Mysql 1

Introduction

- History
- Basic Terminology
- Relational Model versus Relational DBMS
- Domain
- Relation
- Relational Data Integrity

History

- E. F. Codd
 - A relational model of data for large shared data banks (1970)
- Set-Theory model (Childs, 1968)
- Goals of the RM (relational model)
 - high degree of data independence
 - theory of basic problems e.g. consistency, redundancy (normalization)
 - use of set-oriented data manipulation

Terminology

Common Term Relational Term

Table Relation

Record or Row Tuple

Field or Column Attribute

Number of columns Degree

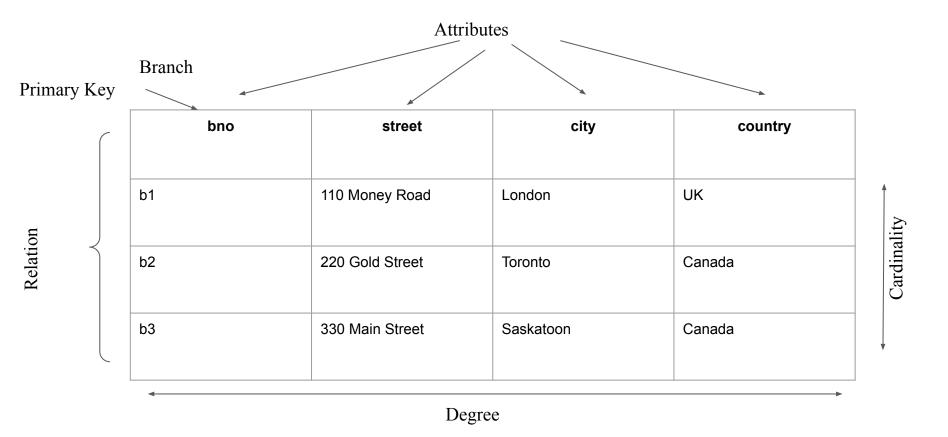
Number of rows Cardinality

Identifier Primary Key

Pool of legal values Domain

RM versus Relational DBMS

- There are differences between RM and relational DBMS
 - RM no duplicate tuples
 - Domains
 - etc...



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Domain

- Pool of values (like a type)
- Values are atomic
- Domains have names
 - unique in database
 - are set of values
 - different domains can overlap
- e.g. {white, red, blue, green }

- Relations consist of
 - heading & body
- Heading
 - fixed set of attributes
 - e.g. { <A1:D1>, <A2:D2>, <AN:DN>}
- Body
 - set of tuples
 - e.g. { T1, T2, TN }
 - Tuple
 - { <A1:V1>, <A2:V2> <An:VN> }

Heading for relation S

• Tuple

- We interpret as humans the relation
- Difference between relation and table
 - relation -> abstract
 - relation -> linked to set theory
 - table -> concrete thing
 - table -> allows duplicate tuples

- Table -> suggest order
 - » tupel
 - » attributes

But there is no order of tuples or attributes!

- Degree (arity)

 - are domains relations?
 - -NO
 - relation dynamic
 - domain static
- Cardinality
 - rows

Operations

- Relation -> Relation
- Operations
 - Selection
 - Projection
 - Cartesian Product
 - Union
 - Intersection
 - Difference
 - Join
 - natural join
 - semi join
 - left outer join
 - Division

Properties of Relations

• Relation has a name that is distinct from other names

• Each cell of the relations contains only a single value

Each attribute has a distinct name

• The values of an attribute are from the same domain

• The order of attributes is not important

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Normalized Relations

- Normalization
 - A technique for producing a set of relations with desirable properties, given the data requirements of the enterprise
- Different types of NF 1,2,3,4,5
 - 1NF
 - A relation in which the intersection of each row and column contains only one value
 - 2NF
 - Relation is in 1NF and non key attributes are fully dependent on key
 - 3NF
 - Relation is in 2NF and all non-key attributes are not dependent on other non-key attribute

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Relational Keys

Super Key

An attribute or a set of attributes that uniquely identify a tuple within a relation

Candidate Key

 A super key such that no proper subset is a super key within the relation

Primary key

The candidate key that is selected to identify tuples uniquely within the relation
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Relational Data Integrity

- Candidate keys
 - uniqueness property
 - irreducibility property
- Why candidate keys?
 - Reference tupleCMPT 353 (c) Ralph Deters 2025
 - _
- Primary key
 - a special candidate key
- Foreign key
 - referential integrity (no unmatched foreign keys)

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MySQL

docker-compose.yml

```
services:
mysql1:
  image: mysql:8.0
  container_name: db1
  volumes:
    - /Users/ralph/classes/353/w6/dbfiles:/var/lib/mysql
  environment:
    MYSQL_ROOT_PASSWORD: admin_xxx
    MYSQL_DATABASE: my_database
    MYSQL_USER: user1
    MYSQL_PASSWORD: user1_xxx
```

```
node1:
build: .
container_name: nodejs1
ports:
    - "80:8080"
volumes:
    - /Users/ralph/classes/353/w6:/usr/src/app
depends_on:
    mysql1:
    condition: service_started
stdin_open: true
tty: true
```

docker exec -it <container_name_or_id> mysql -u <username> -p
docker exec -it db1 bash
mysql -u root -p

Create table

```
show databases;

create database db1;

drop database db1;

create database db1;

use db1;
```

```
CREATE TABLE mytable
id int unsigned NOT NULL auto_increment,
username varchar(100) NOT NULL,
email varchar(100) NOT NULL,
PRIMARY KEY (id)
```

describe mytable;

ALTER TABLE mytable ADD COLUMN phone_number VARCHAR(20);

ALTER TABLE mytable drop COLUMN phone_number;

Insert

INSERT INTO mytable (username, email) VALUES ("myuser", "myuser@example.com");

Select

SELECT C1, C2, C3 FROM table where C1 = "1";

select * from mytable;

Update

UPDATE table_name SET column1 = value1, column2 = value2, column3 = value3, ... WHERE condition;

UPDATE mytable SET username="myuser_1" WHERE id=1;