Sorting Algorithms Summary

In-place: means modifying list to be sorted [as opposed to returning new sorted list]. **Stable:** means the order of elements with equal values is preserved.

Lower bound on Comparion-based Sorting

For a comparion-based algorithm, the expected number of comparions is $\Omega(n\bullet\log(n))$.

Merge Sort

Sort list by merging sorted sub-lists, reduces total number of comparisons needed.

```
1. Divide list into equally sized halves until 1
element per sub-list
2. Sort sub-list [recursively]
3. Merge sub-list using comparator
```

- Comparison-based
- **Not in-place**
- **Stable**
- **Runs in O(n•log(n)) time**

// performs at most n•log(n) comparisons

Heap Sort

Sort by traversing down a heap tree.

```
1. Create a heap from list
```

```
    2. Delete the root node [in list, a[n-1]]
       - Now root is in a[n], since n--
    3. Heapify [make sure heap property is preserved]
    4. Repeat steps 2-4 until no more elements to sort
```

- Comparison-based
- **In-place**
- **Not stable**
- **Runs in O(n•log(n)) time**

`// performs at most 2n•log(n) + O(n) comparisons`

Quick Sort

Sort using a randomly selected value as partition point, sorting sub-lists. Since random selection, might choose worst value [ideally middle value].

```
    1. Randomly select value
    2. Add all values less than to left sub-list, otherwise
    add to right sub-list
    3. Repeat until 1 element in sub-list
```

- Comparison-based
- **In-place**
- **Not stable**
- **Runs in O(n•log(n)) [expected] time**

`// performs at most 2n•log(n) + O(n)`

Comparison-based Algorithms

	Comparisons	In-place	Stable
Merge Sort	*n•log(n)* [worst-case]	no	yes
Heap Sort	*1.38n•log(n) + O(n)* [worst-case]	yes	no

| Quick Sort | $2n \cdot log(n) + O(n)$ [expected] | yes | no |