

How are AVL Trees constructed?

### AVL TREE INVARIANTS:

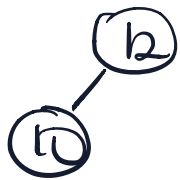
- + Search Tree Invariant
- + Height invariant.

### Q2 from tutes

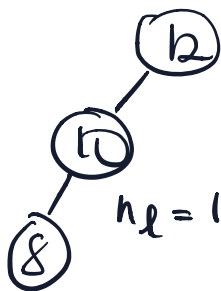
"Show how an AVL tree would be constructed if the following values were inserted in order"

1.  height = 1

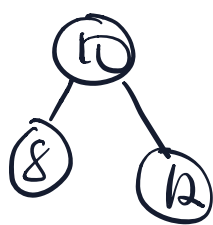
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2.   $|h_l - h_r| > 1$  ?  
 $h_l = 1$   $h_r = 0$

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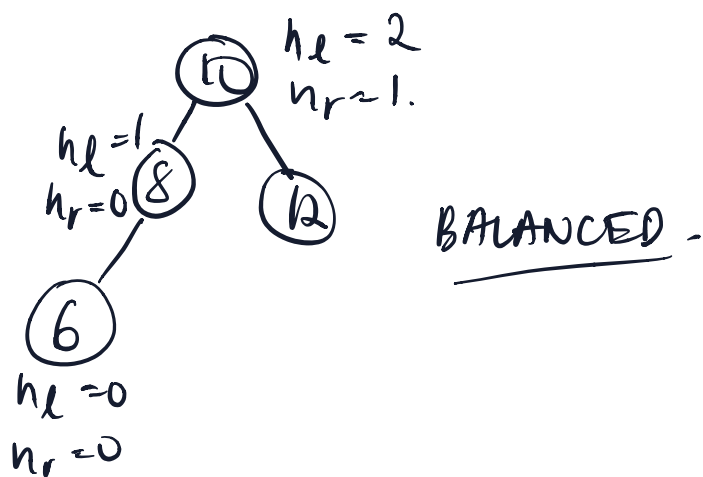
3.   $h_l = 2$   $h_r = 0$   
 $h_l = 1$   $h_r = 0$  ✓  
 $h_l = 0$   $h_r = 0$  ✓

rotate  
right.  
(counter left heavy)

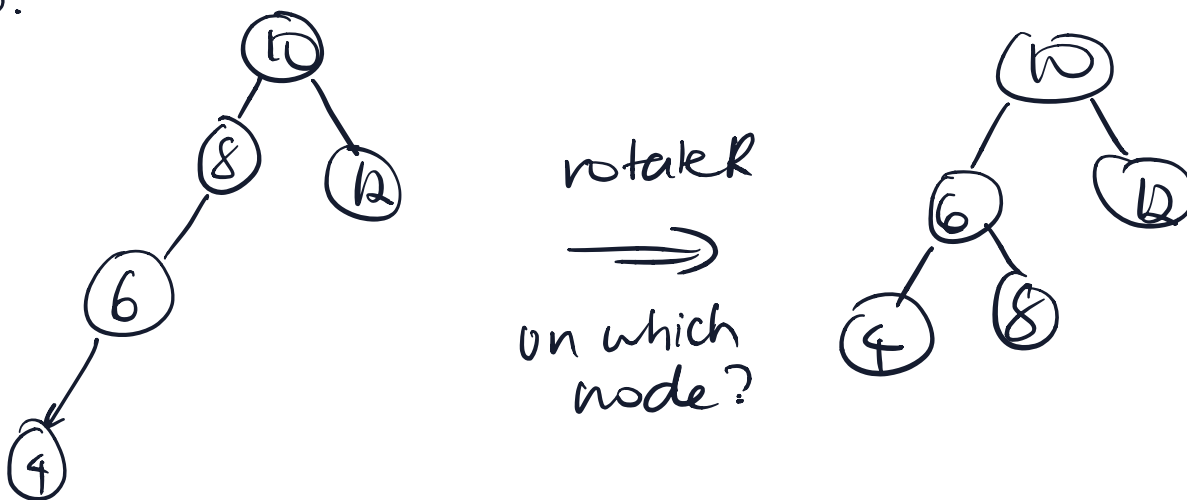


We only ever do one rotate and then it's a balanced tree, why is that?

4.



5.



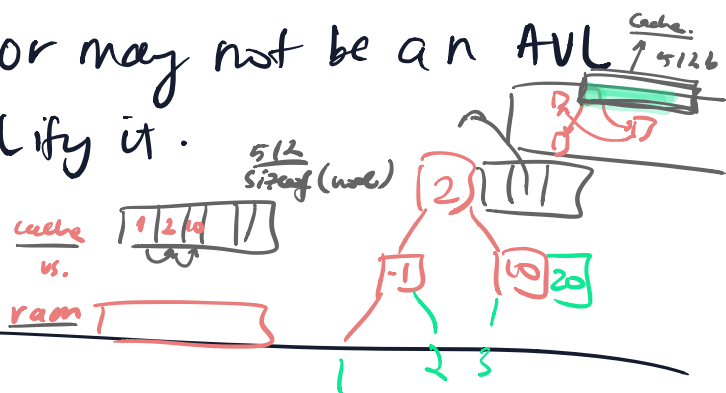
6. ② balanced.

We only ever do a maximum of one rotate / node, and then we move on. How can we guarantee that it's balanced at the end.

Develop an algorithm that checks if a tree is an AVL tree.

> What does it mean to be an AVL tree?

Say we have a BST that may or may not be an AVL tree. Develop an algorithm to AVLify it.

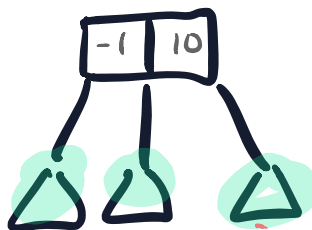


## 2-3-4 TREES.

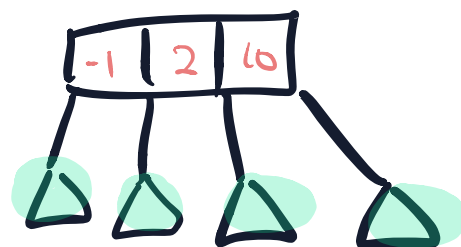
2 NODES



3 NODES



4 NODES



Invariants:

- Every non-leaf node is a 2/3 node
- All leaves are on the same level.

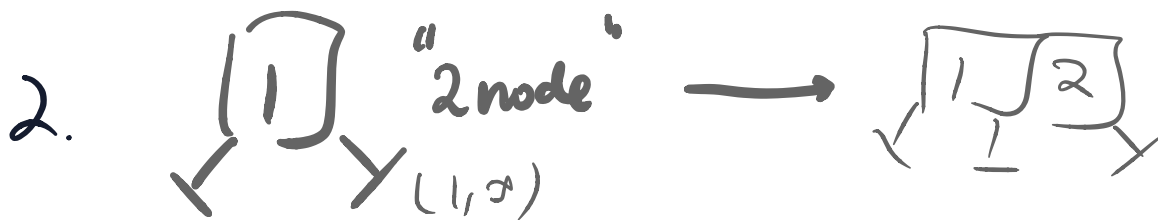
Why? Checkout insert.

(1) Find the position that you'd like to insert the value  $x$ .

(2) If we have overflow remove original middle node and push up. Insert our element as a leaf.

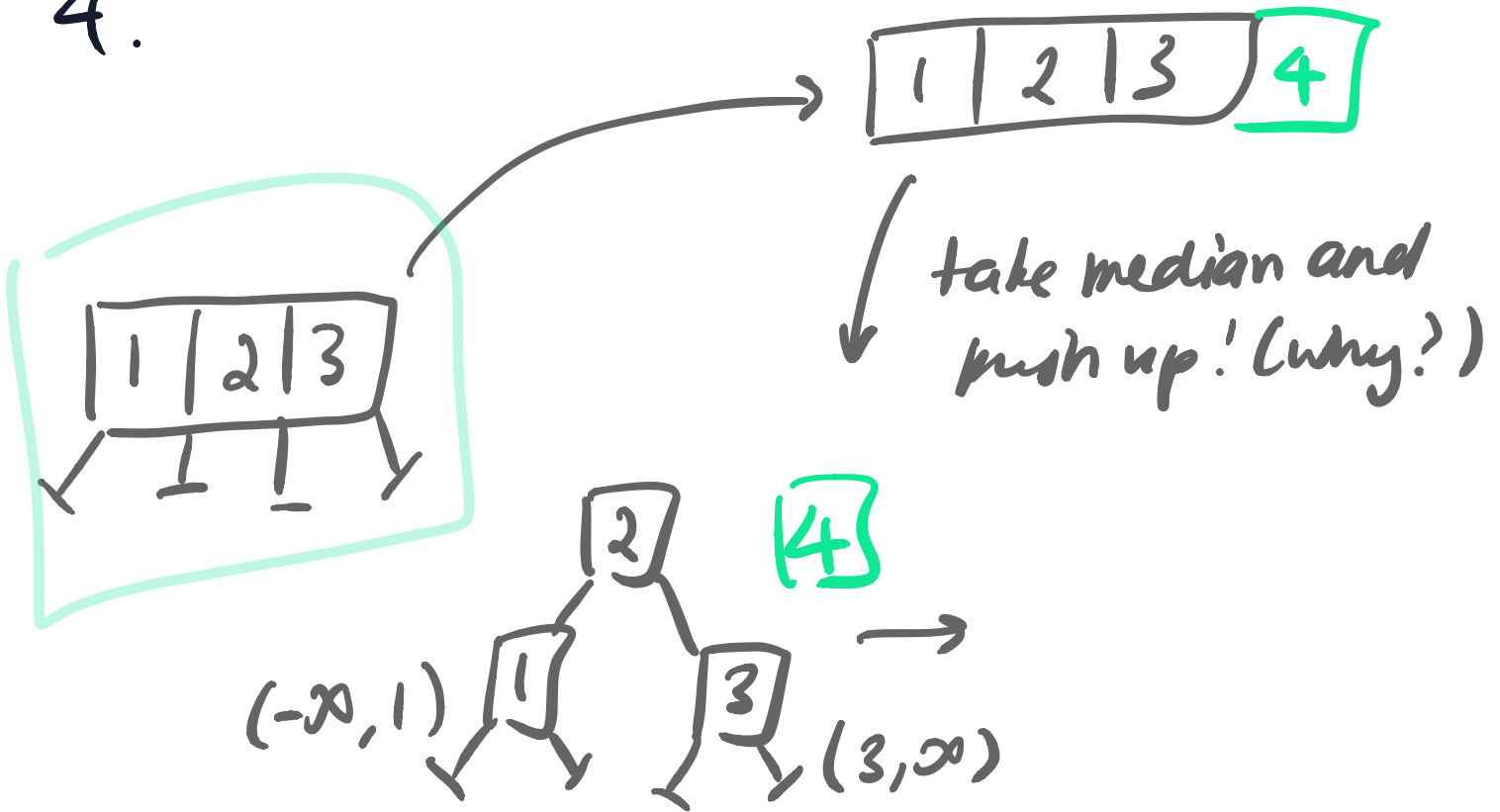
(3) Repeat overflow check for parent, doing steps (2-3) on parent.

If root, the median node is the new root.

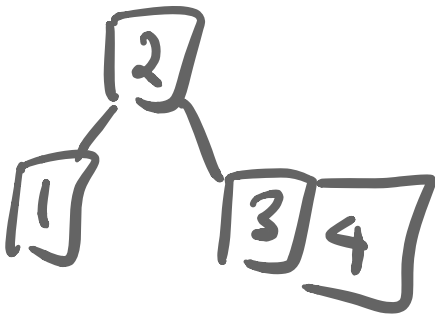


whoh!

4.



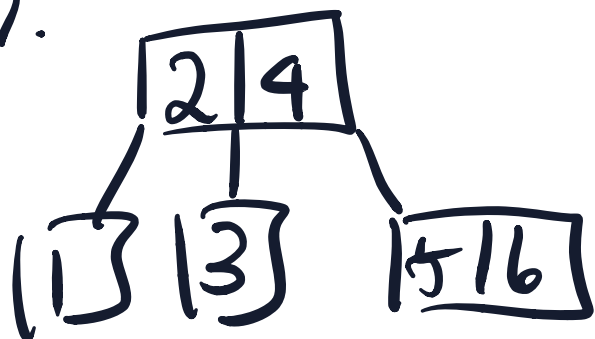
5.



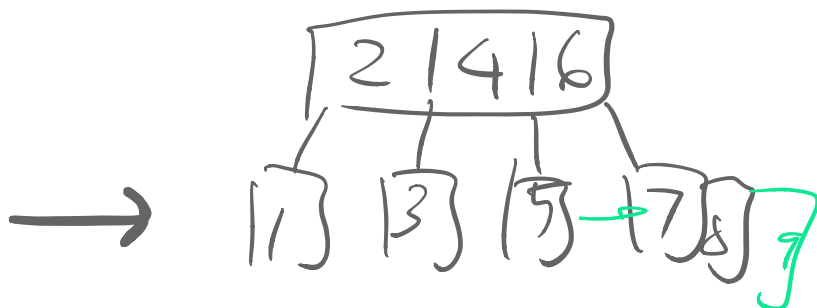
6.



7.



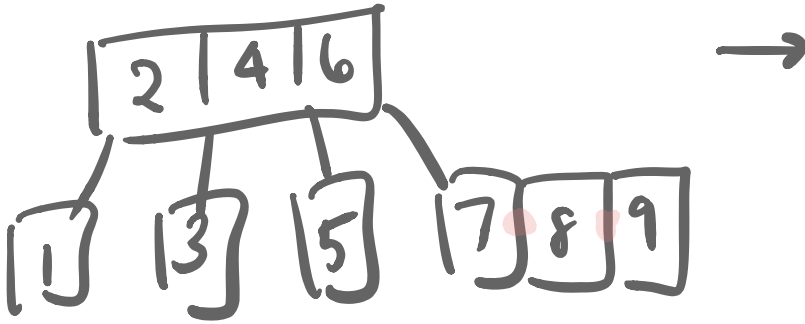
8.



9.

10.

10



search for ① ↖

in interval  $[-\infty, 4]$ ? YES!

in interval  $[-\infty, 2]$ ? YES!

in interval  $[-\infty, 1]$ ? YES!

3 comp.

search for ⑦

can check if 2 node  
or explicit range check

1/  $[-\infty, 4]$ ?  $\times$   $[4, \infty)$ ?  $\checkmark$

2/  $[-\infty, 6]$   $\times$   $[6, 7]$   $\checkmark$ . 3 checks.

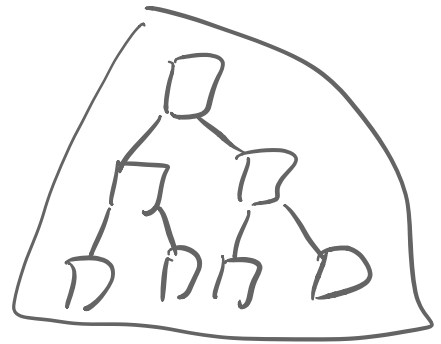
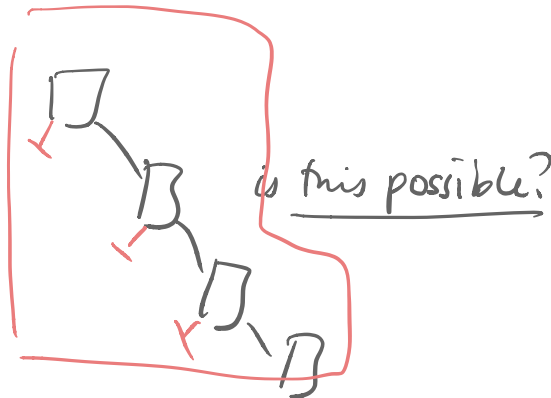
$$\lfloor \log_2(10) \rfloor = 3.$$

3/7

4/8

Some math to derive the worst case time complexities.

- ① Suppose I had a B tree with  $n$  nodes, then in order to find worst case (find the worst possible height of this tree w.r.t.  $n$ ).
- ② Each node has at most 4 children or at ~~worst~~ 2 children.
- ③ The worst case is when each node has 2 children. (why?)



- ④ If every node has exactly 2 children in worst case  $\Rightarrow O(\log_2(n))$  height.  
(balanced tree).

ALL LEAVES SAME LEVEL :

push one up  $\rightarrow$  leaves go one down the tree

$\rightarrow$  subtrees go down one level only by worst case.