



Operating Systems CS F372

Synchronization

BIJU K RAVEENDRAN

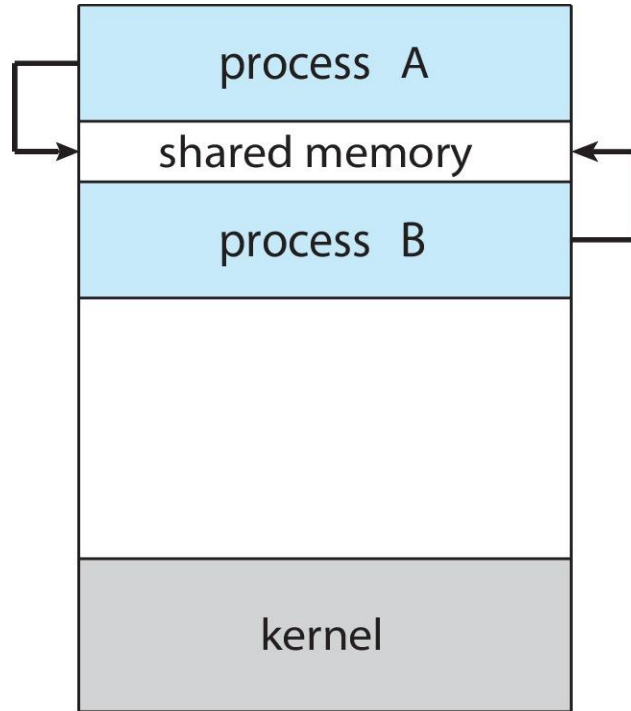


Inter Process Communication

- Processes within a system may be **independent** or **cooperating**
- Cooperating process can affect or be affected by other processes, including sharing data
- Reasons for cooperating processes:
 - Information sharing, Computation speedup, Modularity, Convenience
- Cooperating processes need **inter process communication (IPC)**
 - Shared memory, Message passing

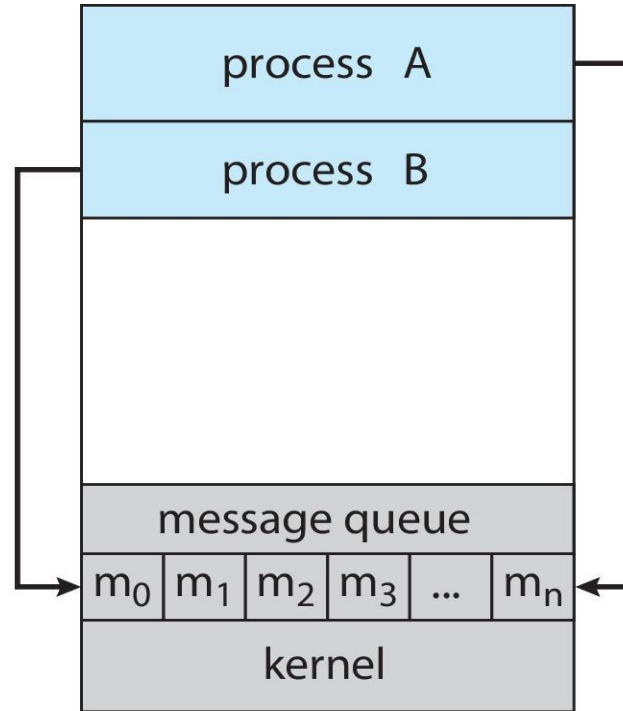
Communication Models

(a) Shared memory.



(a)

(b) Message passing.



(b)

Message Passing

- Message system – processes communicate with each other without resorting to shared variables
 - Very useful in Distributed systems
- IPC facility provides two operations:
 - **send**(*message*) – message size fixed/variable]
 - **receive**(*message*)
- If P and Q wish to communicate, they need to:
 - establish a *communication link* between them
 - exchange messages via send/receive



Communication

- Direct
 - Processes must name each other explicitly:
 - **send** (P , *message*), **receive**(Q , *message*)
 - [exactly one] link [unidirectional or bi-directional] is associated with exactly one pair of processes
- Indirect
 - Mailbox [also referred to as port] with unique ID
 - Processes can communicate only if they share a mailbox
 - A link [uni /bi directional] may be associated with many processes
 - Each pair of processes may share several communication links, each link corresponds to a mail box
 - Operations: create new mailbox, send/receive messages, destroy mailbox

Synchronization

- Message passing may be either blocking or non-blocking
- **Blocking** is considered **synchronous**
 - **Blocking send:** The sender blocks until the message is received
 - **Blocking receive:** The receiver blocks until a message is available
- **Non-blocking** is considered **asynchronous**
 - **Non-blocking:** The sender sends the message and continue
 - **Non-blocking:** The receiver receives a valid message or null



Synchronization

- If a message is lost or if process fails before sending , receive process is permanently blocked in case of blocking receive
- In case of non blocking receive, if the process executes receive before message is sent the message will be lost
- Allow process to test whether a message is waiting before issuing a receive primitive.
- Receive can also test for arrival before issuing receive.



Buffering

- Queue of messages attached to the link; implemented in one of three ways
 1. Zero capacity – 0 messages
Sender must wait for receiver
 2. Bounded capacity – finite length of n messages
Sender must wait if link full
 3. Unbounded capacity – infinite length
Sender never waits





Producer – Consumer Problem


- Paradigm for cooperating processes, *producer* process produces information that is consumed by a *consumer* process
 - *unbounded-buffer* places no practical limit on the size of the buffer
 - *bounded-buffer* assumes that there is a fixed buffer size

Producer – Consumer Problem

$in \leftarrow 0, out \leftarrow 0$

PRODUCER

CONSUMER



```
while (1) {  
    // Produce item;  
    Buffer[in] = item;  
    in = in + 1;  
}
```

```
while (1) {  
    while (in == out);  
    item = Buffer[out];  
    out = out + 1;  
}
```

Bounded Buffer – Shared Memory Solution

- Shared data

```
#define BUFFER_SIZE 10
```

```
typedef struct {
```

```
    ...
```


```
} item;
```

```
item buffer[BUFFER_SIZE];
```

```
int in = 0;
```

```
int out = 0;
```

Producer – Consumer Problem



```
/* Producer */
while (true) {
    /* Produce an item */
    while (((in+1)% BUFFER_SIZE) == out);
        /* do nothing -- no free buffers */

    buffer[in] = item;
    in = (in + 1) % BUFFER_SIZE;
}

/*Consumer */
while (true) {
    while (in == out); // do nothing -- nothing to consume
    // remove an item from the buffer
    item = buffer[out];
    out = (out + 1) % BUFFER_SIZE;
}
```




```
/* Producer */
```

```
while (true) {
```

```
    /* produce an item and put in nextProduced */
```

```
    while (count == BUFFER_SIZE); // do nothing
```

```
    buffer [in] = nextProduced;
```

```
    in = (in + 1) % BUFFER_SIZE;
```

```
    count++;
```

```
}
```

```
/* Consumer */
```

```
while (true) {
```

```
    while (count == 0); // do nothing
```

```
    nextConsumed = buffer[out];
```

```
    out = (out + 1) % BUFFER_SIZE;
```

```
    count--;
```

```
    /* consume the item in nextConsumed */
```

```
}
```