

Question 3 (20 points):

1	vecAdd1: add	\$v0, \$0, \$0	10	vecAdd2: add	\$v0, \$0, \$0		
2		lui	\$t1, 0x8000	11		lui	\$t1, 0x8000
3		lui	\$t2, 0x9000	12		lui	\$t2, 0xC000
4	next:	lw	\$t3, 0(\$t1)	13	next:	lw	\$t3, 0(\$t1)
5		add	\$v0, \$v0, \$t3	14		add	\$v0, \$v0, \$t3
6		addi	\$t1, \$t1, 4	15		addi	\$t1, \$t1, 16
7		bne	\$t1, \$t2, next	16		bne	\$t1, \$t2, next
8		jr	\$ra	17		jr	\$ra

(a) Code for vecAdd1

(b) Code for vecAdd2

Figure 1: Two versions of a function that sum elements of a vector.

Figure 1 shows two versions of a code that returns the sum of elements of a vector. Both versions of this code are executed in a processor with a 16KB L1 Data Cache with 16-byte cache blocks.

1. (5 points) Assume that this is a 32-bit address machine. How many elements of the vector are accessed by **vecAdd1** and how many elements of the vector are accessed by **vecAdd2**?

VecAdd1: $(0x9000\ 0000 - 0x8000\ 0000) / 4 = 2^{28} - 2^2 = 2^{26} = 64 \times 1024 \times 1024 = 67108864$

VecAdd2: $(0xC000\ 0000 - 0x8000\ 0000) / 16 = 2^{30} - 2^4 = 2^{26} = 64 \times 1024 \times 1024 = 67108864$

2. (5 points) If the L1 Data Cache is directly mapped, what is the hit ratio for the L1 Data Cache for **vecAdd1** and for **vecAdd2**?

There is not temporal data reuse in either program. Thus, all the hits are because of the size of the cache block.

For **vecAdd1** every fourth data access will be a miss. Therefore the hit ratio is 75%.

vecAdd2 access every fourth element of the array and thus, its hit ratio is 0%.

3. (5 points) What is the effect in the hit ratios if the L1 Data Cache retains the same total data storage of 16KB but is made two-way set associative?

The hit ratio does not change because there are no conflict misses in the direct mapped cache. Higher associativity has no effect on cold misses.

4. (5 points) Assume that this machine has a 32-bit address bus. If the L1 Data Cache is two-way set associative, how many bits are used for each of the following components of a data cache access?

- Offset: 4 bits
- Index: $\frac{16KB}{16 \times 2} = 512 \text{ sets} \Rightarrow 9 \text{ bits}$
- Tag: $32 - 9 - 4 = 19 \text{ bits}$