Question 1: Inference in Markov Networks

- a. Computing P(A|b1) using variable elimination
 - i. Factors are f0(D,A), f1(A,B), f2(B,C), f3(C,D)

f0		
d0	a0	100
d0	al	1
d1	a0	1
d1	al	100
fl		
a0	b0	30
a0	b1	5
al	b0	1
al	b1	10
f2		
b0	c0	100
b0	cl	1
b1	c0	1
b1	cl	100
f3		
c0	d0	1
c0	d1	100
cl	d0	100
cl	d1	1

ii. Then we update the observed value B=b1, so f1 and f2 changes to f4 and f5 as follows:

f4	
a0	5
al	10
f 5	
c0	1
cl	100

iii. Now, specify elimination ordering D,C:

$$\textstyle \sum_D f_0(D,A) f_3(C,D) \sum_C f_4(A) f_5(C)$$

iv. Combine f0 and f3 to form f6(D,A,C):

f6			
d0	a0	c0	100
d0	a0	cl	10000
d0	al	c0	1
d0	al	cl	100
d1	a0	c0	100
d1	a0	cl	1
d1	al	c0	10000
d1	al	cl	100

v. Sum out D to form f7(A,C):

f 7		
a0	c0	200
a0	cl	10001
al	c0	10001
al	cl	200

vi. Eliminate C by combining f5(C) and f7(A,C), then summing out on C:

f8		
a0	c0	200
a0	cl	1000100
al	c0	10001
al	cl	20000

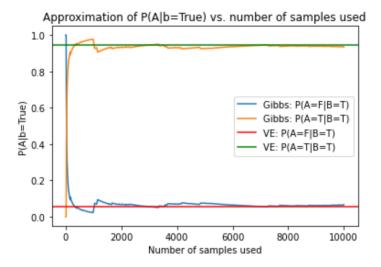
f 9	
a0	1000300
al	30001

vii. Combine f4(A) and f9(A) to form f10(A):

f10	
a0	5001500
al	300010

viii. Normalize f10(A) to get P(A|b1):

		-
a0	0.94341	
al	0.05659	



```
def prob_ab(A, B):
    if not A and not B:
        return 30
    elif not A and B:
        return 5
    elif A and not B:
        return 1
    else:
        return 10
def prob bc(B, C):
    if not B and not C:
        return 100
    elif not B and C:
        return 1
    elif B and not C:
        return 1
    else:
        return 100
def prob cd(C, D):
    if not C and not D:
        return 1
    elif not C and D:
        return 100
    elif C and not D:
        return 100
    else:
        return 1
def prob_da(D, A):
    if not D and not A:
        return 100
    elif not D and A:
        return 1
    elif D and not A:
        return 1
    else:
        return 100
```

```
def random_bool(v1, v2):
   p = v1 / (v1 + v2)
   return random.uniform(0, 1) <= p</pre>
def next_variable(curr_variable):
    variables = ['A', 'C', 'D']
    return random.choice(variables)
def sample_variable(curr_variable, A, B, C, D):
   if curr_variable == 'A':
        A0 = prob_ab(0, B) * prob_da(D, 0)
        A1 = prob_ab(1, B) * prob_da(D, 1)
        A = random\_bool(A1, A0)
   elif curr_variable == 'D':
        D0 = prob_cd(C, 0) * prob_da(0, A)
        D1 = prob_cd(C, 1) * prob_da(1, A)
        D = random\_bool(D1, D0)
   elif curr_variable == 'C':
        C0 = prob_cd(0, D) * prob_bc(B, 0)
        C1 = prob_cd(1, D) * prob_bc(B, 1)
        C = random\_bool(C1, C0)
   return A, B, C, D
```

```
A1 count = 0
A0_count = 0
num\_samples = 10000
curr_variable = 'A'
(A, B, C, D) = (True, True, True, True)
pa0_givenb1 = []
pa1_givenb1 = []
for i in range(1,num_samples+1):
    (A, B, C, D) = sample_variable(curr_variable, A, B, C, D)
    if A:
        A1_count += 1
    else:
        A0_count += 1
    curr_variable = next_variable(curr_variable)
    pa0_givenb1.append(A0_count / i)
    pa1_givenb1.append(A1_count / i)
```

Question 2: Resolution in Propositional Logics

$$\frac{KB}{(B \lor C) \land \neg A}$$

$$\neg (B \to A)$$

$$A \to (\neg B \lor C)$$

$$A \lor E$$

$$C$$

- a. $KB \models \neg (A \leftrightarrow C)$
 - Convert second and third clauses to CNF

$$\neg (B \to A) \equiv \neg (\neg B \lor A) \equiv \neg A \land B$$
$$A \to (\neg B \lor C) \equiv \neg A \lor \neg B \lor C$$

- ii. Simplify KB to CNF
 - Conjunction of second and first clause:
 - $[\neg A \land B] \land [(B \lor C) \land \neg A] \equiv \neg A \land B$
 - Conjunction of result from above with fourth clause:
 - $[\neg A \land B] \land [A \lor E] \equiv \neg A \land B \land E$
 - Conjunction of result from above with fifth clause:
 - $[\neg A \land B \land E] \land C \equiv \neg A \land B \land E \land C$
 - Conjunction of result from above with third clause:
 - $[\neg A \land B \land E \land C] \land [\neg A \lor \neg B \lor C] \equiv \neg A \land B \land E \land C$
 - Conjunction of all clauses in $KB \equiv \neg A \land B \land E \land C$
- iii. Apply resolution algorithm, where $KB \equiv \neg A \land B \land E \land C$, $\alpha \equiv \neg (A \leftrightarrow C)$
 - Simplify alpha: $(A \leftrightarrow C) \equiv (A \to C) \land (C \to A) \equiv (\neg A \lor C) \land (\neg C \lor A)$
 - New = []
 - Clauses = $CNF(KB \land \alpha) \equiv \neg A \land B \land E \land C \land (\neg A \lor C) \land (\neg C \lor A)$
 - Resolvents = $\neg A \land (\neg C \lor A) \equiv \neg A \land \neg C$
 - i. New = new U resolvents = $\neg A \land \neg C$
 - Resolvents = $C \land (\neg C \lor A) \equiv A \land C$
 - i. New = new U resolvents = $\neg A \land \neg C \land A \land C$
 - Remaining resolvents are redundant, and are left out for brevity
 - New clauses (A and C) were created, so starting the second loop, clauses = clauses U new
 - Clauses = $\neg A \land B \land E \land C \land (\neg A \lor C) \land (\neg C \lor A) \land \neg C \land A$
 - Resolvents = $\neg A \land A \equiv \emptyset$, return true
 - It can also be shown that: resolvents = $\frac{\neg C \land C \equiv \emptyset}{}$, in which we also return true
- iv. The algorithm returns true. Therefore the query is entailed, i.e. $KB \models \neg (A \leftrightarrow C)$.
- b. Since KB entails $\neg (A \leftrightarrow C)$, no interpretation exists in which KB is true and $\neg (A \leftrightarrow C)$ is false.

Question 3: WalkSAT

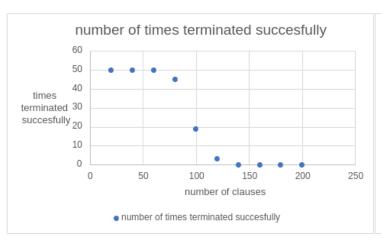
Used: https://www.cs.rochester.edu/u/kautz/walksat/

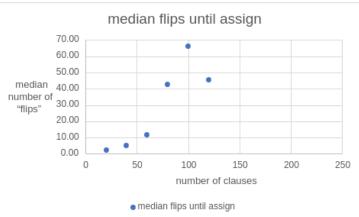
How many times WalkSAT terminates successfully for each of these sets of 50 problems

The median number of "flips" (i.e. steps) for each of these sets of 50 problems, considering only the cases when WalkSAT terminates successfully

The table and figures below are determined by fixing N (number of variables) at 20 and generating 50 random 3SAT problems for each value of C/N from 1 to 10 (so, with number of clauses ranging from 20, 40, ..., 200).

Number of clauses	Median number of flips until assigned	Number of times terminated successfully
20	2.25	50
40	5.15	50
60	11.60	50
80	42.80	45
100	66.40	19
120	45.30	3
140		0
160		0
180		0
200		0





Source code: https://github.com/brucecui97/cpsc422A3Q3

Question 4: Semantics of First-Order Logic

Interpretations for a first-order logic specification that includes: 10 constant symbols, 2 ternary predicates, 2 binary predicates, 10 unary predicates Assume the world contains 10 objects and they are all assigned unique names.

10 objects and 10 symbols:

- Assuming unique names, each object is assigned 1 symbol, so there are 10! interpretations 10 unary predicates:
 - 10 objects, each can be 0 or 1, so there are 2^10 interpretations for 1 predicate
 - With 10 predicates, we have (2^10)^10 interpretations

2 binary predicates:

- 10*10 = 100 possibilities, each one can be 0 or 1, so there are 2^100 interpretations for 1 predicate
- With 2 predicates, we have (2^100)^2 interpretations

2 ternary predicates:

- 10*10*10 = 1000 possibilities, each one can be 0 or 1, so there are 2^1000 interpretations for 1 predicate
- With 2 predicates, we have (2^1000)^2 interpretations

TOTAL:

 $10! * 2^{10^{10}} * 2^{100^{2}} * 2^{1000^{2}}$ interpretations