

Stutter-invariance and equivalences in logic and automata

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Bachelor thesis Mathematics

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Abstract

Schrijf een samenvatting van hoogstens een halve bladzijde waarin je kort uitlegt wat je hebt gedaan. De samenvatting schrijf je als laatste. Je mag er vanuit gaan dat je docent de lezer is.

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1. Inleiding

Vertel iets over je artikel, noem je vraagstelling. Richt je tekst op medestudenten.
(Richtlijn 2 bladzijden).

2. Proving equivalence

2.1. Monadic Second Order logic and Deterministic parity automaton

As proven earlier in this thesis there is a functional (expressive?) equivalence between Deterministic parity automaton and non-deterministic Muller automaton. For this proof I will obtain an equivalence between Monadic Second Order logic and Non-deterministic Muller Automaton.

2.1.1. Lemma 12.18 (Een muller automoot naar MSO)

Theorem 2.1. *There is a effective procedure transforming a Non-deterministic Muller automata \mathcal{A} to an equivalent Monadic Second Order logic formula $\varphi_{\mathcal{A}}$. That is for a ω -word w we have $w \models \varphi_{\mathcal{A}}$ if and only iff \mathcal{A} accepts w*

Proof. Let $\mathcal{A} = (Q, \Sigma, q_I, \Delta, \mathcal{F})$ be a non-Deterministic Muller Automata (conform ..). Now I will define a monadic second order formula to describe the workings of this automata. At first we see that the acceptance condition there *exists* a succesfull run already hints to the start of this second order formula. We can define a run as following: for every state we define a monadic predicate R_q consisting of word positions that are in state q . Now define $\bar{R} = (R_q)_{q \in Q}$. Now we will describe the working of this automaton. First define the notion of a state

$$\text{State}_q(x) := x \in R_q \wedge \bigwedge_{q' \in Q \setminus \{q\}} \neg x \in R_{q'}.$$

Now we check if the R_q indeed form a partition of Q ,

$$\text{Part} := \forall x(\text{sing}(x) \rightarrow \bigvee_{q \in Q} \text{State}_q(x)).$$

To program the initial condition we define:

$$\text{Init} := \exists x(x \in R_{q_I} \wedge \forall y(\text{sing}(y) \rightarrow x \leq y)).$$

Next we want to model the transition relation in our formula

$$\text{Trans} := \forall x \forall y \left((\text{sing}(x) \wedge \text{sing}(y) \wedge Sxy) \rightarrow \bigvee_{(q,a,q') \in \Delta} (\text{State}_q(x) \wedge x \in P_a \wedge \text{State}_{q'}(y)) \right)$$

The only thing left is to express the acceptance condition of this automata. For the non-deterministic Muller automata this constst of the fact that there exists a run ρ (which we now defined) where $\inf(\rho) \in \mathcal{F}$. So we need to define the inf set and express that this is a Muller set. So first to express the fact that a state q occurs infinitely often in the run:

$$\text{InfOcc}_q := \exists Q(Q \subseteq R_q \wedge \text{Inf}(Q)).$$

And to express the fact that this is indeed a Muller set:

$$\text{Muller} := \bigvee_{F \in \mathcal{F}} \left(\bigwedge_{q \in F} \text{InfOcc}_q \wedge \bigwedge_{q \notin F} \neg \text{InfOcc}_q \right)$$

Now we can define

$$\varphi_A := \exists \bar{R} (\text{Part} \wedge \text{Init} \wedge \text{Trans} \wedge \text{Muller})$$

This gives the translation. Now I need to prove that these are indeed equivalent. (Maar is een beetje per constructie toch???) \square

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3. Differentiëren en integreren

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4. Conclusie

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Bibliography

- [1] Erich Gradel, Wolfgang Thomas, and Thomas Wilke. *Automata, logics, and infinite games: a guide to current research*, volume 2500. Springer, 2003.

Populaire samenvatting

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A. Lineaire Algebra

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