

# *CI-1220 Ensamblador y Microprocesadores*

## Memory Layout and Buffer Overflows

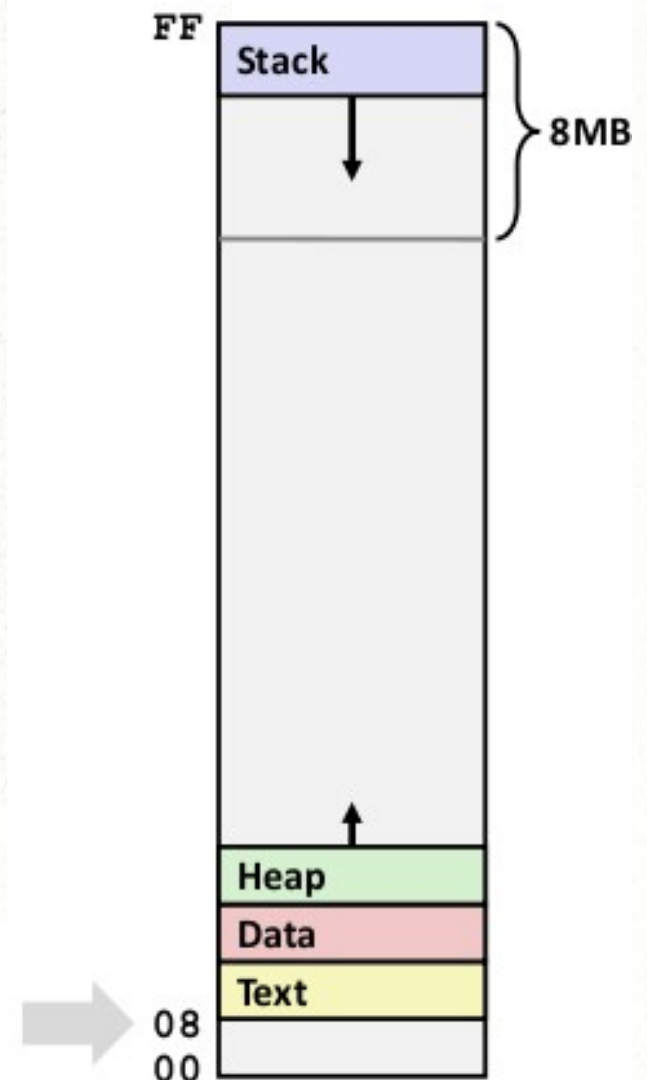
Prof. Ricardo Villalón

# IA32/Linux Memory Layout

*not drawn to scale*

- **Stack**
  - Runtime stack (8MB limit)
- **Heap**
  - Dynamically allocated storage
  - When call `malloc()`, `calloc()`, `new()`
- **Data**
  - Statically allocated data
  - E.g., arrays & strings declared in code
- **Text**
  - Executable machine instructions
  - Read-only

Upper 2 hex digits  
= 8 bits of address



# Memory Allocation Example

*not drawn to scale*

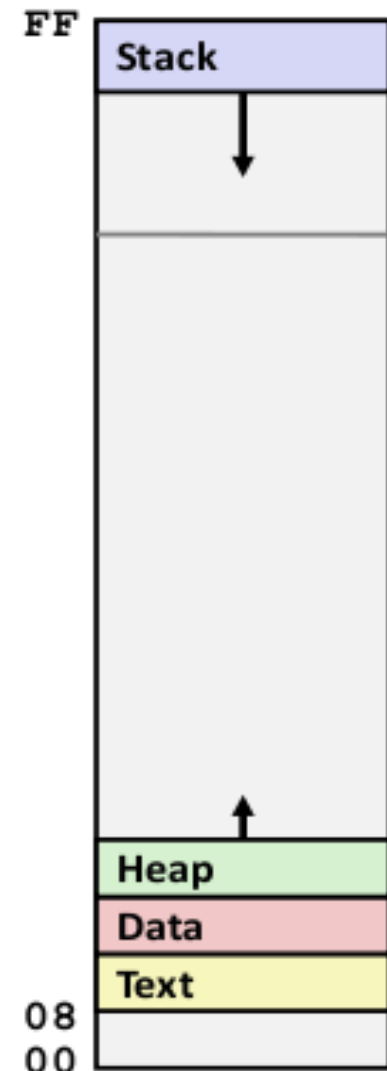
```
char big_array[1<<24]; /* 16 MB */
char huge_array[1<<28]; /* 256 MB */

int beyond;
char *p1, *p2, *p3, *p4;

int useless() { return 0; }

int main()
{
    p1 = malloc(1 << 28); /* 256 MB */
    p2 = malloc(1 << 8); /* 256 B */
    p3 = malloc(1 << 28); /* 256 MB */
    p4 = malloc(1 << 8); /* 256 B */
    /* Some print statements ... */
}
```

*Where does everything go?*



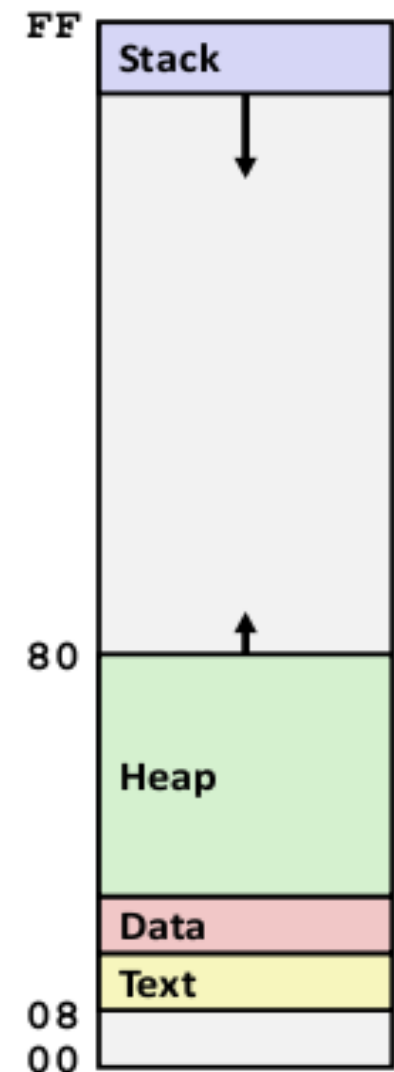
# IA32 Example Addresses

*not drawn to scale*

*address range  $\sim 2^{32}$*

<code>\$esp</code>	<code>0xffffbcd0</code>
<code>p3</code>	<code>0x65586008</code>
<code>p1</code>	<code>0x55585008</code>
<code>p4</code>	<code>0x1904a110</code>
<code>p2</code>	<code>0x1904a008</code>
<code>&amp;p2</code>	<code>0x18049760</code>
<code>beyond</code>	<code>0x08049744</code>
<code>big_array</code>	<code>0x18049780</code>
<code>huge_array</code>	<code>0x08049760</code>
<code>main()</code>	<code>0x080483c6</code>
<code>useless()</code>	<code>0x08049744</code>
<code>final malloc()</code>	<code>0x006be166</code>

`malloc()` is dynamically linked  
address determined at runtime





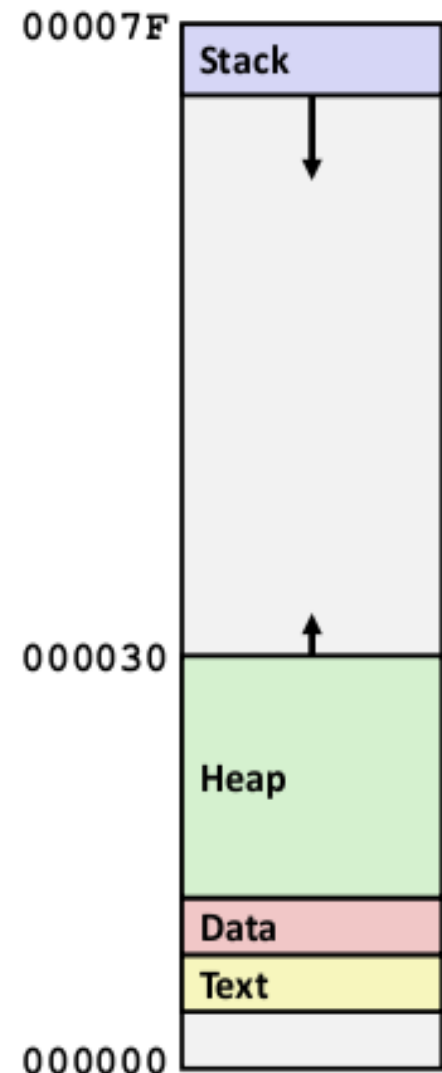
# X86-64 Example Addresses

*not drawn to scale*

*address range  $\sim 2^{47}$*

\$rsp	0x7fffffff8d1f8
p3	0x2aaabaadd010
p1	0x2aaaaadc010
p4	0x000011501120
p2	0x000011501010
&p2	0x000010500a60
beyond	0x000000500a44
big_array	0x000010500a80
huge_array	0x000000500a50
main()	0x000000400510
useless()	0x000000400500
final malloc()	0x00386ae6a170

`malloc()` is dynamically linked  
address determined at runtime



# *Security Problems with Memory Management*

# *Problems with String Library Code*

## ■ Implementation of Unix function gets ()

```
/* Get string from stdin */
char *gets(char *dest)
{
    int c = getchar();
    char *p = dest;
    while (c != EOF && c != '\n') {
        *p++ = c;
        c = getchar();
    }
    *p = '\0';
    return dest;
}
```

- No way to specify limit on number of characters to read

## ■ Similar problems with other Unix functions

- **strcpy**: Copies string of arbitrary length
- **scanf**, **fscanf**, **sscanf**, when given **%s** conversion specification

# Vulnerable Buffer Code

```
/* Echo Line */  
void echo()  
{  
    char buf[4]; /* Way too small! */  
    gets(buf);  
    puts(buf);  
}
```

```
int main()  
{  
    printf("Type a string:");  
    echo();  
    return 0;  
}
```

```
unix>./bufdemo  
Type a string:1234567  
1234567
```

```
unix>./bufdemo  
Type a string:12345678  
Segmentation Fault
```

```
unix>./bufdemo  
Type a string:123456789ABC  
Segmentation Fault
```



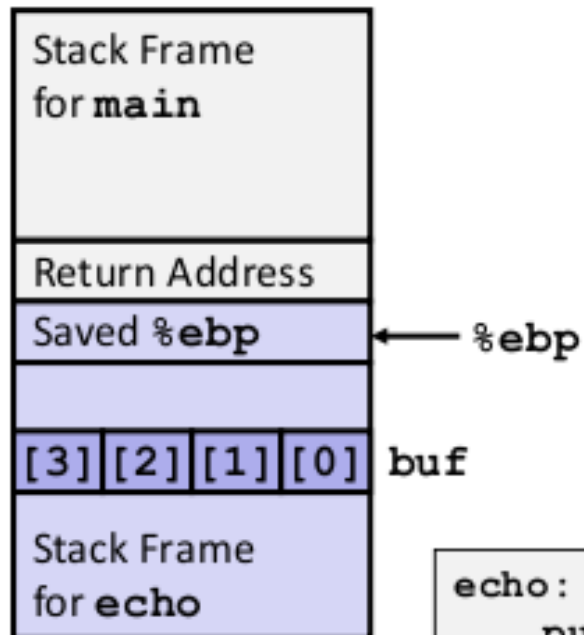
# Buffer Overflow Disassembly

```
080484f0 <echo>:
80484f0: 55                push    %ebp
80484f1: 89 e5            mov     %esp, %ebp
80484f3: 53              push    %ebx
80484f4: 8d 5d f8        lea     0xffffffff8(%ebp), %ebx
80484f7: 83 ec 14        sub     $0x14, %esp
80484fa: 89 1c 24        mov     %ebx, (%esp)
80484fd: e8 ae ff ff ff  call    80484b0 <gets>
8048502: 89 1c 24        mov     %ebx, (%esp)
8048505: e8 8a fe ff ff  call    8048394 <puts@plt>
804850a: 83 c4 14        add     $0x14, %esp
804850d: 5b              pop     %ebx
804850e: c9              leave
804850f: c3              ret

80485f2: e8 f9 fe ff ff  call    80484f0 <echo>
80485f7: 8b 5d fc        mov     0xffffffffc(%ebp), %ebx
80485fa: c9              leave
80485fb: 31 c0            xor     %eax, %eax
80485fd: c3              ret
```

# Buffer Overflow Stack

*Before call to gets*



```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

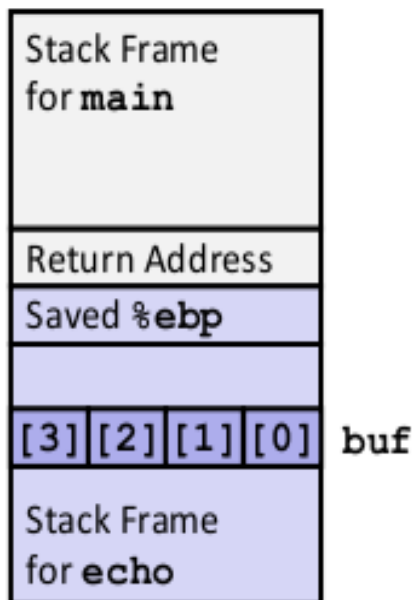
```
echo:
    pushl %ebp                # Save %ebp on stack
    movl  %esp, %ebp
    pushl %ebx                # Save %ebx
    leal  -8(%ebp), %ebx      # Compute buf as %ebp-8
    subl  $20, %esp           # Allocate stack space
    movl  %ebx, (%esp)        # Push buf on stack
    call  gets                # Call gets
    . . .
```

# Buffer Overflow Stack Example

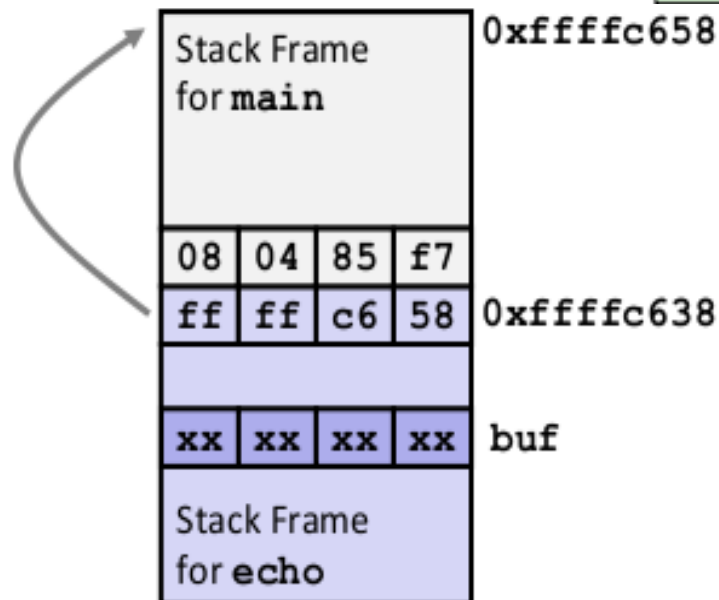
```

unix> gdb bufdemo
(gdb) break echo
Breakpoint 1 at 0x8048583
(gdb) run
Breakpoint 1, 0x8048583 in echo ()
(gdb) print /x $ebp
$1 = 0xffffc638
(gdb) print /x *(unsigned *)$ebp
$2 = 0xffffc658
(gdb) print /x *((unsigned *)$ebp + 1)
$3 = 0x80485f7
    
```

*Before call to gets*



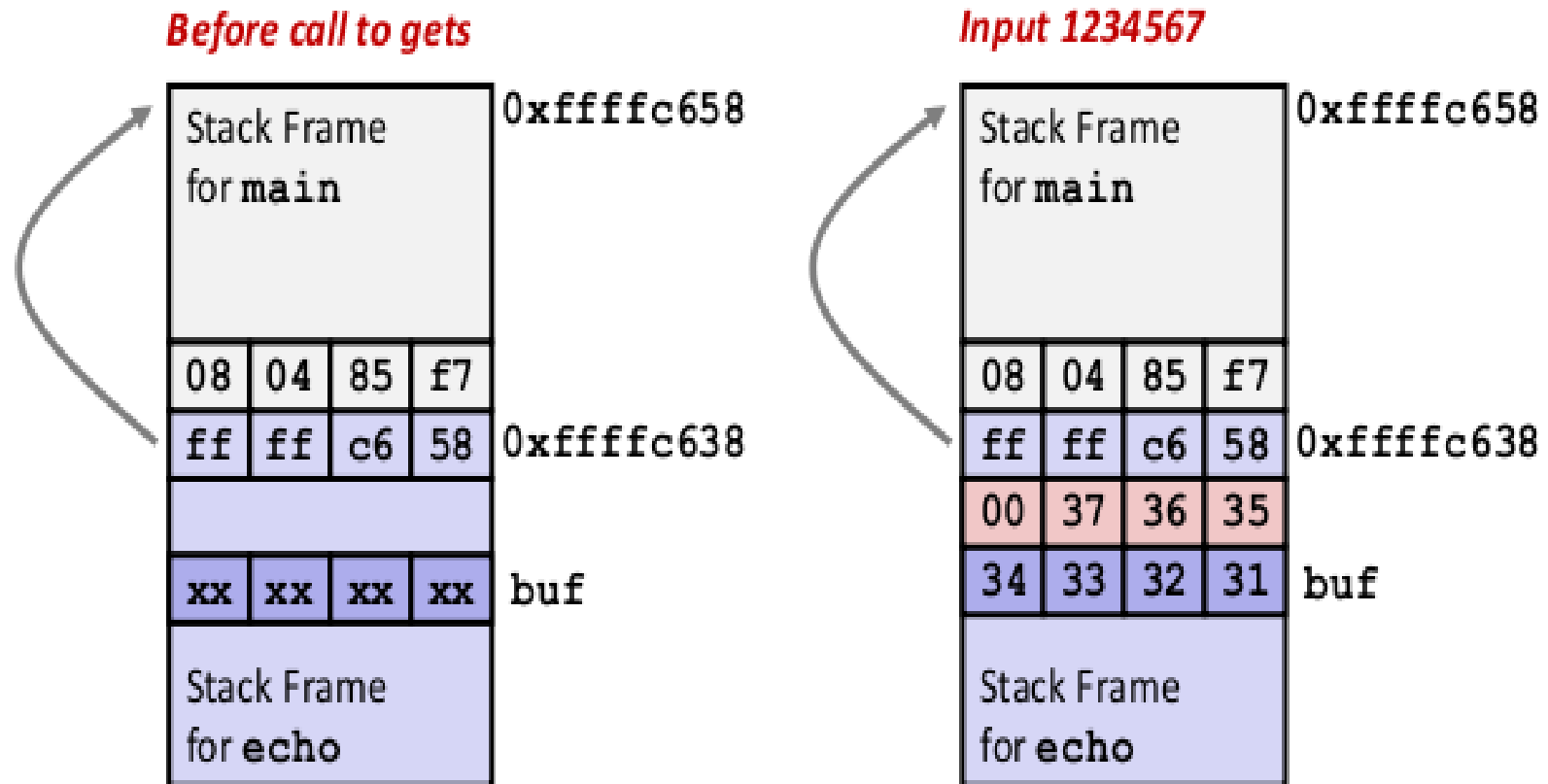
*Before call to gets*



80485f2: call 80484f0 <echo>

80485f7: mov 0xffffffffc(%ebp), %ebx # Return Point

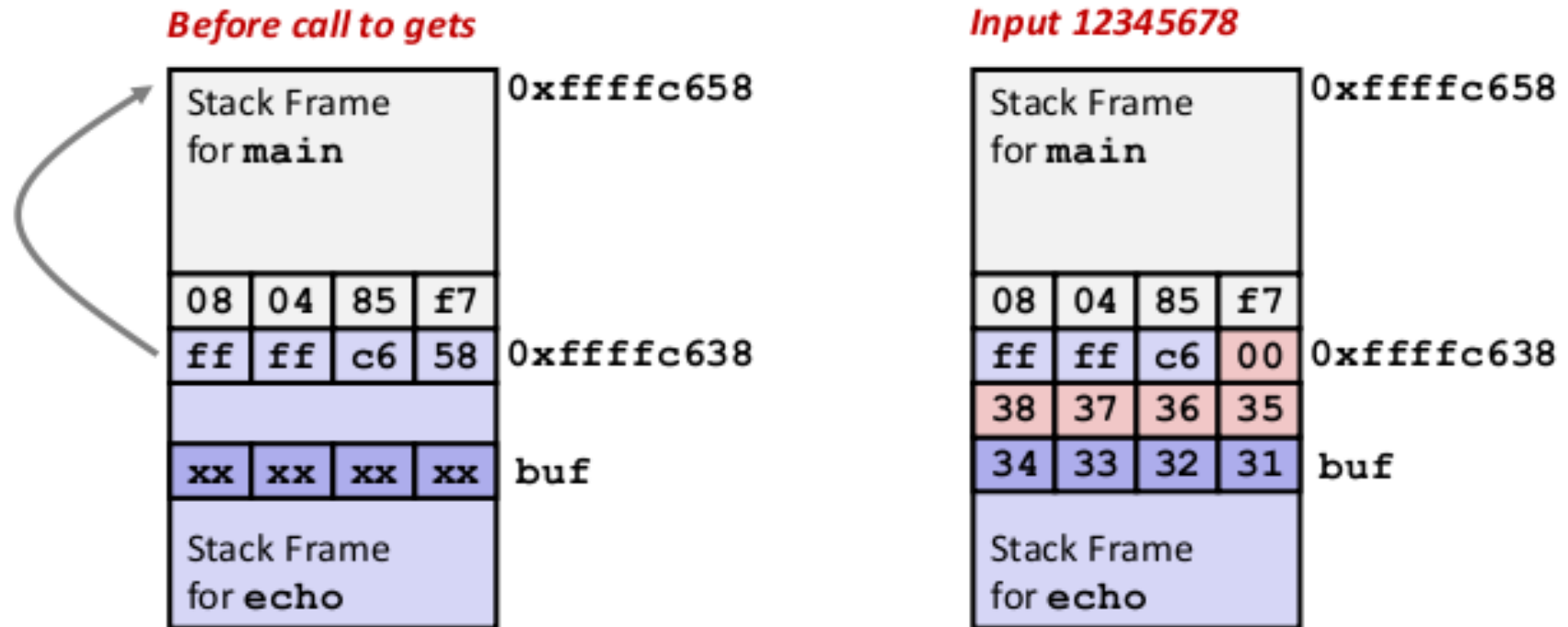
# Buffer Overflow Example #1



Overflow buf, but no problem



# Buffer Overflow Example #2

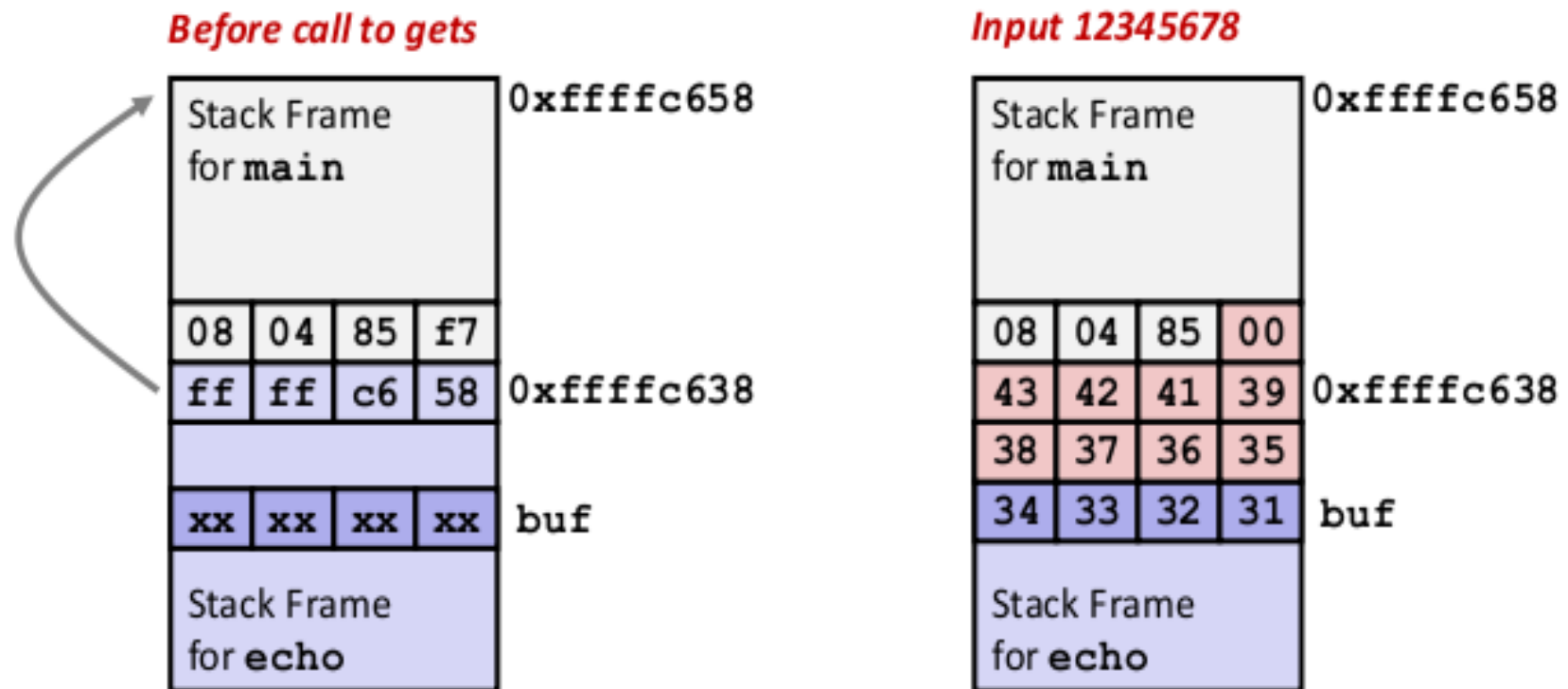


**Base pointer corrupted**

```

. . .
804850a: 83 c4 14  add    $0x14,%esp  # deallocate space
804850d: 5b        pop     %ebx    # restore %ebx
804850e: c9        leave   # movl %ebp, %esp; popl %ebp
804850f: c3        ret     # Return
    
```

# Buffer Overflow Example #3



**Return address corrupted**

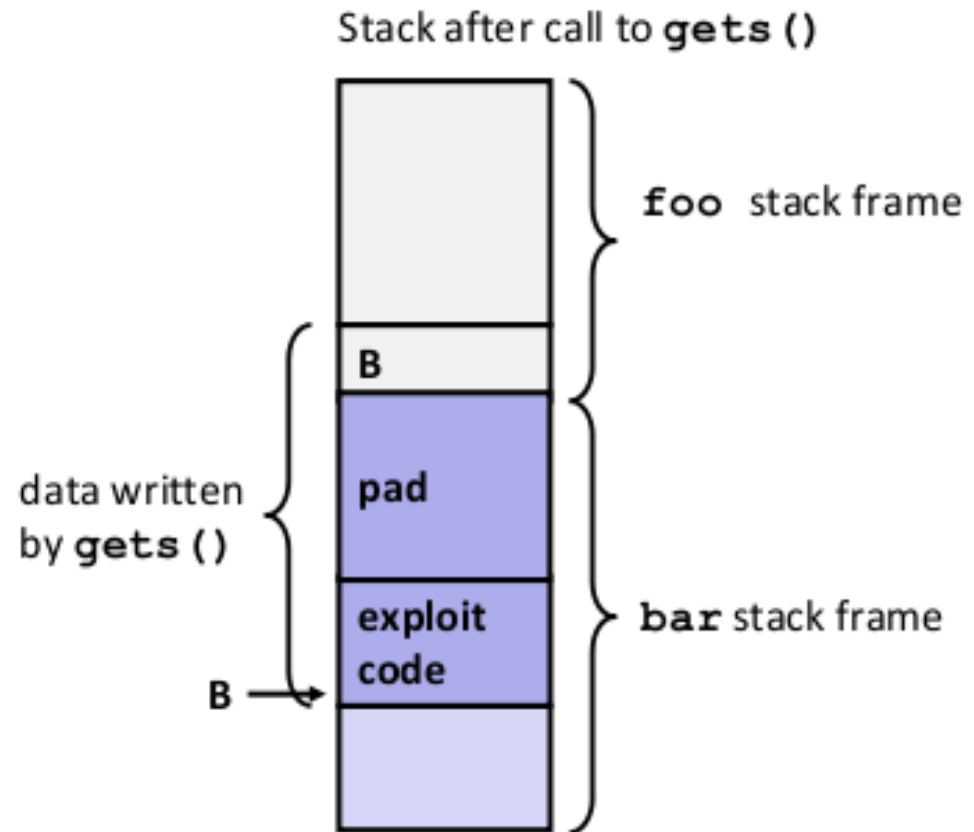
```
80485f2: call 80484f0 <echo>
80485f7: mov 0xffffffffc(%ebp),%ebx # Return Point
```

# Malicious Use of Buffer Overflow

```
void foo() {  
    bar();  
    ...  
}
```

return  
address  
A

```
int bar() {  
    char buf[64];  
    gets(buf);  
    ...  
    return ...;  
}
```



- Input string contains byte representation of executable code
- Overwrite return address with address of buffer
- When `bar()` executes `ret`, will jump to exploit code



# *Exploits Based on Buffer Overflow*

- *Buffer overflow bugs allow remote machines to execute arbitrary code on victim machines*
- **Internet worm**
  - Early versions of the finger server (fingerd) used **gets ()** to read the argument sent by the client:
    - **finger droh@cs.cmu.edu**
  - Worm attacked fingerd server by sending phony argument:
    - **finger "exploit-code padding new-return-address"**
    - exploit code: executed a root shell on the victim machine with a direct TCP connection to the attacker.



# *Avoiding Overflow Vulnerability*

```
/* Echo Line */  
void echo()  
{  
    char buf[4]; /* Way too small!  
*/  
    fgets(buf, 4, stdin);  
    puts(buf);  
}
```

- **Use library routines that limit string lengths**
  - **fgets** instead of **gets**
  - **strncpy** instead of **strcpy**
  - Don't use **scanf** with **%s** conversion specification
    - Use **fgets** to read the string
    - Or use **%ns** where **n** is a suitable integer

# *System Level Protections*

## ■ Randomized stack offsets

- At start of program, allocate random amount of space on stack
- Makes it difficult for hacker to predict beginning of inserted code

## ■ Nonexecutable code segments

- In traditional x86, can mark region of memory as either “read-only” or “writeable”
  - Can execute anything readable
- Add explicit “execute” permission

```
unix> gdb bufdemo
(gdb) break echo

(gdb) run
(gdb) print /x $ebp
$1 = 0xffffc638

(gdb) run
(gdb) print /x $ebp
$2 = 0xffffbb08

(gdb) run
(gdb) print /x $ebp
$3 = 0xffffc6a8
```

# *Worms and Viruses*

- **Worm: A program that**
  - Can run by itself
  - Can propagate a fully working version of itself to other computers
  
- **Virus: Code that**
  - Add itself to other programs
  - Cannot run independently
  
- **Both are (usually) designed to spread among computers and to wreak havoc**