Organization of Digital Computers Lab EECS 112L

Lab 2 - Single Cycle RISC-V datapath and control

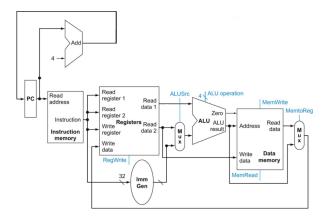
University of California, Irvine

Due date: February 26, 2018

Winter 2018

1 RISC-V datapath

System-verilog code for RISC-V processor is given to you. According to the module definition, RISC-V is the top module which includes Controller, ALUController and Datapath. The provided code for datapath implements the following figure:



RISC-V base instruction formats are as follows extracted from [1]:

31	25	24 2	0 19	15 14 1	2 11 7	6	0
funct	7	rs2	rs1	funct3	rd	opcode	R-type
•					•		
i	mm[11:0]		rs1	funct3	rd	opcode	I-type
imm[1]	1:5]	rs2	rs1	funct3	imm[4:0]	opcode	S-type
		imm[31:12	2]		rd	opcode	U-type

The RISC-V ISA keeps the source rs1 and rs2 and destination register rd in the same position in all the formats. Also, there are a further two variants of the instruction formats (B/J) based on the handling of immediate. RISC-V base instruction formats including immediate variants are represented below:

31 30 25 funct7	5 24 21 20 rs2	19 rs1	5 14 12 funct3	2 11 8 7	6 0 opcode	R-type
imm[1	1:0]	rsl	funct3	rd	opcode	I-type
imm[11:5]	rs2	rs1	funct3	imm[4:0]	opcode	S-type
imm[12] imm[10:5]	rs2	rs1	funct3	imm[4:1] imm[11]	opcode	B-type
	imm[31:12]		rd	opcode	U-type	
imm[20] imm[1	0:1] imm[11]	imm	19:12]	rd	opcode	J-type

To add all the instructions of different types to the datapath, the first step is to modify the Immediate Generator unit. According to the following figure, Imm Gen produces the 32-bit output differently according to the type of instruction. The current code for the Imm Gen only generates output for I-type and S-type instructions.

31	30	20	19 1	2	11	10 5	4	1 1	0	
		— inst[3	1] —			inst[30:25] i	nst[24:21]	inst[20]	I-immediate
		— inst[3	1] —			inst[30:25		inst[11:8]	inst[7]	S-immediate
	_	inst[31] —		ir	ıst[7]	inst[30:25		inst[11:8]	0	B-immediate
inst[31]	ins	st[30:20]	inst[19:12]			_	0 -	_		U-immediate
_	- inst[3	1] —	inst[19:12]	in	st[20]	inst[30:25] i	nst[24:21]	0	J-immediate

The following table has the values control signals for some of the instructions which might be useful for you.

Input or output	Signal name	R-format	ld	sd	beq
Inputs	I[6]	0	0	0	1
	I[5]	1	0	1	1
	I[4]	1	0	0	0
	I[3]	0	0	0	0
	I[2]	0	0	0	0
	l[1]	1	1	1	1
	I[0]	1	1	1	1
Outputs	ALUSrc	0	1	1	0
	MemtoReg	0	1	Х	Х
	RegWrite	1	1	0	0
	MemRead	0	1	0	0
	MemWrite	0	0	1	0
	Branch	0	0	0	1
	ALUOp1	1	0	0	0
	ALUOp0	0	0	0	1

2 Integer Computational Instructions

Integer computational instructions are either encoded as register-immediate operations using the I-type format or as register-register operations using the R-type format. No integer computational instructions cause arithmetic exceptions.

3 Integer Register-Immediate Instructions

ADDI adds the sign-extended 12-bit immediate to register rs1. Arithmetic overflow is ignored.

SLTI (set less than immediate) places the value 1 in register rd if register rs1 is less than the sign-extended immediate when both are treated as signed numbers, else 0 is written to rd. SLTIU is similar but compares the values as unsigned numbers.

ANDI, ORI, XORI are logical operations that perform bitwise AND, OR, and XOR on register rs1 and the sign-extended 12-bit immediate and place the result in rd.

	31	20 19	15 1	4 12	11 7	6	0
1	imm[11:0]	rs	1	funct3	$^{\mathrm{rd}}$	opcode	
	12	5		3	5	7	
	I-immediate[11:0]	sr	c A	ADDI/SLTI[U]	dest	OP-IMM	
	Limmediate[11:0]	or	c A	NDL/ORL/XO	RI doet	OP-IMM	

Shifts by a constant are encoded as a specialization of the I-type format. The operand to be shifted is in rs1, and the shift amount is encoded in the lower 5 bits of the I-immediate field. The right shift type is encoded in a high bit of the I-immediate. SLLI is a logical left shift (zeros are shifted into the lower bits); SRLI is a logical right shift (zeros are shifted into the upper bits); and SRAI is an arithmetic right shift (the original sign bit is copied into the vacated upper bits).

31	25 24 20	19 15	14 12	11 7	6 0
imm[11:5]	imm[4:0]	rs1	funct3	rd	opcode
7	5	5	3	5	7
0000000	shamt[4:0]	src	SLLI	dest	OP-IMM
0000000	shamt[4:0]	src	SRLI	dest	OP-IMM
0100000	chamt[4:0]	erc	SRAI	dest	OP-IMM

LUI (load upper immediate) is used to build 32-bit constants and uses the U-type format. LUI places the U-immediate value in the top 20 bits of the destination register rd, filling in the lowest 12 bits with zeros. AUIPC (add upper immediate to pc) is used to build pc-relative addresses and uses the U-type format. AUIPC forms a 32-bit offset from the 20-bit U-immediate, filling in the lowest 12 bits with zeros, adds this offset to the pc, then places the result in register rd.



4 Integer Register-Register Instructions

RV32I defines several arithmetic R-type operations. All operations read the rs1 and rs2 registers as source operands and write the result into register rd. The funct7 and funct3 fields select the type of operation.

ADD and SUB perform addition and subtraction respectively. Over ows are ignored. SLT and SLTU perform signed and unsigned compares respectively. OR, and XOR perform bitwise logical operations. SLL, SRL, and SRA perform logical left, logical right, and arithmetic right shifts on the value in register rs1 by the shift amount held in the lower 5 bits of register rs2.

31	25	24 20) 19	15 14	12	11	7 6	0
funct7		rs2	rs1		funct3	rd	opcode	
7		5	5		3	5	7	
0000000		src2	src1	AD	D/SLT/SLT	U dest	OP	
0000000		src2	src1	AN	D/OR/XOR	dest	OP	
0000000		src2	src1		SLL/SRL	dest	OP	
0100000		src2	src1	5	SUB/SRA	dest	OP	

5 Control Transfer Instructions

5.1 Unconditional Branches

The jump and link (JAL) instruction uses the J-type format, where the J-immediate encodes a signed offset in multiples of 2 bytes. The offset is sign-extended and added to the pc to form the jump target address. Plain unconditional jumps (assembler pseudo-op J) are encoded as a JAL with rd=x0.

The indirect jump instruction JALR (jump and link register) uses the I-type encoding. The target address is obtained by adding the 12-bit signed I-immediate to the register rs1, then setting the least-significant bit of the result to zero. The address of the instruction following the jump (pc+4) is written to register rd.

31	30	21 20	19	12	11	7 6		0
imm[20]	imm[10:1]	imm[1	1] imn	ı[19:12]	$_{\mathrm{rd}}$		opcode	
1	10	1		8	5		7	
	offset[2	20:1]			dest		$_{ m JAL}$	
31		20 19	15	14 12	11	7 6		0
	imm[11:0]		rs1	funct3	rd		opcode	
	12			3	5		7	
	offset[11:0]						JALR	

5.2 Conditional Branches

All branch instructions use the B-type instruction format. The 12-bit B-immediate encodes signed offsets in multiples of 2, and is added to the current pc to give the target address. The conditional branch range is 4 KiB.

Branch instructions compare two registers. BEQ and BNE take the branch if registers rs1 and rs2 are equal or unequal respectively. BLT and BLTU take the branch if rs1 is less than rs2, using signed and unsigned comparison respectively. BGE and BGEU take the branch if rs1 is greater than or equal to rs2, using signed and unsigned comparison respectively.

	31	30 25	24 20	19 15	14 12	11	8 7	6	0
	imm[12]	imm[10:5]	rs2	rs1	funct3	imm[4:1]	imm[11]	opcode	
١	1	6	5	5	3	4	1	7	
	offset	[12,10:5]	src2	src1	BEQ/BNE	offset[1	1,4:1]	BRANCH	
	offset	[12,10:5]	src2	src1	BLT[U]	offset[1	1,4:1]	BRANCH	
	offset	[12,10:5]	src2	src1	BGE[U]	offset[1	1,4:1	BRANCH	

6 Load and Store Instructions

The provided RISC-V processor implements Load and Store instructions. Load and store instructions transfer a value between the registers and memory. Loads are encoded in the I-type format and stores are S-type. The effective byte address is obtained by adding register rs1 to the sign-extended 12-bit offset. Loads copy a value from memory to register rd.

Stores copy the value in register rs2 to memory. The LW instruction loads a 32-bit value from memory into rd. LH loads a 16-bit value from memory, then sign-extends to 32-bits before storing in rd. LHU loads a 16-bit value from memory but then zero extends to 32-bits before storing in rd. LB and LBU are defined analogously for 8-bit values. The SW, SH, and SB instructions store 32-bit, 16-bit, and 8-bit values from the low bits of register rs2 to memory.

The table below summarizes the newly added instructions with the desired opcodes.

		Dase Instr	uction 5			
	imm[31:12]		rd	0110111	LUI	
	imm[31:12]			rd	0010111	AUIPC
	m[20 10:1 11 1:			rd	1101111	JAL
imm[11:		rs1	000	rd	1100111	JALR
imm[12 10:5]	rs2	rs1	000	imm[4:1 11]	1100011	BEQ
imm[12 10:5]	rs2	rs1	001	imm[4:1 11]	1100011	BNE
imm[12 10:5]	rs2	rs1	100	imm[4:1 11]	1100011	BLT
imm[12 10:5]	rs2	rs1	101	imm[4:1 11]	1100011	BGE
imm[12 10:5]	rs2	rs1	110	imm[4:1 11]	1100011	BLTU
imm[12 10:5]	rs2	rs1	111	imm[4:1 11]	1100011	BGEU
imm[11:		rs1	000	rd	0000011	LB
imm[11:		rs1	001	rd	0000011	LH
imm[11:		rs1	010	rd	0000011	LW
imm[11:		rs1	100	rd	0000011	LBU
imm[11:		rs1	101	rd	0000011	LHU
imm[11:5]	rs2	rs1	000	imm[4:0]	0100011	SB
imm[11:5]	rs2	rs1	001	imm[4:0]	0100011	SH
imm[11:5]	rs2	rs1	010	imm[4:0]	0100011	SW
imm[11:	0]	rs1	000	rd	0010011	ADDI
imm[11:		rs1	010	rd	0010011	SLTI
imm[11:		rs1	011	rd	0010011	SLTIU
imm[11:	0]	rs1	100	rd	0010011	XORI
imm[11:		rs1	110	rd	0010011	ORI
imm[11:	0]	rs1	111	rd	0010011	ANDI
0000000	shamt	rs1	001	rd	0010011	SLLI
0000000	shamt	rs1	101	rd	0010011	SRLI
0100000	shamt	rs1	101	rd	0010011	SRAI
0000000	rs2	rs1	000	rd	0110011	ADD
0100000	rs2	rs1	000	rd	0110011	SUB
0000000	rs2	rs1	001	rd	0110011	SLL
0000000	rs2	rs1	010	rd	0110011	SLT
0000000	rs2	rs1	011	rd	0110011	SLTU
0000000	rs2	rs1	100	rd	0110011	XOR
0000000	rs2	rs1	101	rd	0110011	SRL
0100000	rs2	rs1	101	rd	0110011	SRA
0000000	rs2	rs1	110	rd	0110011	OR
0000000	rs2	rs1	111	rd	0110011	AND

RV32I Base Instruction Set

7 Material to be submitted

- You should submit a report including how different instructions have be to implemented. The document contains all the necessary modifications in the datapath to add all the instructions.
- Modify the code to implement all the instructions. You can download the updated code for RISC-V from Github.

References

[1] Waterman, Andrew, et al., *The RISC-V Instruction Set Manual*, volume I: User-level ISA, version 2.0, EECS Department, University of California, Berkeley, Tech. Rep. UCB/EECS-2014-54 (2014).