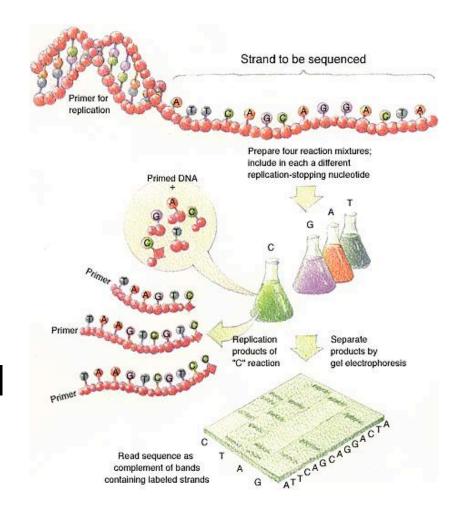
Eulerian graph

DNA Sequencing

- Shear DNA into millions of small fragments
- Read 500 700
 nucleotides at a
 time from the small
 fragments (Sanger
 method)



Fragment Assembly

- Computational Challenge: assemble individual short fragments (reads) into a single genomic sequence ("superstring")
- Until late 1990s the shotgun fragment assembly of human genome was viewed as intractable problem

Shortest superstring problem

- Problem: Given a set of strings, find a shortest string that contains all of them
- Input: Strings s_1, s_2, \ldots, s_n
- Output: A string s that contains all strings
 s₁, s₂,..., s_n as substrings, such that the
 length of s is minimized
- Complexity: NP—hard
- Note: this formulation does not take into account sequencing errors

Shortest superstring problem: example

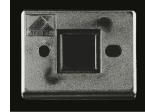
```
The Shortest Superstring problem
Set of strings: {000, 001, 010, 011, 100, 101, 110, 111}
Concatenation
              000 001 010 011 100 101 110 111
Superstring
                       010
                    110
                 011
Shortest
             0001110100
superstring
               001
                   111
                      101
                          100
```

Sequencing by Hybridization (SBH): History

1988: SBH is suggested as an First an alternative sequencing method.

First microarray prototype (1989)

 1991: Light directed polymer synthesis developed by Steve Fodor and colleagues. First commercial DNA microarray prototype w/16,000 features (1994)



500,000 features per chip **(2002)**

 1994: Affymetrix develops first 64-kb DNA microarray



How SBH works

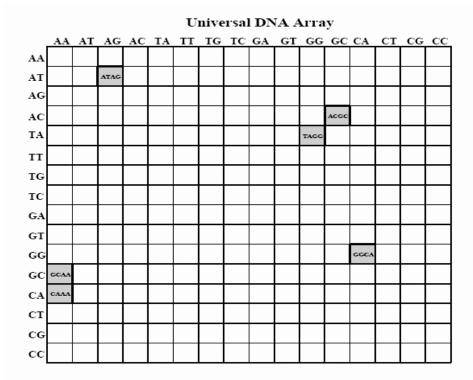
- Attach all possible DNA probes of length / to a flat surface, each probe at a distinct and known location. This set of probes is called the DNA array.
- Apply a solution containing fluorescently labeled DNA fragment to the array.
- The DNA fragment hybridizes with those probes that are complementary to substrings of length / of the fragment.

How SBH works

 Using a spectroscopic detector, determine which probes hybridize to the DNA fragment to obtain the *I*—mer composition of the target DNA fragment.

 Apply the combinatorial algorithm (below) to reconstruct the sequence of the target DNA fragment from the *I* – mer composition.

Hybridization on DNA Array



DNA target TATCCGTTT (complement of ATAGGCAAA) hybridizes to the array of all 4-mers:

ATAGGCAAA ATAG TAGG AGGC GGCA GCAA

I-mer composition

- Spectrum (s, l) unordered multiset of all possible (n l + 1) l-mers in a string s of length n
- The order of individual elements in Spectrum (s, l) does not matter
- For s = TATGGTGC all of the following are equivalent representations of Spectrum (s, 3):
 {TAT, ATG, TGG, GGT, GTG, TGC}
 {ATG, GGT, GTG, TAT, TGC, TGG}

{TGG, TGC, TAT, GTG, GGT, ATG}

I-mer composition

- Spectrum (s, l) unordered multiset of all possible (n l + 1) l-mers in a string s of length n
- The order of individual elements in Spectrum (s, I) does not matter
- For s = TATGGTGC all of the following are equivalent representations of Spectrum (s, 3):
 {TAT, ATG, TGG, GGT, GTG, TGC}
 {ATG, GGT, GTG, TAT, TGC, TGG}
 {TGG, TGC, TAT, GTG, GGT, ATG}
- We usually choose the lexicographically maximal representation as the canonical one.

Observations: different sequences may have the same spectrum

Different sequences may have the same spectrum:

```
Spectrum(GTATCT,2)=
Spectrum(GTCTAT,2)=
{AT, CT, GT, TA, TC}
```

The SBH problem

Goal: Reconstruct a string from its *I*-mer composition

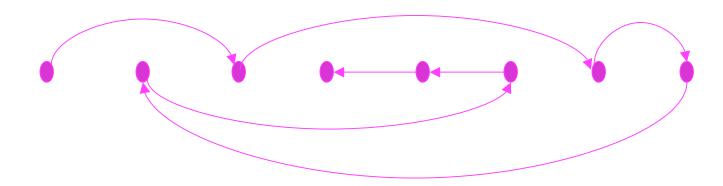
 Input: A set S, representing all I-mers from an (unknown) string s

Output: String s such that Spectrum (s,l) =

SBH: Hamiltonian path approach

S = { ATG AGG TGC TCC GTC GGT GCA CAG }

H ATG AGG TGC TCC GTC GGT GCA CAG



ATGCAGGTCC

Path visited every VERTEX once

SBH: Hamiltonian path approach

A more complicated graph:

```
S = { ATG TGG TGC GTG GGC GCA GCG CGT }
```

SBH: Hamiltonian path approach

 $S = \{ATG TGG TGC GTG GGC GCA GCG CGT\}$

Path 1:

ATGCGTGGCA

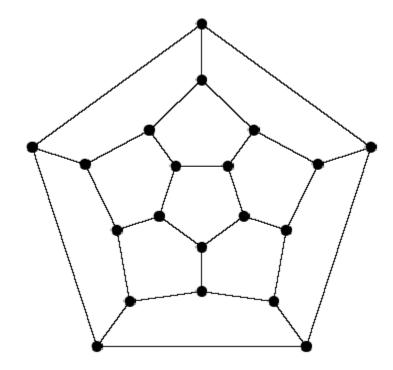
Path 2:

ATGGCGTGCA

Hamiltonian cycle problem

 Find a cycle that visits every *vertex* exactly once

NP-complete

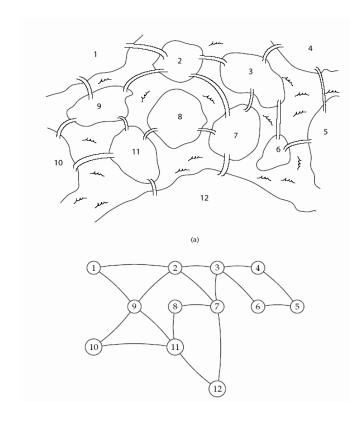


Game invented by Sir William Hamilton in 1857

Eulerian cycle problem

 Find a cycle that visits every edge exactly once

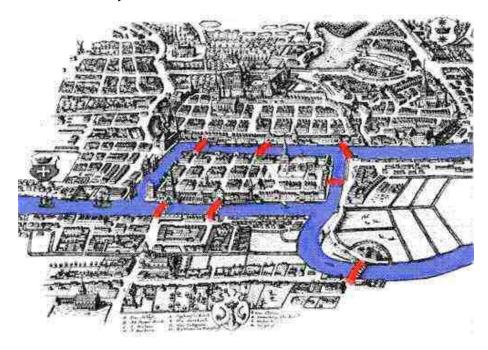
Linear time



More complicated Königsberg

The Bridge Obsession Problem

Find a tour crossing every bridge just once *Leonhard Euler*, 1735



Bridges of Königsberg

Euler Theorem

 A graph is balanced if for every vertex the number of incoming edges equals to the number of outgoing edges:

$$in(v) = out(v)$$

• Theorem: A connected graph is Eulerian if and only if each of its vertices is balanced.

Euler Theorem: proof

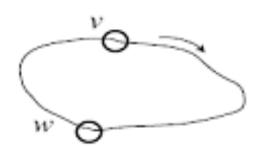
Eulerian → balanced
 for every edge entering v (incoming edge)
 there exists an edge leaving v (outgoing edge). Therefore

$$in(v) = out(v)$$

Balanced → Eulerian
 ???

Algorithm for Constructing an Eulerian Cycle

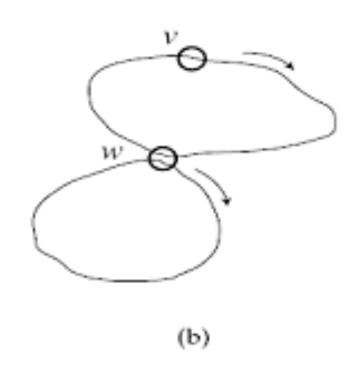
Start with an arbitrary vertex v and form an arbitrary cycle with unused edges until a dead end is reached. Since the graph is Eulerian this dead end is necessarily the starting point, i.e., vertex v.



(a)

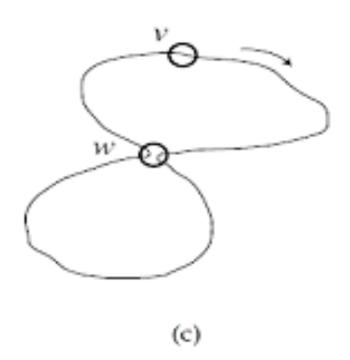
Algorithm for Constructing an Eulerian Cycle (cont'd)

If cycle from (a) above is not an Eulerian cycle, it must contain a vertex w, which has untraversed edges. Perform step (a) again, using vertex w as the starting point. Once again, we will end up in the starting vertex w.



Algorithm for Constructing an Eulerian Cycle

c. Combine the cycles from (a) and (b) into a single cycle and iterate step (b).



Euler Theorem: extension

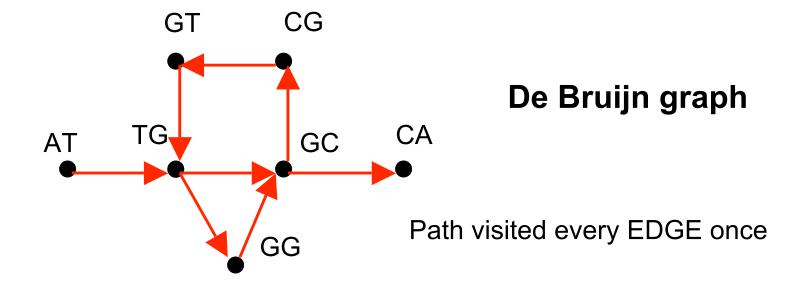
 Theorem: A connected graph has an Eulerian path if and only if it contains at most two semi-balanced vertices and all other vertices are balanced.

SBH: Eulerian path approach

S = { ATG, TGC, GTG, GGC, GCA, GCG, CGT }

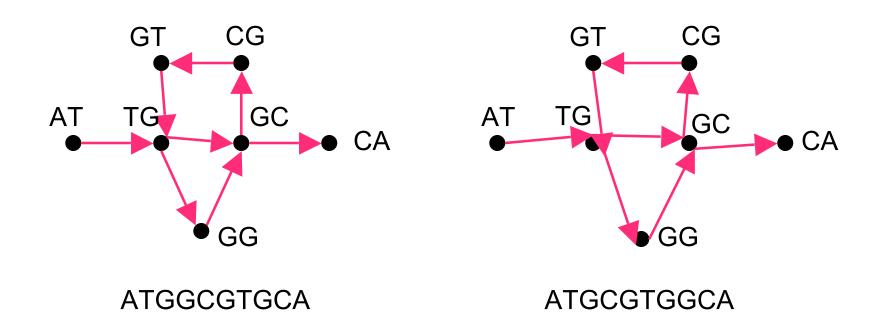
Vertices correspond to (l-1) – mers : $\{AT, TG, GC, GG, GT, CA, CG\}$

Edges correspond to I – mers from S



SBH: multiple solutions

S = { AT, TG, GC, GG, GT, CA, CG } corresponds to two different paths:



Representing alternative SBH solutions by Eulerian graphs

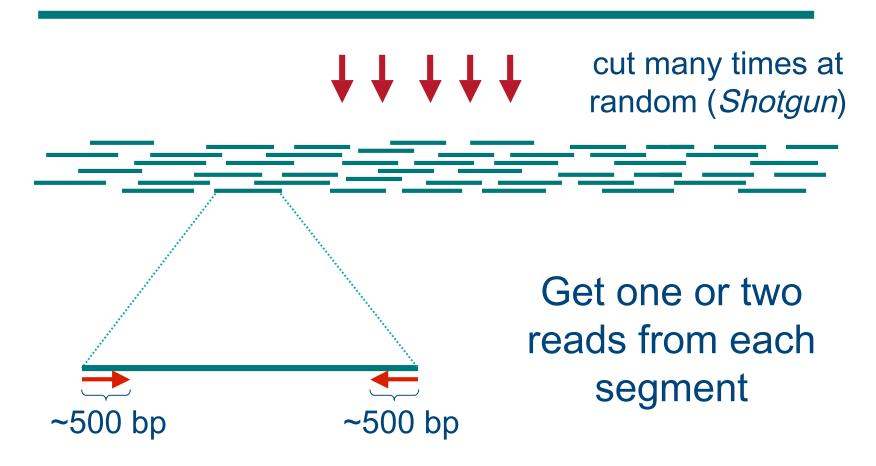
- # of solutions ←→ # of Eulerian cycles
 - BEST theorem

Obtains all solutions from one solution ->
 Cassette transformations

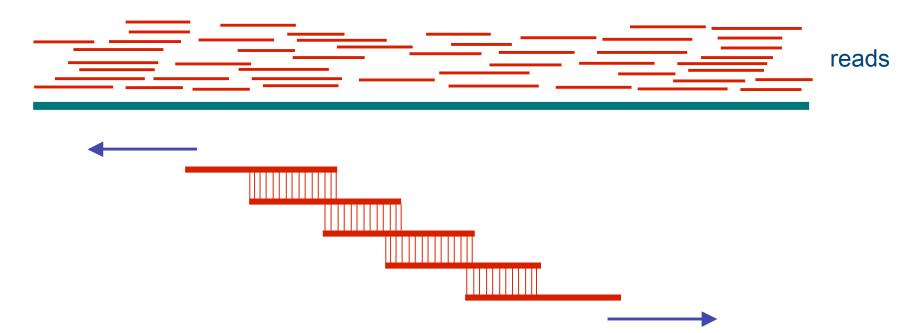
Reads assembly

Shotgun Sequencing

genomic segment



Fragment Assembly



Cover region with ~7-fold redundancy

Overlap reads and extend to reconstruct the original genomic region

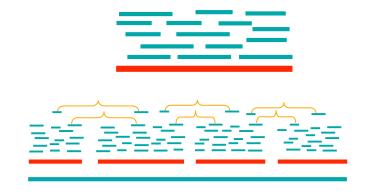
Overlap-Layout-Consensus

Assemblers: ARACHNE, PHRAP, CAP, TIGR, CELERA

Overlap: find potentially overlapping reads



Layout: merge reads into contigs and contigs into supercontigs



Consensus: derive the DNA sequence and correct read errors

..ACGATTACAATAGGTT..

Overlap

- Find the best match between the suffix of one read and the prefix of another
- Due to sequencing errors, need to use dynamic programming to find the optimal overlap alignment
- Apply a filtration method to filter out pairs of fragments that do not share a significantly long common substring

Overlapping Reads

- Sort all k-mers in reads $(k \sim 24)$
- Find pairs of reads sharing a k-mer
- Extend to full alignment throw away if not >95% similar



Overlapping Reads and Repeats

- A *k*-mer that appears N times, initiates N² comparisons
- For an Alu that appears 10⁶ times → 10¹² comparisons too much

Solution:

Discard all k-mers that appear more than $t \times \text{Coverage}$, $(t \sim 10)$

Finding Overlapping Reads

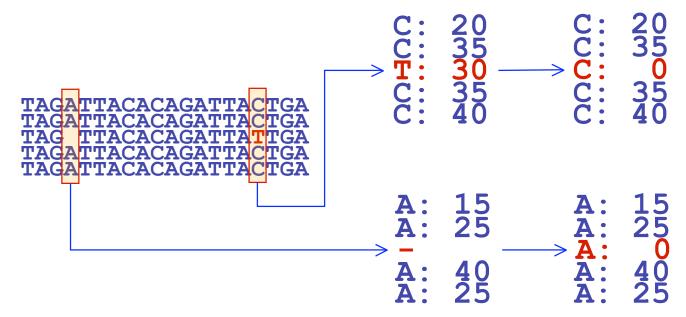
Create local multiple alignments from the overlapping reads

TACATTACACACATTACTCA
TAGATTACACACATTACTCA
TAG TTACACACATTATTCA
TAGATTACACACACATTACTCA
TAGATTACACACACATTACTCA
TAGATTACACACACATTACTCA
TAG TTACACACACATTACTCA
TAGATTACACACACATTACTCA

Finding Overlapping Reads

(cont'd)

Correct errors using multiple alignment

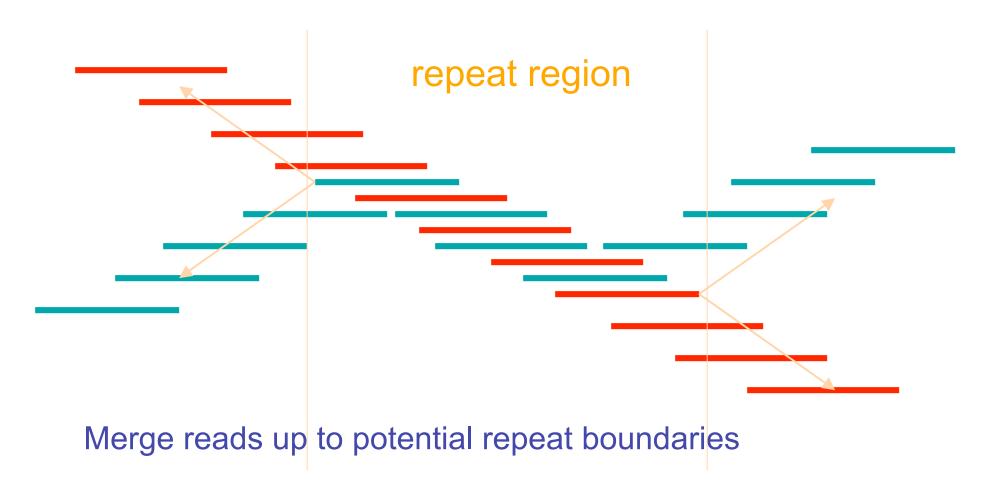


- Score alignments
- Accept alignments with good scores

Layout

- Repeats are a major challenge
- Do two aligned fragments really overlap, or are they from two copies of a repeat?
- Solution: repeat masking hide the repeats!!!
- Masking results in high rate of misassembly (up to 20%)
- Misassembly means alot more work at the finishing step

Merge Reads into Contigs



Repeats, Errors, and Contig Lengths

- Repeats shorter than read length are OK
- Repeats with more base pair differencess than sequencing error rate are OK
- To make a smaller portion of the genome appear repetitive, try to:
 - Increase read length
 - Decrease sequencing error rate

Error Correction

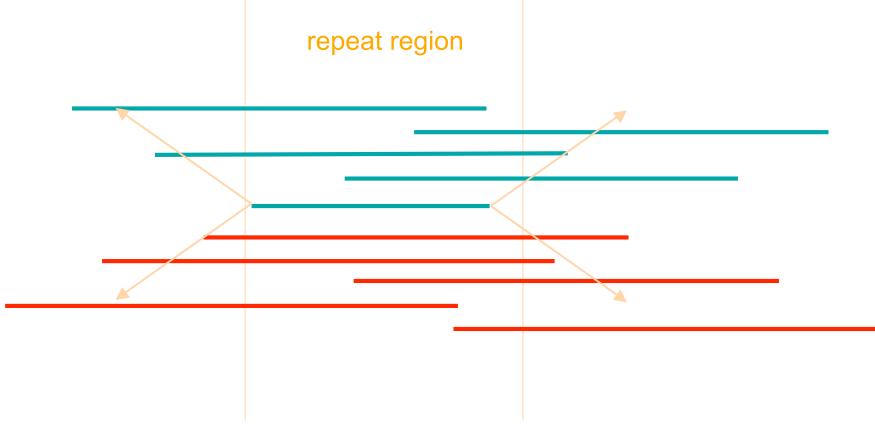
Role of error correction:

Discards ~90% of single-letter sequencing errors

decreases error rate

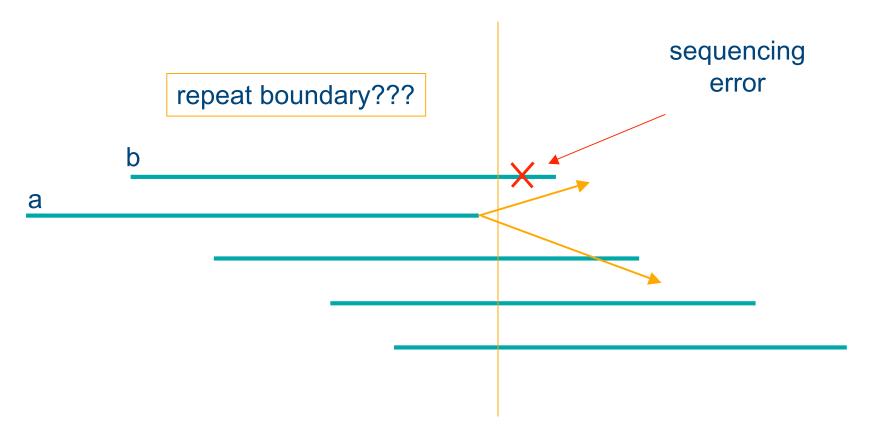
- ⇒ decreases effective repeat content
- ⇒ increases contig length

Merge Reads into Contigs (cont'd)



- Ignore non-maximal reads
- Merge only maximal reads into contigs

Merge Reads into Contigs (cont'd)



Ignore "hanging" reads, when detecting repeat boundaries

Merge Reads into Contigs (cont'd)



Insert non-maximal reads whenever unambiguous

Link Contigs into Supercontigs



Normal density

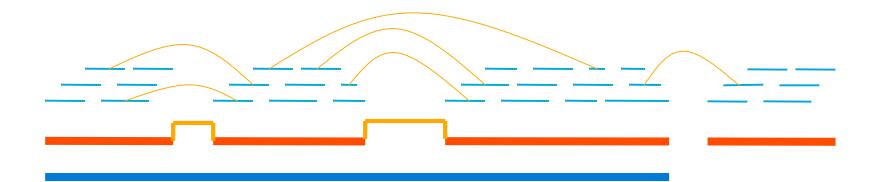
Too dense: Overcollapsed?

Inconsistent links: Overcollapsed?

Link Contigs into Supercontigs (cont'd)

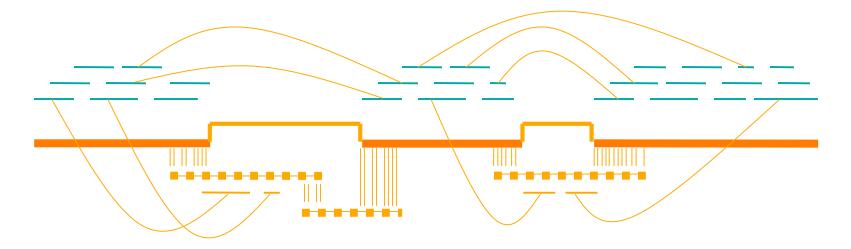
Find all links between unique contigs

Connect contigs incrementally, if ≥ 2 links

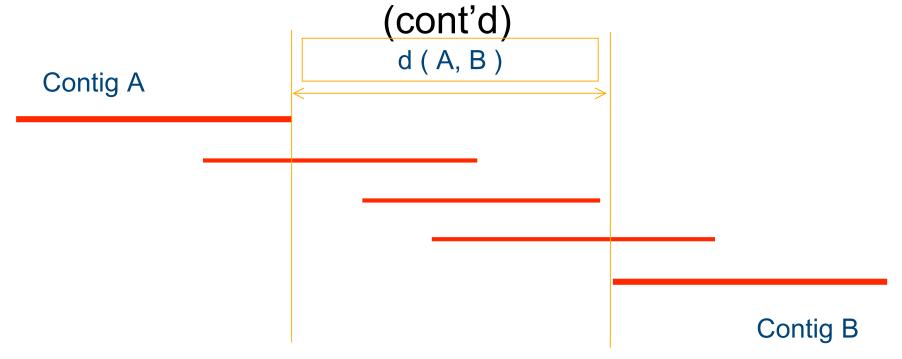


Link Contigs into Supercontigs (cont'd)

Fill gaps in supercontigs with paths of overcollapsed contigs



Link Contigs into Supercontigs



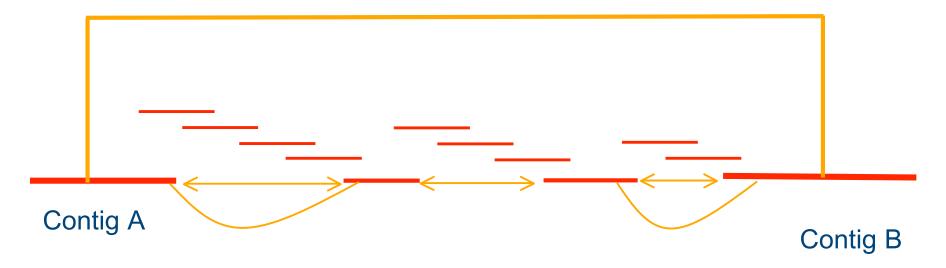
Define G = (V, E)

V := contigs

E := (A, B) such that d(A, B) < C

Reason to do so: Efficiency; full shortest paths cannot be computed

Link Contigs into Supercontigs (cont'd)



Define T: contigs linked to either A or B

Fill gap between A and B if there is a path in G passing only from contigs in T

Consensus

 A consensus sequence is derived from a profile of the assembled fragments

 A sufficient number of reads is required to ensure a statistically significant consensus

Reading errors are corrected

Derive Consensus Sequence



Derive multiple alignment from pairwise read alignments

Derive each consensus base by weighted voting

Problems with the shotgun approach

