algorithm even if we only know f upto a normalising constant (as is often the case in Bayesian statistics - see chapter 4).

#### Generalising the Rejection Idea

If the support of f is not finite, then bounding it within a rectangle will not work. Instead of using a box to bound the density f(x) (ie requiring f(x) < m for some constant m) we can use a function m(x) such that  $f(x) \le m(x)$  for all x.

Suppose the larger bounding set is

$$\mathcal{L} = \{ (y, u) : 0 < u < m(y) \}$$

then all we require is that simulation of a uniform from  $\mathcal L$  is feasible. Note

- ightharpoonup The closer m is to f the more efficient our algorithm.
- ▶ Because  $m(x) \ge f(x)$ , m cannot be a probability density. We write

$$m(x) = Mg(x)$$
 where  $\int m(x)dx = \int Mg(x)dx = M$ 

for some density g.

#### Example: Sampling from a beta distribution

Consider sampling from  $X \sim \text{Beta}(\alpha, \beta)$  for  $\alpha, \beta > 1$  which has pdf

$$f(x) = \frac{\Gamma(\alpha + \beta)}{\Gamma(\alpha)\Gamma(\beta)} x^{\alpha - 1} (1 - x)^{\beta - 1} \quad 0 < x < 1.$$

We note

$$f(x) \propto f_1(x) = x^{\alpha - 1} (1 - x)^{\beta - 1}$$
  $0 < x < 1$ 

and that  $M = \sup_{0 < x < 1} x^{\alpha - 1} (1 - x)^{\beta - 1}$  occurs at  $x = \frac{\alpha - 1}{\alpha + \beta - 2}$  (mode) and hence

$$M = \frac{(\alpha - 1)^{\alpha - 1}(\beta - 1)^{\beta - 1}}{(\alpha + \beta - 2)^{\alpha + \beta - 2}}.$$

The rejection algorithm is

- 1. Generate  $Y \sim U(0,1)$  and  $U \sim U(0,M)$ .
- 2. If  $U \leq f_1(Y) = Y^{\alpha-1}(1-Y)^{\beta-1}$  then let X = Y (accept) else go to 1 (reject).

#### Generalising the Rejection Idea II

This suggests a more general implementation of the fundamental theorem:

**Corollary:** Let  $X \sim f(x)$  and let g(x) be a density function that satisfies  $f(x) \leq Mg(x)$  for some constant  $M \geq 1$ . Then, to simulate  $X \sim f$ , it is sufficient to generate

$$Y \sim g$$
 and  $U|Y = y \sim U(0, Mg(y))$ 

and set 
$$X = Y$$
 if  $U < f(Y)$ .

#### **Proof:**

$$\mathbb{P}(X \in A) = \mathbb{P}(Y \in A | U \le f(Y))$$

$$= \frac{\int_{A} \int_{0}^{f(y)} \frac{du}{Mg(y)} g(y) dy}{\int \int_{0}^{f(y)} \frac{du}{Mg(y)} g(y) dy}$$

$$= \int_{A} f(y) dy$$

## The Rejection Algorithm

The rejection algorithm is usually stated in a slightly modified form:

#### Rejection Algorithm

If g is such that f/g is bounded, so there exists M such that  $Mg(x) \ge f(x)$  for all x then

- 1. Generate Y from density g, and U from U(0,1).
- 2. If  $U \le f(Y)/Mg(Y)$  set X = Y. Otherwise, return to step 1.

produces simulations from f

We keep sampling new Y and U until the condition is satisfied.

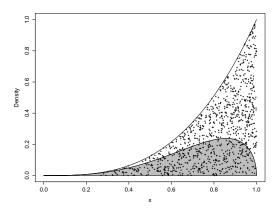
Exercise: Convince yourself that these two descriptions of the rejection algorithm are the same.

How to simulate Y with pdf  $g(y) = \alpha y^{\alpha-1}$ ?

- We note that the cdf of Y is  $G(y) = y^{\alpha}$ , 0 < y < 1.
- ▶ Therefore we can use inversion. Let  $Z \sim U(0,1)$  then solve  $Z = G(Y) = Y^{\alpha}$  and so  $Y = Z^{\frac{1}{\alpha}}$ .

#### Full algorithm is:

- 1. Generate  $U \sim U(0,1)$  and  $Z \sim U(0,1)$ . Let  $Y = Z^{\frac{1}{\alpha}}$ .
- 2. If  $U \leq (1-Y)^{\beta-1}$  then set X=Y else go to 1.



## Example: Sampling from a beta distribution revisited

Use rejection to sample from  $X \sim \text{Beta}(\alpha, \beta)$ . Let  $g(y) = \alpha y^{\alpha-1}$ , 0 < y < 1, then

$$\frac{f_1(x)}{g(x)} = \frac{(1-x)^{\beta-1}}{\alpha}$$
 is bounded if and only if  $\beta \ge 1$ 

Then 
$$M = \sup_{x} \left\{ \frac{f_1(x)}{g(x)} \right\} = \frac{1}{\alpha}$$
 occurs at  $x = 0$ .

- 1. Simulate Y with pdf  $g(y) = \alpha y^{\alpha-1}$ , 0 < y < 1 and  $U \sim U(0, 1)$ .
- 2. If  $U \le \frac{f_1(Y)}{Mg(Y)} = \frac{(1-Y)^{\beta-1}}{\left(\frac{1}{\alpha}\right)\alpha} = (1-Y)^{\beta-1}$  then set X = Y else go to 1.

## Efficiency of the rejection method

Each time we generate a (Y, U) pair,

$$\operatorname{Prob}(\operatorname{Reject}) = P\big(U \geq f(Y)/Mg(Y)\big) = 1 - \frac{1}{M}, \quad \operatorname{Prob}(\operatorname{Accept}) = \frac{1}{M}.$$

The number of tries until we accept Y is a geometric random variable with expectation M.

Note that M here must be calculated with the normalised density f, i.e.,  $M = \sup \frac{f(x)}{g(x)}$ .

If we used an unnormalised density  $f_1(x)$ , where  $\int f_1(x) dx = c$ , so that  $f(x) = \frac{1}{c} f_1(x)$ , then if we used

$$M = \sup \frac{f_1(x)}{g(x)}$$

the acceptance rate is

$$\mathbb{P}(\text{Accept}) = \frac{c}{M}$$

For maximum efficiency, we want M as small as possible, i.e.  $\sup f(x)/g(x)$  as small as possible. This means finding a g that

- (a) we can sample from efficiently, and
- (b) mimics f as closely as possible.

There are many good generators based on rejection from a well-chosen g.

#### Truncated distributions

Suppose we wish to sample X from the following distribution:

$$f_X(x) \propto \begin{cases} g_X(x) & \text{for } x \in A \\ 0 & \text{otherwise} \end{cases}$$

where  $g_X(x)$  is a known density that we can sample from, e.g.  $g_X(x)$  is the N(0,1) density, and  $A = [0, \infty)$ .

$$f_X(x) = \begin{cases} k g_X(x) & \text{for } x \in A \\ 0 & \text{otherwise} \end{cases}$$

where k is a normalising constant, given by

$$k^{-1} = \int_A g_X(x) dx$$

#### Rejection Example III

Let  $\theta$  have von Mises distribution with pdf

$$f(\theta) = \frac{\exp(k\cos\theta)}{2\pi I(k)} \quad 0 < \theta < 2\pi \quad (k \ge 0)$$

where I(k) is the normalising constant.

Let 
$$f_1(\theta) = \frac{1}{2\pi} \exp(k \cos \theta), 0 < \theta < 2\pi.$$

$$Y \sim U(0, 2\pi)$$
 so that  $g(y) = \frac{1}{2\pi}$ ,  $0 < y < 2\pi$ .

Then

$$M = \sup_{\theta} \left\{ \frac{f_1(\theta)}{g(\theta)} \right\} = \sup_{\theta} \{ \exp(k \cos \theta) \} = \exp k.$$

Let  $U \sim U(0, 1)$ .

If

$$U \le \frac{f_1(Y)}{Mg(Y)} = \frac{\exp(k\cos Y)}{2\pi \cdot \frac{1}{2\pi} \cdot \exp k} = \exp\left(k(\cos Y - 1)\right)$$

we accept  $\theta = Y$  otherwise reject.

$$f_X(x) \propto \begin{cases} g_X(x) & \text{for } x \in A \\ 0 & \text{otherwise} \end{cases}$$

Consider using rejection method to sample X from  $f_X(x)$ . We sample Y from the full (non-truncated) density  $g_X(x)$ .

$$\frac{f_X(x)}{g_X(x)} = \begin{cases} k & \text{if } x \in A \\ 0 & \text{otherwise} \end{cases}$$

So 
$$M = \sup_{x} \frac{f_X(x)}{g_X(x)} = k$$
.

Rejection algorithm: sample u from U[0,1] and y from  $g_Y(y)$ , and accept X = y if  $u \leq \frac{f_X(y)}{M g_Y(y)}$ .

But since

$$\frac{f_X(x)}{M g_X(x)} = \begin{cases} \frac{f_X(x)}{k g_X(x)} = 1 & \text{if } x \in A \\ 0 & \text{otherwise} \end{cases}$$

we will always have  $u \le \frac{f_X(y)}{M\,g_Y(y)}$  if  $y \in A$ , and  $u \ge \frac{f_X(y)}{M\,g_Y(y)}$  if  $y \notin A$ .

So we don't need to sample u. Can just do

- 1. generate y from  $g_Y(y)$
- 2. if  $y \in A$ , accept X = y
- 3. otherwise, return to step 1.

As usual, acceptance probability will be high if M is small, i.e  $\int_A g_Y(y) dy$  is near 1. So if the truncated region is large, rejection sampling will be inefficient.

## Sequential methods

We can obviously write

$$f(\mathbf{x}) = f_1(x_1)f_2(x_2|x_1)f_3(x_3|x_1,x_2)\dots$$

So we can first generate  $X_1$  from  $f_1$ . Then for that given value of  $X_1$ , generate  $X_2$  from  $f_2$ , and so on.

#### 3.5 Multivariate generators

Now suppose we want to generate a random vector  $\mathbf{X} = (X_1, \dots, X_p)$  from density  $f(\mathbf{x})$ . We can note the following simple points.

1. If the elements of X are to be independent, i.e.

$$f(\mathbf{x}) = f_1(x_1) f_2(x_2) \dots f_p(x_p),$$

then we can separately generate  $X_1$  from  $f_1$ ,  $X_2$  from  $f_2, \ldots, X_p$  from  $f_p$  using different uniforms.

- 2. Inversion no longer works as the theorem can't be generalised.
- 3. Rejection does work. If we can generate from g(x) (and g may be a product of independent components) and find  $M \ge \sup_{x} \frac{f(x)}{g(x)}$  and otherwise reject.

#### Example

REPLACE THIS EXAMPLE WITH Normal inverse gamma instead.

We wish to sample  $\{\theta, \phi\}$  from the density function

$$f(\theta, \phi) \propto \phi^3 \exp\{-1(1+5\phi)\theta^2 + 40\phi\theta - 81\phi\}.$$

Firstly, it can be shown that the marginal density of  $\phi$  is

$$f(\phi) \propto \phi^3 (1+5\phi)^{-1/3} \exp\left\{-\phi \left(1 + \frac{80}{1+5\phi}\right)\right\}.$$

Additionally, the conditional distribution of  $\theta | \phi$  is normal:

$$\theta | \phi \sim N\left(\frac{20\phi}{1+5\phi}, \frac{1}{2(1+5\phi)}\right).$$

To sample from  $f(\theta, \phi)$  we could first sample  $\phi$  from  $f(\phi)$ , using rejection sampling, then generate Z from N(0, 1), and finally set

$$\theta = \frac{20\phi}{1 + 5\phi} + \frac{1}{\sqrt{2(1 + 5\phi)}}Z.$$

#### Multivariate normal distributions

Consider generating **X** from  $N(\mathbf{m}, V)$ , for some non-diagonal matrix V, given a sample of independent standard normal random variables  $Z_1, Z_2, \ldots$ 

One technique involves the use of the **Cholesky square root** of the matrix V. For any (symmetric, square) positive definite matrix V, we can find a square root U, such that  $U^TU = V$ .

## 3.6 Importance sampling

In order to estimate an integral of the form  $\int h(x)f(x)dx$  we find that it is sometimes better to generate values not from the distribution f(x), but instead from some other distribution g(x) and to then account for this by using a weighting. This is the idea behind importance sampling. To introduce the idea we consider a simple example.

Multiple solutions for U, but upper triangular matrix U known as the Cholesky square root. To find the Cholesky square root of a matrix V in  $\mathbb{R}$ , type chol(V).

Define  $\mathbf{Z} \sim N(\mathbf{0}, I)$  (equal dimension to  $\mathbf{X}$ ), consider  $\mathbf{Y} = \mathbf{m} + U^T \mathbf{Z}$ . Then have  $\mathbf{Y}$  normally distributed, (each element of  $\mathbf{Y}$  is linear combination of the elements of  $\mathbf{Z}$ ), and

$$E(\mathbf{m} + U^T \mathbf{Z}) = \mathbf{m},$$

$$Var(\mathbf{m} + U^T \mathbf{Z}) = U^T I U = V,$$

(with I the identity matrix, the variance of  $\mathbf{Z}$ ). Hence to generate  $\mathbf{X}$ , we generate independent standard normal random variables  $\mathbf{Z}$ , and then transform them by  $\mathbf{m} + U^T \mathbf{Z}$  to obtain  $\mathbf{X}$ .

## Example of Monte Carlo/Importance Sampling

Let X be Cauchy  $f(x) = \frac{1}{\pi(1+x^2)}, -\infty < x < \infty$ . Let  $\theta = P(X > 2) = I = \int_2^\infty \frac{1}{\pi(1+x^2)} dx$  (= 0.1476). Use Monte Carlo Methods to estimate  $\theta$ .

(i) Generate n Cauchy variates,  $X_1, \ldots, X_n$ . Let  $Y_1$  be the number that are greater than 2,  $Y_1 = \sum \mathbb{I}_{X_i > 2}$ . Then  $Y_1 \sim B(n, \theta)$  so that

$$E(Y_1) = n\theta, \quad V(Y_1) = n\theta(1 - \theta)$$

$$\hat{\theta}_1 = \frac{Y_1}{n}$$

$$E(\hat{\theta}_1) = \frac{E(Y_1)}{n} = \frac{n\theta}{n} = \theta$$

and

$$V(\hat{\theta}_1) = \frac{V(Y_1)}{n^2} = \frac{n\theta(1-\theta)}{n^2} = \frac{\theta(1-\theta)}{n} = \frac{0.126}{n}.$$

## Example of Monte Carlo/Importance Sampling - II

# (ii) Note that $\theta = \frac{1}{2}P(|X| > 2)$ - we want to use this to reduce the variance of our estimator $\hat{\theta}$ .

Generate n Cauchy variates.

Let  $Y_2$  be the number that are greater than 2 in modulus then  $Y_2 \sim B(n, 2\theta)$ 

and 
$$\hat{\theta}_2 = \frac{1}{2} \frac{\hat{Y_2}}{n}$$

$$\implies E(\hat{\theta}_2) = \frac{1}{2} \frac{E(Y_2)}{n} = \frac{1}{2} \cdot \frac{n2\theta}{n} = \theta$$

and

$$V(\hat{\theta}_2) = \frac{V(Y_2)}{2^2 n^2} = \frac{n2\theta(1-2\theta)}{2^2 n^2} = \frac{\theta(1-2\theta)}{2n} = \frac{0.052}{n}.$$

## Example of Monte Carlo/Importance Sampling - IV

We can see that

$$\mathbb{E}(\hat{\theta}_3) = \frac{1}{2} - \frac{1}{n} \sum_{i=1}^n \int_0^2 \frac{2}{\pi(1+x^2)} dx = \frac{1}{2} - \mathbb{P}(0 < X < 2)$$

where  $X \sim \text{Cauchy}$ , so that it too is an unbiased estimator.

The variance of  $\hat{\theta}_3$  is Var(h(U))/n and we can see that

$$\mathbb{E}h(U) = \int_0^2 h(x) \frac{1}{2} dx = 0.5 - 0.1475 = 0.3525$$

$$\mathbb{E}h(U)^2 = \int_0^2 h(x)^2 \frac{1}{2} dx = \int_0^2 \frac{2}{\pi^2 (1 + x^2)^2} dx$$

$$= \frac{1}{\pi^2} \left[ \frac{x}{x^2 + 1} + \tan^{-1}(x) \right]_0^2 = 0.1527$$

Hence  $Var(h(x)) = 0.1527 - 0.3525^2 = 0.02851$  and thus

$$Var(\hat{\theta}_3) = \frac{0.02851}{n}$$

## Example of Monte Carlo/Importance Sampling - III

(iii) The relative inefficiency of these methods is due to generation of values outside the domain of interest  $[2, \infty)$ . Alternatively note we can write

$$\theta = \frac{1}{2} - \int_0^2 \frac{1}{\pi(1+x^2)} \, \mathrm{dx}.$$

This integral can be considered the expectation of  $h(X) = \frac{2}{\pi(1+x^2)}$  where  $X \sim U[0,2]$  as the density of U[0,2] is g(x) = 1/2.

An alternative method of evaluation of  $\theta$  is therefore

$$\hat{\theta}_3 = \frac{1}{2} - \frac{1}{n} \sum_{i=1}^n h(U_i)$$

where  $U_i \sim U[0,2]$ .

## Example of Monte Carlo/Importance Sampling - V

(iv) Finally, note that another possibility is to note that if  $y = \frac{1}{x}$ 

$$\theta = \int_{+2}^{\infty} \frac{1}{\pi(1+x^2)} dx = \int_{0}^{\frac{1}{2}} \frac{y^{-2} dy}{\pi(1+y^{-2})} = \int_{0}^{\frac{1}{2}} h(y) dy.$$

This can be seen as the expectation of  $h(X) = \frac{X^{-2}}{2\pi(1+X^{-2})}$  where  $X \sim U[0, \frac{1}{2}]$ . We can estimate this as

$$\hat{\theta}_4 = \frac{1}{n} \sum_{i=1}^n h(U_i)$$

where  $U_1, ..., U_n \sim U[0, 1/2]$ .

Again, we have  $\mathbb{E}\hat{\theta_4} = \theta$  and now

$$\mathbb{E}h(U)^{2} = \int_{0}^{1/2} h(x)^{2} \cdot 2dx = \frac{1}{4\pi^{2}} \left[ \frac{x}{x^{2} + 1} + \tan^{-1}(x) \right]_{0}^{1/2} = 0.02188$$
Hence  $\mathbb{V}ar(\hat{\theta}_{4}) = \frac{0.02188 - 0.1476^{2}}{2} = \frac{0.0000955}{2}$ 

#### Summary of Example

We found 4 unbiased estimators of  $\theta$ , each with a different variance.

$$Var(\hat{\theta}_1) = \frac{0.126}{n} \qquad Var(\hat{\theta}_2) = \frac{0.052}{n}$$
$$Var(\hat{\theta}_3) = \frac{0.02851}{n} \qquad Var(\hat{\theta}_4) = \frac{0.0000955}{n}$$

The best estimator is the one with the smallest variance, namely  $\hat{\theta}_4$ .

Compared with  $\hat{\theta}_1$ , the evaluation of  $\hat{\theta}_4$  requires  $\sqrt{(0.126/0.000955)} \approx 36$  times fewer simulations to achieve the same precision.

By carefully considering our simulation method we can hope to get more accurate estimates.

Estimate  $\hat{\theta}_2$  and  $\hat{\theta}_4$  are both types of importance sampling.

#### Some comments:

- ▶ g(x) is called the importance function, and  $w(X_i)$  are called the importance weights.
- ► The sum (1) will converge for the same reasons the Monte Carlo sum does.
- ▶ Notice that this sum is valid for any choice of the distribution g, as long as  $supp(f) \subseteq supp(g)$ .
- ▶ This is a very general representation that expresses the fact that a given integral is not intrinsically associated with a given distribution.
- $\triangleright$  Because very little restriction is put on the choice g, we can choose a distribution which is easy to sample from, and one which gives nice properties for the sum.

#### Importance Sampling

Consider calculating the integral

$$I = \mathbb{E}_f h(X) = \int h(\boldsymbol{x}) f(\boldsymbol{x}) d\boldsymbol{x}.$$

#### Importance sampling

Let  $X_1, X_2, \ldots, X_n$  be independently and identically distributed random variables with common density g(x).

Define  $w(\mathbf{x}) = f(\mathbf{x})/g(\mathbf{x})$ , so that

$$\mathbb{E}_g\{h(\boldsymbol{X}_i)w(\boldsymbol{X}_i)\} = \int h(\boldsymbol{x})w(\boldsymbol{x})g(\boldsymbol{x}) \; \mathrm{d}\boldsymbol{x} = \int h(\boldsymbol{x})f(\boldsymbol{x}) \; \mathrm{d}\boldsymbol{x} = \mathrm{I}.$$

Therefore

$$\hat{I} = \frac{1}{n} \sum_{i=1}^{n} w(\boldsymbol{X}_i) h(\boldsymbol{X}_i)$$
(1)

is an unbiased estimator of I.

## Cauchy example revisited

We can now understand the estimator  $\hat{\theta}_4$  in the Cauchy example. Recall that we want to estimate

$$\mathbb{E}\mathbb{I}_{X>2} = \int h(x)f(x)\mathrm{dx}$$

where  $h(x) = \mathbb{I}_{x>2}$  and  $f(x) = \frac{1}{\pi(1+x^2)}$ .

Noticing that for large x, f(x) is similar to the density

$$g(x) = 2/x^2 \text{ for } x > 2.$$

suggests g() might be a good importance density. We can sample from g by letting  $X_i = 1/U_i$  where  $U_i \sim U[0, \frac{1}{2}]$  (inversion method). Thus our estimator is

$$\hat{\theta} = \frac{1}{n} \sum_{i=1}^{n} h(x_i) \frac{f(x_i)}{g(x_i)} = \frac{1}{n} \sum_{i=1}^{n} \frac{x_i^2}{2\pi (1 + x_i^2)}$$
$$= \frac{1}{n} \sum_{i=1}^{n} \frac{u_i^{-2}}{2\pi (1 + u_i^{-2})} = \hat{\theta}_4$$

#### The variance of the estimator

Since the  $X_i$ s are iid,  $\operatorname{Var}(\hat{I}) = \frac{\sigma^2}{n}$ , where

$$\sigma^{2} = \operatorname{Var}_{g}\{h(\boldsymbol{X})w(\boldsymbol{X})\} = \mathbb{E}\{h(\boldsymbol{X})^{2}w(\boldsymbol{X})^{2}\} - \mathbb{E}\{h(\boldsymbol{X})w(\boldsymbol{X})\}^{2}$$

$$= \int h(\boldsymbol{x})^{2}w(\boldsymbol{x})^{2}g(\boldsymbol{x}) d\boldsymbol{x} - I^{2}$$

$$= \int \frac{h(\boldsymbol{x})^{2}f(\boldsymbol{x})^{2}}{g(\boldsymbol{x})} d\boldsymbol{x} - I^{2} \quad \text{since} \quad g(\boldsymbol{x}) = \frac{f(\boldsymbol{x})}{w(\boldsymbol{x})}.$$

We do not of course know  $\sigma^2$  in practice, but we can see that  $\hat{I}$  will be a better estimator if we can make  $w(\mathbf{X})$  less variable. Our objective, therefore, is to find a distribution  $g(\mathbf{x})$  that we know how to obtain independent samples from, and which mimics  $h(\mathbf{x})f(\mathbf{x})$  as closely as possible.

#### Unnormalised densites

Suppose we only know f upto a normalising constant, i.e., we know

$$f(x) = \frac{f_1(x)}{c}$$
 where  $c = \int f_1(x) dx$ 

We can still use importance sampling

Importance sampling with unnormalised densites Let  $X_1, X_2, ..., X_n$  be independently and identically distributed random variables with common density g(x). Define  $\tilde{w}(x) = f_1(x)/g(x)$ . Estimate I by

$$\hat{I} = \frac{\sum_{i=1}^{n} \tilde{w}(\boldsymbol{X}_i) h(\boldsymbol{X}_i)}{\sum_{i=1}^{n} \tilde{w}(\boldsymbol{X}_i)}$$

Alternatively, we can write this as

$$\hat{I} = \sum_{i=1}^{n} w_i h(\boldsymbol{X}_i)$$
 where  $w_i = \frac{\tilde{w}(X_i)}{\sum \tilde{w}(X_i)}$ 

#### Optimal choice of g

**Theorem** The choice of  $g = g^* = \frac{|h(x)|f(x)}{\int |h(z)|f(z)dz}$  minimises the variance of the estimator (1).

**Proof** We've seen that it is sufficient to minimise

$$\int \frac{h^2(\boldsymbol{x})f^2(\boldsymbol{x})}{g(\boldsymbol{x})} d\boldsymbol{x} = \mathbb{E}_{g}\left(\frac{h^2(X)f^2(X)}{g^2(X)}\right)$$

and using Jensen's inequality we can see that

$$\mathbb{E}_{g}\left(\frac{h^{2}(X)f^{2}(X)}{g^{2}(X)}\right) \ge \left(\mathbb{E}_{g}\left[\frac{|h(X)|f(X)}{g(X)}\right]\right)^{2}$$
$$= \left(\int |h(x)|f(x)dx\right)^{2}$$

and that this lower bound is achieved by choosing  $g = g^*$ . NB: We won't be able to calculate  $g^*$ ! But the theorem suggests

that choosing q to look like hf will be a good choice.

 $\frac{1}{n}\sum \tilde{w}(\boldsymbol{X}_i)$  is an unbiased estimator of c as

$$\mathbb{E}_g \tilde{w}(X) = \int \frac{f_1(x)}{g(x)} g(x) dx = \int f_1(x) dx = c.$$

When we use unnormalised densities,  $\hat{I}$  is a biased estimator of I, however it is possible to prove that we still have  $\hat{I} \to I$  almost surely as  $n \to \infty$ .

This will be important when we use importance sampling to estimate Bayesian quantities.

#### Effective sample size

How variable the weights are tells us how efficient our choice of g is. s

In the best case, where g = f, then  $\tilde{w}(X) = 1$  so that  $w_i = \frac{1}{n}$ , which is the case in plain Monte Carlo. In this case  $\mathbb{V}ar(w(X)) = 0$ .

If f and g are very different, then the weights will be very variable, and we can find that one or two particles  $(X_i)$  dominate the sum.

We often calculate the **effective sample size** 

$$ESS = \frac{1}{\sum w_i^2}$$

- ▶ In the best case,  $w_i = \frac{1}{n}$  and ESS= n so we have an effective sample size equal to the true sample size.
- ▶ The worst case is when one of the  $w_i = 1$  and all the others are equal to zero. Then ESS= 1, i.e., we effectively have only a single sample.

We want to choose g so that the ESS is large.

#### Antithetic variables - II

We need to find two estimators which are negatively correlated. This can be done as follows:

- If  $U \sim U[0,1]$  then  $1 U \sim U[0,1]$  also.
- ▶ If F is the distribution function of X then  $X_1 = F^{-1}(U)$ and  $X_2 = F^{-1}(1 - U)$  are both distributed according to F
- ▶ and  $\mathbb{C}ov(X_1, X_2) < 0$ .

#### Proof (non-examinable):

Let  $h(u) = F^{-1}(u)$ . Then h(u) is a non-decreasing function. We need to show

$$\mathbb{E}h(U)h(1-U) \le (\mathbb{E}h(U))^2$$

Let  $Q = \mathbb{E}h(U)$ . The since h is non-decreasing on [0,1]

$$h(0) \le Q \le h(1)$$

#### 3.7 Variance reduction techniques

#### Antithetic variables

The method of antithetic variables uses two correlated estimators and combines them to get an estimator with a lower variance (i.e. a better estimator).

Suppose we have two different estimators  $\hat{\theta}_1$  and  $\hat{\theta}_2$  of  $\theta$ ,

- ▶ with the same mean and variance
- ▶ but which are negatively correlated

Define  $\hat{\theta}_3 = \frac{1}{2}(\hat{\theta}_1 + \hat{\theta}_2)$ . Then

$$Var(\hat{\theta}_3) = \frac{1}{4}(Var(\hat{\theta}_1) + Var(\hat{\theta}_2) + 2Cov(\hat{\theta}_1, \hat{\theta}_2))$$
$$= \frac{1}{2}(Var(\hat{\theta}_1) + Cov(\hat{\theta}_1, \hat{\theta}_2))$$
$$< \frac{1}{2}Var(\hat{\theta}_1)$$

This is twice the cost of computing  $\hat{\theta}_1$  but the variance is more than halved!

Let 
$$f(y) = \int_0^y h(1-x)dx - Qy$$
 on  $[0,1]$   
Then  $f(0) = f(1) = 0$  and

$$f'(y) = h(1-y) - Q$$

is also a non-increasing function.

Since  $f'(0) = h(1) - Q \ge 0$  and  $f'(1) = h(0) - Q \le 0$  we must have

$$f(u) > 0$$
 on  $[0, 1]$ 

Therefore

$$0 \le \int_0^1 f(y)h'(y)dy = [fh]_0^1 - \int f'h(y)dy$$
$$= -\int_0^1 f'(y)h(y)dy$$

Therefore

$$\int_0^1 f'(y)h(y) = \int_0^1 h(y)(h(1-y) - Q)dy = \int_0^1 h(y)h(1-y)dy - Q^2 \le 0$$

Hence  $\int_0^1 h(y)h(1-y)dy \leq Q^2$  as required.

#### Cauchy Example Revisited

Above we used

$$\hat{\theta}_3 = \frac{1}{2} - \frac{2}{n} \sum_{i=1}^n \left[ \frac{1}{\pi (1 + u_i^2)} \right]$$

as an estimator of  $\mathbb{P}(X > 2)$  where  $X \sim Cauchy$ .

An estimator with a smaller variance can be found using antithetic variables

$$\frac{1}{2} \left( \frac{1}{2} - \frac{2}{n} \sum_{i=1}^{n} \left[ \frac{1}{\pi (1 + u_i^2)} \right] + \frac{1}{2} - \frac{2}{n} \sum_{i=1}^{n} \left[ \frac{1}{\pi (1 + (2 - u_i)^2)} \right] \right)$$

which gives

$$\hat{\theta}_{antithetic} = \frac{1}{2} - \frac{1}{n} \sum_{i=1}^{n} \left[ \frac{1}{\pi (1 + u_i^2)} + \frac{1}{\pi (1 + (2 - u_i)^2)} \right]$$

The for n=10 we find the variance of  $\hat{\theta}_3$  is  $2.7 \times 10^{-4}$  whereas the variance of  $\hat{\theta}_{antithetic}$  is  $5.5 \times 10^{-6}$  - a substantial improvement.