# Concordia University

Department of Computer Science and Software

Engineering

# **SOEN 331:**

Formal Methods for Software Engineering

# Assignment 2 Richard Badir 40249566 Valentin Gornostaev 40211600

March 26, 2024

#### Problem 2

1. CL-USER 1 > (defvar \*airports\* '((YUL.Montreal)(LCY.London\_UK)(LHR.London\_UK)(MIL.M \*AIRPORTS\*

CL-USER 2 > (print \*airports\*)

((YUL.MONTREAL) (LCY.LONDON\_UK) (LHR.LONDON\_UK) (MIL.MILAN) (SFO.SAN\_FRANCISCO) (SD ((YUL.MONTREAL) (LCY.LONDON\_UK) (LHR.LONDON\_UK) (MIL.MILAN) (SFO.SAN\_FRANCISCO) (SD CL-USER 3 > \*airports\*

((YUL.MONTREAL) (LCY.LONDON\_UK) (LHR.LONDON\_UK) (MIL.MILAN) (SFO.SAN\_FRANCISCO) (SD CL-USER 3 > \*airports\*

CL-USER 4 >

- 2. monitored = dom(airports)
- 3. airports is a variable that holds one to many mappings of objects of type AIRPORT to objects of type CITY.  $dom(airports) = \mathbb{P}$  AIRPORT.  $codomain(airports) = \mathbb{P}$  CITY.  $airports = \mathbb{P}$  AIRPORT $\mapsto$  CITY
- 4. This should be an unordered structe, because order doesn't matter in this scenario, nor does it make sense. Looking up in an unordered structure would also be a lot faster than in an ordered one O(1) vs O(n). Assuming a Hashmap or any variation is not permitted, a set containing tuples of type (AIRPORT, CITY) should be used. There should be a restriction on the domain of this data structure, indicating no head of tuple can be a duplicate of another head. Lists of strict length 2 with types (AIRPORT, CITY) could also be used to replace the tuples.

5.

AirportSystem \_\_\_\_\_ codes:  $\mathbb{P}$  AIRPORT $\it airports: AIRPORT \rightarrow CITY$  --non-injective, surjective  $cities: \mathbb{P} CITY$  $codes = dom \ airports$  $deployed = ran \ airports$ 6. Success \_\_\_  $\Xi Airport System$ response!: MESSAGEresponse! = 'ok'CodeAlreadyInUse \_\_\_\_\_\_  $\Xi Airport System$ code?: AIRPORTresponse!: Message $code? \in dom\ airports$ 

Remove Everything After this line ————

 $response! = 'Code \ already \ in \ use'$ 

# Part 1: Temperature monitoring system with the Z specification

Consider a system called 'TempMonitor' that keeps a number of sensors, where each sensor is deployed in a separate location in order to read the location's temperature. Before the system is deployed, all locations are marked on a map, and each location will be addressed by a sensor. The formal specification of the system introduces the following three types:

 $SENSOR\_TYPE, LOCATION\_TYPE, TEMPERATURE\_TYPE$ 

We also introduce an enumerated type MESSAGE which will assume values that correspond to success and error messages.

Provide a formal specification in Z, with the following operations:

- DeploySensorOK: Places a new sensor to a unique location. You may assume that some (default) temperature is also passed as an argument.
- ReadTemperatureOK: Obtain the temperature reading from a sensor, given the sensor's location.

Provide appropriate success and error schemata to be combined with the definitions above to produce robust specifications for the following interface:

- DeploySensor,
- ReadTemperature.

#### Solution:

```
TempMonitor
deployed': \mathbb{P} SENSOR_TYPE
map: SENSOR\_TYPE \rightarrow LOCATION\_TYPE
                                             --partial bijective
read: SENSOR\_TYPE \rightarrow TEMPERATURE\_TYPE
deployed = dom map
deployed = dom read
DeploySensorOK \_
\Delta TempMonitor
sensor?: SENSOR\_TYPE
location?: LOCATION\_TYPE
temperature?: TEMPERATURE\_TYPE
sensor? \notin deployed
location? \not\in ran map
deployed' = deployed \cup \{sensor?\}
map' = map \cup \{sensor? \mapsto location?\}
read' = read \cup \{sensor? \mapsto temperature?\}
ReadTemperatureOK _____
\Xi TempMonitor
location?: LOCATION\_TYPE
temperature!: TEMPERATURE\_TYPE
location? \in ran map
temperature! = read(map^{-1}(location?))
Success _____
\Xi TempMonitor
response!: MESSAGE
response! = 'ok'
```

SensorAlreadyDeployed \_\_\_\_\_

 $\Xi TempMonitor$ 

 $sensor?: SENSOR\_TYPE$ 

 $\frac{response!: Message}{sensor? \in deployed}$ 

response! = 'Sensor deployed'

LocationAlreadyCovered \_\_\_\_\_

 $\Xi TempMonitor$ 

 $location?: LOCATION\_TYPE$ 

 $\frac{response!: Message}{location? \in ran \ map}$ 

response! = 'Location already covered'

Location Unknown \_\_\_\_\_

 $\Xi TempMonitor$ 

 $location?: LOCATION\_TYPE$ 

response!: Message

 $location? \not\in ran map$ 

response! = 'Location not covered'

DeploySensor =

 $(DeploySensorOK \land Success) \oplus (SensorAlreadyDeployed \lor LocationAlreadyCovered)$ 

 $ReadTemperature \triangleq (ReadTemperatureOK \land Success) \oplus LocationUnknown$ 

### Part 2: A booking system with the Object Z specifica-

#### tion

We introduce the basic types [Person, SeatType]. We also introduce an enumerated type Message which will assume values (feel free to define your own) that correspond to success and error messages. Consider a system to book seats for a theater play. A customer can book a single seat, and a seat can only accommodate a single customer. The booking system keeps a log of the customers that have booked a seat. The system publishes a plan of the theater and it allows customers to access it online and make a booking or cancel a booking.

### Class Booking

Define a formal specification in Object-Z for class *Bookingt* to support the following operations:

- BookOK: Reserves a seat for a given customer.
- CancelOK: Frees a seat for a given customer.

You will also need to provide appropriate success and error schemata to be combined with the definitions above to produce *robust specifications* for the following interface:

- Book, and
- Cancel.

## Class Booking2

Subclassify Booking to introduce class Booking2 that behaves exactly like Booking, while introducing the following operations:

• GetNumberOfCustomers returns the total number of customers who have made a booking.

• ModifyBookingOK assigns an existing customer to a different seat. Provide any additional schema(ta) in order to extend the interface to include a robust operation ModifyBooking.

The extended interface will now include operations

- GetNumberOfCustomers, and
- ModifyBooking.

#### Solution:

The main functionality for class Booking together with the success and error schemata are given below:

```
\_Booking _____
\upharpoonright (Book, Cancel)
  booked: \mathbb{P} Person
  booking: Person \rightarrowtail SeatType
  capacity: \mathbb{N}
  count: \mathbb{N}
  booked = dom\ booking
  capacity > 0
  count \ge 0
  INIT ____
  booked = \emptyset
  capacity = 100
  count = 0
  _ BookOK _____
  \Delta(booking, count)
  customer?: Person
  seat?: SeatType
  customer? \not \in booked
  seat? \not\in ran SeatType
  count < capacity
  booking' = booking \cup \{customer? \mapsto seat?\}
  count' = count + 1
  _ CancelOK _____
  \Delta(booking, count)
  customer?: Person
  customer? \in booked
  count > 0
  booking' = \{customer?\} \lhd booking
  count' = count - 1
```

```
Booking/cont._____
  Success ____
 response! : Message
 response! = 'ok'
 \_CustomerExists
 customer?: Person
 response! : Message
 customer? \in booked
 response! = 'Customer already exists'
 \_CustomerUnknown \_\_
 customer?: Person
 response!: Message
 customer? \not \in booked
 response! = 'Customer unknown'
 \_SeatTaken\_
 seat?: SeatType
 response!: Message
 seat? \in ran\ booking
 response! = 'Seat \ taken'
 \_ TheaterFull \_
 response!: Message
 count = capacity
 response! = 'Theater full'
 _{-} TheaterEmpty _{-}
 response!: Message
 count = 0
 response! = 'Theater empty'
Book \triangleq (BookOK \land Success) \oplus CustomerExists \oplus (SeatTaken \lor TheaterFull)
Cancel = (CancelOK \land Success) \oplus (CustomerUnknown \lor TheaterEmpty)
```