

# FINDING MAX SUM OF CONTIGUOUS SUBARRAY

Let's say we've been given an array of length 5,

-1 -1 4 5 6

Here all the contiguous subarrays are,

-1	0th index	sum = -1
-1 -1	0th,1st index	sum = -2
-1 -1 4	0th,1st,2nd index	sum = 2
-1 -1 4 5	0th,1st,2nd,3rd index	sum = 7
-1 -1 4 5 6	0th,1st,2nd,3rd,4th index	sum = 13
-1	1th index	sum = -1
-1 4	1th,2nd index	sum = 3
-1 4 5	1th,2nd,3rd index	sum = 8
-1 4 5 6	1th,2nd,3rd,4th index	sum = 14
4	2nd index	sum = 4
4 5	2nd,3rd index	sum = 9
4 5 6	2nd,3rd,4th index	sum = 15
5	3rd index	sum = 5
5 6	3rd,4th index	sum = 11
6	4th index	sum = 6

Here, the subarray with maximum sum is 15, which is this subarray

4 5 6

But how are we going to find this?

# BRUTE FORCE

For the brute force method, we need to traverse the array twice, for every subarray. So for this, time needed is  $O(n^2)$ .

```
int ans = 0;
for(int i=0; i<n; i++){
    int sum = 0;
    for(int j=i; j<n; j++){
        sum += arr[j];
    }
    sum = max(ans, sum);
}
```

We need a better approach!!

# KADANE'S ALGORITHM

We need a sum variable and maxi variable.

Sum initialized to 0, and maxi initialized to the 0th element of the array.

```
int sum = 0, maxi = arr[0];
```

Then we'll traverse through the array once and keep updating maxi and sum.

```
for(int i=0; i<n; i++){
    int sum += arr[i];
    maxi = max(sum, maxi);
}
```

And another thing, if the sum becomes negative, we'll set this to zero. Because a negative sum will reduce the sum.

```
int sum = 0, maxi = arr[0];
for(int i=0;i<n;i++){
    int sum +=arr[i];
    maxi = max(sum,maxi);
    if(sum<0){
        sum =0;
    }
}
```

This approach costs  $O(n)$  time and will give us the maximum sum of the contiguous subarrays.

*Note: This approach won't work for negative arrays( an array with all negative values).*