### HW2

Q1

The following grammar is not suitable for a top-down predictive parser. Identify the problem and correct it by rewriting the grammar. Show that your new grammar satisfies the LL(1) condition.

## ORIGINAL GRAMMAR

$$egin{array}{ccccc} L & 
ightarrow & Ra & \ & | & Qba & \ R & 
ightarrow & aba & \ & | & caba & \ & | & Rbc & \ Q & 
ightarrow & bbc & \ & | & bc & \ \end{array}$$

**ISSUE** 

The rule  $R \to Rbc$  is an example of left recursion. There are two choices in the rule  $Q \to bbc|bc$  if the parser is presented with a b.

### **NEW GRAMMAR**

#### **LEFT RECURSION**

## No longer present

### FIRST TERMINAL SETS

$$\mathit{First}(L) = \{a,b,c\} \ \mathit{First}(R) = \{a,c\} \ \mathit{First}(S) = \{a,b\} \ \mathit{First}(Q) = \{b\} \ \mathit{First}(T) = \{b,c\}$$
 
$$\mathit{First}(R) \cap \mathit{First}(Q) = \emptyset$$

Q2

1	A	$\rightarrow$	Ba
2	B	$\rightarrow$	dab
3			Cb
4	C	$\rightarrow$	cB
5		1	Ac

### **LEFT RECURSION**

Yes, this grammar has left recursion. (Rules 1, 3 and 5 applied one after the other is circular. In other words  $A \Rightarrow^+ A\beta$  for  $b \in N \cup T$ .)

#### **REWRITTEN GRAMMAR**

Step 1: Eliminate C

$$egin{array}{ccccc} A & 
ightarrow & Ba \ B & 
ightarrow & dab \ & phantom{\phantom{phantom{phantom{phantom{phantom{phantom{phantom{\phantom{\phantom{\phantom{\phantom{\phantom{\phantom{\phantom{\phantom{\phantom{\phantom{\phantom{\phantom{phantom{\phant$$

# Step 2: Remove circular left recursion

$$egin{array}{cccc} A & 
ightarrow & Ba \ B & 
ightarrow & dab \ & ert & cBb \ & ert & Bacb \end{array}$$

# Step 3: Remove left recursion

$$egin{array}{ccccc} A & 
ightarrow & Ba \ B & 
ightarrow & dabC \ & phantom{\phantom{phantom{phantom{phantom{phantom{phantom{\phan$$

LL(1)

No, since the grammar has left recursion, it is not LL(1).

Q3