

Homework 7: Monomorphic and Polymorphic Type Systems

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Formation

This law defines the $\text{LIST}(\tau)$ type, which represents a list which contains elements of type τ . Note that this means that lists can only have one type τ .

$$\frac{\tau \text{ is a type}}{\text{LIST}(\tau) \text{ is a type}} \text{ (LIST)}$$

Introduction

This law defines a non-empty $\text{LIST}(\tau)$ is of type τ created using cons.

$$\frac{\Gamma_{\xi, \phi, \rho} \vdash x : \tau \quad \Gamma_{\xi, \phi, \rho} \vdash y : \text{LIST}(\tau)}{\Gamma_{\xi, \phi, \rho} \vdash \text{CONS-LIST}(x, y) : \text{LIST}(\tau)} \text{ (CONS LIST)}$$

$$\frac{\tau \text{ is a type}}{\Gamma_{\xi, \phi, \rho} \vdash \text{EMPTY-LIST}(\tau) : \text{LIST}(\tau)} \text{ (EMPTY LIST)}$$

Elimination

null? , car , and cdr are all elimination rules, as they are rules which can "observe" information about the list. null? checks if the list is empty, so it returns a boolean given a list of type τ . car returns the first element of the cons cell, which will be the type of the list τ , so car will return type τ . cdr returns the second element of the cons cell, which in a list, is the rest of the list, so cdr will return type $\text{LIST}(\tau)$.

$$\frac{\Gamma_{\xi, \phi, \rho} \vdash x : \text{LIST}(\tau)}{\Gamma_{\xi, \phi, \rho} \vdash \text{IS-NULL}(x) : \text{BOOL}} \text{ (IS NULL)}$$

$$\frac{\Gamma_{\xi, \phi, \rho} \vdash x : \text{LIST}(\tau)}{\Gamma_{\xi, \phi, \rho} \vdash \text{CAR}(x) : \tau} \text{ (CAR)}$$

$$\frac{\Gamma_{\xi, \phi, \rho} \vdash x : \text{LIST}(\tau)}{\Gamma_{\xi, \phi, \rho} \vdash \text{CDR}(x) : \text{LIST}(\tau)} \text{ (CDR)}$$