

NrOS: Effective Replication and Sharing in an Operating System

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OSDI'21

Presented by Yuchao Qian

Outline

- Background & Overview
- Node Replication
- NrOS Design
- Evaluation
- Conclusion

Background

- Increasing CPU core count
- Non-uniform memory access (NUMA)
- Elaborate concurrent data structures
 - fine-grained locking
 - read-copy-update (RCU)
 - good performance but increased complexity

Background

- Big kernel lock works for microkernel^[1]
 - does not target NUMA
- Multikernel: per-core kernels, communicating via message passing^[2]
 - scales well
 - too much complexity and overhead for hosts with shared mem

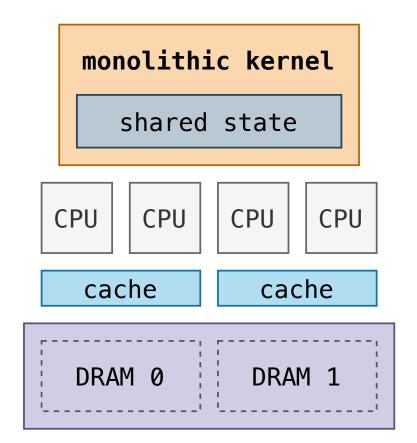
^{1.} For a Microkernel, a big lock is fine. APSys'15.

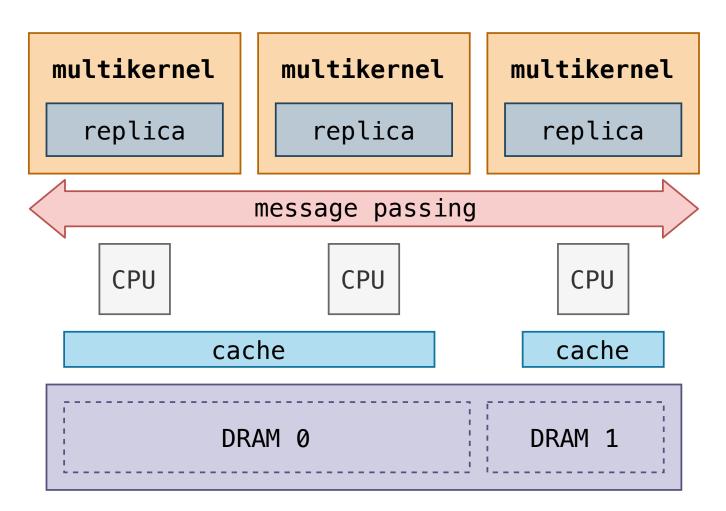
^{2.} The Multikernel: A New OS Architecture for Scalable Multicore Systems. SOSP'09.

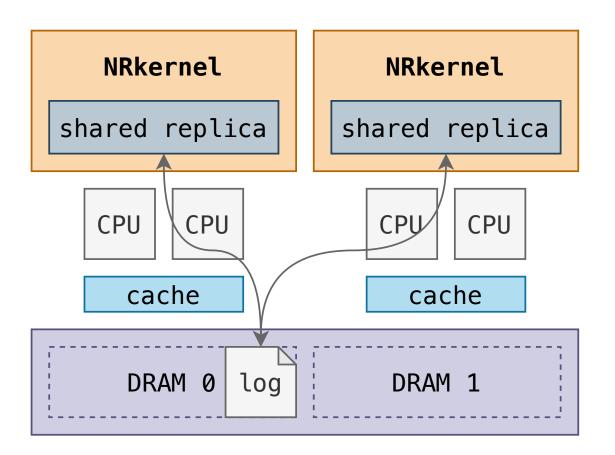
Overview

- NRkernel
 - node replication^[1]
 - kernel state replica per NUMA node
 - read local replica concurrently
 - mutate by shared operation log, serially
- NrOS: an NRkernel

Overview







Outline

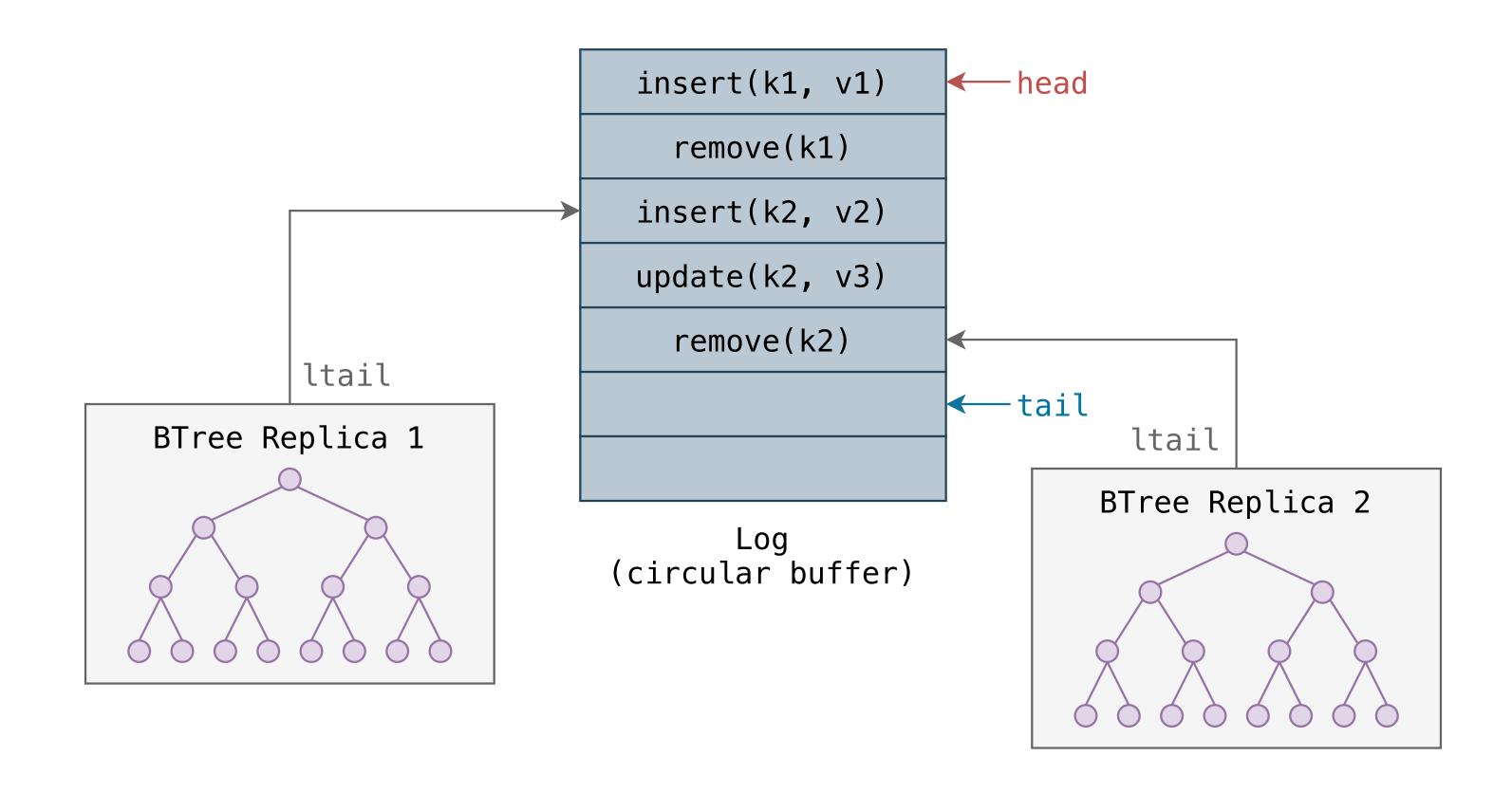
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Node Replication (NR)

Turn a sequential data struture into a linearizable NUMA-aware concurrent data structure.

- Replica on each NUMA node
- Operation log
- Flat combining
- Optimized readers-writer lock

Operation Log

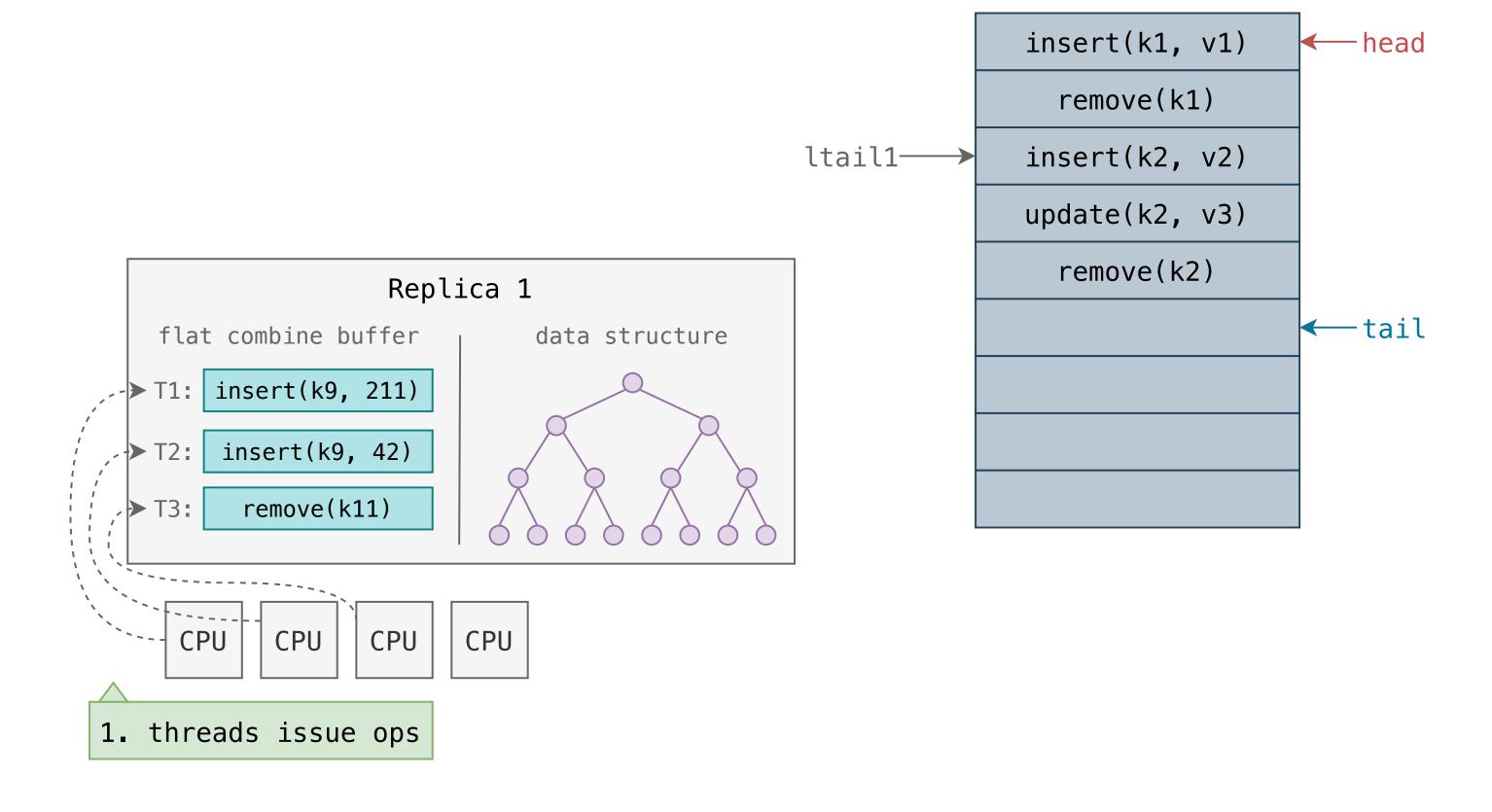


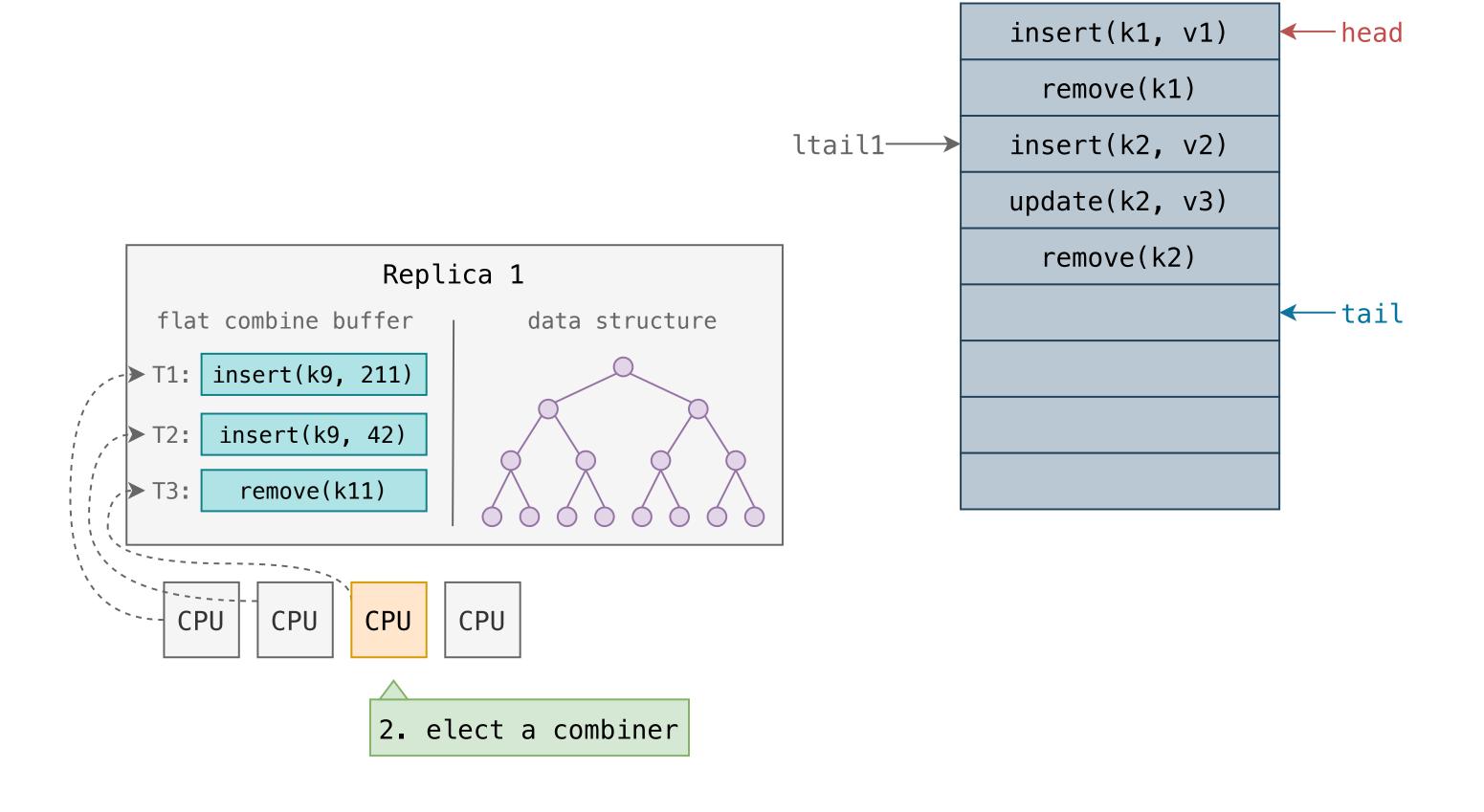
Flat Combining

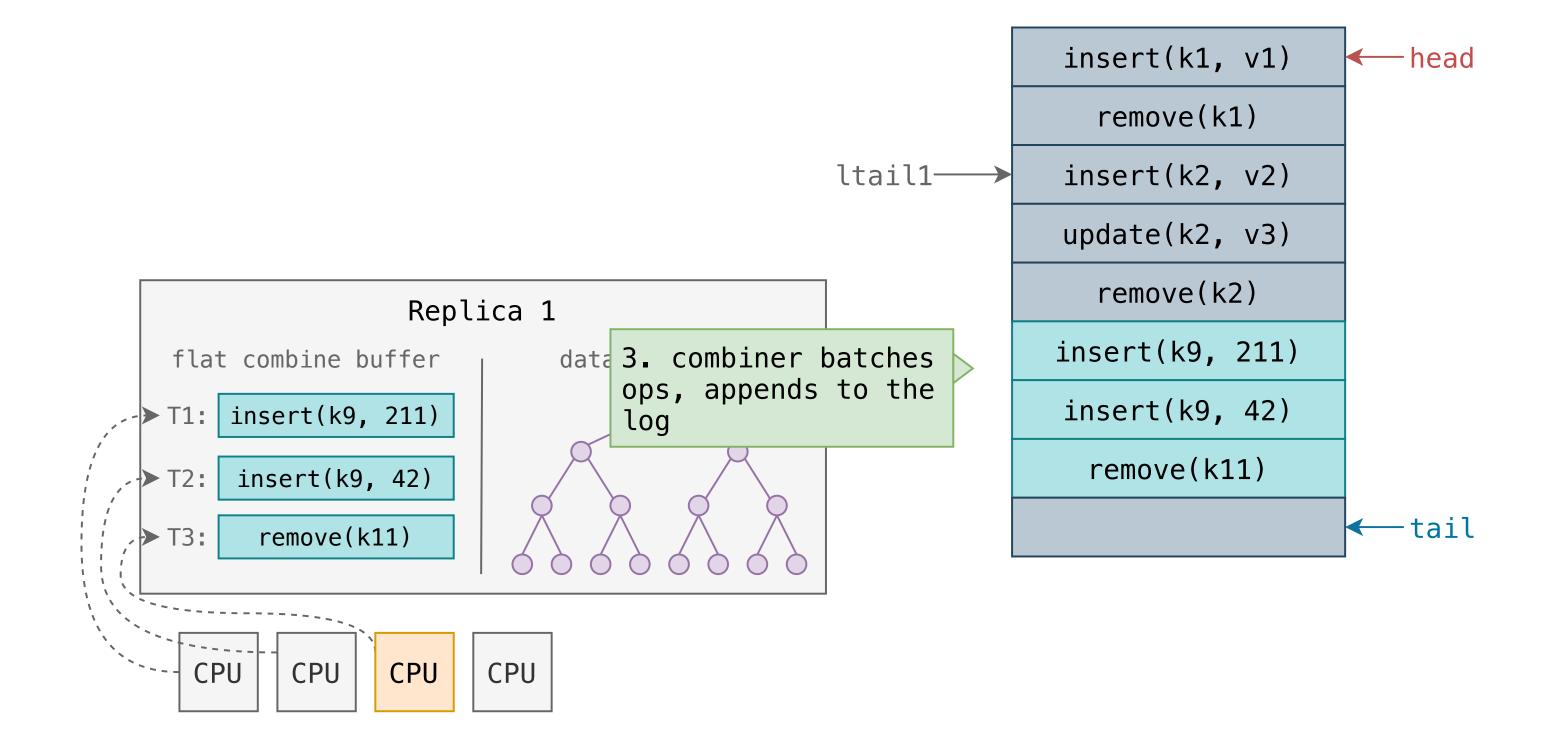
- One combiner per NUMA node
 - batch appending operation logs
 - batch executing mutations
- Lower cost, better cache locality

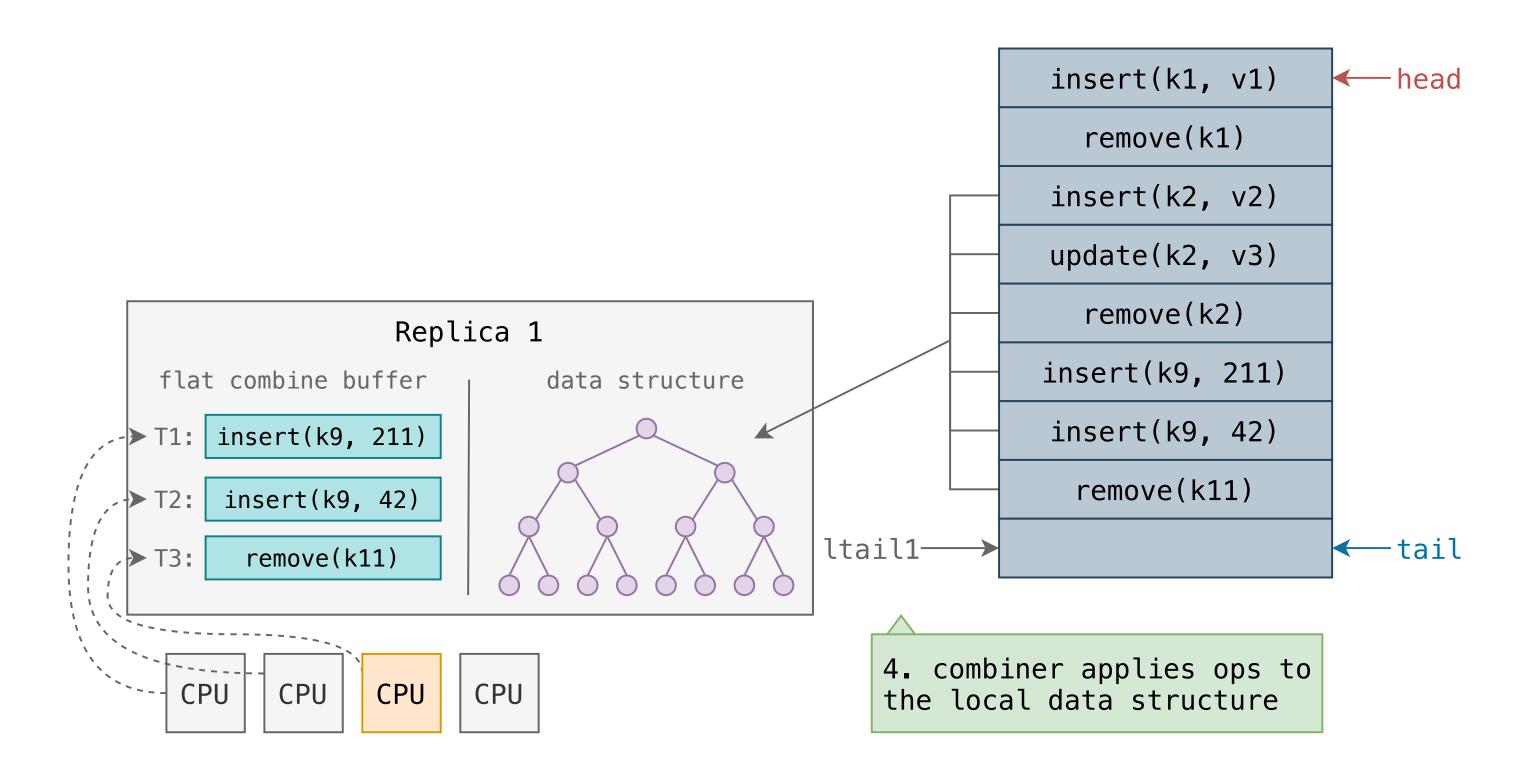
Optimized Readers-Writer Lock

- Protect concurrent read/write on local NUMA node
- Writer-preference
- Split writer lock and combiner lock
 - allow parallel reading and combining









```
struct NrHashMap { storage: HashMap<u64, u64> }
enum WriteOp { Put(u64, u64) }
enum ReadOp { Get(u64) }
impl Dispatch for NrHashMap {
    type ReadOperation = ReadOp;
    type WriteOperation = WriteOp;
    type Response = Option<u64>;
    fn dispatch(&self, op: Self::ReadOperation) → Self::Response {
        match op {
            ReadOp::Get(key) \Rightarrow self.storage.get(&key).map(|v| *v),
    fn dispatch_mut(&mut self, op: Self::WriteOperation) → Self::Response {
        match op {
            WriteOp::Put(key, value) ⇒ self.storage.insert(key, value),
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```
let logsize = 2 * 1024 * 1024;
let log = Log::<<NrHashMap as Dispatch>::WriteOperation>::new(logsize);

let replica1 = Replica::<NrHashMap>::new(&log);
let replica2 = Replica::<NrHashMap>::new(&log);

let tid1 = replica1.register();

let r = replica1.execute(ReadOp::Get(1), tid1);
let r = replica1.execute_mut(WriteOp::Put(1, 1), tid1);
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Problems with NR

- Frequent mutating operations (e.g. FS)
 - mostly independent (commutative)
- Limit scalability

Concurrent Node Replication (CNR)

Turn an *already concurrent data structure* into a NUMA-aware concurrent data structure.

Multiple logs

- assign commutative operations to different logs
- assign conflicting operations to the same log

Multiple combiners per NUMA node

concurrently append and apply operation logs

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NrOS

- Designed around per-NUMA node kernel replicas
- Major subsystems
 - NR-vMem: virtual memory management
 - NR-FS: in-memory file system
 - NR-Scheduler: process management

NR-vMem

- Per-process mapping B-Tree & hardware PT are replicated
- Mutating Ops: `Map`, `Unmap`, `Adjust`
- Non-mutating Ops: `Resolve`
- Problem: Out-of-band read by hardware
 - Map`: page fault → `Resolve`
 - Unmap & Adjust: $IPI \rightarrow update replica \rightarrow TLB flush$

NR-FS

- Entire in-memory FS data structure is replicated
- Problem: POSIX read operations mutate kernel state (e.g. fd offset)
 - only implement `pread`/`pwrite` in kernel
 - fd offset in userspace lib
- Problem: Large amound data in operation log
 - allocate kernel buffer, put only references in log

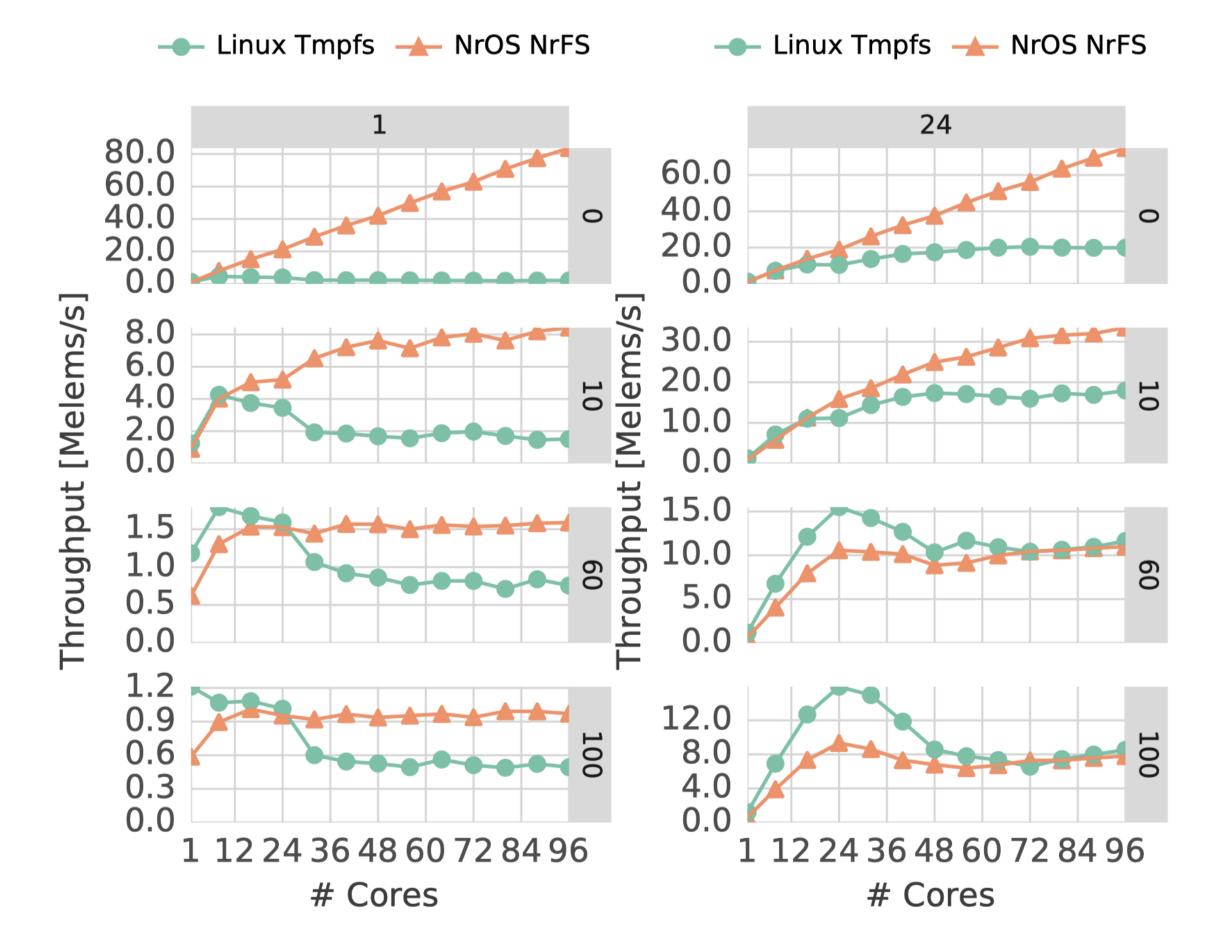
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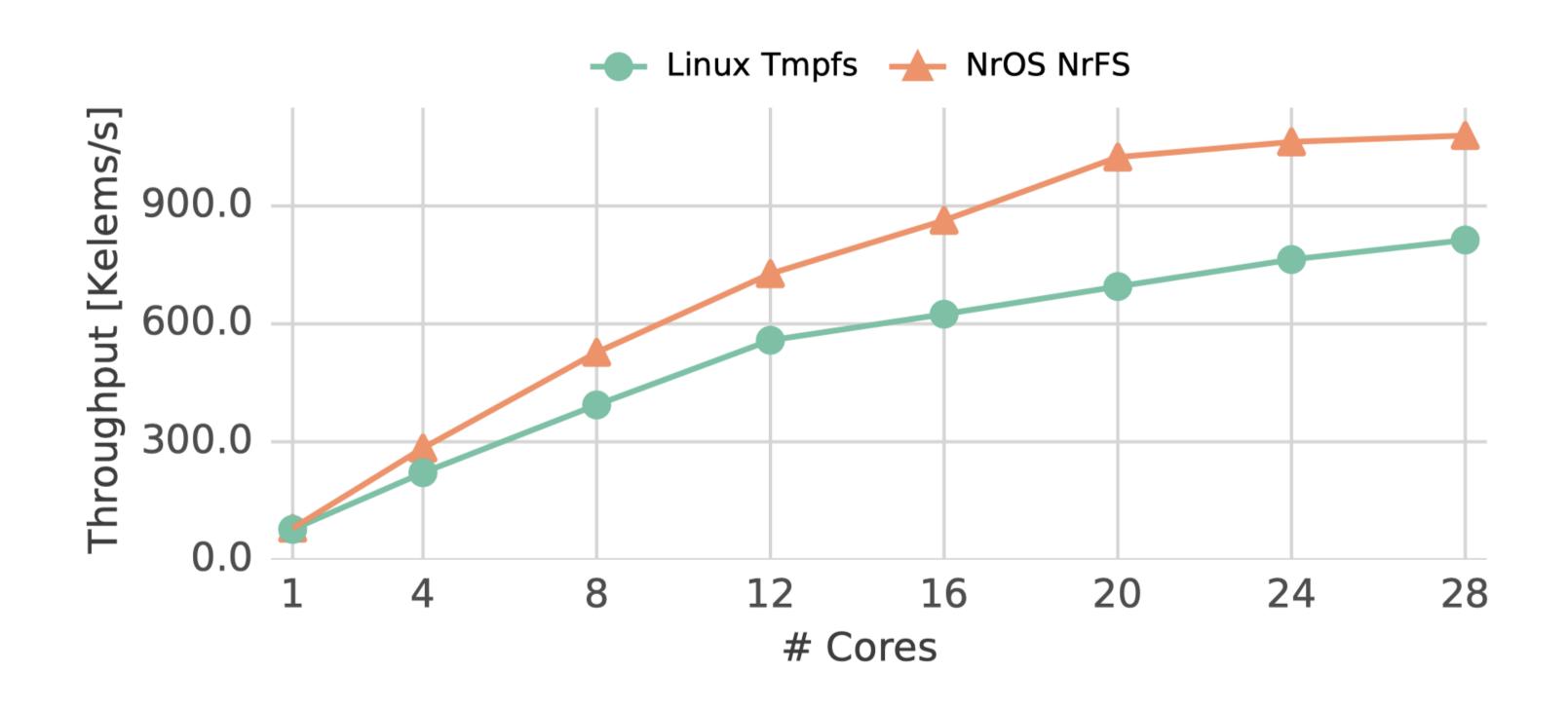
Evaluation Platforms

Name	Memory	Nodes/Cores/Threads
2×14 Skylake	192 GiB	2×14×2 Xeon Gold 5120
4×24 Cascade	1470 GiB	4×24×2 Xeon Gold 6252

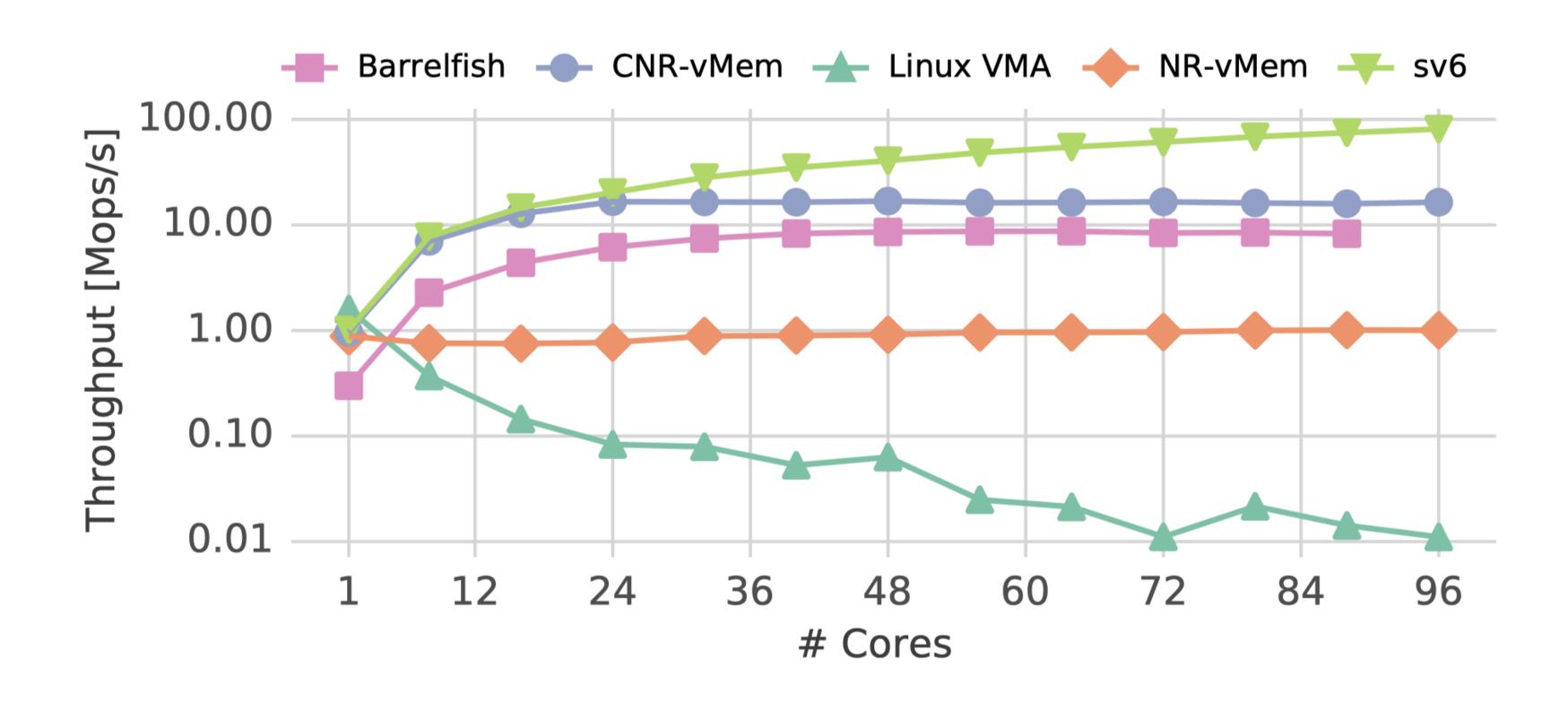
NR-FS Microbenchmark



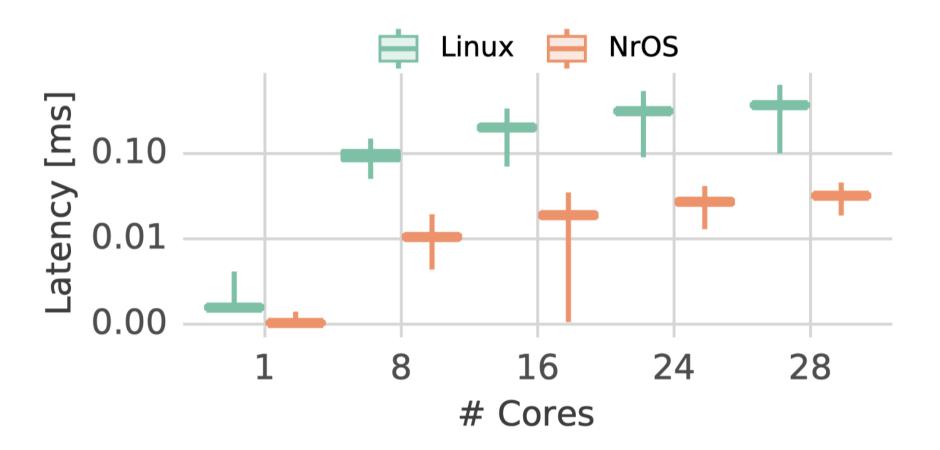
NR-FS LevelDB



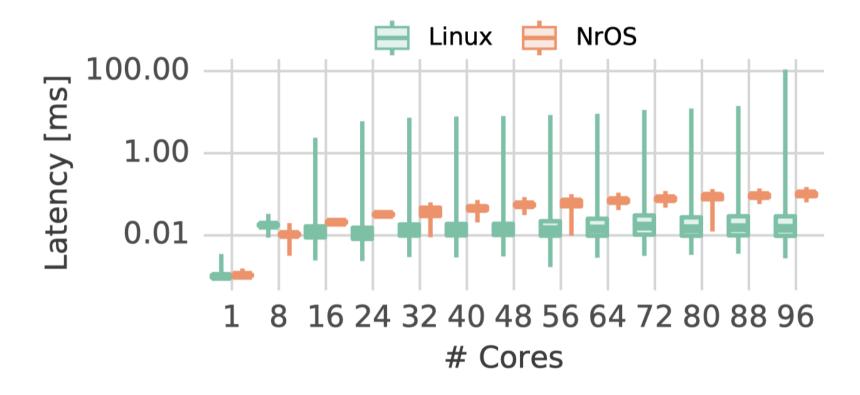
NR-vMem Map Throughput



NR-vMem Map Latency

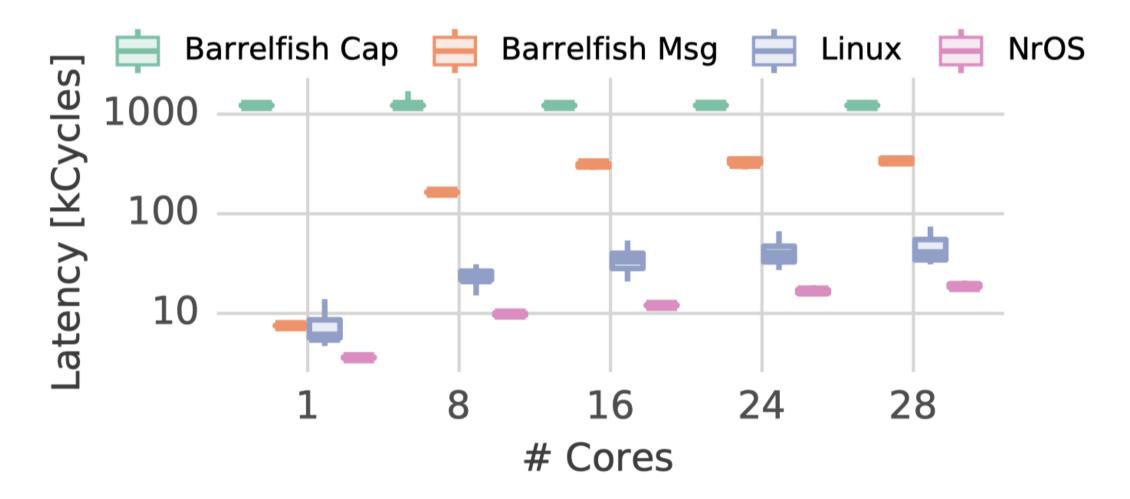


(a) Map latency for 2×14 Skylake.



(b) Map latency for 4×24 Cascade.

NR-vMem Unmap Latency



(c) Unmap latency on 2×14 Skylake.

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NRkernel Principles

- Combining replicated and shared state
- Replica consistency via operation log
- Compiler-enforced memory and concurrency safety

Future Directions

- Relaxing consistency
- Verifying correctness
- Extending NrOS for compute clusters

Thank you!